

# Spectral Interpretations of the Riemann Zeros

Quantum Chaos and the Computational Hardness of the Inverse  
Spectral Problem

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## Abstract

This paper explores the theoretical framework linking the Riemann Hypothesis (RH) to the spectral theory of unitary operators, following the Hilbert-Pólya conjecture. We discuss the geometric realization of the Riemann zeros as eigenvalues of a Frobenius operator on a stable arithmetic bundle. Furthermore, we address the implications for cryptography. We argue that even if a spectral realization exists, the **Fourier Uncertainty Principle** governing the duality between the prime spectrum and the zero spectrum imposes an exponential lower bound on the resolution time required to recover prime factors. This suggests that the validity of RH is consistent with the computational hardness of integer factorization.

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## 1 Introduction: The Hilbert-Pólya Operator

The Hilbert-Pólya conjecture posits that the non-trivial zeros of the Riemann zeta function  $\zeta(s)$  correspond to the eigenvalues of a self-adjoint operator acting on a suitable Hilbert space.

**Conjecture 1.1** (Spectral Existence). *There exists a Hamiltonian operator  $\hat{H}$  acting on a Hilbert space  $\mathcal{H}$  such that the eigenvalues  $E_n$  satisfy:*

$$\frac{1}{2} + iE_n = \rho_n$$

where  $\rho_n$  are the non-trivial zeros of  $\zeta(s)$ .

If such an operator exists and is self-adjoint (or the associated flow is unitary), the Riemann Hypothesis follows automatically, as the eigenvalues of self-adjoint operators are real.

## 2 Geometric Framework

### 2.1 Arakelov Bundles and Stability

In the context of Arakelov geometry, the search for such an operator leads to the study of vector bundles over  $\overline{\text{Spec}(\mathbb{Z})}$ .

**Definition 2.1** (Slope Stability). An Arakelov vector bundle  $\overline{E}$  is said to be semistable if for every subbundle  $\overline{F} \subset \overline{E}$ , the slope satisfies  $\mu(\overline{F}) \leq \mu(\overline{E})$ .

Using the analogy with the Narasimhan-Seshadri theorem (and the Hitchin-Kobayashi correspondence), a stable bundle corresponds to a unitary connection. The Frobenius flow on the moduli space of such bundles is a natural candidate for the operator whose spectrum matches the Riemann zeros.

## 3 Quantum Chaos and the Berry-Keating Model

### 3.1 The Hamiltonian

Berry and Keating proposed a semiclassical Hamiltonian responsible for the Riemann dynamics:

$$\hat{H}_{BK} = \frac{1}{2}(\hat{x}\hat{p} + \hat{p}\hat{x})$$

where  $\hat{x}$  and  $\hat{p}$  are the position and momentum operators.

### 3.2 Semiclassical Density of States

The density of states for this system approximates the asymptotic distribution of the Riemann zeros:

$$N(E) \sim \frac{E}{2\pi} \ln \left( \frac{E}{2\pi e} \right)$$

This provides strong heuristic evidence that the Riemann zeros are spectral in nature, related to a chaotic quantum system with no time-reversal symmetry.

## 4 Computational Implications

A common concern is whether a spectral proof of RH would render RSA cryptography insecure. We argue that physical and analytical constraints preserve the hardness of factorization.

## 4.1 The Inverse Spectral Problem

The problem of factoring an integer  $N$  given the “Riemann spectrum” is an Inverse Spectral Problem. The Explicit Formula of Riemann-Weil establishes a Fourier duality between the zeros  $\rho$  and the prime powers  $\ln p$ :

$$\sum_{\rho} \phi(\gamma) \sim \sum_{p,k} \frac{\ln p}{p^{k/2}} \hat{\phi}(\ln p^k)$$

## 4.2 Fourier Uncertainty as a Computational Bound

To distinguish a specific prime factor  $p$  from nearby integers, one must resolve the “frequency”  $\omega_p = \ln p$  in the spectral density.

- **Resolution Limit:** To resolve a frequency scale of  $\Delta\omega \approx 1/p$  (the spacing between  $\ln p$  and  $\ln(p+1)$ ), the Fourier Uncertainty Principle requires an energy (eigenvalue) range  $E_{max}$  proportional to the inverse of the spacing.

- **Scaling:**

$$E_{max} \sim \frac{1}{\Delta\omega} \sim p$$

- **Complexity:** For an RSA modulus  $N$ , the factors are  $p \approx \sqrt{N}$ . Thus, recovering  $p$  from the spectral data requires summing over zeros up to height  $T \sim \sqrt{N} = e^{\frac{1}{2}\ln N}$ .

**Theorem 4.1** (Conditional Hardness). *Assuming the spectral interpretation holds, the recovery of prime factors from the operator’s spectrum requires evaluating the trace formula up to an energy level that scales exponentially with the bit-length of the key ( $n = \log_2 N$ ).*

$$T_{compute} = O(e^{n/2})$$

## 5 Conclusion

The geometrization of arithmetic suggests that the Riemann Hypothesis is a consequence of the unitary evolution of the underlying arithmetic cohomology. However, this structural order does not imply computational triviality. The **Fourier Uncertainty Principle**—the mathematical dual of the Heisenberg Uncertainty Principle—ensures that extracting specific arithmetic data (primes) from the spectral data remains an exponentially hard problem. Thus, the integrity of RSA cryptography is compatible with the truth of the Riemann Hypothesis.