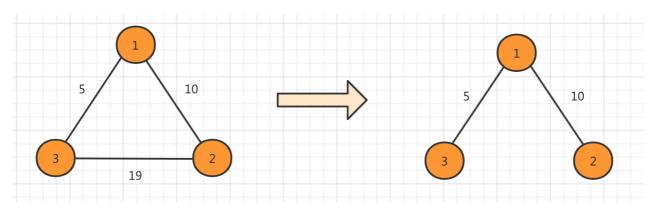
详解: 最小生成树算法

最小生成树(Minimum Spanning Tree, MST)是在一个给定的**无向图**G(V, E)中求一棵树,使得这棵树有图G的所有顶点,且所有边都来自图G中的边,并且满足整棵树的边权之和最小.

最下生成树满足如下性质:

- 最小生成树是树,因此其边数等于顶点数减1,且树内一定没有环;
- 对给定的图, 其最小生成树不唯一, 但边权之和是唯一的;
- 最小生成树是在无向图上生成的,因此其根结点可以是这个图上的任意一个点。题目一般会指出那个点作为生成树的根结点

示例



1 输入
2 3 3
3 1 2 10
4 1 3 5
5 3 2 19
6 输出
7 15

Prim算法

跟Dijkstra算法类似,将顶点分为两个集合,一个是已经在生成树中的顶点集合S,另一个是还未访问的点。

用数组d[]表示各个顶点到集合S的最短距离(只有这里d的含义与Dijkstra算法中d的含义不一样)

```
Prim(G, d[]) {
1
2
       初始化;
3
       for (循环n次) {
 4
          u = \phi d[u]最小的还未被访问的顶点的标号;
          记u已被访问
5
          for (从u出发能到达的所有顶点v){
 6
              if(v + k) 表象 以u为中介点使得v与集合S的最短距离d[v]更优){
8
                 将G[u][v]赋值给v与集合S的最短距离d[v]
9
              }
10
          }
11
12
   }
```

基于邻接矩阵实现的Prim算法

Java

```
import java.util.Arrays;
 2
    import java.util.Scanner;
 3
 4
    public class Main{
5
 6
        private static final int INF = Integer.MAX VALUE;
 7
        public static void main(String[] args) throws InterruptedException {
 8
9
            Scanner input = new Scanner(System.in);
10
            int n = input.nextInt();
            int m = input.nextInt();
11
            int[][] graph = new int[n + 1][n + 1];
12
            for (int i = 0; i < n + 1; i++) Arrays.fill(graph[i], INF);
13
14
            int u, v, w;
15
            for (int i = 0; i < m; i++) {
16
                u = input.nextInt();
17
                v = input.nextInt();
18
                w = input.nextInt();
19
                graph[u][v] = graph[v][u] = w;
2.0
            int start = 1; //指定一个结点作为生成树的根结点
21
            int res = prim(graph, n, 1);
2.2
            System.out.println(res);
23
        }
24
25
        public static int prim(int[][] graph, int n, int start) {
2.6
27
            int[] d = new int[n + 1];
            boolean[] visit = new boolean[n + 1];
28
29
            Arrays.fill(d, INF);
30
            d[start] = 0;
```

```
31
            int ans = 0;
32
            for (int i = 1; i <= n; i++) {
33
                 int u = -1, MIN = INF;
                 for (int j = 1; j \le n; j++) {
34
35
                     if (!visit[j] && d[j] < MIN) {</pre>
36
                         u = j;
37
                         MIN = d[j];
38
                     }
39
                 }
                 if (u == -1) return -1;
40
41
                 visit[u] = true;
                 ans += d[u]; // 将与集合S距离最小的边加入最小生成树
42
43
                 for (int v = 1; v \le n; v ++) {
                     if (!visit[v] && graph[u][v] < INF && graph[u][v] < d[v])
44
45
                         d[v] = graph[u][v];
46
                     }
47
                 }
48
49
            return ans;
50
        }
51
    }
```

C++

```
#include <iostream>
 2
    #include <vector>
 3
    using namespace std;
 4
 5
    const int N = 100;
    const int INF = INT MAX;
 6
 7
 8
    int prim(vector<vector<int>>& graph, int n, int start) {
9
        vector<int> d(n + 1, INF);
        vector<bool> visit(n + 1, false);
10
        d[start] = 0;
11
12
        int ans = 0;
        for (int i = 1; i <= n; i++) {
13
            int u = -1, MIN = INF;
14
15
             for (int j = 1; j \le n; j++) {
16
                 if (!visit[j] && d[j] < MIN) {</pre>
17
                     u = j;
18
                     MIN = d[j];
19
                 }
20
21
             if (u == -1) return -1;
22
            visit[u] = true;
23
             ans += d[u];
             for (int v = 1; v \le n; v ++) {
24
```

```
25
                 if (!visit[v] && graph[u][v] < INF && graph[u][v] < d[v]) {
26
                     d[v] = graph[u][v];
27
                 }
28
             }
29
30
        return ans;
31
32
33
    int main() {
        int n, m;
34
35
        cin >> n >> m;
        vector<vector<int>>> graph(n + 1, vector<int>(n + 1, INF));
36
37
        int u, v, w;
        for (int i = 0; i < m; i++) {
38
39
            cin >> u >> v >> w;
40
             graph[u][v] = graph[v][u] = w;
41
        int start = 1;
42
43
        int res = prim(graph, n, start);
        cout << res << endl;</pre>
45
        return 0;
46
    }
```

基于邻接表实现的Prim算法

Java

```
import java.util.ArrayList;
    import java.util.Arrays;
 3
    import java.util.List;
 4
    import java.util.Scanner;
 6
    public class Main{
7
8
        private static final int INF = Integer.MAX_VALUE;
 9
10
        public static void main(String[] args) throws InterruptedException {
            Scanner input = new Scanner(System.in);
11
            int n = input.nextInt();
12
13
            int m = input.nextInt();
            List<List<Node>> graph = new ArrayList<>(n + 1);
            for (int i = 0; i < n + 1; i++) graph.add(new ArrayList<>());
15
16
            int u, v, w;
            for (int i = 0; i < m; i++) {
17
                u = input.nextInt();
18
                v = input.nextInt();
19
20
                w = input.nextInt();
```

```
21
                 graph.get(u).add(new Node(v, w));
22
                 graph.get(v).add(new Node(u, w));
23
             }
24
             int start = 1;
25
             int res = prim(graph, n, start);
             System.out.println(res);
26
27
        }
28
29
        public static int prim(List<List<Node>>> graph, int n, int start) {
             int[] d = new int[n + 1];
30
31
             boolean[] visit = new boolean[n + 1];
             Arrays.fill(d, INF);
32
33
             d[start] = 0;
             int ans = 0;
34
35
             for (int i = 1; i \le n; i++) {
                 int u = -1, MIN = INF;
36
37
                 for (int j = 1; j \le n; j++) {
38
                     if (!visit[j] && d[j] < MIN) {</pre>
39
                         u = j;
40
                         MIN = d[j];
41
                     }
42
                 }
43
                 if (u == -1) return -1;
44
                 visit[u] = true;
                 ans += d[u];
45
                 for (int k = 0; k < graph.get(u).size(); k++) {
46
                     int v = graph.get(u).get(k).v; //从u能到达的顶点v
47
                     if (!visit[v] && graph.get(u).get(k).dis < d[v]) {</pre>
48
                         d[v] = graph.get(u).get(k).dis;
49
50
                     }
51
                 }
52
             }
5.3
             return ans;
54
        }
55
56
        static class Node{
             int v, dis;
57
             public Node(int v, int dis) {
58
59
                 this.v = v;
                 this.dis = dis;
60
61
62
        }
63
    }
```

C++

```
#include <iostream>
#include <vector>
using namespace std;
```

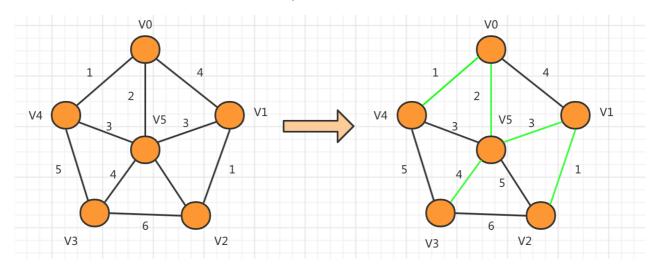
```
5
    const int N = 100;
    const int INF = INT MAX;
 7
8
    struct Node{
9
        int v, dis; //v为边的目标顶点, dis为边权
        Node(int v, int dis): v( v), dis( dis) {}
10
11
    };
12
13
    int prim(vector<vector<Node*>>& graph, int n, int start) {
14
        vector<int> d(n + 1, INF);
        vector<bool> visit(n + 1, false);
15
        d[start] = 0;
16
        int ans = 0;
17
18
        for (int i = 1; i \le n; i++) {
            int u = -1, MIN = INF;
19
            for (int j = 1; j \le n; j++) {
20
21
                if (!visit[j] && d[j] < MIN) {
22
                    u = j;
23
                    MIN = d[j];
2.4
25
            }
            if (u == -1) return -1;
2.6
            visit[u] = true; // 将u放入集合S中
27
            ans += d[u];
28
            // 遍历从u能到达的顶点,更新它们到达集合S的最短距离
29
30
            for (int k = 0; k < graph[u].size(); k++) {
                int v = graph[u][k]->v; //从u能到达的顶点
31
32
                if (!visit[v] && graph[u][k]->dis < d[v]) {</pre>
33
                    d[v] = graph[u][k] -> dis;
34
                }
35
            }
36
        }
37
        return ans;
38
    }
39
40
    int main() {
41
        int n, m;
42
        cin >> n >> m;
43
        vector<vector<Node*>> graph(n + 1);
        int u, v, w;
44
45
        for (int i = 0; i < m; i++) {
            cin >> u >> v >> w;
46
            graph[u].push back(new Node(v, w));
47
            graph[v].push_back(new Node(u, w));
48
49
50
        int start = 1;
51
        int res = prim(graph, n, start);
52
        cout << res << endl;</pre>
```

```
53 return 0;
54 }
```

Kruskal算法

采用边贪心的策略,思想很简单:

- 1. 对所有边权按从小到大进行排序;
- 2. 按边权从小到大测试所有边,如果当前测试边所连接的两个顶点不在同一个连通块中,则把这条边加入当前最小生成树中;否则,将边舍弃;
- 3. 执行步骤2,知道最小生成树中的边数等于总顶点数减1或是测试完所有边时结束。而当结束时如果最小生成树中的边数小于总顶点树减1,说明该图不连通.



```
int kruskal() {
2
      令最小生成树的边权之和为 ans, 最小生成树的当前边数为Num_Edge;
      将所有边按边权从小到大排序;
3
      for (从小到大枚举所有边) {
4
         if (当前测试边的两个端点在不同的连通块中) {
5
6
             将测试边加入最小生成树中;
             ans += 测试边的边权;
7
             最小生成树的当前边数Num Edge加1;
8
9
             当边数等于顶点数减1时结束循环;
10
         }
11
12
      return ans;
13
   }
```

伪代码中有两个细节私护不太直观, 即:

- 1. 如何判断测试边的两个端点在不同的连通块中;
- 2. 如何将测试边加入最小生成树中。

把每个连通块当作一个集合,那么就可以把问题转换为判断两个端点是否在同一个集合中,这正好可以 使用**并查集**。

<u>并查集</u>可以通过查询两个结点所在集合的根结点判断是否来自同一集合;

将测试边加入连通块可以通过将两个端点所在集合合并,也正好利用了并查集的合并特性;

Java

```
import java.util.Arrays;
 2
    import java.util.Comparator;
 3
    import java.util.Scanner;
 4
 5
    public class Main{
        public static int[] father;
 6
 7
        public static void main(String[] args) {
            Scanner input = new Scanner(System.in);
8
9
            int n = input.nextInt();
10
            int m = input.nextInt();
            Edge[] edges = new Edge[m];
11
12
            father = new int[n + 1];
13
            for (int i = 0; i < n; i++) father[i] = i;</pre>
14
            int u, v, w;
            for (int i = 0; i < m; i++) {
15
16
                 u = input.nextInt();
17
                 v = input.nextInt();
18
                 w = input.nextInt();
19
                 edges[i] = new Edge(u, v, w);
20
            }
2.1
            int res = kruskal(edges, n, m);
            System.out.println(res);
2.2
23
        }
24
25
        public static int kruskal(Edge[] edges, int n, int m) {
26
            Arrays.sort(edges, new Comparator<Edge>() {
                 @Override
27
2.8
                 public int compare(Edge o1, Edge o2) {
29
                     if (o1.w > o2.w) return 1;
                     else if (o1.w == o2.w) return 0;
30
31
                     else return -1;
32
                 }
33
            });
            int ans = 0;
34
35
            int numEdges = 0;
36
            for (int i = 0; i < m; i++) {
37
                 int faU = findFather(edges[i].u);
38
                 int faV = findFather(edges[i].v);
39
                 if (faU != faV) {
40
                     ans += edges[i].w;
                     father[faU] = faV;
41
```

```
42
                    numEdges++;
43
                    if (numEdges == n - 1) break;
               }
44
            }
45
46
            if (numEdges != n - 1) return -1;
47
            else return ans;
48
        }
49
50
        public static int findFather(int a) {
51
            int x = a;
            while (x != father[x]) {
52
               x = father[x];
53
54
55
            // 压缩路径
            while (a != father[a]) {
56
57
               int z = a;
58
               a = father[a];
59
                father[z] = x;
60
            }
            return x;
62
        }
63
64
        static class Edge{
            int u, v, w;
65
66
            public Edge(int u, int v, int w) {
67
                this.u = u;
                this.v = v;
68
69
                this.w = w;
70
           }
71
        }
   }
72
```

C++

```
#include <iostream>
    #include <vector>
 3
   using namespace std;
 4
    const int MAXV = 100;
 6
    const int MAXE = 100;
 7
 8
    struct edge{
       int u, v; //边的两个断点
9
       int cost; //边权
10
    }E[MAXE]; //最多有MAXE条边
11
12
13
   bool cmp(edge a, edge b) {
```

```
14
    return a.cost < b.cost;
15
    }
16
    // 并查集部分
17
18
    int father[MAXV]; //并查集数组
    int findFather(int x) { //并查集查询函数
19
2.0
       int a = x;
21
       while (x != father[x]) {
22
           x = father[x];
23
       // 路径压缩
24
        while (a != father[a]) {
25
           int z = a;
26
27
           a = father[a];
28
           father[z] = x;
29
        }
30
       return x;
31
    }
32
    // 返回最小生成树的边权之和,参数n为顶点个数,m为图的边数
33
34
    int kruskal(int n, int m) {
35
       int ans = 0, Num Edge = 0;
36
        for (int i = 0; i < n; i++) {
            father[i] = i;
37
38
        }
39
        sort(E, E + m, cmp);
        for (int i = 0; i < m; i++) { //枚举所有边
40
           int faU = findFather(E[i].u);
41
           int faV = findFather(E[i].v);
42
            if (faU != faV) { // 不在同一个连通块,将边加入最小生成树
43
               father[faU] = faV;
44
45
               ans += E[i].cost;
               Num Edge++;
46
               if (Num Edge == n - 1) break;
47
48
            }
49
        }
        if (Num_Edge != n - 1) return -1; //无法连通返回-1
50
51
        else return ans;
52
    }
53
54
    int main() {
55
       int n, m;
        cin >> n >> m;
56
        int u, v, w;
57
       for (int i = 0; i < m; i++) {
58
59
           cin >> u >> v >> w;
60
           E[i].u = u;
61
           E[i].v = v;
62
           E[i].cost = w;
```