
Algorithm 1 Finding maximal disks following a parallel approach.

Input: Set of points T , maximum distance ϵ and minimum size μ

Output: Set of maximal disks M

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find the set of pairs of points  $P$  in  $T$  which are  $\epsilon$  distance each other
 $C \leftarrow \emptyset$ 
for each  $p_i$  in  $P$  do
    compute disks  $c_i^1$  and  $c_i^2$  of  $p_i$  using  $\epsilon$ 
    add  $c_i^1$  and  $c_i^2$  to  $C$ 
end for
 $D \leftarrow \emptyset$ 
for each  $c_i$  in  $C$  do
    find the set of points  $\rho_i$  which lie  $\epsilon$  distance around  $c_i$ 
    if  $|\rho_i| \geq \mu$  then
        compute centroid  $\varsigma_i$  of the MBR of  $\rho_i$ 
        set  $d_i.center$  as  $\varsigma_i$ 
        set  $d_i.points$  as  $\rho_i$ 
        if  $d_i$  not in  $D$  then
            add  $d_i$  to  $D$  {Prunning duplicate candidates...}
        end if
    end if
end for
build an R-Tree  $disksRT$  using centers in  $D$ 
 $E \leftarrow \emptyset$ 
for each MBR in  $disksRT$  do
    expand MBR to create an expanded MBR  $\varepsilon_i$  using a buffer of  $\epsilon$  distance
    add  $\varepsilon_i$  to  $E$ 
end for
for each  $d_i$  in  $D$  do
    for each  $\varepsilon_j$  in  $E$  do
        if  $d_i.center \cap \varepsilon_j$  then
            add  $d_i$  to  $\varepsilon_j$ 
        end if
    end for
end for
 $M \leftarrow \emptyset$ 
for each  $\varepsilon_i$  in  $E$  do
     $\chi \leftarrow \emptyset$ 
    for each  $d_i$  in  $\varepsilon_i$  do
        add  $d_i.points$  to  $\chi$ 
    end for
    find the set of maximal patterns  $F$  in  $\chi$ 
    for each  $f_i$  in  $F$  do
        compute centroid  $\varsigma_i$  of the items in  $f_i$ 
        if  $\varsigma_i$  is not in the expansion area of  $\varepsilon_i$  then
            set  $m_i.center$  as  $\varsigma_i$ 
            set  $m_i.points$  as  $f_i$ 
            add  $m_i$  to  $M$ 
        end if
    end for
end for
return  $M$ 
```
