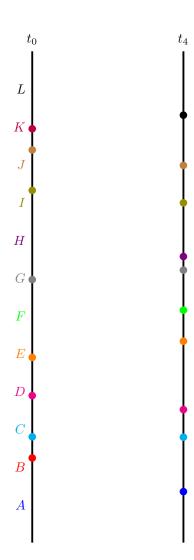
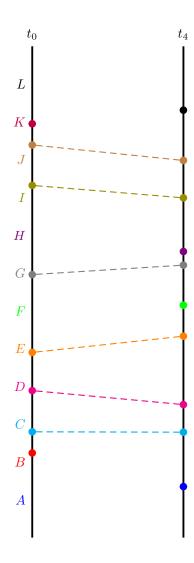
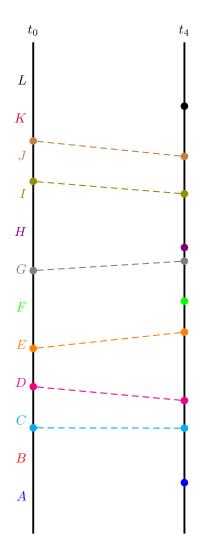
Maximal disks t_0 and t_4

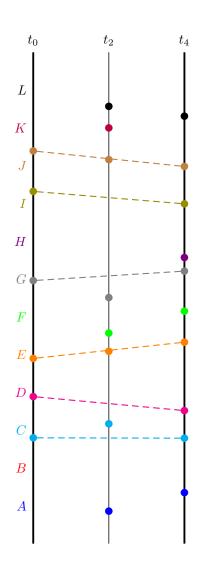


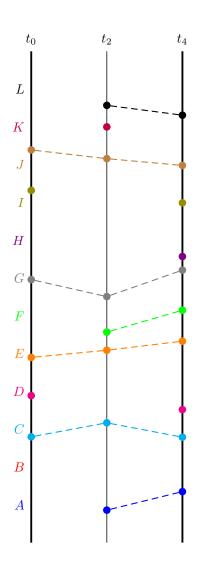
Join t_0 and t_4



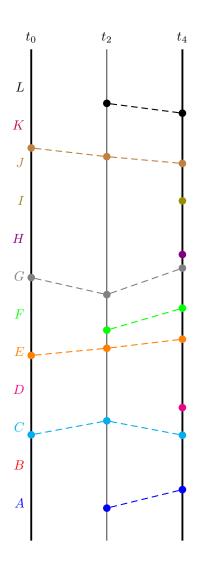
Filter t_0 and t_4

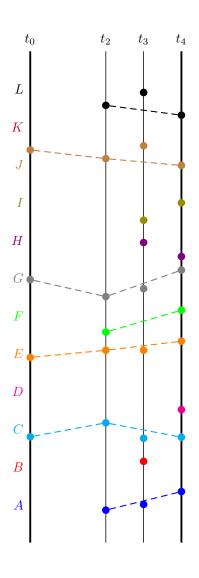


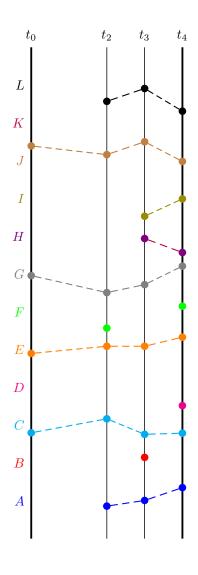




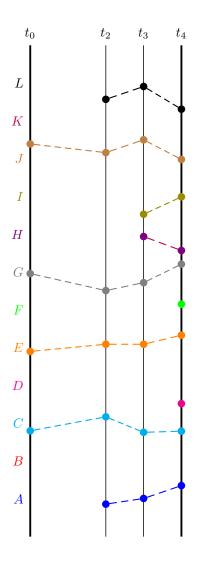
Filter t_2

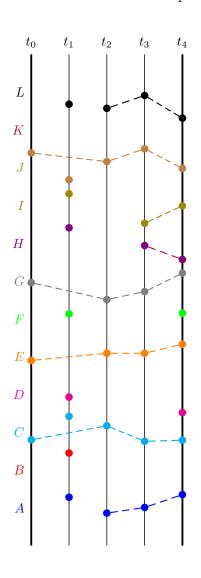


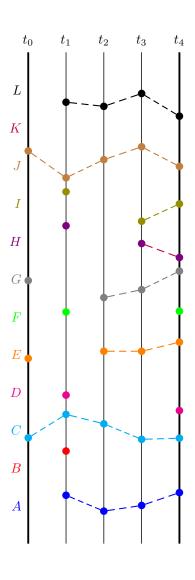




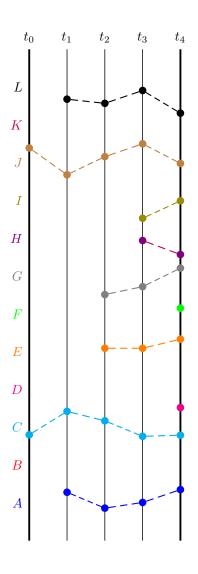
Filter t_3



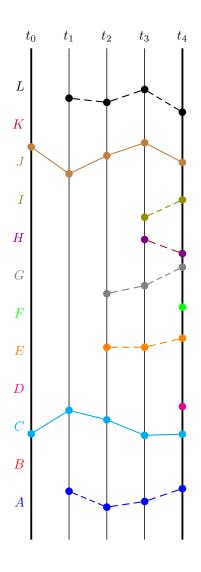




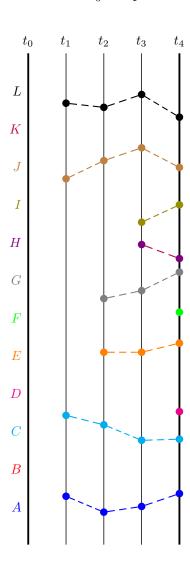
Filter t_1

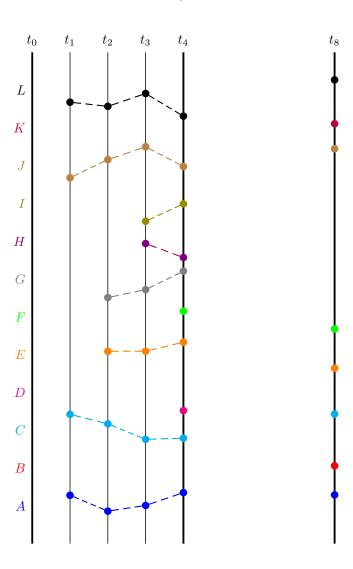


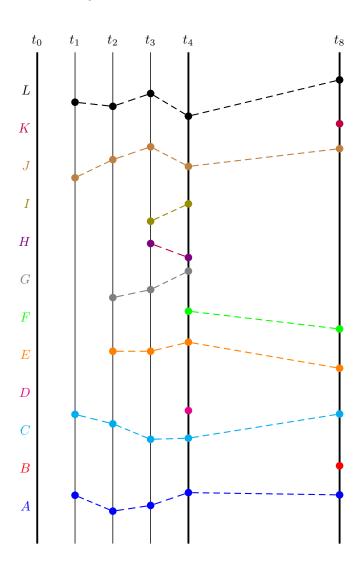
Flocks t_0 - t_4



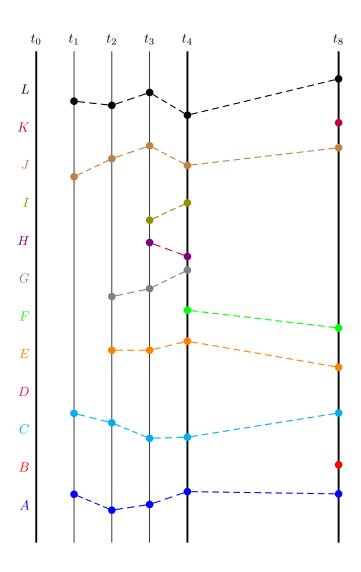
Prune t_0 - t_4

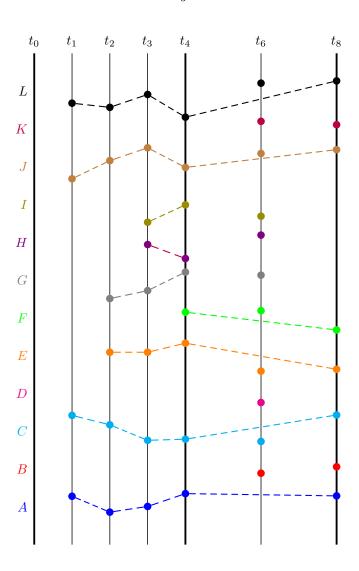


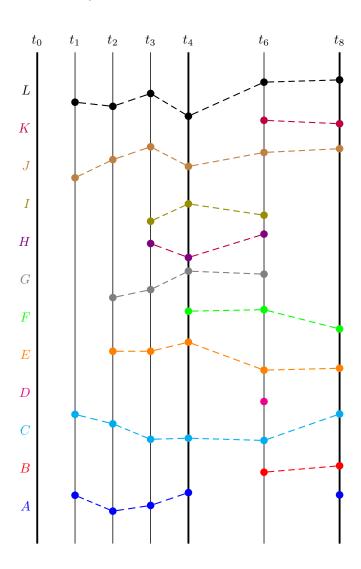




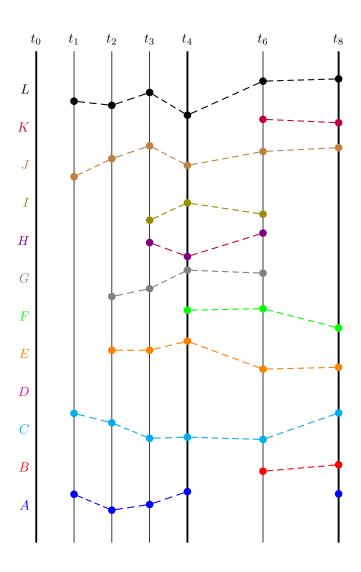
Filter t_8

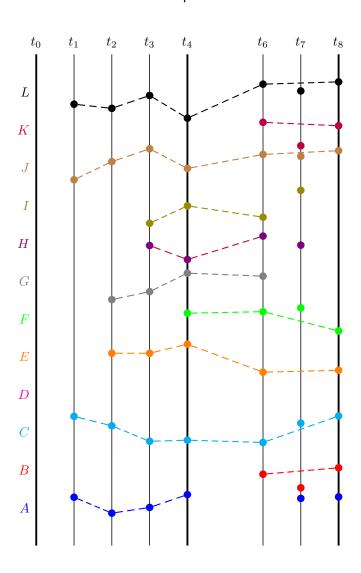


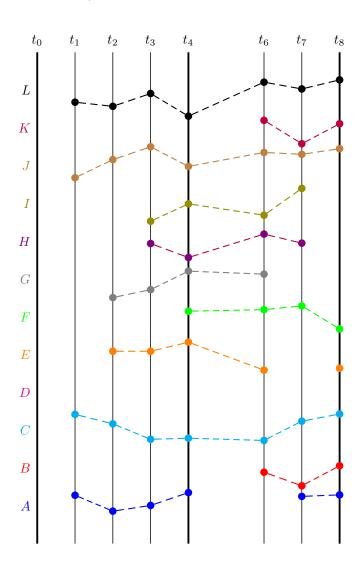




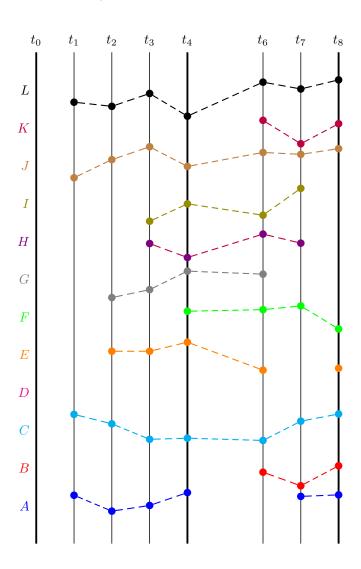
Filter t_6

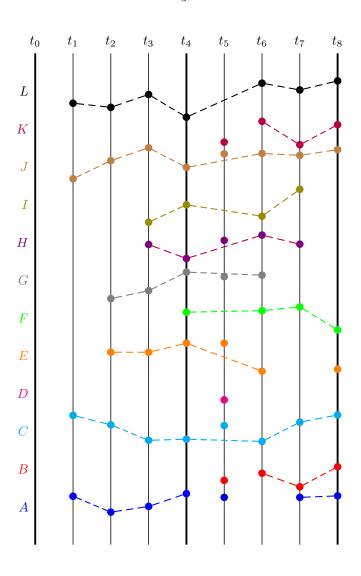


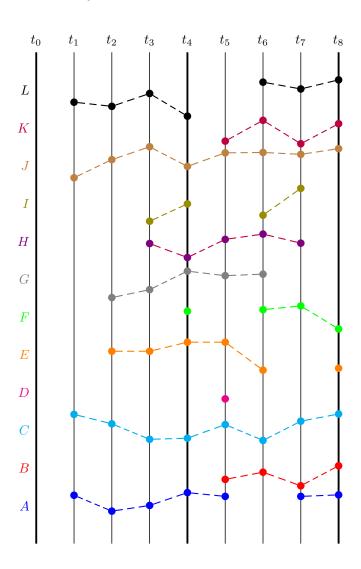




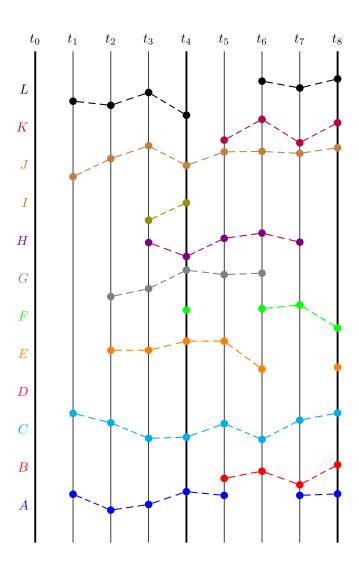
Filter t_7



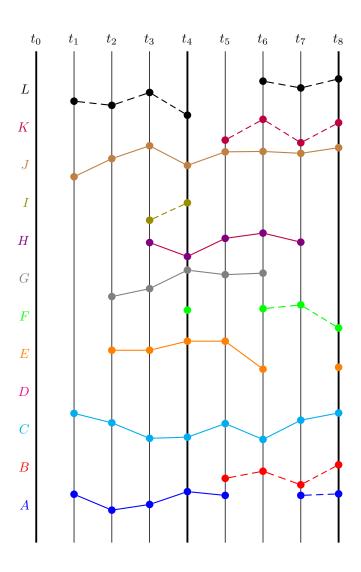




Filter t_5



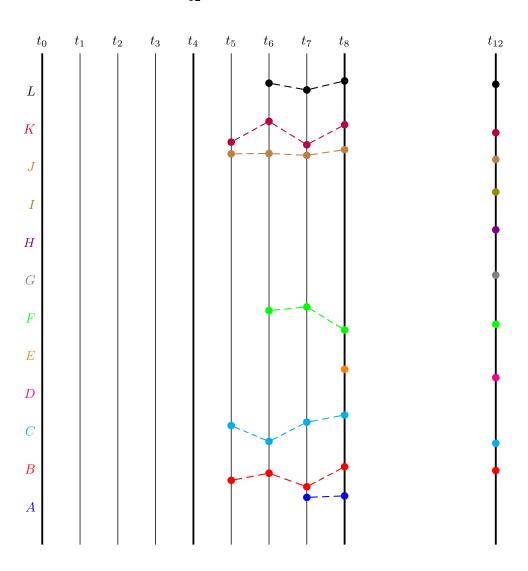
Flocks t_4 - t_8



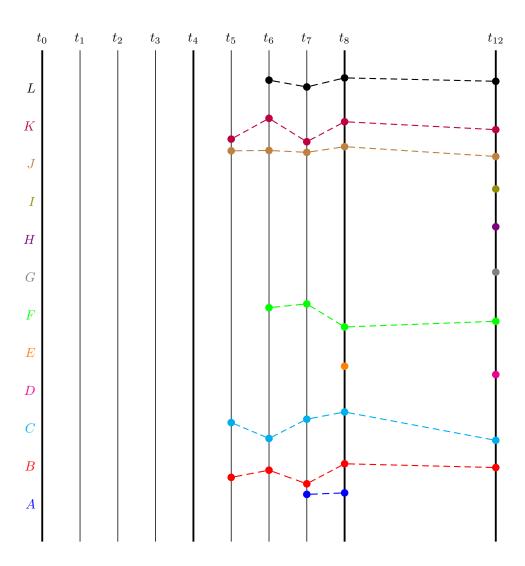
Prune t_4 - t_8

t_0	t	$_{1}$ t	$_{2}$ t	t_{3}	t_{ξ}	t_0	t_{6}	t_8
L						•		
K								
J								
I								
H								
G								
F								
E								
D								
C								
B							.	
A								

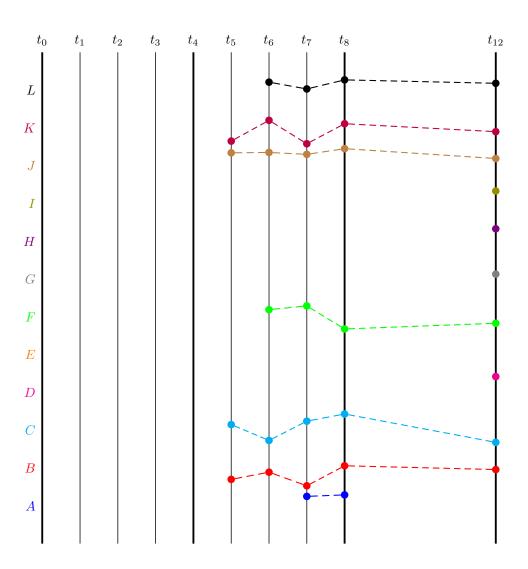
Maximal disks t_{12}



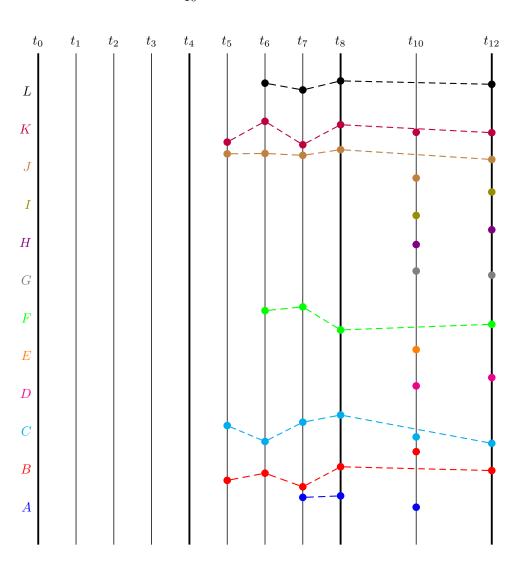
Join t_{12}



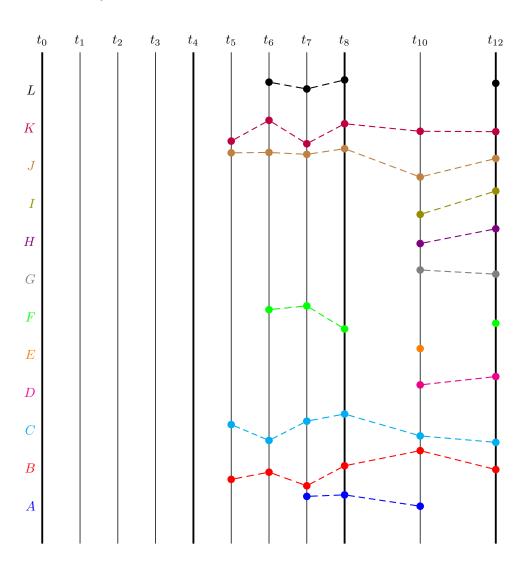
Filter t_{12}



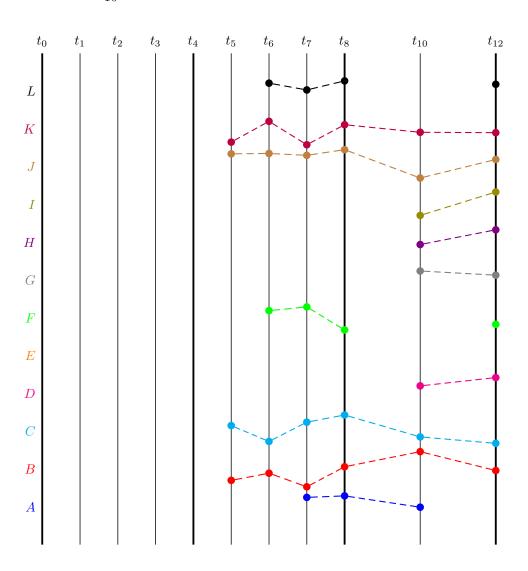
Maximal disks t_{10}



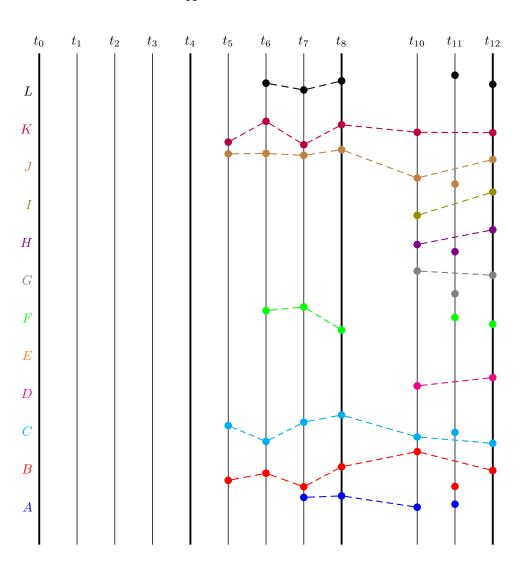
Join t_{10}



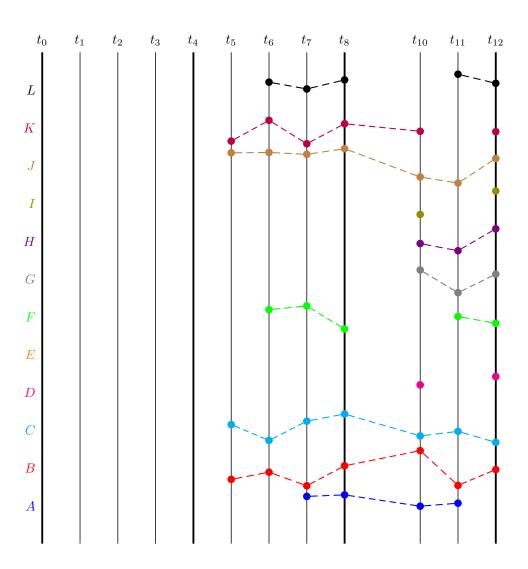
Filter t_{10}



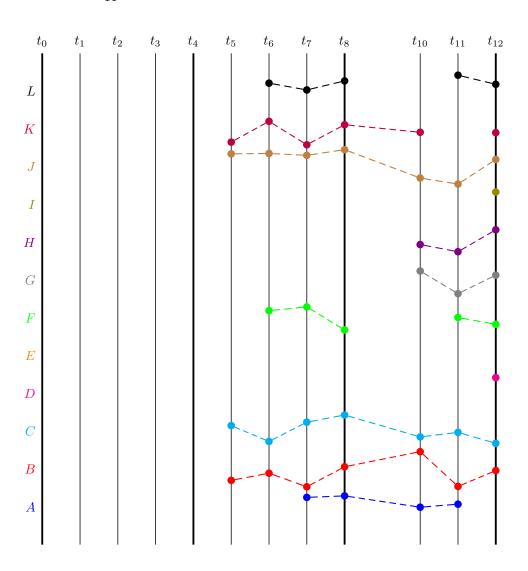
Maximal disks t_{11}



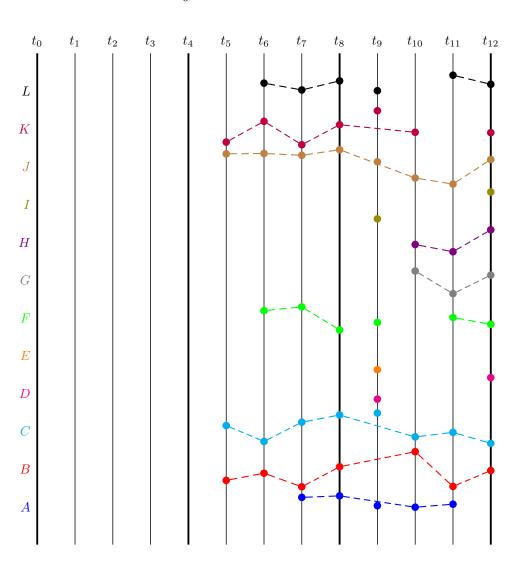
Join t_{11}



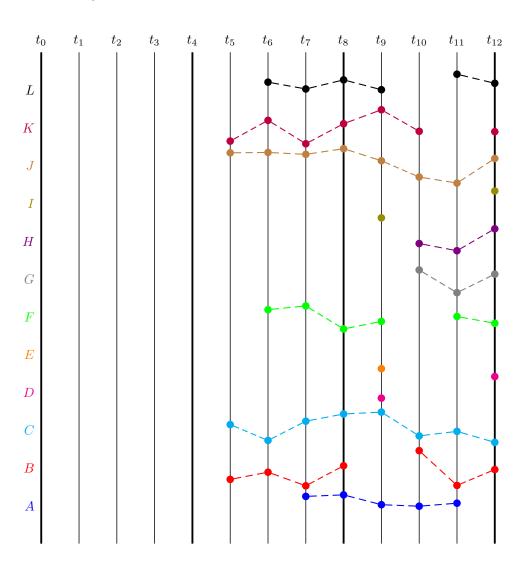
Filter t_{11}



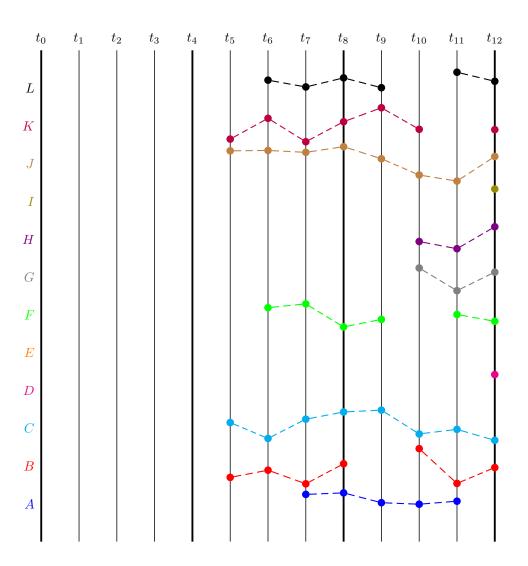
Maximal disks t_9



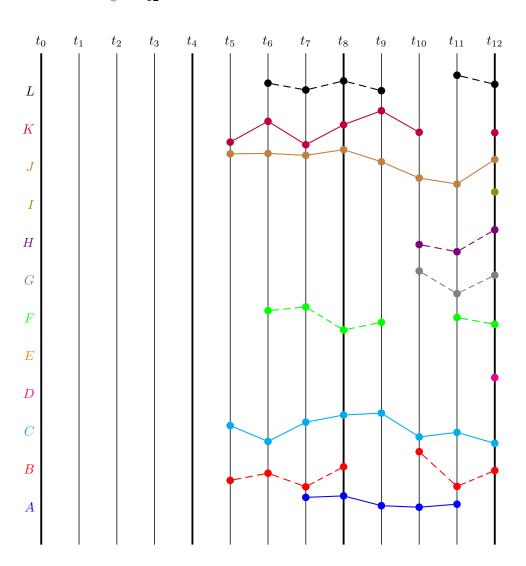
Join t_9



Filter t_9



Flocks t_8 - t_{12}



Prune t_8 - t_{12}

t_0	t_1	t_2	t_3	t_4	t_5	t_6	t_7	t_8	t_9	t_{10}	t_{11}	t_{12}
L												-
K												•
J									•	-		_
I												1
H											-	_1
G										•		-
F												- 🗼
E												
D												
C										\ \ \ \ -		
B											\	_
A											Ĭ	

MergeLast issues

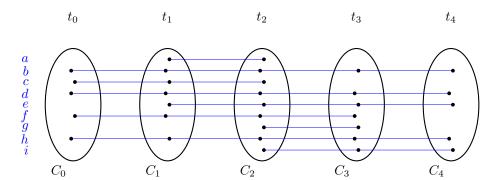
Drawbacks:

- Has to store intermediate points to deal appropriately with 'holes' in flocks.
- Join, subset elimination and consecutive checking at each timestamp.
- Expensive operations to maintain set of candidate flocks.

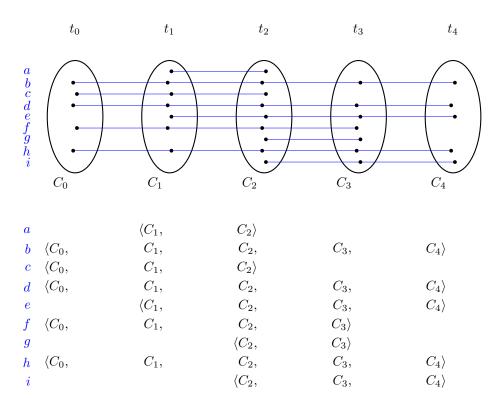
However, it has some interesting features to keep in mind:

- Performs an early pruning of candidate flocks which do not touch the borders.
- Keeping tracking of candidates which touch the borders ease a parallel implementation.

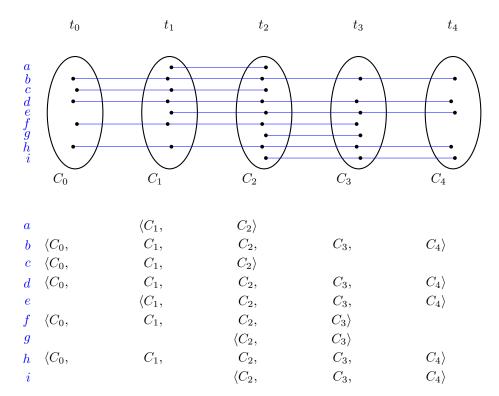
From trajectories to transactions



From trajectories to transactions



From trajectories to transactions



If we apply a Maximal Pattern (MP) algorithm over the new transactions...

$$MP = \langle C_0,$$

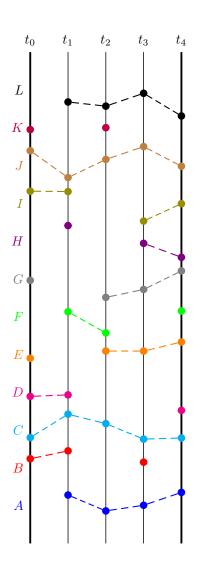
 C_1 ,

 C_2 ,

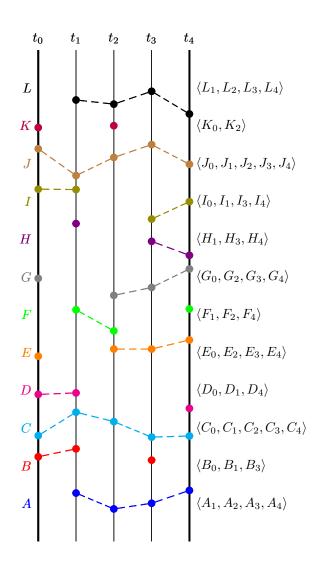
 C_3 ,

 $C_4\rangle$: $3(min_sup \ge \mu)$

MP finding per window...



MP finding per window...



MP finding per window...

Each trajectory is associated with just the maximal disks it touches. MP algorithms returns sets of disks which are visited by the same trajectories. If they happen in consecutive order, it is a flock.

Pros:

- Do not perform distance join at each timestamp.
- Although still have to deal with consecutive checking, it is done just at the end of the window.
- It deals with subset elimination.

Cons:

 Overlapping disks could introduce false flocks. It will require an additional filter at the end of the window.

Some reading...

- B. Negrevergne, A. Termier, J.-F. Mhaut, and T. Uno, Discovering closed frequent itemsets on multicore: Parallelizing computations and optimizing memory accesses, in High Performance Computing and Simulation (HPCS), 2010 International Conference on, 2010, pp. 521528.
- M. Kirchgessner, Mining and ranking closed itemsets from large-scale transactional datasets, Universit Grenoble Alpes, 2016.
- S. Cong, J. Han, and D. Padua, Parallel mining of closed sequential patterns, in Proceedings of the eleventh ACM SIGKDD international conference on Knowledge discovery in data mining, 2005, pp. 562567.