Update

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Outline

Algorithms

Performance evaluation



Algorithms

- ► Split the problem in two stages:
 - 1. Find maximal disks at each timestamp (MaximalFinder) and
 - 2. Join maximal disks between adjancent timestamps (FlockFinder)
- ► Pseudocode for both algorithms available online: MaximalFinder¹ and FlockFinder².



¹https://tinyurl.com/y741ld5k

²https://tinyurl.com/yac26guf

Maximal finder overall steps

- 1. Indexing points...
- 2. Getting pairs...
- 3. Computing centers...
- 4. Indexing centers...
- 5. Getting disks...
- 6. Filtering disks $< \mu$...
- 7. Prunning duplicate candidates...
- 8. Indexing candidates...
- 9. Getting expansions...
- 10. Finding maximal disks.



Flock finder

- 1. Set of disks for t_i ...
- 2. Set of disks for $t_{i+\delta}$...
- 3. Joining timestams...
- 4. Checking internal timestamps.



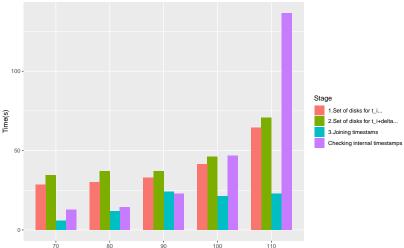
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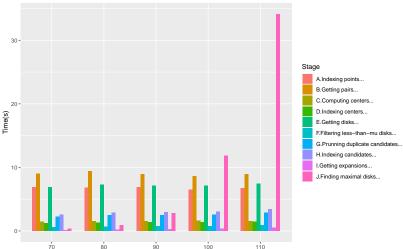


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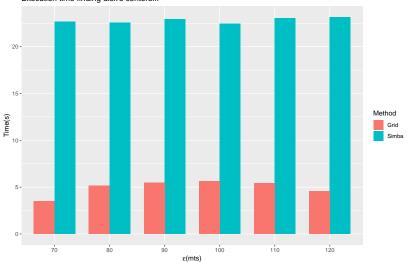
Performance





Performance





Bottlenecks

1. In flock finder:

▶ Checking internal timestamps: When merge last approach prunes enough points it works as expected but large amount of intermediate points have huge impact.

2. In maximal finder:

► Finding maximal disks: Even the new implementation is more stable, the most costly operation is removing duplicates and redundant disks.

3. Overall:

- ▶ Online approach requires indexing at each timestamp.
- Simba indexing is slow.

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- 1. Explore alternatives in Simba³
 - ► There are QuadTree and KDTree partitioners but they are not fully-integrated as indeces. (QuadTree only support 2D.)
 - ► For partitioning, RTree is already faster than QuadTree and KDTree in 2D and 3D datasets.
- 2. Grid indexing
 - Include the Grid partitioner in Simba and work on its index integration.
 - ► Implement distance join by my own.

