

### 2.1 Pizza or Pasta

#### Module ALICE:

Init:

```

custPizzaBuffer =  $\emptyset$ 
custPastaBuffer =  $\emptyset$ 
pizzaBuffer =  $\emptyset$ 
pastaBuffer =  $\emptyset$ 

```

```

upon  $\langle \text{Bob}, \text{getJob} \mid \text{NULL} \rangle$ 
  while (pizzaBuffer ==  $\emptyset$ ) {}
  dish  $\leftarrow$  pizzaBuffer
  trigger  $\langle \text{Bob}, \text{prepareDish} \mid \text{dish} \rangle$ 

```

```

upon  $\langle \text{Carole}, \text{getJob} \mid \text{NULL} \rangle$ 
  while (pastaBuffer ==  $\emptyset$ ) {}
  dish  $\leftarrow$  pastaBuffer
  trigger  $\langle \text{Carole}, \text{prepareDish} \mid \text{dish} \rangle$ 

```

```

upon  $\langle \text{Customer}, \text{orderDish} \mid \text{customer}, \text{menu} \rangle$  do
  if (menu.sortOf(Pizza))
    custPizzaBuffer = custPizzaBuffer  $\cup$  customer
    pizzaBuffer = pizzaBuffer  $\cup$  menu
  else
    custPastaBuffer = custPastaBuffer  $\cup$  customer
    pastaBuffer = pastaBuffer  $\cup$  menu

```

```

upon  $\langle \text{Bob}, \text{returnDish} \mid \text{menu} \rangle$  do
  customer  $\leftarrow$  custPizzaBuffer
  custPizzaBuffer = custPizzaBuffer  $\setminus$  customer
  trigger  $\langle \text{Alice}, \text{serveDish} \mid \text{customer}, \text{menu} \rangle$ 

```

```

upon  $\langle \text{Carole}, \text{returnDish} \mid \text{menu} \rangle$  do
  customer  $\leftarrow$  custPastaBuffer
  custPastaBuffer = custPastaBuffer  $\setminus$  customer
  trigger  $\langle \text{Alice}, \text{serveDish} \mid \text{customer}, \text{menu} \rangle$ 

```

#### Module BOB:

Init:

```

bobPizzaBuffer =  $\emptyset$ 
triggerCounter = 0

```

```

while (TRUE)
  if ( $| \text{bobPizzaBuffer} | + \text{triggerCounter} < 3$ )
    trigger  $\langle \text{Bob}, \text{getJob} \mid \text{NULL} \rangle$ 
    triggerCounter ++
  if ( $| \text{bobPizzaBuffer} | \neq 0$ )
    dish  $\leftarrow$  bobPizzaBuffer
    prepare the dish
    bobPizzaBuffer = bobPizzaBuffer  $\setminus$  dish
    trigger  $\langle \text{Bob}, \text{returnDish} \mid \text{dish} \rangle$ 

```

```

upon  $\langle \text{Bob}, \text{prepareDish} \mid \text{dish} \rangle$ 
  bobPizzaBuffer = bobPizzaBuffer  $\cup$  dish
  triggerCounter -- = 1

```

#### Module CAROLE:

Init:

```

carolePastaBuffer =  $\emptyset$ 
triggerCounter = 0

```

```

while (TRUE)
  if ( $| \text{carolePastaBuffer} | + \text{triggerCounter} < 7$ )
    trigger  $\langle \text{Carole}, \text{getJob} \mid \text{NULL} \rangle$ 
    triggerCounter ++
  if ( $| \text{carolePastaBuffer} | \neq 0$ )
    dish  $\leftarrow$  carolePastaBuffer
    prepare the dish
    carolePastaBuffer = carolePastaBuffer  $\setminus$  dish
    trigger  $\langle \text{Carole}, \text{returnDish} \mid \text{dish} \rangle$ 

```

```

upon  $\langle \text{Carole}, \text{prepareDish} \mid \text{dish} \rangle$ 
  carolePastaBuffer = carolePastaBuffer  $\cup$  dish
  triggerCounter -- = 1

```

## 2.2 Safety and liveness

- (a) *If some general attacks at time  $t$ , then the other general attacks at the same time.*  
This is a **safety property**, because
- (b) *If  $m_2$  arrives after time  $t$ , then General A attack after General B.*  
This is a **safety property**, because this statement ensures that General A will attack if and only if it receives the message delivered by  $m_2$  and not just out of thin air.
- (c) *Eventually, General B will attack.*  
This is a **liveness property**, because
- (d) *If messengers  $m_1$  and  $m_2$  are not intercepted, then eventually both generals attack.*  
This is a **liveness property**, because
- (e) *If  $m_1$  and  $m_2$  are not intercepted, then eventually both generals attack at time  $t$ .*  
This is a **mixture**, because we have the safety property that both messenger will arrive in time - so before time  $t$  - and the liveness property that eventually both attack at the same time  $t$ .

## 2.3 Unreliable clocks

- (a) *Find two examples, where timing issues lead to safety violations.*
- (b) *Find two examples, where timing issues lead to liveness violations.*  
As a first example, we can take the case from exercise 2.2. In there it could happen that  $m_2$  is delivering the message after the time  $t$ . Therefore General A will attack after General B which violates the liveness property.