# **Exercise 8**

#### 8.1 Total-order broadcast using consensus (4pt)

Consider Algorithm 6.1 (Consensus-Based Total-Order Broadcast) and observe that the payload messages proposed to consensus are represented as a set.

- a) What happens if the messages decided by consensus are not sorted deterministically before *tob-delivering* them?
- b) Suppose one leaves out the sorting of decided payload messages. How can one modify the algorithm so that every process nevertheless outputs the same ordered sequence of messages?

### 8.2 Atomic register as a replicated state machine (4pt)

One of the simplest state machines is the register abstraction discussed earlier. Implement an (N, N)-atomic register using a primitive for total-order broadcast.

The replicated state is the current value of the register and the relevant commands are write(v) and  $read() \rightarrow v$ . Use Module 4.3 [CGR11] and show that the implemented register satisfies the respective properties.

## 8.3 Replicated register with local read (2pt)

The atomic-register implementation of Ex. 8.2 sends every operation through total-order broadcast. Describe a modification of the protocol, where *reads* are executed locally, without being *tob-broadcast* to all processes. One often finds this optimization in practical systems. Does the resulting register still satisfy *atomicity*? Either argue why the property holds or demonstrate an execution that violates atomicity.

#### References

[CGR11] C. Cachin, R. Guerraoui, and L. Rodrigues, *Introduction to reliable and secure distributed programming (Second Edition)*, Springer, 2011.