Solution for exercise 7

7.1 Worst-case latency of eager reliable broadcast (2pt)

The *latency* of a broadcast protocol is defined as the maximum number of point-to-point message transmissions taken by the protocol, measured from when the sender broadcasts a payload message until every receiver has delivered it.

Consider the *eager reliable broadcast* protocol [CGR11, Algorithm 3.3]. What is its maximum latency? Describe a corresponding execution.

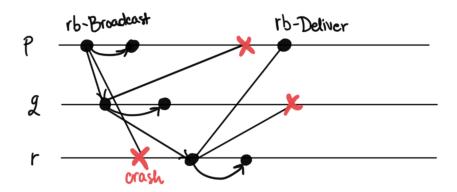


Figure 1. An example of the worst case execution of eager reliable broadcast.

7.2 Uniform reliable broadcast in the fail-stop model (4pt)

Algorithm 3.4 [CGR11, Sec. 3.4.2] implements uniform reliable broadcast using a perfect failure detector \mathcal{P} .

- a) What may happen if \mathcal{P} does not satisfy its *strong completeness* property?
- b) What may happen if \mathcal{P} violates its *strong accuracy* property?

In each case, explain whether this violates liveness or safety.

Solution.

- a) Consider a system of three processes: p, q, and r. Suppose that p urb-broadcasts a message m. If strong completeness is not satisfied then p might never urb-deliver m if either of q and r crashes and p never detects their crash or beb-delivers m from them. Process p might wait indefinitely for them to relay m and the algorithm may violate the validity property which is a liveness property.
- b) Assume that $strong\ accuracy$ is violated and p falsely suspects q and r to have crashed. Process p eventually urb-delivers m. Assume that p crashes afterward. It might be the case that q and r have never beb-delivered m and have no way of knowing about m. Hence, they will not urb-deliver m, violating the $uniform\ agreement$ property which is a liveness property.

7.3 FIFO broadcast from FIFO links (4pt)

An implementation of reliable broadcast with FIFO (first-in, first-out) order is found in Algorithm 3.12 [CGR11]. It runs on top of a reliable broadcast primitive. Here we consider a different implementation of FIFO reliable broadcast, without underlying reliable broadcast.

The usual specification of perfect point-to-point links (Module 2.3) does not impose any ordering. Assume, additionally, that the links between all servers are FIFO, that is, messages are delivered over each link *per process* in the order they were sent. A specification of FIFO links is available in Module 2.11 (Section 2.9.) Run the *basic broadcast* (Algorithm 3.1) using these FIFO links instead of perfect point-to-point links; this gives a FIFO best-effort broadcast (FIFO-BEB).

Now run the *eager reliable broadcast* protocol (Algorithm 3.3) with this FIFO-BEB (and the FIFO links). Does this implement FIFO-order reliable broadcast? Justify your answer.

Solution. Yes. Note that the *eager reliable broadcast* protocol uses best-effort broadcast and satisfies the properties of reliable broadcast. If it uses FIFO-BEB then its *FIFO delivery* property ensures that all DATA messages transmitted from one process to another one are delivered in FIFO order. Since a process *rb*-delivers every payload message immediately after receiving it for the first time through a DATA message, multiple messages from a specific sending process are also *rb*-delivered in the same order in which they were *rb*-broadcast. This, the modified implementation of reliable broadcast also satisfies the FIFO delivery order property.

References

[CGR11] C. Cachin, R. Guerraoui, and L. Rodrigues, *Introduction to reliable and secure distributed programming (Second Edition)*, Springer, 2011.