

## Exercise 4

### 4.1 Deterministic libraries (3pt)

A *deterministic* program is one that uses no random choices. Suppose  $\mathcal{L}_1$  and  $\mathcal{L}_2$  are two deterministic libraries with a common interface. Show that either  $\mathcal{L}_1 \equiv \mathcal{L}_2$ , or else  $\mathcal{L}_1$  &  $\mathcal{L}_2$  can be distinguished with advantage 1.

### 4.2 Hash-function collisions (2pt)

Consider an ideal hash function  $H : \{0, 1\}^* \rightarrow \{0, 1\}^\lambda$ , which can be modeled as follows:

$\mathcal{L}_{\text{idealhash}}$
$T := [ ]$
QUERY( $x \in \{0, 1\}^*$ ):
if $T[x]$ is undefined:
$T[x] \leftarrow \{0, 1\}^\lambda$
return $T[x]$

- For output length  $\lambda = 40$ , estimate the probability of finding a collision if one computes  $10^6$  hashes on arbitrary, distinct inputs.
- For  $\lambda = 256$  (SHA-256), estimate the number of hashes we need to compute to find a collision with probability  $\frac{1}{2}$ .

### 4.3 Salt (2pt)

Many computer systems use passwords for user authentication. It is common practice that the hash of a password is stored on a server instead of the clear-text password. This prevents that passwords are exposed directly in case the server is compromised.

The number of words in a language like German or English is estimated to be a few hundred thousand, but the active vocabulary is smaller than 100'000 words. Many users select a password from the vocabulary of their language and typically add a digit or two. To be concrete, assume below that every user selects a password randomly among  $2^{20}$  (roughly 1'000'000) words.

- Give an estimate on the number of users that are needed such that the password database contains two equal hashed passwords with probability  $\frac{1}{2}$  or more. Assume that the hash function is perfect, that is with no collisions.
- Salt* is a random data that enters the hash calculation of the password, but which is stored together with the password hash by the server. Hence, a typical password file contains an entry of the form

$$(username, salt, h)$$

for each user, where  $h = H(\textit{salt}||\textit{password})$  represents a hashed and “salted” password. Assume  $\textit{salt}$  is a random 256-bit value. What is now the estimated minimum number of users such that the password database contains two entries that have the same  $h$  values with probability  $\frac{1}{2}$  or more? Assume here that the hash function is perfect, i.e., that no collisions occur.