Distributed AlgorithmsLectures

1 Introduction - February 19, 2020

1.1 Defining Dependable Systems

QUOTES:

A distributed system is a system where a computer of which you did not know it exists can prevent you from getting your job done. - Leslie LAMPORT

There is perhaps a market for maybe five computers in the world. - TJ WATSON

 $FAULT \rightarrow ERROR \rightarrow FAILURE$

- Train delayed because of tree has fallen on the tracks
- Travelers reach destination too late
- Alice misses her exam

	FAULT	Error	FAILURE
Train:	Tree fallen	no train	delay for passengers
Journey:	Train delay	delay	reached destination 2h after intention
Exam:	arrival 2h late	missed time-slot	repeat exam

FAULT: cause of failure

ERROR: internal state of system, not according to specification

FAILURE: observable deviation of specification

FAULT examples:

- timing
- cables
- power supply
- messages lost
- data loss (solved with RAIDs)

1.1.1 How to make systems tolerate faults

- PREVENTION
- TOLERANCE
 - Replication/Redundancy
 - Recovery
- REMOVAL
- FORECASTING/PREDICTION

 $SAFETY \neq SECURITY$

SAFETY is connected to loss of live/material due to accidents

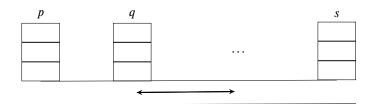
SECURITY is connected to malicious intent

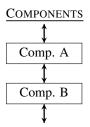
1.1.2 Defining distributed computation

Processes
$$\Pi = \{p, q, r, s \dots\}$$

 $\mid \Pi \mid = N$







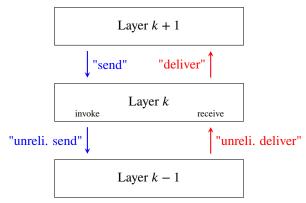
EVENTS for Component c:

$$\langle c, event \mid param_1, param_2 \dots \rangle$$

$$\frac{\text{upon } \langle c, ev_1 \mid param_1 \rangle \underline{\text{do}}}{\text{do something trigger } \langle b, domore \mid p \rangle}$$

$$\text{upon } \langle b, domore \mid p \rangle \underline{\text{do}}$$

1.1.3 Layered modules



Events either travel:

- upwards (red): indication - downwards (blue): request

Events on a given layer may be:

- input events (IN)
- output events (OUT)

Distributed Algorithms Lectures

1.1.4 Module Jobhandler

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Events:
   Request: \langle jh, handle \mid job \rangle
   Indication: \langle jh, confirm \mid job \rangle
Properties:
   Every job submitted for handling is eventually confirmed.
Implementation (synchronized) JOBHANDLER
State
upon \langle jh, handle \mid job \rangle do
    "process job"
   trigger \langle jh, confirm \mid job \rangle
upon ...
upon ...
Implementation (asynchronized) JOBHANDLER
State
   \overline{bu}f \leftarrow \emptyset
upon \langle jh, handle \mid job \rangle \underline{do}
   \overline{buf} \leftarrow buf \cup \{job\}
   trigger \langle jh, confirm \mid job \rangle
upon buf \neq \emptyset do
   \overline{job} \leftarrow \text{some element of } buf
    "process job"
   buf \leftarrow buf \setminus \{job\}
```

1.2 Concurrency and Replication in Distributed Systems

2 2nd Lecture - February 20, 2020