

Chapter 4

String Matching Algorithm

string-matching algorithms, sometimes called **string-searching algorithms**, are an important class of string algorithms that try to find a place where one or several strings (also called patterns) are found within a larger string or text.

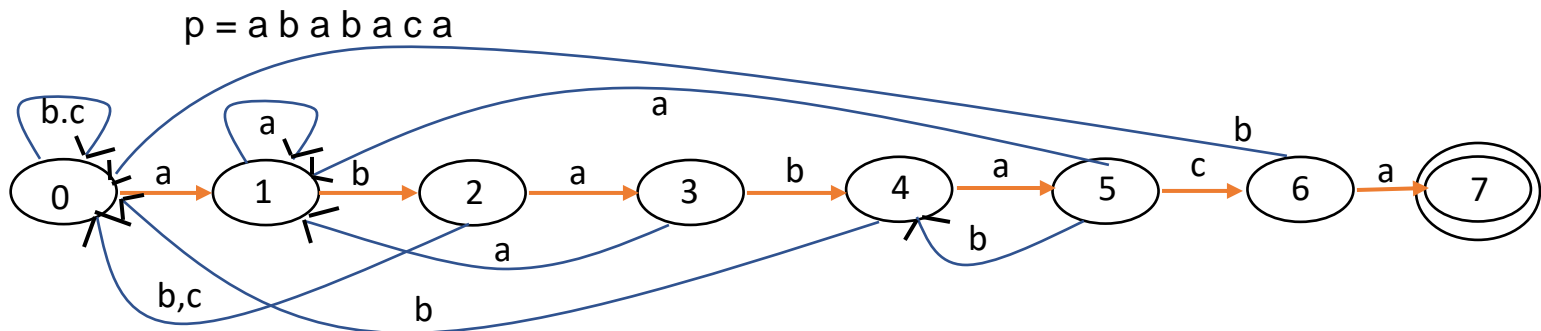
T → Text

aabbcdabbca

P → bca : Pattern

Algorithm

```
n ← length( T )
m ← length( P )
for ( s = 0; s <= ( n-m); s++)
    if ( p[1..m] = T[ s+1 ... s+m ]
        print “ pattern found”
    end if
end for
```



Transition Table

State	A	b	c
0	1	0	0
1	1	2	0
2	3	0	0
3	1	4	0
4	5	0	0
5	1	4	6
6	7	0	0
7	1/0	2/0	0/0

With loop without loop

Find_Shift_Key (T, R, P)

// T: text, R: Transition Table, P: Pattern

$n \leftarrow \text{length}(T)$

$q \leftarrow 0$ // state number

for ($i = 0$; $i \leq n$; $i++$)

$q = R[q, T[i+1]]$;

 if ($q == m$) // m: Final state

$s = i - m$;

 print " Pattern is at position ",s

 end if

end for

Example:

$T = a b a c a b a a c b a b a b a c a b a \dots$

$P = a b a b a c a$

$q = 0$

$q = R[0, a] = 1$

$q = R[1, b] = 2$

$q = R[2, a] = 3$

$q = R[3, c] = 0$

$q = R[0, a] = 1$

$q = R[1, b] = 2$

$q = R[2, a] = 3$

$q = R[3, a] = 1$

$q = R[1, c] = 0$

$q = R[0, b] = 0$

$q = R[0, a] = 1$

$q = R[1, b] = 2$

$q = R[2, a] = 3$

$q = R[3, b] = 4$

$q = R[4, a] = 5$

$q = R[5, c] = 6$

$q = R[6, a] = 7$

$q = R[7, b] = 2$

$q = R[2, a] = 3$