

# **Trees and Traversals**

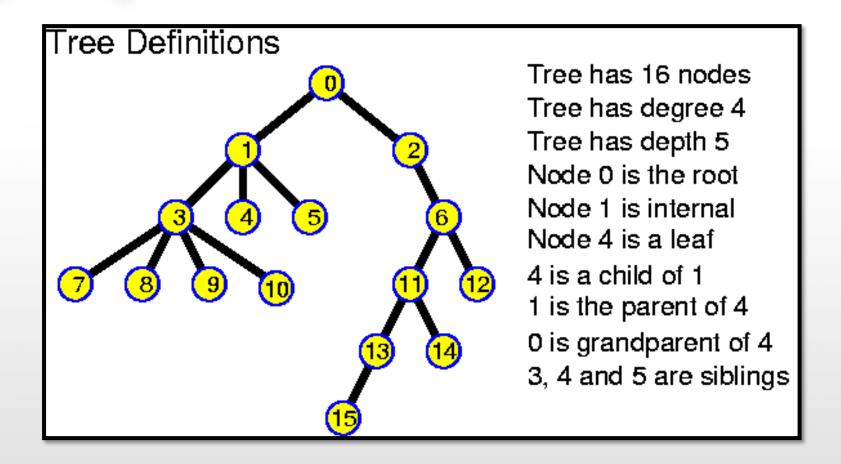
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COMP242

# Tree



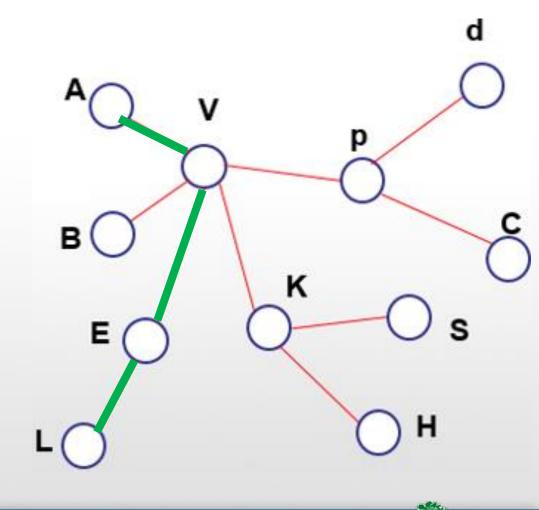
## Motivation



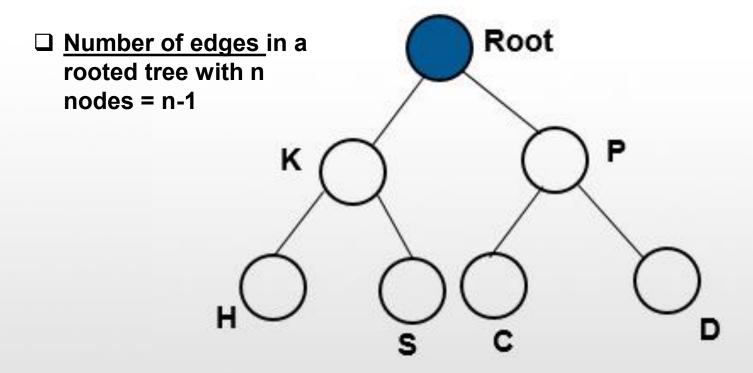
### Tree

- ☐ <u>Tree:</u> Set of nodes and edges that connect them.
- Exactly one path between any two nodes.
- □ <u>Path</u>: connected sequence of edges.

A **path** is a sequence of nodes such that the next node in the sequence is a child of the previous {A,V,E,L}



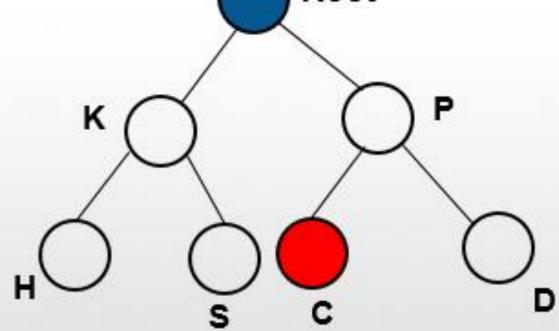
□ Rooted tree: One distinguished node is called the root.



□ Every node C, except root, has one parent P, the first node on path form c to the root.

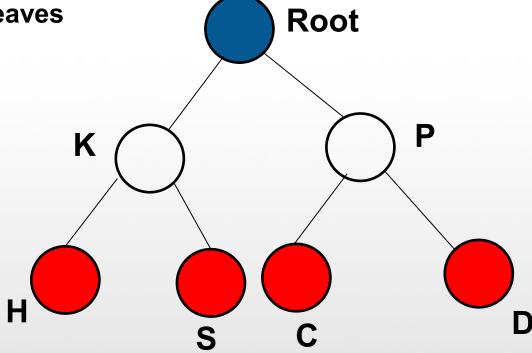
☐ if p is parent of C,then C is P's child ☐ Root

☐ Root has no parents



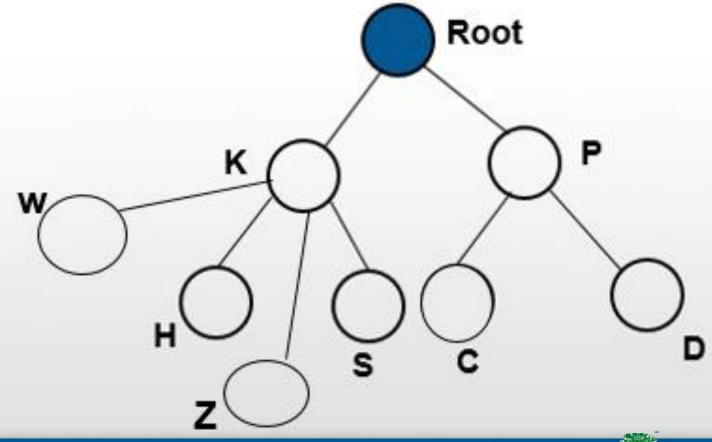
**Leaf**: Node with no children.

Example. H,S,C, and D are leaves in our tree

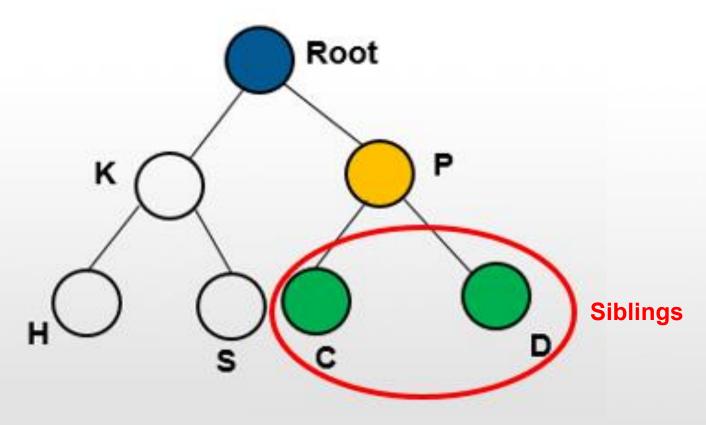


Leaf node also called external node, all other nodes are internal

☐ A node can have any number of children

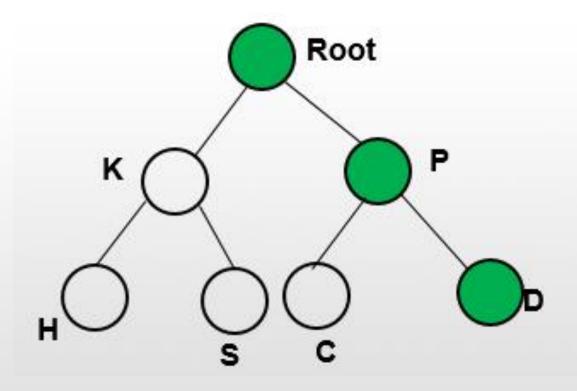


☐ Siblings : Nodes with same parent.

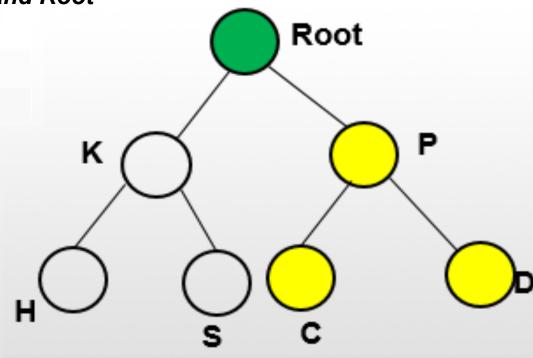


P is parent of C, and P is parent of D

- □ Ancestors of a node D : nodes on path from D to root, including D, D's parent, D's grandparent,...root (included).
- ☐ If P is <u>ancestor</u> of D, then D is <u>descendant</u> of P



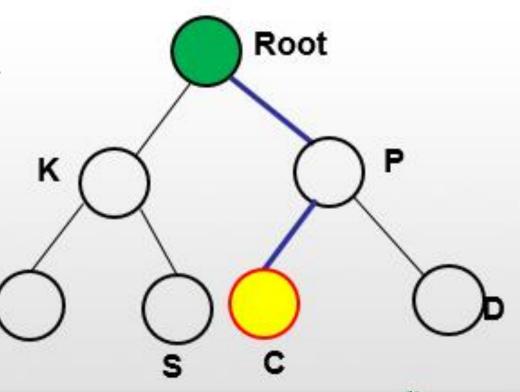
Example: Descendants of P are P,C, and D
Ancestors of H are H,K, and Root



□ Length of path: number of edges in path.

#### **Example:**

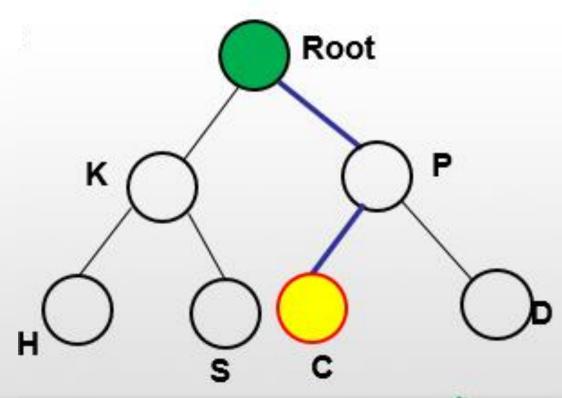
- → path from c to itself, the length is zero (empty path or no path)
- □ path from c to p, length is 1 (one edge in path)
- □ path from c to root, length is 2



□ Depth of node n is length of path from n to root.

EX: (Depth of root is zero)

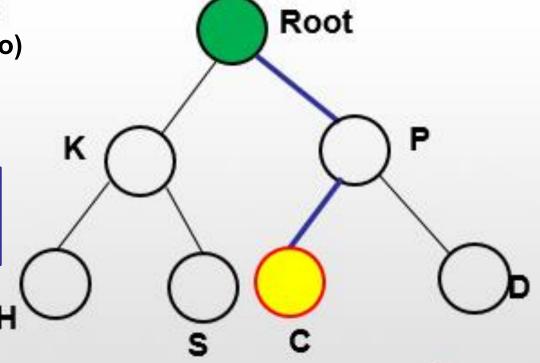
EX: (Depth of C is 2)



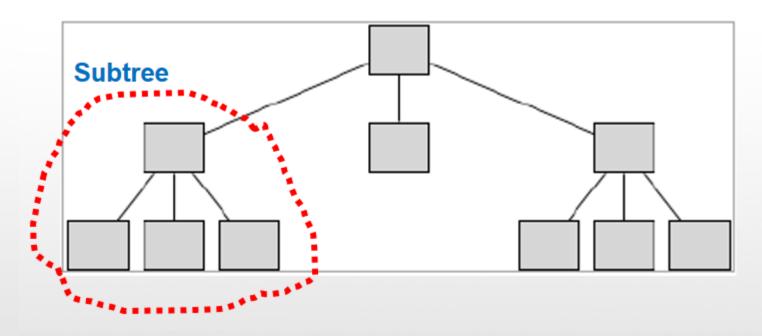
Height of node n is the length of path from n to its deepest descendant

Examples:
(Height of any leaf node is zero)
(Height of Root node is 2)
(Height of P node is 1)

Height of a tree= height of the root (The longest path length from the root to a leaf)



☐ <u>Degree</u>: the maximum number of possible children for every node in the tree.



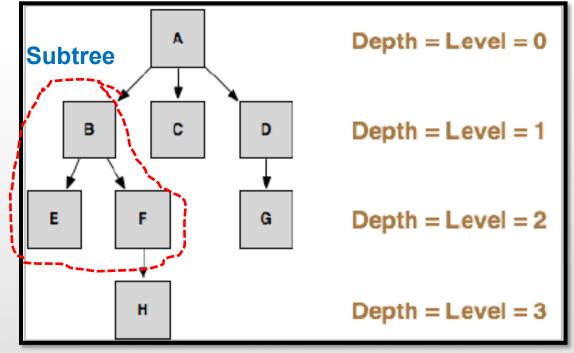
The height of tree is 2 and the degree 3

Node level & node depth: is the path length from the root

- The root is level 0 and depth 0
- Other nodes depth is 1 + depth of parent

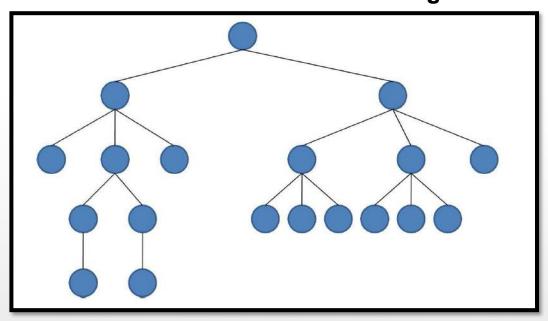
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depth(A)=0
depth(B)=1
depth(E)=2
depth(H)=3
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height(A)=3 height(F)=1 height(C)=0 height(E)=0 height(D)=1 height(B)=2



#### Rooted Tree: H.W

#### You have one week to do the following

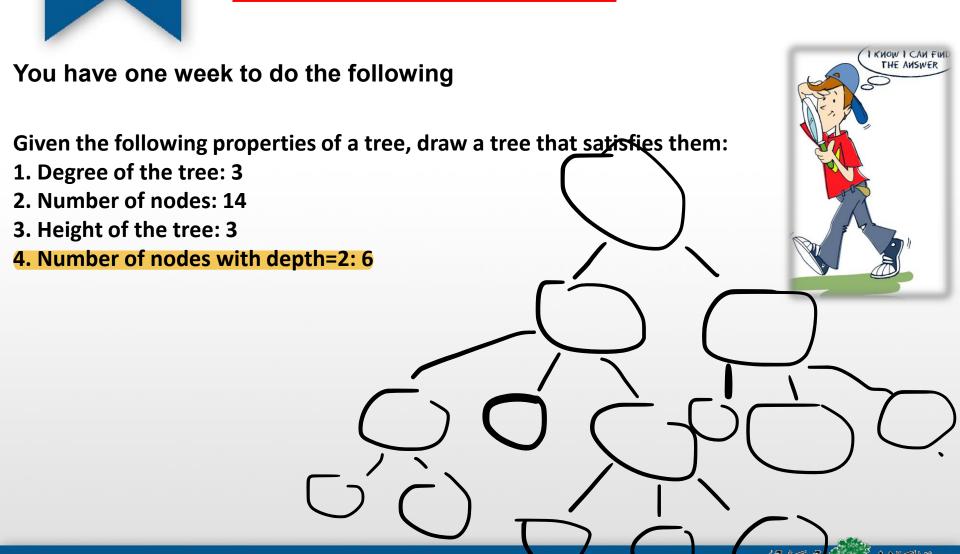




**□** Explain the values of the main characteristics of the tree shown in the figure.

NOTE: These characteristics are grade(degree) of the tree, height, number of nodes, external and internal nodes.

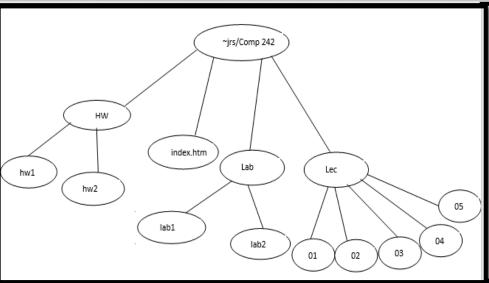
### Rooted Tree: H.W

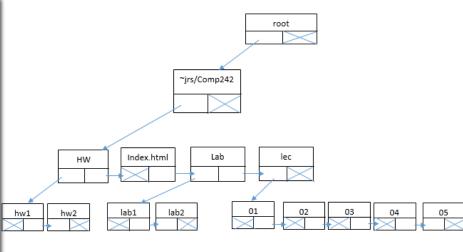


**Abdallah Karakr** 

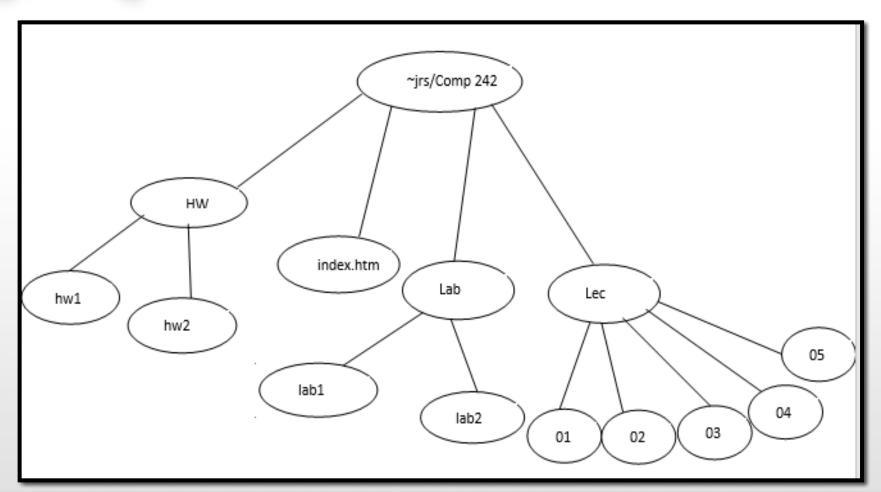
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### Implementation:Rooted Tree

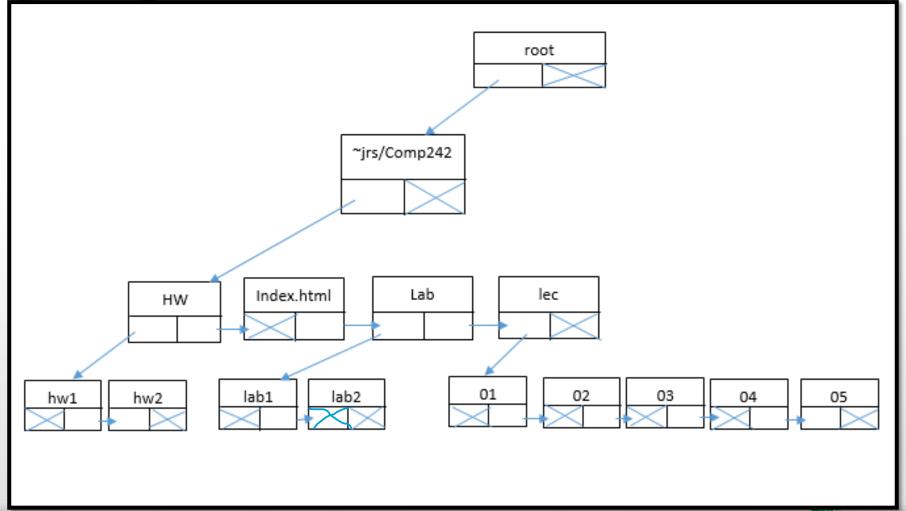




# Implementation:Rooted Tree



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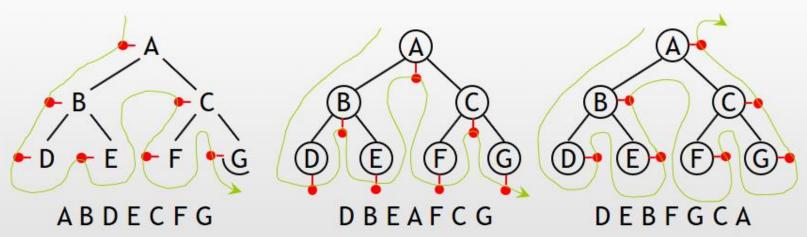


- ☐ A traversal is a manner of visiting each node in a tree once.
- What you do when visiting any particular node depends on the application; for instance, you might print a node's value, or perform some calculation upon it.
  There are several different traversals, each of which orders the nodes differently

□ Preorder: visits nodes as root → left → right

☐ Inorder: visits nodes as left → root → right

□ Postorder: visits nodes as left → right → root



preorder





Preorder

Inorder

Postorder

# Tree Traversal: More Details

#### Preorder traversal

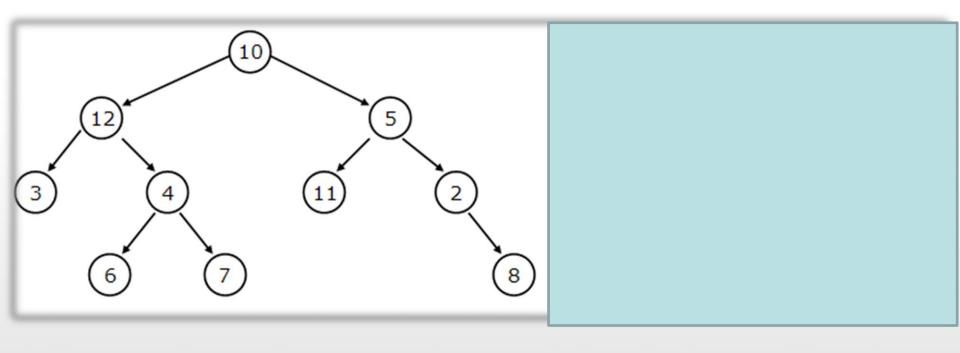
Let T be a tree with root r and subtrees  $T_1, T_2, ..., T_n$ . In Preorder traversal, we visit the root r first, then traverse the subtree  $T_1$  in preorder, then traverse the subtree  $T_2$  in preorder, and so on up to the traversal of the subtree  $T_n$  in preorder.

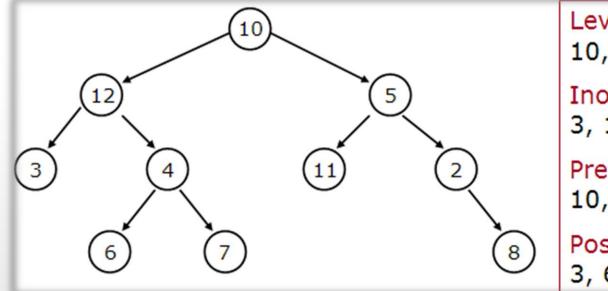
#### Inorder traversal

Let T be a tree with root r and subtrees  $T_1, T_2, ..., T_n$ . In an Inorder traversal, we traverse the subtree  $T_1$  in inorder, then we visit the root r, then traverse the subtree  $T_2$  in inorder, and so on up to the traversal of the subtree  $T_n$  in inorder.

#### Postorder traversal

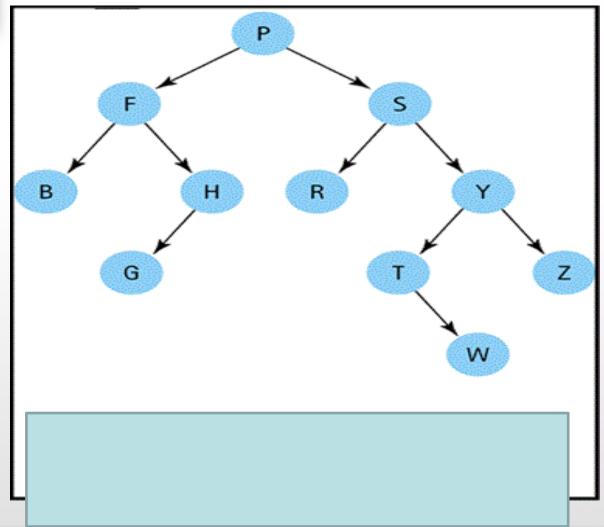
Let T be a tree with root r and subtrees  $T_1, T_2, ..., T_n$ . In a Postorder traversal, we traverse the subtree  $T_1$  in postorder, then traverse the subtree  $T_2$  in postorder, and so on up to the traversal of the subtree  $T_n$  in postorder, and finally we visit the root r.

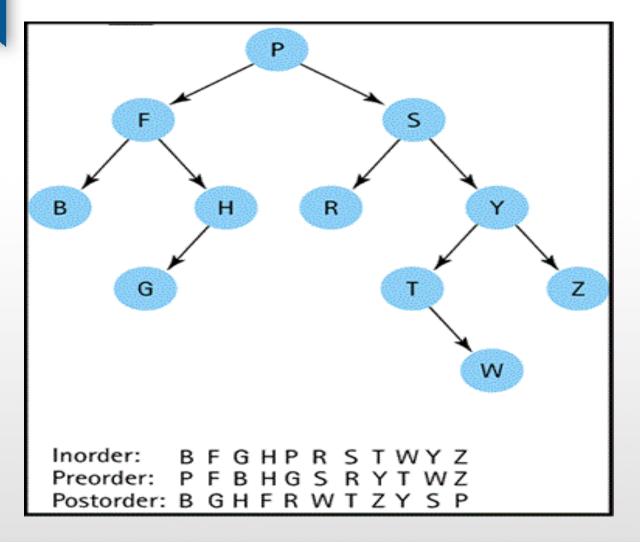


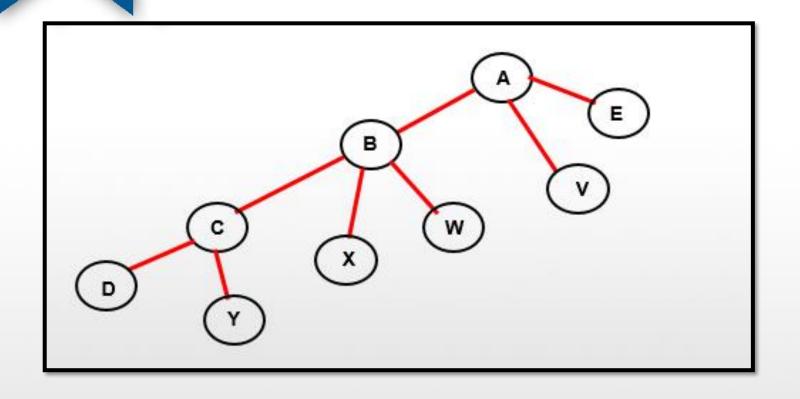


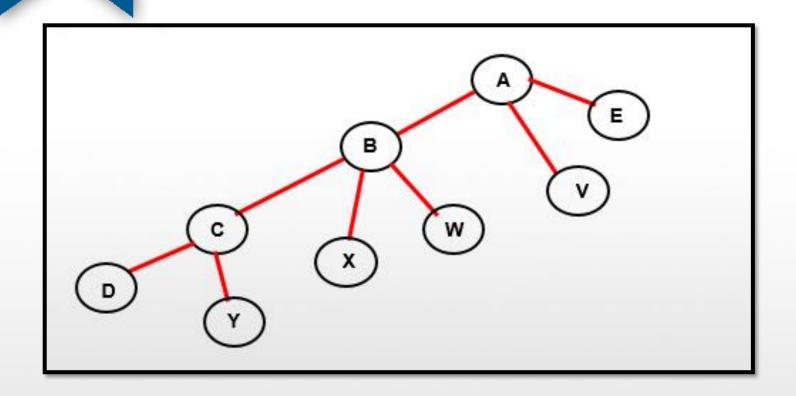
Levelorder tree traversal 10, 12, 5, 3, 4, 11, 2, 6, 7, 8 Inorder tree traversal 3, 12, 6, 4, 7, 10, 11, 5, 2, 8 Preorder tree traversal 10, 12, 3, 4, 6, 7, 5, 11, 2, 8

Postorder tree traversal 3, 6, 7, 4, 12, 11, 8, 2, 5, 10









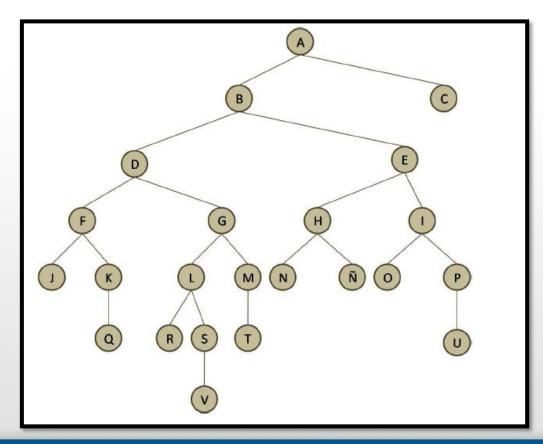
Pre-order: A,B,C,D,Y,X,W,V,E In-order: D,C,Y,B,X,W,A,V,E Post-order: D,Y,C,X,W,B,V,E,A

### Rooted Tree: H.W

You have one week to do the following

Given the following tree, write the pre-order, in-order and post-order

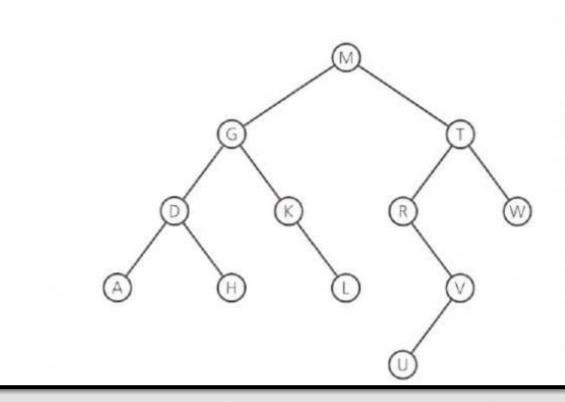
traversals.





# Extra Exercises

Carrano, 4th edition, Chapter 10, Exercise 2: What are the preorder, inorder, and postorder traversals of the following binary tree:



# Question?



"Success is the sum of small efforts, repeated day in and day out."
Robert Collier





- 1. Algorithms and Data Structures, Julian Moreno Schneider et al.
- 2. Fundamentals of Data Structures in C, Ellis Horowitz et al.
- 3. Data Structures and Problem Solving with C++: Walls and Mirrors
- 4. Analysis of algorithms robert Sedgewick
- 5. Prof. Sin-Min Lee Lecture Notes
- 6. Prof. Evan Korth Lecture Notes