2019-02-14 Skip Lists

Thursday, February 14, 2019

- o Finding, adding, and removing elements requires you to start at the front
- In last class, we used a "last accessed" pointer to speed up iteration

Downside of linked lists is that they don't support random access

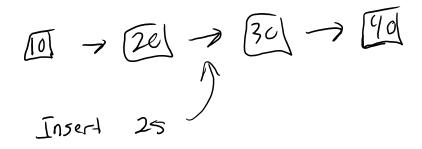
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- Made iteration *much* faster
- Downside: if we try to retrieve something physically before our "last access", we again have to start at the front of the list -> back to being slow
- The observation that linked lists become much faster with additional "cache pointers" is not very new.
 - The question then become: how many cache pointers do we need? And where should they be placed within the linked list?
- Suggestion: place pointer 1/2 way through LL
 - o Problem: where is 1/2 way? What happens when 1/2 way moves
 - Backwards is particularly bad because we'd have to reset from the front (SLOW!)
- Solution to the problem of "where to place pointers" is addressed in the Skip List data structure
 - Create "layers" of pointers.
 - Each layer points to data that is spatially farther away
 - Placement of pointers in these layers is handled randomly

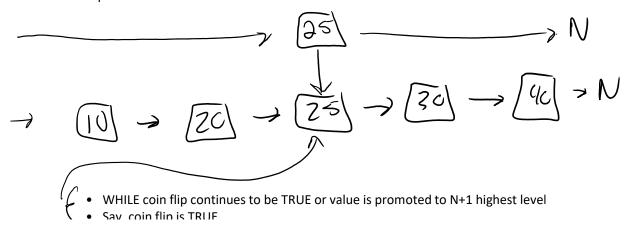
General Idea of Skip Lists

- Unlike normal linked lists, Skip Lists *must* be sorted
- On an insert of new data, flip a coin.
 - o If TRUE, turn this element into a pointer at layer N.
 - While coin flip continues to be true, add element to successively higher levels of pointers

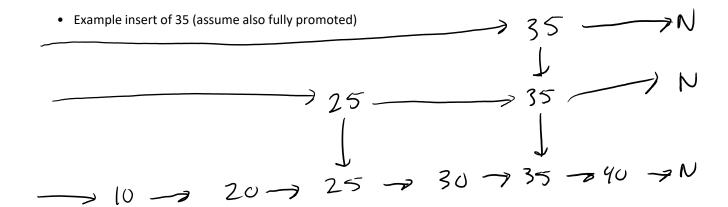
Graphical Example



• Step #1: Insert 25 as normal



• WHILE coin flip continues to be TRUE or value is promoted to N+1 highest level



- Performing operations on a Skip List start at the top layer and work down as needed
- Insert 15, with a single promotion

- If our base layer has N elements, on average how many elements does the layer immediately above it have?
 - o Each layer will, on average, have 1/2 the number of elements