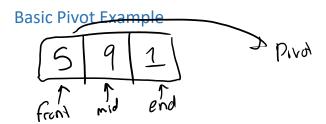
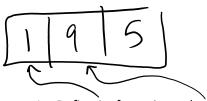
2:59 PM

## Quicksort

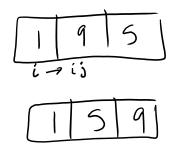
- General idea: find a "pivot." Put items smaller than the pivot to the "left" of the pivot and larger items to the "right" of the pivot.
- Recursively do this on the subarray bound by the pivot
- There are several ways that we might select a pivot.
- Richard Weiss suggests "best of 3" approach:
  - Pick middle-most value between front element, middle element, and end element
  - We do this in order to avoid QuickSort's worst-case N^2 runtime



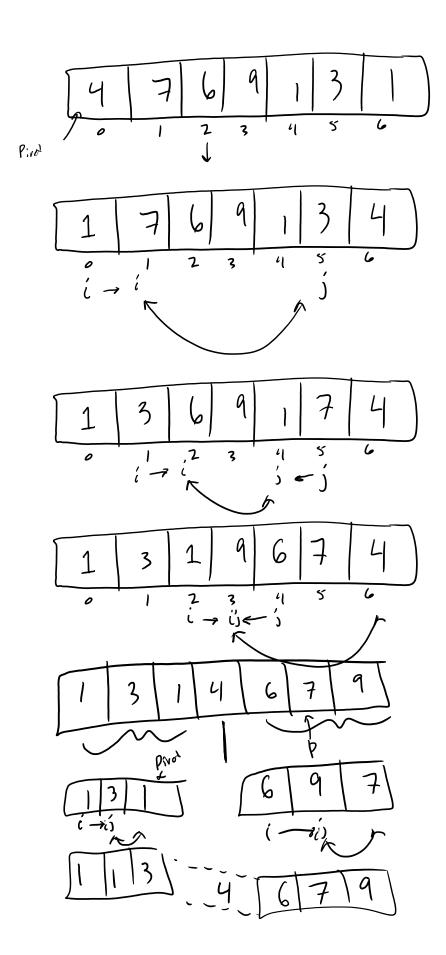
• Step #1: move pivot out of the way (swap with last element)



- 1. Define i = front; j = end 1
- 2. WHILE numbers[i] < pivot AND i < j
  - a. Increment i
- 3. WHILE numbers[j] > pivot AND i < j
  - a. Decrement j
- 4. If i != j:
  - a. Swap numbers[i] with numbers[j]
  - b. Go back to #1
- 5. Otherwise, loop is done. Swap numbers[end] with numbers[i]
  - 6. Repeat process recursively, splitting data on index i



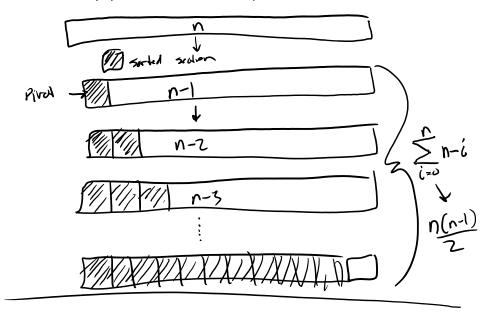
Another example



## **Analysis of Quicksort**

• Consider if we were to always choose the worst possible pivot

o Always pick smallest item in sequence



- Consider a quicksort that always picks the perfect pivot
  - o Perfect pivot is the median value

