Guru Nanak Institutions Technical Campus (Autonomous)

School of Engineering & Technology

DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Operating Systems Lab
Lab Manual [Subject Code: 22PC0CS05]
For the Academic year 2023-24

II B Tech Semester-II [CSE]



Guru Nanak Institutions Technical Campus(Autonomous)

Ibrahimpatnam, R R District – 501 506 (T.S.)



Department of Computer Science & Engineering

LAB MANUAL FOR THE ACADEMIC YEAR 2023-24

SUB : Operating Systems Lab

SUB CODE : 22PC0CS11

SEMESTER : II

STREAM : CSE

Document No : GNITC/CSE/

OS/R22Date of Issue: 15-02-2024

Prepared : Dr.Geeta Tripathi / Dr.Harish Kundra/ Mr. Md.Sirajul

Huque/ Mr. E. Kiran Kumar /Mr.Kumar Gaurav

PROGRAMMER: Chitti

VENUE : CSE-003 & 004

BLOCK : CSE

Verified by: Mr. B. Lalu

Authorized by Dr. Geeta Tripathi HOD-CSE

INDEX

S.No	Contents	Page.no
1	Lab Objective	4
2	Lab outcomes	5
3	Introduction About Lab	6
4	Guidelines to students a). Standard Operating Procedure (SOP) b). General guidelines	8-9
5	List of experiments as per the Atonoums curriculum	10
6	List of Additional Experiments	11
7	Text Books / Reference Books	11
8	Content of the experiment. Objective of the Experiment Hardware & Software Requirements Pre- requisite Flow chart/ algorithm with inputs and outputs Program Output Conclusion	12-107

1. OPERATING SYSTEMS LAB:

LAB OBJECTIVE

Upon successful completion of this Lab the student will be able to:

- To provide an understanding of the design aspects of operating system concepts through simulation
- Introduce basic Unix commands, system call interface for process management, interprocess communication and I/O in Unix

2. OPERATING SYSTEMS LAB:

LAB OUTCOME

Upon successful completion of this Lab the student will be able to:

- Simulate and implement operating system concepts such as scheduling, deadlock management, file management and memory management.
- Able to implement C programs using Unix system calls

3. INTRODUCTION ABOUT LAB

There are 60 systems (Compaq Presario) installed in this Lab. Their configurations are as follows:

Processor : Intel core i5

RAM : 4 GB

Hard Disk : 500 GB

Mouse : Optical Mouse

Software:

❖ All systems are configured with Linux Ubuntu or Windows 7 as per their lab requirement. This is very useful for students because they are familiar with different Operating Systems so that they can execute their programs in any programming environments.

- **❖ Software installed:** Turbo C/C++
- ❖ Systems are provided for students in the 1:1 ratio.
- Systems are assigned numbers and same system is allotted for students when they do the lab.

29. A. STANDARD OPERATING PROCEDURE - SOP

- a) Explanation on today's experiment by the concerned faculty using OHP/PPT/white Board covering the following aspects:60 mins.
 - 1) Name of the experiment/Aim
 - 2) Software/Hardware required
 - 3) Description about the program
 - 4) C Program code
- b) Writing of C programs by the student

30mins.

c) Compiling and execution of the program

90mins.

30. Writing of the experiment in the Observation Book:

The students will write the today's experiment in the Observation book as per the following format:

- a) Name of the experiment/Aim
- b) Software/Hardware required
- c) Source Progam
- d) Results for the written code
- e) Viva-Voce Questions and Answers
- f) Errors observed (if any) during compilation/execution
- g) Signature of the Faculty

4. Guide Lines to Students in Lab

Students are advised to maintain discipline and follow the guidelines given below:

- Keep all your bags in the racks and carry the observation book and record book.
- Mobile phones/pen drives/ CDs are not allowed in the labs.
- Maintain proper dress code along with ID Card
- Occupy the computers allotted to you and maintain the discipline.
- Student must submit the record with the last week experiment details and observation book with the brief of the present experiment.
- Read the write up of the experiment given in the manual.
- Students must use the equipment with care. Any damage is caused student is punishable
- After completion of every experiment, the observation notes to be shown to the lab in charge and after correction the record must be updated and submit to the lab in charge for
 correction.
- Lab marks are given on Continuous Evaluation Basis as per GNITC(A)guidelines
- If any student is absent for any lab, they need to be complete the same experiment in the free time before attending next lab session.

Steps to perform experiments in the lab by the student

Step1: Students have to write the Date, aim, Software and Hardware requirements for the scheduled experiment in the observation book.

Step2: Students have to listen and understand the experiment explained by the faculty and note down the important points in the observation book.

Step3: Students need to write procedure/algorithm in the observation book.

Step4: Analyze and Develop/implement the logic of the program by the student in respective platform

Step5: After approval of logic of the experiment by the faculty then the experiment has to be executed on the system.

Step6: After successful execution, the results have to be recorded in the observation book and shown to the lab in charge faculty..

Step7: Students need to attend the Viva-Voce on that experiment and write the same in the observation book.

Step8: Update the completed experiment in the record and submit to the concerned faculty incharge.

Instructions to maintain the record

- Before staring of the first lab session students must buy the record book and bring the same to the lab.
- Regularly (Weekly) update the record after completion of the experiment and get it corrected with concerned lab in-charge for continuous evaluation.
- In case the record is lost, inform on the same day to the faculty in charge and submit the new record within 2 days for correction.
- If record is not submitted in time or record is not written properly, the record evaluation marks (5M) will be reduced accordingly.

Awarding the marks for day to day evaluation:

Total marks for day to day evaluation is 15 Marks as per JNTUH. These 15 Marks are distributed as:

Record	5 Marks
Exp setup/program written	5 Marks
and execution	
Result and Viva-Voce	5 Marks

Allocation of Marks for Lab Internal

Total marks for lab internal are 40 Marks as per GNITC (AUTONOMOUS).

These 40 Marks are distributed as:

Average of day to day evaluation marks: 10 Marks

Lab Mid exam: 10 Marks

Viva Marks: 10 Marks

Additional lab project: 10 Marks

Allocation of Marks for Lab External

Total marks for lab External are 60 Marks as per GNITC AUTONOMOUS).

These 60 Marks are distributed as:

Program Written: 20 Marks

Program Execution and Result: 20 Marks

Viva-Voce: 10 Marks

Record: 10 Marks

5. List of Lab Exercises:

S. No	Name of the experiment	Page No
1	Write C programs to simulate the following CPU Scheduling algorithms a) FCFS b) SJF c) Round Robin d) priority	
2	Write programs using the I/O system calls of UNIX/LINUX operating system (open, read, write, close, fcntl, seek, stat, opendir, readdir)	
3	Write a C program to simulate Bankers Algorithm for Deadlock Avoidance and Prevention	
4	Write a C program to implement the Producer – Consumer problem using Semaphores using UNIX/LINUX system calls.	
5	Write C programs to illustrate the following IPC mechanisms a) Pipes b) FIFOs c) Message Queues d) Shared Memory	
6	Write C programs to simulate the following memory management techniques a) Paging b) Segmentation	
7	Write C programs to simulate Page replacement policies a) FCFS b) LRU c) Optimal	

6. List of additional experiments for the semester

S. No	Name of the experiment	
1		

TEXT BOOKS:

- 1. Operating System Principles- Abraham Silberchatz, Peter B. Galvin, Greg Gagne 7th Edition, John Wiley
- 2. Advanced programming in the Unix environment, W.R. Stevens, Pearson education.

REFERENCE BOOKS:

- 1. Operating Systems Internals and Design Principles, William Stallings, Fifth Edition–2005, Pearson Education/PHI
- 2. Operating System A Design Approach-Crowley, TMH.
- 3. Modern Operating Systems, Andrew S Tanenbaum, 2nd edition, Pearson/PHI
- 4. UNIX Programming Environment, Kernighan and Pike, PHI/Pearson Education
- 5. UNIX Internals: The New Frontiers, U. Vahalia, Pearson Education

7. Content of Lab Experiments

WEEK-1

Write C programs to simulate the following CPU Scheduling algorithms.

FCFS SJF Round Robin Priority

a) FCFS (First Come First Serve)

Aim: Write a C program to implement the various process scheduling mechanisms such as FCFS scheduling.

Algorithm:

- 1: Start the process
- 2: Accept the number of processes in the ready Queue
- 3: For each process in the ready Q, assign the process id and accept the CPU burst time
- 4: Set the waiting of the first process as '0' and its burst time as its turn around time
- 5: for each process in the Ready Q calculate
- a. Waiting time for process(n)= waiting time of process (n-1) + Burst time of process(n-1)
- b. Turnaround time for Process(n)= waiting time of Process(n)+ Burst time for process(n)
- 6: Calculate
- a. Average waiting time = Total waiting Time / Number of process
- b. Average Turnaround time = Total Turnaround Time / Number of process
- 7: Stop the process

Program:

```
#include<stdio.h>
int main()
  int bt[20],p[20],wt[20],tat[20],i,j,n,total=0,pos,temp;
  float avg_wt,avg_tat;
  printf("Enter number of process:");
  scanf("%d",&n);
  printf("\nEnter Burst Time:\n");
  for(i=0;i< n;i++)
     printf("p % d:",i+1);
     scanf("%d",&bt[i]);
                     //contains process number
     p[i]=i+1;
wt[0]=0;
             //waiting time for first process will be zero
//calculate waiting time
  for(i=1;i< n;i++)
     wt[i]=0;
     for(j=0;j< i;j++)
       wt[i]+=bt[j];
     total+=wt[i];
  }
  avg_wt=(float)total/n;
                            //average waiting time
  total=0;
  printf("\nProcess\t Burst Time \tWaiting Time\tTurnaround Time");
for(i=0;i< n;i++)
  {
     tat[i]=bt[i]+wt[i]; //calculate turnaround time
     total+=tat[i];
    printf("\np\%d\t\t \%d\t\t \%d\t\t, \np[i], \nt[i], \nt[i], \nt[i]);
  }
  avg_tat=(float)total/n; //average turnaround time
  printf("\n\nAverage Waiting Time=%f",avg_wt);
  printf("\nAverage Turnaround Time=%f\n",avg_tat);
}
```

Output:

```
[188r1a0501@localhost ~]$ vi w1.c
[188r1a0501@localhost ~]$ gcc w1.c
[188r1a0501@localhost ~]$ ./a.out
Enter number of process:3
Enter Burst Time:
p1:3
p2:4
p3:2
                               Waiting Time
Process
           Burst Time
                                                Turnaround Time
p1
p2
р3
Average Waiting Time=3.333333
Average Turnaround Time=6.333333
[188r1a0501@localhost ~]$ [
```

b) SJF (Shortest Job First)

Aim: Write a C program to implement the various process scheduling mechanisms such as SJF Scheduling.

Algorithm:

- 1: Start the process
- 2: Accept the number of processes in the ready Queue
- 3: For each process in the ready Q, assign the process id and accept the CPU burst time 4: Start the Ready Q according the shortest Burst time by sorting according to lowest to highest burst time.
- 5: Set the waiting time of the first process as '0' and its turnaround time as its burst time.
- 6: For each process in the ready queue, calculate
- (a) Waiting time for process(n)= waiting time of process (n-1) + Burst time of process(n-1)
- (b) Turn around time for Process(n)= waiting time of Process(n)+ Burst time for process(n)
- 7: Calculate
- (c) Average waiting time = Total waiting Time / Number of process
- Average Turnaround time = Total Turnaround Time / Number of process
 Stop the process

Program:

```
#include<stdio.h>
int main()
  int bt[20],p[20],wt[20],tat[20],i,j,n,total=0,pos,temp;
  float avg_wt,avg_tat;
  printf("Enter number of process:");
  scanf("%d",&n);
  printf("\nEnter Burst Time:\n");
  for(i=0;i< n;i++)
     printf("p%d:",i+1);
     scanf("%d",&bt[i]);
     p[i]=i+1;
                     //contains process number
  }
//sorting burst time in ascending order using selection sort
  for(i=0;i< n;i++)
     pos=i;
     for(j=i+1;j< n;j++)
       if(bt[j]<bt[pos])</pre>
          pos=j;
     }
     temp=bt[i];
     bt[i]=bt[pos];
     bt[pos]=temp;
     temp=p[i];
     p[i]=p[pos];
     p[pos]=temp;
wt[0]=0;
                //waiting time for first process will be zero
  //calculate waiting time
  for(i=1;i<n;i++)
  {
     wt[i]=0;
```

```
for(j=0;j< i;j++)
       wt[i]+=bt[j];
     total+=wt[i];
                            //average waiting time
  avg_wt=(float)total/n;
  total=0;
  printf("\nProcess\t Burst Time \tWaiting Time\tTurnaround Time");
for(i=0;i< n;i++)
  {
     tat[i]=bt[i]+wt[i];
                          //calculate turnaround time
     total+=tat[i];
     printf("\np\%d\t\t \%d\t\t \%d\t\t \%d\t\t\t\fi], tt[i], tt[i], tt[i], tt[i]);
  }
  avg_tat=(float)total/n; //average turnaround time
  printf("\n\nAverage Waiting Time=%f",avg_wt);
  printf("\nAverage Turnaround Time=%f\n",avg tat);
}
```

output

```
[188r1a0501@localhost ~]$ vi wla.c
[188r1a0501@localhost ~]$ gcc wla.c
[188r1a0501@localhost ~]$ ./a.out
Enter number of process:3
Enter Burst Time:
p1:5
p2:3
p3:7
Process
            Burst Time
                                Waiting Time
                                                 Turnaround Time
p2
р1
p3
                                                         15
Average Waiting Time=3.666667
Average Turnaround Time=8.666667
[188r1a0501@localhost ~]$
```

c) Round Robin

Aim: Write a C program to implement the various process scheduling mechanisms such as Round Robin Scheduling.

Algorithm

- 1: Start the process
- 2: Accept the number of processes in the ready Queue and time quantum (or) time slice
- 3: For each process in the ready Q, assign the process id and accept the CPU burst time
- 4: Calculate the no. of time slices for each process where
- No. of time slice for process(n) = burst time process(n)/time slice
- 5: If the burst time is less than the time slice then the no. of time slices =1.
- 6: Consider the ready queue is a circular Q, calculate
- (a) Waiting time for process(n) = waiting time of process(n-1)+ burst time of process(n-1) + the time difference in getting the CPU from process(n-1)
- (b) Turn around time for process(n) = waiting time of <math>process(n) + burst time of process(n) + the time difference in getting CPU from <math>process(n).
- 7: Calculate
- (a) Average waiting time = Total waiting Time / Number of process
- (b) Average Turnaround time = Total Turnaround Time / Number of process Step 8: Stop the process

Program:

```
#include<stdio.h>
main()
{
     int st[10],bt[10],wt[10],tat[10],n,tq;
     int i,count=0,swt=0,stat=0,temp,sq=0;
     float awt, atat;
     printf("enter the number of processes");
     scanf("%d",&n);
     printf("enter the burst time of each process /n");
     for(i=0;i<n;i++)
printf(("p%d",i+1);
scanf("%d",&bt[i]);
          st[i]=bt[i];
     }
     printf("enter the time quantum");
     scanf("%d",&tq);
     while(1)
     {
         for(i=0,count=0;i<n;i++)
          {
               temp=tq;
               if(st[i]==0)
                    count++;
                    continue;
               }
               if(st[i]>tq)
               st[i]=st[i]-tq;
```

```
else
                if(st[i]>=0)
                     temp=st[i];
                     st[i]=0;
                sq=sq+temp;
                tat[i]=sq;
           }
          if(n==count)
          break;
     }
     for(i=0;i< n;i++)
     {
          wt[i]=tat[i]-bt[i];
          swt=swt+wt[i];
          stat=stat+tat[i];
     }
     awt=(float)swt/n;
     atat=(float)stat/n;
     printf("process no\t burst time\t waiting time\t turnaround time\n");
     for(i=0;i<n;i++)
     printf("\%d\t\t\%d\t\t\%d\t\t\%d\t\t\%d\t\t\%d\t\t\%d\t\t\%d\t\t\%d\t,i+1,bt[i],wt[i],tat[i]);
     printf("avg wt time=%f,avg turn around time=%f",awt,atat);
}
```

Output:

```
[188rla0501@localhost ~]$ vi wld.c
[188rla0501@localhost ~]$ gcc wld.c
[188rla0501@localhost ~]$ ./a.out
enter the number of processes 3
enter the burst time of each process
pl 7
p2 2
p3 8
enter the time quantum 2
process no burst time waiting time turnaround time
1 7 8 15
2 2 2 4
3 8 9 17
avg wt time=6.333333, avg turn around time=12.000000[188rla0501@localhost ~]$
```

d) Priority

Aim: Write a C program to implement the various process scheduling mechanisms such as Priority Scheduling.

Algorithm:

- 1: Start the process
- 2: Accept the number of processes in the ready Queue
- 3: For each process in the ready Q, assign the process id and accept the CPU burst time 4: Sort the ready queue according to the priority number.
- 5: Set the waiting of the first process as '0' and its burst time as its turn around time
- 6: For each process in the Ready Q calculate
- (e) Waiting time for process(n)= waiting time of process (n-1) + Burst time of process(n-1)
- (f) Turn around time for Process(n)= waiting time of Process(n)+ Burst time for process(n)
- 7: Calculate
- (g) Average waiting time = Total waiting Time / Number of process
- (h) Average Turnaround time = Total Turnaround Time / Number of process Step 8: Stop the process

Program:

```
#include<stdio.h>
int main()
  int bt[20],p[20],wt[20],tat[20],pri[20],i,j,k,n,total=0,pos,temp;
  float avg_wt,avg_tat;
  printf("Enter number of process:");
  scanf("%d",&n);
  printf("\nEnter Burst Time:\n");
  for(i=0;i<n;i++)
  {
     printf("p%d:",i+1);
     scanf("%d",&bt[i]);
     p[i]=i+1;
                     //contains process number
printf(" enter priority of the process ");
for(i=0;i< n;i++)
p[i] = i;
//printf("Priority of Process");
printf("p%d ",i+1);
scanf("%d",&pri[i]);
}
for(i=0;i< n;i++)
for(k=i+1;k< n;k++)
if(pri[i] > pri[k])
{
 temp=p[i];
 p[i]=p[k];
 p[k]=temp;
 temp=bt[i];
bt[i]=bt[k];
bt[k]=temp;
temp=pri[i];
pri[i]=pri[k];
pri[k]=temp;
}
```

14

```
wt[0]=0;
                   //waiting time for first process will be zero
  //calculate waiting time
  for(i=1;i<n;i++)
     wt[i]=0;
    for(j=0;j< i;j++)
       wt[i]+=bt[j];
     total+=wt[i];
  }
                          //average waiting time
avg_wt=(float)total/n;
  total=0;
  printf("\nProcess\t Burst Time \tPriority \tWaiting Time\tTurnaround Time");
  for(i=0;i<n;i++)
     tat[i]=bt[i]+wt[i];
                          //calculate turnaround time
     total+=tat[i];
    printf("\np\%d\t\t\ \%d\t\t\ \%d\t\t\ \%d\t\t\ \%d\t\t\); p[i],bt[i],pri[i],wt[i],tat[i]);
  }
  avg_tat=(float)total/n;
                            //average turnaround time
  printf("\n\nAverage Waiting Time=%f",avg_wt);
  printf("\nAverage Turnaround Time=%f\n",avg_tat);
}
```

Output:

```
[188r1a0501@localhost ~]$ vi wlb.c
[188r1a0501@localhost ~]$ gcc wlb.c
[188r1a0501@localhost ~]$ ./a.out
Enter number of process:3

Enter Burst Time:
01:5
02:6
03:7
enter priority of the process pl 1
02 3
03 2

Process Burst Time Priority Waiting Time Turnaround Time
00 5 1 0 5
02 7 2 5 12
01 6 3 12 18

Average Waiting Time=5.666667

Average Turnaround Time=11.666667
```

WEEK-2

Write programs using the I/O system calls of UNIX/LINUX operating system (open, read, write, close, fcntl, seek, stat, opendir, readdir)

Aim: C program using open, read, write, close system calls

Theory:

There are 5 basic system calls that Unix provides for file I/O.

1. Create: Used to Create a new empty file

Syntax: int creat(char *filename, mode_t mode) filename: name of the file which you want to create

mode: indicates permissions of new file.

2. **open**: Used to Open the file for reading, writing or both.

Syntax: int open(char *path, int flags [, int mode]);

Path: path to file which you want to use

flags: How you like to use

O_RDONLY: read only, O_WRONLY: write only, O_RDWR: read and write, O_CREAT: create file if it doesn't exist, O_EXCL: prevent creation if it already exists

3. **close**: Tells the operating system you are done with a file descriptor and Close the file which pointed by fd.

Syntax: int close(int fd);

fd:file descriptor

4. **read:** From the file indicated by the file descriptor fd, the read() function reads cnt bytes of input into the memory area indicated by buf. A successful read() updates the access time for the file.

Syntax: int read(int fd, char *buf, int size);

fd: file descripter

buf: buffer to read data from

cnt: length of buffer

5. **write**: Writes cnt bytes from buf to the file or socket associated with fd. cnt should not be greater than INT_MAX (defined in the limits.h header file). If cnt is zero, write() simply returns 0 without attempting any other action.

Syntax: int write(int fd, char *buf, int size);

fd: file descripter

buf: buffer to write data to

cnt: length of buffer

^{*}File descriptor is integer that uniquely identifies an open file of the process.

Algorithm

- 1. Star the program.
- 2. Open a file for O_RDWR for R/W,O_CREATE for creating a file ,O_TRUNC for truncate a file.
- 3. Using getchar(), read the character and stored in the string[] array.
- 4. The string [] array is write into a file close it.
- 5. Then the first is opened for read only mode and read the characters and displayed it and close the file.
- 6. Stop the program.

Program

```
#include<sys/stat.h>
#include<stdio.h>
#include<fcntl.h>
#include<sys/types.h>
int main()
  int n,i=0;
  int f1,f2;
  char c,strin[100];
  f1=open("data",O_RDWR|O_CREAT|O_TRUNC);
  while((c=getchar())!='\n')
  {
    strin[i++]=c;
  }
  strin[i]='\0';
  write(f1,strin,i);
  close(f1);
  f2=open("data",O_RDONLY);
  read(f2,strin,0);
  printf("\n%s\n",strin);
  close(f2);
 return 0;
}
```

Output:

Hai Hai

b) Aim: C program using lseek

Theory:

lseek is a system call that is used to change the location of the read/write pointer of a file descriptor. The location can be set either in absolute or relative terms.

Syntax : off_t lseek(int fildes, off_t offset, int whence); int fildes : The file descriptor of the pointer that is going to be moved. off_t offset : The offset of the pointer (measured in bytes).

int whence: Legal values for this variable are provided at the end which are

SEEK_SET (Offset is to be measured in absolute terms), SEEK_CUR (Offset is to be measured relative to the current location of the pointer), SEEK_END (Offset is to be measured relative to the end of the file)

Algorithm:

- 1. Start the program
- 2. Open a file in read mode
- 3. Read the contents of the file
- 4. Use Iseek to change the position of pointer in the read process
- 5. Stop

Program:

```
#include<stdio.h>
#include <unistd.h>
#include <fcntl.h>
#include <sys/types.h>
int main()
     int file=0;
     if((file=open("testfile.txt",O_RDONLY)) < -1)
         return 1;
     char buffer[19];
     if(read(file,buffer,19)!=19) return 1;
     printf("%s\n",buffer);
    if(lseek(file,10,SEEK SET) < 0) return 1;
     if(read(file,buffer,19)!= 19) return 1;
     printf("%s\n",buffer);
    return 0;
}
```

Output:

```
[188rla0501@localhost ~]$ v1 testfile.txt
[188rla0501@localhost ~]$ cat testfile.txt
sdfghjkkl;mnbbvrtyuijnnb
ggtyujjg

[188rla0501@localhost ~]$ gcc w2c.c
[188rla0501@localhost ~]$ ./a.out
sdfghjkkl;mnbbvrtyu
mnbbvrtyuijnnb
ggty
```

c) **Aim:** C program using opendir(), closedir(), readdir()

Theory:

The following are the various operations using directories

1. Creating directories.

Syntax : int mkdir(const char *pathname, mode_t mode);

- 2. The 'pathname' argument is used for the name of the directory.
- 3. Opening directories

Syntax : DIR *opendir(const char *name);

4. Reading directories.

Syntax: struct dirent *readdir(DIR *dirp);

5. Removing directories.

Syntax: int rmdir(const char *pathname);

6. Closing the directory.

Syntax: int closedir(DIR *dirp);

7. Getting the current working directory.

Syntax: char *getcwd(char *buf, size_t size);

Algorithm:

- 1. Start the program
- 2. Print a menu to choose the different directory operations
- 3. To create and remove a directory ask the user for name and create and remove the same respectively.
- 4. To open a directory check whether directory exists or not. If yes open the directory .If it does not exists print an error message.
- 5. Finally close the opened directory.
- 6. Stop

Program:

```
#include<stdio.h>
#include<fcntl.h>
#include<dirent.h>
main()
char d[10]; int c,op; DIR *e;
struct dirent *sd;
printf("**menu**\n1.create dir\n2.remove dir\n 3.read dir\n enter ur choice");
scanf("%d",&op);
switch(op)
case 1: printf("enter dir name\n"); scanf("%s",&d);
c=mkdir(d,777);
if(c==1)
printf("dir is not created");
else
printf("dir is created"); break;
case 2: printf("enter dir name\n"); scanf("%s",&d);
c=rmdir(d);
if(c==1)
printf("dir is not removed");
else
printf("dir is removed"); break;
case 3: printf("enter dir name to open");
scanf("%s",&d);
e=opendir(d);
if(e==NULL)
printf("dir does not exist"); else
printf("dir exist\n"); while((sd=readdir(e))!=NULL) printf("%s\t",sd->d_name);
closedir(e);
break;
}
```

Output:

```
[188r1a0501@localhost f]$ gcc w2e.c
[188r1a0501@localhost f]$ ./a.out
**menu**
1.create dir
2.remove dir
3.read dir
enter ur choice1
enter dir name
d
dir is created[188r1a0501@localhost f]$ ls
a.out a.txt d w2d.c w2e.c
```

WEEK -3

Write a C program to simulate Bankers Algorithm for Deadlock Avoidance and Prevention

a) Aim

Write a C program to simulate the Bankers Algorithm for Deadlock Avoidance.

Data structures

- 1. n- Number of process, m-number of resource types.
- 2. Available: Available[j]=k, k instance of resource type Rj is available.
- 3. Max: If max [i, j]=k, Pi may request at most k instances resource Rj.
- 4. Allocation: If Allocation [i, j]=k, Pi allocated to k instances of resource Rj
- 5. Need: If Need[I, j]=k, Pi may need k more instances of resource type Rj,
- 6. Need [I, j] =Max [I, j]-Allocation [I, j];

Safety Algorithm

- 1. Work and Finish be the vector of length m and n respectively, Work=Available and Finish[i] =False.
- 2. Find an i such that both
- 3. Finish[i] =False
- 4. Need<=Work
- 5. If no such I exist go to step 4.
- 6. work=work+Allocation, Finish[i] =True;
- 7. If Finish [1] =True for all I, then the system is in safe state.

Resource request algorithm

- 1. Let Request i be request vector for the process Pi, If request i=[j]=k, then process Pi wants k instances of resource type Rj.
- 2. If Request<=Need I go to step 2. Otherwise raise an error condition.
- 3. If Request<=Available go to step 3. Otherwise Pi must since the resources are available.
- 4. Have the system pretend to have allocated the requested resources to process Pi by modifying the state as follows;
- 5. Available=Available-Request I;
- 6. Allocation I = Allocation + Request I;
- 7. Need i=Need i-Request I;

If the resulting resource allocation state is safe, the transaction is completed and process Pi is allocated its resources. However, if the state is unsafe, the Pi must wait for Request i and the old resource-allocation state is restore.

Algorithm:

- 1. Start the program.
- 2. Get the values of resources and processes.
- 3. Get the avail value.
- 4. After allocation find the need value.
- 5. Check whether it is possible to allocate.
- 6. If it is possible then the system is in safe state.
- 7. Else system is not in safety state.
- 8. If the new request comes then check that the system is in safety.
- 9. Or not if we allow the request.
- 10. Stop the program.

Program:

```
#include<stdio.h>
int main ()
 int allocated[15][15], max[15][15], need[15][15], avail[15], tres[15],
  work[15], flag[15];
 int pno, rno, i, j, prc, count, t, total;
 count = 0;
 //clrscr();
 printf ("\n Enter number of process:");
 scanf ("%d", &pno);
 printf ("\n Enter number of resources:");
 scanf ("%d", &rno);
 for (i = 1; i \le pno; i++)
   flag[i] = 0;
 printf ("\n Enter total numbers of each resources:");
 for (i = 1; i \le rno; i++)
  scanf ("%d", &tres[i]);
 printf ("\n Enter Max resources for each process:");
 for (i = 1; i \le pno; i++)
  {
```

```
printf ("\n for process %d:", i);
   for (j = 1; j \le rno; j++)
     scanf ("%d", &max[i][j]);
  }
 printf ("\n Enter allocated resources for each process:");
 for (i = 1; i \le pno; i++)
   printf ("\n for process %d:", i);
   for (j = 1; j \le rno; j++)
     scanf ("%d", &allocated[i][j]);
printf ("\n available resources:\n");
 for (j = 1; j \le rno; j++)
   avail[j] = 0;
   total = 0;
   for (i = 1; i \le pno; i++)
      total += allocated[i][j];
    avail[j] = tres[j] - total;
   work[j] = avail[j];
   printf (" %d \t", work[j]);
   }
 do
  {
    for (i = 1; i \le pno; i++)
      for (j = 1; j \le rno; j++)
         need[i][j] = max[i][j] - allocated[i][j];
     }
```

```
printf ("\n Allocated matrix
                                            need");
                                   Max
    for (i = 1; i \le pno; i++)
     {
      printf ("\n");
      for (j = 1; j \le rno; j++)
         printf ("%4d", allocated[i][j]);
      printf ("|");
      for (j = 1; j \le rno; j++)
         printf ("%4d", max[i][j]);
      printf ("|");
      for (j = 1; j \le rno; j++)
         printf ("%4d", need[i][j]);
     }
    prc = 0;
    for (i = 1; i \le pno; i++)
      if (flag[i] == 0)
         prc = i;
         for (j = 1; j \le rno; j++)
            if (work[j] < need[i][j])
              prc = 0;
              break;
if (prc != 0)
        break;
```

```
if (prc != 0)
      printf ("\n Process %d completed", i);
      count++;
      printf ("\n Available matrix:");
      for (j = 1; j \le rno; j++)
        work[j] += allocated[prc][j];
        allocated[prc][j] = 0;
        \max[prc][j] = 0;
        flag[prc] = 1;
        printf (" %d", work[j]);
 while (count != pno && prc != 0);
 if (count == pno)
  printf ("\nThe system is in a safe state!!");
  printf ("\nThe system is in an unsafe state!!");
return 0;
}
```

Output:

```
[188rla0501@localhost ~]$ vi dp.c
[188rla0501@localhost ~]$ gcc dp.c
[188r1a0501@localhost ~]$ ./a.out
 Enter number of process:5
 Enter number of resources:3
 Enter total numbers of each resources:10 5
 Enter Max resources for each process:
 for process 1:7
 for process 2:3
 for process 3:9
 for process 4:2
 for process 5:4
 Enter allocated resources for each process:
 for process 1:0
 for process 2:2
 for process 3:3
 for process 4:2
 for process 5:0
```

```
available resources:
Allocated matrix
                        Max
                                  need
  0
          01
                        31
          0 |
                        2|
                             1
          2|
                        2|
  2
          11
                        2|
                                 1
                                     1
          2|
               4
                    3
                        31
                             4
Process 2 completed
Available matrix:
                     5
                             2
Allocated matrix
                        Max
                                  need
  0
      1
          0 [
                        3|
                                 4
                                     0
          0 [
                        0 [
          21
                        21
                             6
  2
          11
               2
                    2
                        21
                             4
          21
               4
                        3|
Process 4 completed
Available matrix:
                         4
                        Max
Allocated matrix
                                  need
                        3 I
          01
                                 4
  0
          01
                        0 |
          21
                        21
                             6
          01
                        01
          2|
               4
                    3
                        3|
                             4
Process 1 completed
Available matrix:
Allocated matrix
                        Max
                                  need
  0
          01
                        0 [
                                 0
          0 I
                        0 [
  3
                             6
          21
                        21
          0 I
                        01
                             4
          21
               4
                        3|
Process 3 completed
Available matrix:
                        10
                             5
                                  5
Allocated matrix
                            Max
                                       need
  0
       0
            0 |
                  0
                            0 [
                                      0
                                           0
  0
       0
            01
                                           0
                  0
                       0
                           0 [
                                 0
                                      0
                       0
                                           0
       0
            01
                  0
                           01
                                 0
                                      0
                                      0
                                           0
            01
                  0
                       0
                           01
                       3
                           3|
            2|
                  4
                                 4
                                      3
                                           1
Process 5 completed
Available matrix:
                        10
```

b) Aim

Write a C program to simulate Bankers Algorithm for Deadlock Prevention

Algorithm:

- 1. Start
- 2. Attacking Mutex condition: never grant exclusive access. but this may not be possible for several resources.
- 3. Attacking preemption: not something you want to do.
- 4. Attacking hold and wait condition: make a process hold at the most 1 resource at a time.make all the requests at the beginning. All or nothing policy. If you feel, retry. eg. 2-phase locking 34
- 5. Attacking circular wait: Order all the resources. Make sure that the requests are issued in the correct order so that there are no cycles present in the resource graph. Resources numbered 1 ... n. Resources can be requested only in increasing order. ie. you cannot request a resource whose no is less than any you may be holding.
- 6. Stop

Program:

```
#include<stdio.h>
int \max[10][10], alloc [10][10], need [10][10], avail [10], i, i, p, r, finish [10] = \{0\}, flag=0;
main()
{
printf("\n SIMULATION OF DEADLOCK PREVENTION \n ");
printf("Enter no. of processes, resources\n");
scanf("%d%d",&p,&r);
printf("Enter allocation matrix");
for(i=0;i< p;i++)
for(j=0;j< r;j++)
scanf("%d",&alloc[i][j]);
printf("\n enter max matrix");
for(i=0;i<p;i++) /*reading the maximum matrix and availale matrix*/
for(j=0;j<r;j++)
scanf("%d",&max[i][j]);
printf(" \n enter available matrix");
for(i=0;i<r;i++)
scanf("%d",&avail[i]);
```

```
for(i=0;i< p;i++)
for(j=0;j<r;j++)
need[i][j]=max[i][j]-alloc[i][j];
fun(); /*calling function*/
if(flag==0)
{if(finish[i]!=1)
printf("\n Failing :Mutual exclusion");
for(j=0;j<r;j++)
{ /*checking for mutual exclusion*/
if(avail[j]<need[i][j])</pre>
avail[j]=need[i][j];
}fun();
printf("\n By allocating required resources to process %d dead lock is prevented ",i);
printf("\n lack of preemption");
for(j=0;j<r;j++)
if(avail[j]<need[i][j])</pre>
avail[j]=need[i][j];
alloc[i][j]=0;
}
fun();
printf("\n dead lock is prevented by allocating needed resources");
printf(" \n failing:Hold and Wait condition ");
for(j=0;j<r;j++)
{ /*checking hold and wait condition*/
if(avail[j]<need[i][j])</pre>
avail[j]=need[i][j];
fun();
printf("\n AVOIDING ANY ONE OF THE CONDITION, U CAN PREVENT DEADLOCK");
}
}
fun()
while(1)
for(flag=0,i=0;i< p;i++)
                                                                                                  32
```

```
{
    if(finish[i]==0)
    {
        for(j=0;j<r;j++)
        {
            if(need[i][j]<=avail[j])
            continue;
            else
            break;
        }
        if(j==r)
        {
            for(j=0;j<r;j++)
            avail[j]+=alloc[i][j];
        flag=1;
        finish[i]=1;
        }
    }
}</pre>
```

Output:

```
[188r1a0501@localhost ~]$ vi dpl.c
[188r1a0501@localhost ~]$ gcc dpl.c
[188r1a0501@localhost ~]$ ./a.out

SIMULATION OF DEADLOCK PREVENTION
Enter no. of processes, resources
3
2
Enter allocation matrix4
5
3
4
5
6
1
enter max matrix4
3
4
5
6
1
enter available matrix2
5
Failing :Mutual exclusion
By allocating required resources to process 3 dead lock is prevented lack of preemption
dead lock is prevented by allocating needed resources
failing:Hold and Wait condition
AVOIDING ANY ONE OF THE CONDITION, U CAN PREVENT DEADLOCK[188r1a0501@localhost ~]$
```

WEEK-4

Write a C program to implement the Producer – Consumer problem using semaphores using UNIX/LINUX system calls.

Aim:

Write a C program to implement the Producer – Consumer problem using semaphores using UNIX/LINUX system calls.

Algorithm:

- 1. The Semaphore mutex, full & empty are initialized.
- 2. In the case of producer process
- 3. Produce an item in to temporary variable.

 If there is empty space in the buffer check the mutex value for enter into the critical section.

 If the mutex value is 0, allow the producer to add value in the temporary variable to the buffer.
- 4. In the case of consumer process
 - i) It should wait if the buffer is empty
 - ii) If there is any item in the buffer check for mutex value, if the mutex==0, remove item from buffer
 - iii) Signal the mutex value and reduce the empty value by 1.
 - iv) Consume the item.
- 5. Print the result

Program:

```
#include<stdio.h>
#include<stdlib.h>
int mutex = 1, full = 0, empty = 3, x = 0;
int main ()
{
 int n;
 void producer ();
 void consumer();
 int wait (int);
 int signal (int);
 printf ("\n1.Producer\n2.Consumer\n3.Exit");
 while (1)
  {
   printf ("\nEnter your choice:");
   scanf ("%d", &n);
    switch (n)
       {
       case 1:
        if ((mutex == 1) && (empty != 0))
          producer ();
        else
          printf ("Buffer is full!!");
        break;
       case 2:
        if ((mutex == 1) && (full != 0))
          consumer ();
          printf ("Buffer is empty!!");
        break;
       case 3:
        exit (0);
        break;
       }
  }
 return 0;
```

36

```
}
int wait (int s)
 return (--s);
int signal (int s)
 return (++s);
void producer ()
 mutex = wait (mutex);
 full = signal (full);
 empty = wait (empty);
 x++;
 printf ("\nProducer produces the item %d", x);
 mutex = signal (mutex);
void consumer ()
 mutex = wait (mutex);
 full = wait (full);
 empty = signal (empty);
 printf ("\nConsumer consumes item %d", x);
 x--;
 mutex = signal (mutex);
```

Output:

```
[188r1a0501@localhost ~]$ vi pc.c
[188r1a0501@localhost ~]$ gcc pc.c
[188r1a0501@localhost ~]$ ./a.out
1.Producer
2.Consumer
3.Exit
Enter your choice:1
Producer produces the item 1
Enter your choice:1
Producer produces the item 2
Enter your choice:1
Producer produces the item 3 Enter your choice:2
Consumer consumes item 3
Enter your choice:2
Consumer consumes item 2
Enter your choice:2
Consumer consumes item 1
Enter your choice:2
Buffer is empty!!
Enter your choice:3
[188r1a0501@localhost ~]$
```

Week: 5

Write C programs to illustrate the following IPC mechanisms

Aim: Write C programs to illustrate the following IPC mechanisms

ALGORITHM:

- 1. Start the program.
- 2. Declare the variables.
- 3. Read the choice.
- 4. Create a piping processing using IPC.
- 5. Assign the variable lengths6. "strcpy" the message lengths.
- 7. To join the operation using IPC.
- 8. Stop the program

Program: (PIPE PROCESSING)

```
#include <unistd.h>
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
#define MSG_LEN 64
int main()
int result;
int fd[2];
char message[MSG_LEN];
char recvd_msg[MSG_LE];
result = pipe (fd);
//Creating a pipe//fd[0] is for reading and fd[1] is for writing
if (result < 0)
{
perror("pipe ");
exit(1);
strncpy(message,"Linux World!! ",MSG_LEN); result=write(fd[1],message,strlen(message)); if (result
< 0){
perror("write"); exit(2);
strncpy(message,"Understanding ",MSG_LEN);
result=write(fd[1],message,strlen(message));
if(result < 0)
perror("write");
exit(2);
strncpy(message,"Concepts of ",MSG_LEN);
result=write(fd[1],message,strlen(message));
if (result <0)
perror("write");
exit(2);
}
strncpy(message,"Piping ", MSG_LEN);
```

```
result=write(fd[1],message,strlen(message));
if (result < 0)
{
    perror("write");
    exit(2);
}
result=read(fd[0],recvd_msg,MSG_LEN);
if (result < 0)
{
    perror("read");
    exit(3);
}

printf("%s\n",recvd_msg); return 0;
}</pre>
```

a) FIFO

Program:

```
#include <stdio.h>
#include <stdlib.h>
#include <sys/stat.h>
#include <unistd.h>
#include linux/stat.h>
#define FIFO_FILE
                       "MYFIFO"
int main(void)
    FILE *fp;
    char readbuf[80];
    /* Create the FIFO if it does not exist */
     umask(0);
     mknod(FIFO_FILE, S_IFIFO|0666, 0);
     while(1)
         fp = fopen(FIFO_FILE, "r");
         fgets(readbuf, 80, fp);
         printf("Received string: %s\n", readbuf);
```

40

```
fclose(fp);
}
return(0);
}
```

```
#include <stdio.h>
#include <stdlib.h>
#define FIFO_FILE
                       "MYFIFO"
int main(int argc, char *argv[])
    FILE *fp;
    if (argc!=2) {
         printf("USAGE: fifoclient [string]\n");
         exit(1);
     }
    if((fp = fopen(FIFO\_FILE, "w")) == NULL) {
         perror("fopen");
         exit(1);
    fputs(argv[1], fp);
    fclose(fp);
    return(0);
}
```

```
C Program for Message Queue (Writer Process)
#include <stdio.h>
#include <sys/ipc.h>
#include <sys/msg.h>
// structure for message queue
struct mesg_buffer {
  long msg_type;
  char msg_text[100];
} message;
int main()
{
  key_t key;
  int msgid;
 // ftok to generate unique key
  key = ftok("progfile", 65);
 // msgget creates a message queue
  // and returns identifier
  msgid = msgget(key, 0666 | IPC_CREAT);
  message.mesg\_type = 1;
  printf("Write Data : ");
  gets(message.mesg_text);
  // msgsnd to send message
  msgsnd(msgid, &message, sizeof(message), 0);
  // display the message
  printf("Data send is : %s \n", message.mesg_text);
  return 0;
C Program for Message Queue (Reader Process)
#include <stdio.h>
#include <sys/ipc.h>
#include <sys/msg.h>
// structure for message queue
struct mesg_buffer {
  long mesg_type;
```

```
char mesg_text[100];
} message;
int main()
  key_t key;
  int msgid;
  // ftok to generate unique key
key = ftok("progfile", 65);
// msgget creates a message queue and returns identifier
  msgid = msgget(key, 0666 | IPC_CREAT);
  // msgrcv to receive message
  msgrcv(msgid, &message, sizeof(message), 1, 0);
  // display the message
  printf("Data Received is: %s \n",
            message.mesg_text);
  // to destroy the message queue
  msgctl(msgid, IPC_RMID, NULL);
  return 0;
}
C Program for Message Queue (Reader Process)
#include <stdio.h>
#include <sys/ipc.h>
#include <sys/msg.h>
// structure for message queue
struct mesg_buffer {
  long mesg_type;
  char mesg_text[100];
} message;
int main()
{
  key_t key;
  int msgid;
```

```
// ftok to generate unique key
 key = ftok("progfile", 65);
 // msgget creates a message queue
 // and returns identifier
 msgid = msgget(key, 0666 | IPC_CREAT);
 // msgrcv to receive message
 msgrcv(msgid, &message, sizeof(message), 1, 0);
 // display the message
 printf("Data Received is: %s \n",
          message.mesg_text);
 // to destroy the message queue
 msgctl(msgid, IPC_RMID, NULL);
 return 0;
}
OUTPUT: Thus the Piping process using IPC program was executed and verified successfully
       [sree@localhost ~]$ cc pp.c
       [sree@localhost ~]$ ./a.out
       Enter string:1
       05
       er
       tingEnter 1 array elementz:1
       The string length=1
       Sum=0[sree@localhost ~]$ er
       bash: er: command not found
       [sree@localhost ~]$ ating1
       bash: ating1: command not found
       [sree@localhost ~]$ gedit pp.c
       [sree@localhost ~]$ cc pp.c
       [sree@localhost ~]$ ./a.out
       Linux World!!!
       [sree@localhost ~]$ gedit pp.c
       [sree@localhost ~]$ cc pp.c
       [sree@localhost ~]$ ./a.out
       Linux World!! Understanding Concepts of Piping ,
       [sree@localhost ~]$
```

Week: 6

Aim: Write C programs to simulate the following memory management techniques

a) Paging

AIM: To write a C program to implement memory management using paging technique.

ALGORITHM:

Step1: Start the program.

Step2: Read the base address, page size, number of pages and memory unit.

Step3: If the memory limit is less than the base address display the memory limit is less than limit.

Step4: Create the page table with the number of pages and page address.

Step5: Read the page number and displacement value.

Step6: If the page number and displacement value is valid, add the displacement value with the address corresponding to the page number and display the result.

Step7: Display the page is not found or displacement should be less than page size.

Step8: Stop the program.

Program:

```
#include<stdio.h>
#include<conio.h>
main()
{

int ms, ps, nop, np, rempages, i, j, x, y, pa, offset; int s[10], fno[10][20];
printf("\nEnter the memory size -- ");
scanf("%d",&ms);
printf("\nEnter the page size -- ");
scanf("%d",&ps);
nop = ms/ps;

printf("\nThe no. of pages available in memory are -- %d ",nop);
printf("\nEnter number of processes -- ");
scanf("%d",&np);
```

```
rempages = nop; for(i=1;i \le np;i++)
{
printf("\nEnter no. of pages required for p[%d]-- ",i);
scanf("%d",&s[i]);
if(s[i] >rempages)
{
printf("\nMemory is Full");
break;
rempages = rempages - s[i];
printf("\nEnter pagetable for p[%d] --- ",i);
for(j=0;j< s[i];j++)
scanf("%d",&fno[i][j]);
}
printf("\nEnter Logical Address to find Physical Address ");
printf("\nEnter process no. and pagenumber and offset -- ");
scanf("%d %d %d",&x,&y, &offset);
if(x>np \parallel y>=s[i] \parallel offset>=ps)
printf("\nInvalid Process or Page Number or offset");
else
pa=fno[x][y]*ps+offset;
printf("\nThe Physical Address is -- %d",pa);
}
getch();
}
```

OUTPUT:

```
Inter the memory size -- 1000

Inter the page size -- 200

The no. of pages available in memory are -- 5

Inter number of processes -- 2

Inter no. of pages required for p[1]-- 20

Memory is Full
Inter Logical Address to find Physical Address
Inter process no. and pagenumber and offset -- 1

Inter Physical Address is -- 5

Inter Physical Address is -- 5
```

b) Segmentation

Aim: To write a C program to implement memory management using segmentation

Algorithm:

Step1: Start the program.

Step2: Read the base address, number of segments, size of each segment, memory limit.

Step3: If memory address is less than the base address display "invalid memory limit".

Step4: Create the segment table with the segment number and segment address and display it.

Step5: Read the segment number and displacement.

Step6: If the segment number and displacement is valid compute the real address and display the same. Step7: Stop the program.

Program:

```
#include<stdio.h>
#include<conio.h>
struct list
{
int seg;
int base;
int limit:
struct list *next;
} *p;
void insert(struct list *q,int base,int limit,int seg)
if(p==NULL)
{
p=malloc(sizeof(Struct list));
p->limit=limit;
p->base=base;
p->seg=seg;
p->next=NULL;
else
while(q->next!=NULL)
Q=q->next;
Printf("yes")
}
q->next=malloc(sizeof(Struct list));
q->next ->limit=limit;
q->next->base=base;
q->next ->seg=seg;
```

```
q->next ->next=NULL;
int find(struct list *q,int seg)
while(q->seg!=seg)
q=q->next;
return q->limit;
int search(struct list *q,int seg)
while(q->seg!=seg)
q=q->next;
return q->base;
main()
{
p=NULL;
int seg,offset,limit,base,c,s,physical;
printf("Enter segment table/n");
printf("Enter -1 as segment value for termination\n");
do
printf("Enter segment number");
scanf("%d",&seg);
if(seg!=-1)
printf("Enter base value:");
scanf("%d",&base);
printf("Enter value for limit:");
scanf("%d",&limit);
insert(p,base,lmit,seg);
}
while(seg!=-1)
printf("Enter offset:");
```

```
scanf("%d",&offset);
printf("Enter bsegmentation
number:");scanf("%d",&seg);
c=find(p,s
eg);
s=search(
p,seg);
if(offset<c
)
physical=s+offset;
printf("Address in physical memory %d\n",physical);
else
printf("error");
OUTPUT:
Enter segment table
Enter -1 as segmentation value for
terminationEnter segment number:1
Enter base
value:2000 Enter
value for limit:100
Enter segment
number:2Enter
base value:2500
Enter value for
limit:100
Enter segmentation
number:-1Enter
offset:90
Enter segment number:2
Address in physical memory 2590
Enter segment table
Enter -1 as segmentation value for
terminationEnter segment number:1
Enter base
value:2000 Enter
```

value for limit:100

```
Enter segment
number:2Enter
base value:2500
Enter value for
limit:100
Enter segmentation
number:-1Enter
offset:90
Enter segment number:1
Address in physical
memory 20
```

7. Write C programs to simulate Page replacement policies a) FCFS b) LRU c) Optimal

AIM: To Simulate FIRST IN FIRST OUT Page Replacement Algorithm

Theory:

a) FIFO (First in First Out) algorithm: FIFO is the simplest page replacement algorithm, the idea behind this is, "Replace a page that page is oldest page of main memory" or "Replace the page that has been in memory longest". FIFO focuses on the length of time a page has been in the memory rather than how much the page is being used.

PROGRAM:

```
#include<stdio.h>
#include<conio.h>
int i,j,nof,nor,flag=0,ref[50],frm[50],pf=0,victim=-1;
void main()
clrscr();
printf("\n \t\t\t FIFI PAGE REPLACEMENT ALGORITHM");
printf("\n Enter no.of frames....");
scanf("%d",&nof);
printf("Enter number of Pages.\n");
scanf("%d",&nor);
printf("\n Enter the Page No...");
for(i=0;i< nor;i++)
scanf("%d",&ref[i]);
printf("\nThe given Pages are:");
for(i=0;i< nor;i++)
printf("%4d",ref[i]);
for(i=1;i \le nof;i++)
frm[i]=-1;
printf("\n");
for(i=0;i< nor;i++)
flag=0;
```

```
printf("\n\t page no %d->\t",ref[i]);
for(j=0;j< nof;j++)
if(frm[j]==ref[i])
flag=1;
break;
}}
if(flag==0)
pf++;
victim++;
victim=victim%nof;
frm[victim]=ref[i];
for(j=0;j< nof;j++)
printf("%4d",frm[j]);
}
printf("\n\n\t\t No.of pages faults...%d",pf);
getch();
```

OUTPUT:

VIVA QUESTIONS:

- 1. Define FIFO?
- 2. Define page?
- 3. Define Frame?
- 4. Write advantages and disadvantages of FIFO?

AIM: To Simulate LEAST RECENTLY USED Page Replacement Algorithm

Theory:

b) LRU (Least Recently Used): the criteria of this algorithm is "Replace a page that has been used for the longest period of time". This strategy is the page replacement algorithm looking backward in time, rather than forward.

PROGRAM:

#include<stdio.h>
#include<conio.h>

```
main()
int i, j, k, min, rs[25], m[10], count[10], flag[25], n, f, pf=0, next=1;
clrscr();
printf("Enter the length of reference string -- ");
scanf("%d",&n);
printf("Enter the reference string -- ");
for(i=0;i<n;i++)
scanf("%d",&rs[i]);
flag[i]=0;
printf("Enter the number of frames -- ");
scanf("%d",&f);
for(i=0;i<f;i++)
count[i]=0;
m[i]=-1;
printf("\nThe Page Replacement process is -- \n");
for(i=0;i<n;i++)
for(j=0;j< f;j++)
if(m[j]==rs[i])
flag[i]=1;
38
count[j]=next;
next++;
if(flag[i]==0)
if(i < f)
m[i]=rs[i];
count[i]=next;
next++;
}
else
min=0;
for(j=1;j< f;j++)
if(count[min] > count[j])
min=j;
m[min]=rs[i];
count[min]=next;
next++;
}
pf++;
for(j=0;j< f;j++)
printf("%d\t", m[j]);
if(flag[i]==0)
```

```
printf("PF No. -- %d" , pf);
printf("\n");
}
printf("\nThe number of page faults using LRU are %d",pf);
getch();
}
```

OUTPUT

```
exam1@localhost:~
                                                                    [exam1@localhost ~]$ vi lru.c
[exam1@localhost ~]$ cc lru.c
[exam1@localhost ~]$ ./a.out
Enter the length of reference string -- 8
Enter the reference string -- 2
Enter the number of frames -- 4
The Page Replacement process is --
                    -1 PF No. -- 1
2
                             PF No. -- 2
                              PF No.
                      5
                              PF No.
       3
                              PF No.
0
                              PF No.
                       5
o
               2
                       5
                               PF No. --
                               PF No.
The number of page faults using LRU are 8[exam1@localhost ~]$
```

AIM: To Simulate LEAST FREQUENTLY USED Page Replacement Algorithm

Theory:

c) OPTIMAL Page Replacement:

Optimal page replacement algorithm says that if page fault occurs then that page should be removed that will not be used for maximum time in future.

It is also known as clairvoyant replacement algorithm or Bélády's optimal page replacement policy.

PROGRAM:

```
#include<stdio.h>
int main()
{
   int no_of_frames, no_of_pages, frames[10], pages[30], temp[10], flag1, flag2, flag3, i, j, k, pos,
max, faults = 0;
   printf("Enter number of frames: ");
   scanf("%d", &no_of_frames);
```

```
printf("Enter number of pages: ");
scanf("%d", &no_of_pages);
printf("Enter page reference string: ");
for(i = 0; i < no\_of\_pages; ++i){
  scanf("%d", &pages[i]);
}
for(i = 0; i < no\_of\_frames; ++i){
  frames[i] = -1;
}
for(i = 0; i < no\_of\_pages; ++i){
  flag1 = flag2 = 0;
  for(j = 0; j < no\_of\_frames; ++j){
     if(frames[j] == pages[i]){
         flag1 = flag2 = 1;
         break;
       }
  }
  if(flag1 == 0){
     for(j = 0; j < no\_of\_frames; ++j){
       if(frames[j] == -1){
          faults++;
          frames[j] = pages[i];
          flag2 = 1;
          break;
     }
  }
  if(flag2 == 0)
     flag3 = 0;
     for(j = 0; j < no\_of\_frames; ++j){
       temp[i] = -1;
       for(k = i + 1; k < no\_of\_pages; ++k){
          if(frames[j] == pages[k]){
             temp[j] = k;
             break;
          }
        }
     for(j = 0; j < no\_of\_frames; ++j){
       if(temp[j] == -1)
          pos = j;
```

```
flag3 = 1;
          break;
       }
     }
     if(flag3 ==0){
       max = temp[0];
       pos = 0;
       for(j = 1; j < no\_of\_frames; ++j){
          if(temp[j] > max){
             max = temp[j];
             pos = j;
          }
       }
     }
     frames[pos] = pages[i];
     faults++;
  }
  printf("\n");
  for(j=0;\,j < no\_of\_frames;\, +\!+j)\{
     printf("%d\t", frames[j]);
  }
}
printf("\n\nTotal Page Faults = %d", faults);
return 0;
```