

Algorithm – Search and Sort (1)

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Review

- Class vs. Class object
- Method vs. Function
- Object-oriented programming
 - Encapsulation
 - Abstraction
 - Inheritance
 - Polymorphism
 - This is a CS course dedicated for OOP...

Why Search?

- Searching is a fundamental part of programming, especially in data science
- There are **massive** amount of data in the world and you want to find data you are interested
- You should find data that you want, **efficiently**

Linear Search – Algorithm

- Find if a target value exists in a list
- To do this, search from the first item to the last item sequentially (**linear search**)
- If the target value exists, return the index where the value first occurs
- Otherwise, return -1

Linear Search – Algorithm

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target = 4

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

Linear Search – Algorithm

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- If the target value exists, return the index where the value first occurs
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Linear Search – Algorithm

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target = 4




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- Find if a target value exists in a list
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- If the target value exists, return the index where the value first occurs
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target = 4



index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

Linear Search – Algorithm

- Find if a target value exists in a list
- To do this, search from the first item to the last item sequentially (**linear search**)
- **If the target value exists, return the index where the value first occurs**
- Otherwise, return -1

target = 4

*Found!
Return 6*

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

Linear Search – Algorithm

- Find if a target value exists in a list
- To do this, search from the first item to the last item sequentially (**linear search**)
- If the target value exists, return the index where the value **first occurs**
- Otherwise, return -1

target = 4

Found!
Return 6

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

Ignored...

Linear Search – Algorithm

- Find if a target value exists in a list
- To do this, search from the first item to the last item sequentially (**linear search**)
- If the target value exists, return the index where the value first occurs
- **Otherwise, return -1**

target = 1

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

Linear Search – Algorithm

- Find if a target value exists in a list
- To do this, search from the first item to the last item sequentially (**linear search**)
- If the target value exists, return the index where the value first occurs
- Otherwise, return -1

target = 1

	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	No target! Return -1
index	0	1	2	3	4	5	6	7	8	9	10	11	
values	5	-2	0	100	-6	7	4	9	-7	50	4	3	

Algorithm vs. Programming

Algorithm: A recipe for computers to follow (logical steps)

Program: An instruction set in programming languages for a computer to understand and put an algorithm to practice

There can be ***many different*** ways to implement (program)
a ***single*** algorithm!

Linear Search – Impl (1): While Loop

- `def linear_search_while(L: list, value: Any) -> int:`
- `i = 0`
- `while i < len(L) and L[i] != value:`
- `i = i + 1`
- `if i == len(L):`
- `return -1`
- `else:`
- `return i`

Linear Search – Impl (2): While Loop with Sentinel

- The while loop version needs to do ($i < \text{len}(L)$) every time
 - Because we need to know when the loop reaches the end of the list
- How can we remove this? – by using **sentinel** at the end of the list!



Linear Search – Impl (2): While Loop with Sentinel

- `def linear_search_sentinel(L: list, value: Any) -> int:`
- `L.append(value) # Add the sentinel`
- `i = 0`
- `while L[i] != value: # This condition is enough!`
- `i = i + 1`
- `L.pop() # Remove the sentinel`
- `if i == len(L):`
- `return -1`
- `else:`
- `return i`



Caveat

Some people do not like modifying the input list because it could be dangerous and possibly incur errors

Linear Search – Impl (3): For Loop

- `def linear_search_for(L: list, value: Any) -> int:`
- `for i in range(len(L)):`
- `if L[i] == value:`
- `return i`
- `return -1`
- Simple code, no complex conditions
- But some people dislike returning in the middle of a loop
- We have learnt three types of linear search, among which you can choose according to your taste 😊

Linear Search – Time Complexity

- How to measure time spent for an algorithm?
 - `import time`
 - `t_start = time.perf_counter()`
 - **<<Your Algorithm>>**
 - `t_end = time.perf_counter()`
 - `return (t_end - t_start) * 1000.0` # the unit becomes milliseconds

Linear Search – Time Complexity (10 M items)

- When the value is located at the end of the list, it takes more time (**linear increase**)
 - This is why the algorithm is called **linear** search!
- Built-in list.index is the **fastest**
 - Python program is notoriously slow since every line of code needs to pass through the Python **interpreter** at run time

Case	while	sentinel	for	list.index
First	0.01	0.01	0.01	0.01
Middle	1261	697	515	106
Last	2673	1394	1029	212

*What if the list is **sorted**?
Can we do anything better?*

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

*What if the list is **sorted**?
Can we do anything better?*

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Motivation

- Linear search does work for a sorted list, but does **NOT** take advantage of the fact that it is sorted

target = 4

*Return 5
(Takes similar)*



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Motivation

- Linear search does work for a sorted list, but does **NOT** take advantage of the fact that it is sorted

target = 100

*Return 11
Takes very long*

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search - Idea

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Algorithm

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

target = 9

start = 0
end = 11

$mid = (0+11)//2$
 $4 < target$



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Algorithm

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

target = 9

$start = 0$
 $end = 11$

$mid = (0+11)//2$
 $4 < target$



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Algorithm

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

target = 9

start = 6
end = 11

$mid = (6+11)//2$
 $7 < target$



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Algorithm

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

target = 9

start = 9
end = 11

$mid = (6+11)//2$
 $7 < target$



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Algorithm

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

target = 9

start = 9
end = 11

$mid = (9+11)//2$
 $50 \geq target$



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Algorithm

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

target = 9

start = 9
end = 9

$mid = (9+11)//2$
 $50 \geq target$

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Algorithm

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

target = 9

start = 9
end = 9

$mid = (9+9)//2$
 $9 \geq target$



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Algorithm

- Idea: Evaluate the **middle** of the sorted list and removes half of candidate entries
- Linear search: one evaluation removes **one** candidate entry
- Binary search: one evaluation removes **half** of candidate entries

target = 9

start = 9
end = 8

Cannot proceed further!

*If there is target in L,
L[start] must be the target!*



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

Binary Search – Code

- `def binary_search(L: list, v: Any) -> int:`
- `start, end = 0, len(L) - 1`
- `while start != end + 1:`
- `mid = (start+end) // 2`
- `if L[mid] < v:`
- `start = mid + 1`
- `else:`
- `end = mid - 1`
- `if start < len(L) and L[start] == v:`
- `return start`
- `else:`
- `return -1`

Binary Search – Time Complexity (10 M items)

- Linear search
 - Time delay is proportional to $\text{len}(L)$
- Binary search
 - Time delay is proportional to $\log_2^{\text{len}(L)}$
- A good example why **sorting** is useful!
 - But remember that sorting is **NOT** free either. It also takes non-negligible time...

Case	list.index	binary_search
First	0.007	0.02
Middle	105	0.02
Last	211	0.02 (WoW!)

Summary

- Linear search
 - Evaluate the **first** item and cut the **one** evaluated item
 - Time proportional to **len(L)**
 - Applicable to **any** list
- Binary search
 - Evaluate the **middle** item and cut the **half**
 - Time proportional to $\log_2^{\text{len}(L)}$
 - Applicable to a **sorted** list

Let's move onto sorting

Why Sorting?

- People often want to see numerous items sorted!
 - Midterm score, sports...
 - Dictionary
- Sorting helps searching
 - Binary search

투수 순위 | 타자 순위

순위	선수	타율	경기수	타수	안타	2루타	3루타	홈런	타점	득점	도루	볼넷	삼진	출루율	장타율	>
1	페르난데스 (두산)	0.369	92	374	138	21	0	16	76	72	0	37	30	0.431	0.553	
2	토마스 (KT)	0.354	89	359	127	28	1	31	84	76	0	32	89	0.409	0.696	
3	김현수 (LG)	0.350	93	369	129	26	1	20	80	67	0	37	32	0.407	0.588	
4	이정후 (키움)	0.349	95	373	130	36	4	14	74	63	8	35	29	0.404	0.579	
5	손아섭 (롯데)	0.346	86	327	113	25	0	6	57	64	3	43	38	0.418	0.477	
6	강진성 (NC)	0.339	74	251	85	17	0	12	54	40	5	11	28	0.371	0.550	
7	이명기 (NC)	0.337	84	297	100	14	2	1	34	52	8	26	49	0.391	0.407	
8	오재일 (두산)	0.336	75	289	97	20	0	11	55	40	0	29	56	0.395	0.519	
9	최형우 (KIA)	0.333	87	318	106	20	1	11	57	56	0	46	58	0.423	0.506	
10	박민우 (NC)	0.326	70	264	86	17	2	4	28	47	8	22	24	0.375	0.451	
11	나성범 (NC)	0.323	83	341	110	25	2	25	78	74	0	35	101	0.396	0.628	
12	배정대 (KT)	0.321	89	327	105	22	3	9	41	55	16	42	85	0.401	0.489	
13	김상수 (삼성)	0.320	73	259	83	17	1	3	25	48	8	39	36	0.422	0.429	
14	정훈 (롯데)	0.319	59	235	75	14	1	7	39	48	6	30	52	0.400	0.477	
15	조용호 (KT)	0.315	81	254	80	11	0	0	17	49	8	36	51	0.403	0.358	
16	구자욱 (삼성)	0.313	69	262	82	15	0	8	43	38	12	30	58	0.390	0.462	
17	박해민 (삼성)	0.310	81	294	91	12	3	7	33	49	15	17	48	0.349	0.442	
18	강백호 (KT)	0.309	74	285	88	17	1	15	51	53	3	34	57	0.383	0.533	
19	마창도 (롯데)	0.309	87	311	96	22	1	7	47	46	10	30	38	0.370	0.453	
20	정수빈 (두산)	0.305	89	311	95	11	5	2	35	54	9	28	37	0.365	0.392	

Then, how can we sort a list?

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

Selection Sort – Idea

- Find the minimum value of the unsorted list and swap it with the leftmost entry

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

1-st iteration

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

unsorted

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

1-st iteration

Minimum in [0:11]



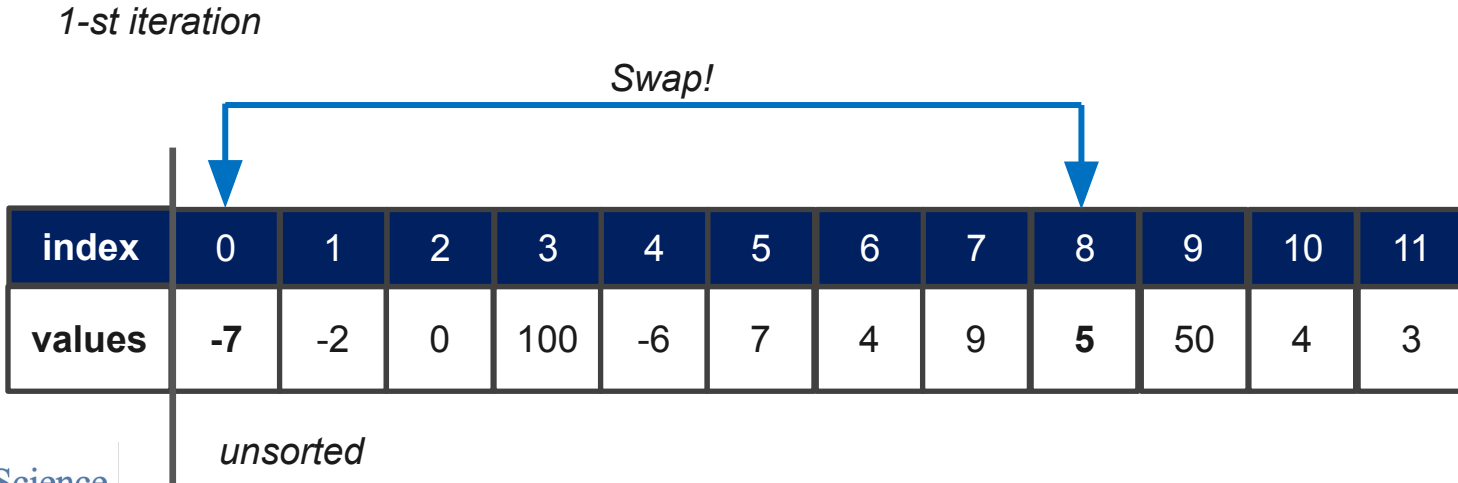
The diagram illustrates the first iteration of the Selection Sort algorithm. A horizontal array of 12 elements is shown, with indices 0 through 11. The values are 5, -2, 0, 100, -6, 7, 4, 9, -7, 50, 4, and 3. A vertical line is positioned between index 0 and index 1, separating the sorted portion (empty) from the unsorted portion. A blue arrow points down to the value -7 at index 8, which is the minimum in the unsorted range [0:11]. The word 'unsorted' is written below the array, starting from index 1.

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

unsorted

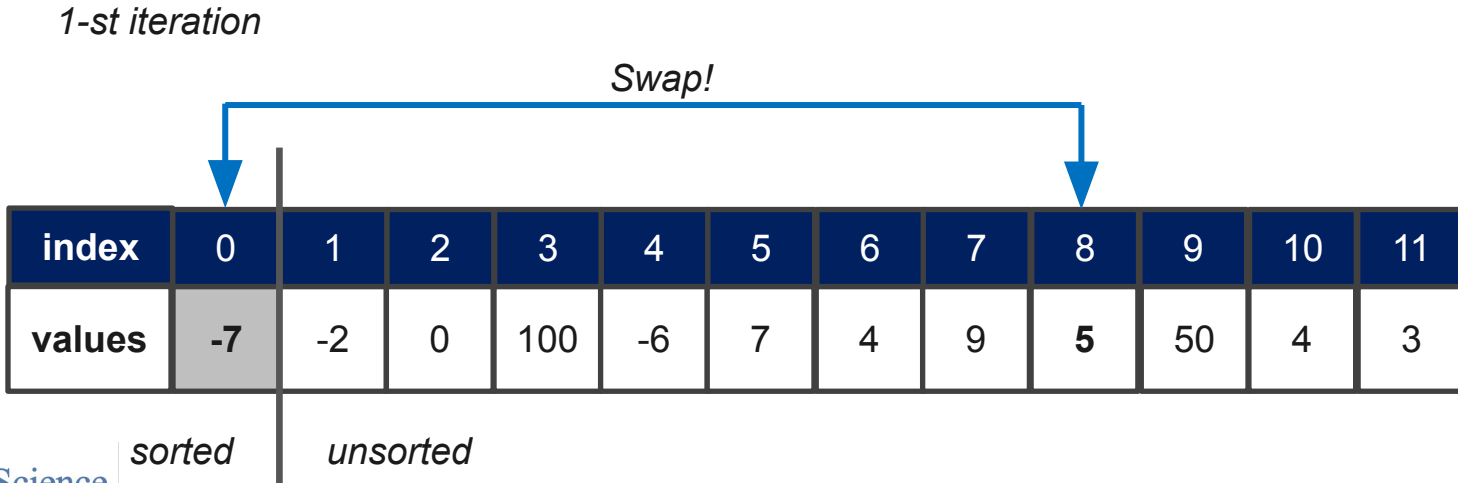
Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry



Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry



Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

2-nd iteration

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-2	0	100	-6	7	4	9	5	50	4	3


sorted *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

2-nd iteration

Minimum in [1:11]

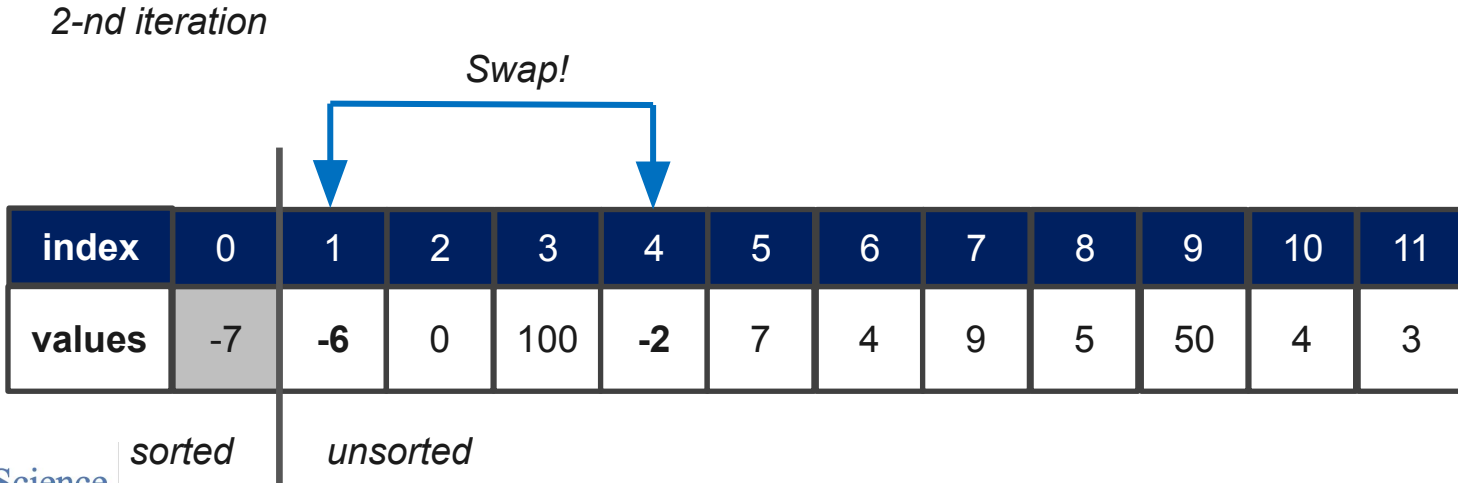


index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-2	0	100	-6	7	4	9	5	50	4	3

sorted *unsorted*

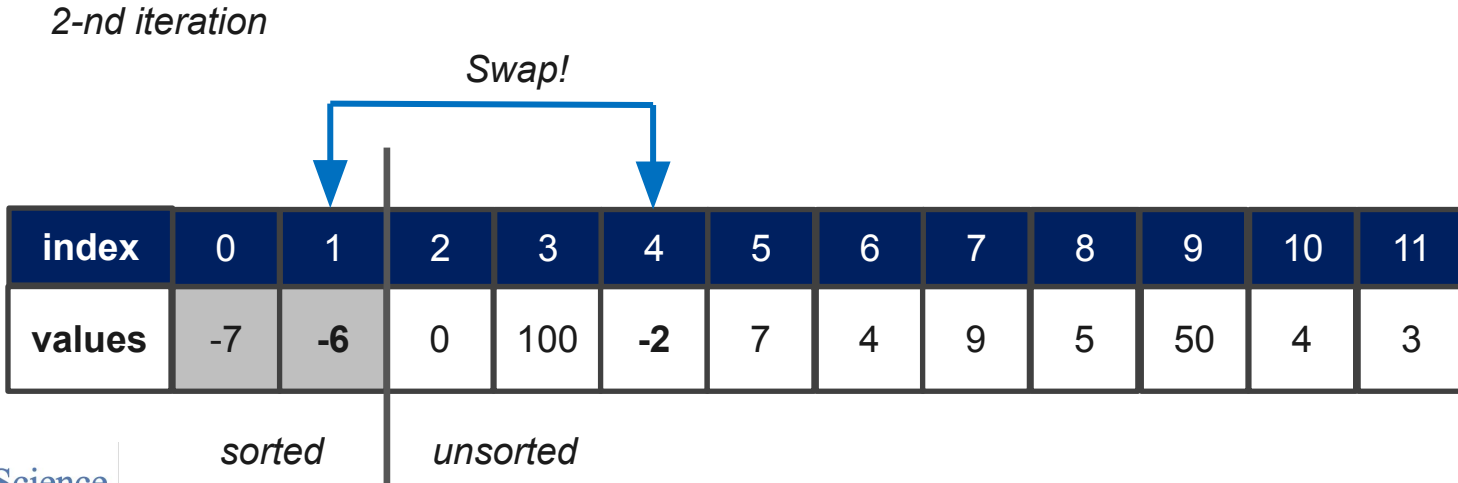
Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry



Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry



Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

3-rd iteration

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	0	100	-2	7	4	9	5	50	4	3


sorted *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

3-rd iteration

Minimum in [2:11]



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	0	100	-2	7	4	9	5	50	4	3


sorted | *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

3-rd iteration

Swap!



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	100	0	7	4	9	5	50	4	3


sorted *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

3-rd iteration

Swap!



The diagram shows a horizontal array of 12 cells. A vertical line is placed between the 3rd and 4th cells. A blue bracket with arrows at both ends spans from the 3rd cell to the 4th cell, with the text 'Swap!' centered above it. The first three cells (indices 0, 1, 2) are shaded gray and labeled 'sorted' below. The remaining nine cells (indices 3-11) are white and labeled 'unsorted' below.

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	100	0	7	4	9	5	50	4	3

sorted *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

4-th iteration

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	100	0	7	4	9	5	50	4	3

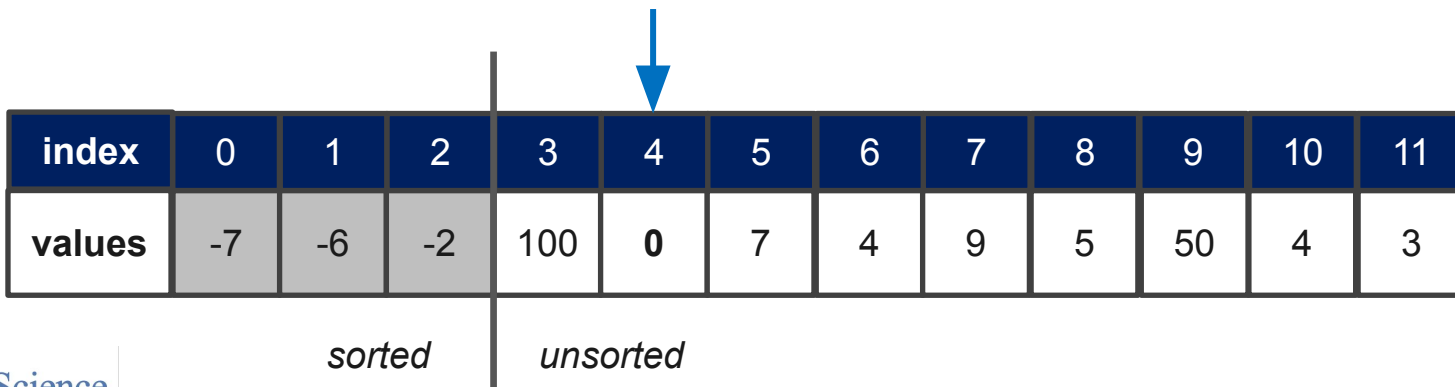
sorted | *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

4-th iteration

Minimum in [3:11]



The diagram illustrates the 4th iteration of the Selection Sort algorithm. A horizontal array of 12 elements is shown, indexed from 0 to 11. A vertical line separates the array into a 'sorted' region (indices 0-2) and an 'unsorted' region (indices 3-11). The values are: -7, -6, -2, 100, 0, 7, 4, 9, 5, 50, 4, 3. The minimum value in the unsorted region, 0 at index 4, is highlighted with a blue arrow pointing to it from the text 'Minimum in [3:11]' above. The sorted region is shaded gray, and the unsorted region is white.

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	100	0	7	4	9	5	50	4	3


sorted | unsorted

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

4-th iteration

Swap!



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	100	7	4	9	5	50	4	3


sorted *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

4-th iteration

Swap!



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	100	7	4	9	5	50	4	3

sorted | *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

5-th iteration

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	100	7	4	9	5	50	4	3


sorted | *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

5-th iteration

Minimum in [4:11]



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	100	7	4	9	5	50	4	3

sorted | *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

5-th iteration

Swap!

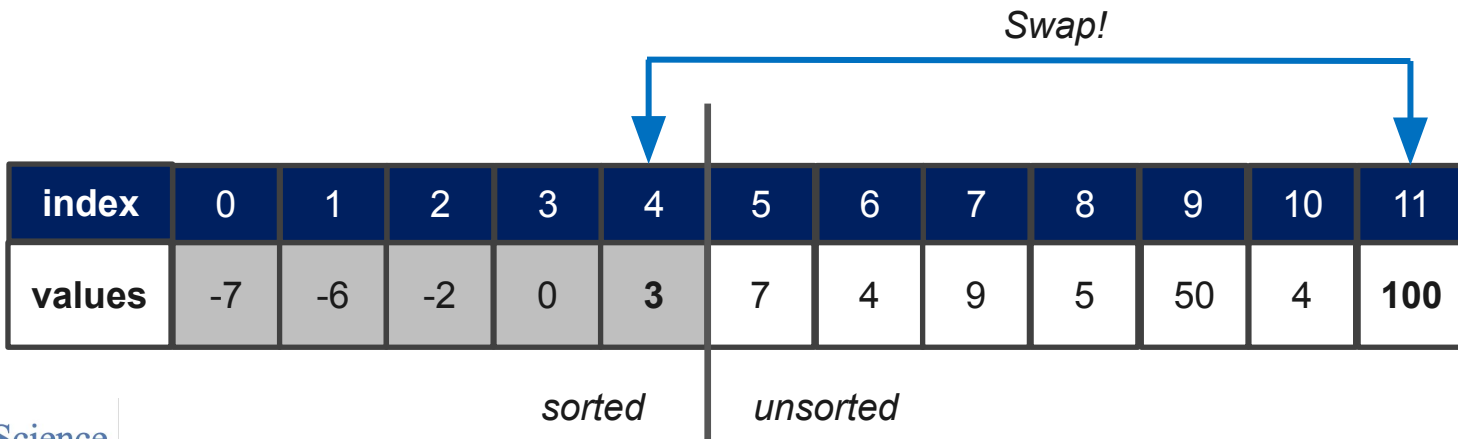
index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	7	4	9	5	50	4	100

sorted | *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

5-th iteration



The diagram illustrates the 5th iteration of the Selection Sort algorithm. A horizontal array of 12 elements is shown, with indices 0 to 11. A vertical line separates the array into a 'sorted' region (indices 0-4) and an 'unsorted' region (indices 5-11). The values in the sorted region are -7, -6, -2, 0, and 3. The values in the unsorted region are 7, 4, 9, 5, 50, 4, and 100. A blue arrow labeled 'Swap!' points from the value 100 at index 11 to the value 3 at index 4, indicating the swap of the minimum element in the unsorted region with the first element of the unsorted region.

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	7	4	9	5	50	4	100

sorted | *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

Repeat the procedure 12 times!

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	7	4	9	5	50	4	100

sorted | *unsorted*

Selection Sort – Algorithm

- Find the minimum value of the unsorted list and swap it with the leftmost entry

Repeat the procedure 12 times!

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

sorted

Selection Sort – Code

- `def selection_sort(L: list) -> None:`
- `for i in range(len(L)):`
- *# Find the index of the smallest item in L[i:]: **smallest***
- `L[i], L[smallest] = L[smallest], L[i] # swap`

Selection Sort – Code

- `def selection_sort(L: list) -> None:`
- `for i in range(len(L)):`
- `smallest = find_min(L, i)`
- `L[i], L[smallest] = L[smallest], L[i]` *# swap*

Selection Sort – Code

- `def find_min(L: list, start_idx: int) -> int:`
- `smallest = start_idx` *# (1) Initialize smallest*
- `for i in range(start_idx+1, len(L)):` *# (2) Update smallest*
- `if L[i] < L[smallest]:`
- `smallest = i`
- `return smallest` *# (3) Return the final value*

Selection Sort – Code (in one function)

- `def selection_sort(L: list) -> None:`
- `for i in range(len(L)):`
- `smallest = i`
- `for j in range(i+1, len(L)):`
- `if L[j] < L[smallest]:`
- `smallest = j`
- `L[i], L[smallest] = L[smallest], L[i] # swap`

Selection Sort – Time Complexity

- At i-th iteration, its inner loop (func **find_min**) needs to look up (N+1-i) items
 - When $N = \text{len}(L)$
- $N + (N-1) + (N-2) + \dots + 1 = N(N+1)/2$

Let's see another sorting algorithm called
Insertion Sort

Insertion Sort – Idea

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

1-st iteration



index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3

sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

1-st iteration

	<i>Insert location</i>		<i>Current target</i>									
												
index	0	1	2	3	4	5	6	7	8	9	10	11
values	5	-2	0	100	-6	7	4	9	-7	50	4	3
	<i>sorted</i>		<i>unsorted</i>									

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

1-st iteration

Insert

↓

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	5	0	100	-6	7	4	9	-7	50	4	3


sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

1-st iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	5	0	100	-6	7	4	9	-7	50	4	3

sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

2-nd iteration



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	5	0	100	-6	7	4	9	-7	50	4	3

sorted | *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

2-nd iteration


	<i>Insert location</i>		<i>Current target</i>									
												
index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	5	0	100	-6	7	4	9	-7	50	4	3
	<i>sorted</i>		<i>unsorted</i>									

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

2-nd iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	0	5	100	-6	7	4	9	-7	50	4	3


sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

2-nd iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	0	5	100	-6	7	4	9	-7	50	4	3

sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

3-rd iteration

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	0	5	100	-6	7	4	9	-7	50	4	3

sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

3-rd iteration

Insert location *Current target*

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	0	5	100	-6	7	4	9	-7	50	4	3


sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

3-rd iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	0	5	100	-6	7	4	9	-7	50	4	3


sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

3-rd iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	0	5	100	-6	7	4	9	-7	50	4	3

sorted | *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

4-th iteration



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-2	0	5	100	-6	7	4	9	-7	50	4	3

sorted | *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

4-th iteration


		<i>Insert location</i>				<i>Current target</i>							
													
index		0	1	2	3	4	5	6	7	8	9	10	11
values		-2	0	5	100	-6	7	4	9	-7	50	4	3
		<i>sorted</i>				<i>unsorted</i>							

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

4-th iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-6	-2	0	5	100	7	4	9	-7	50	4	3


sorted | *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

4-th iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-6	-2	0	5	100	7	4	9	-7	50	4	3

sorted | *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

5-th iteration



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-6	-2	0	5	100	7	4	9	-7	50	4	3

sorted | *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

5-th iteration


	<i>Insert location</i>				<i>Current target</i>							
												
index	0	1	2	3	4	5	6	7	8	9	10	11
values	-6	-2	0	5	100	7	4	9	-7	50	4	3
	<i>sorted</i>					<i>unsorted</i>						

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

5-th iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-6	-2	0	5	7	100	4	9	-7	50	4	3


sorted *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

5-th iteration

Insert



index	0	1	2	3	4	5	6	7	8	9	10	11
values	-6	-2	0	5	7	100	4	9	-7	50	4	3

sorted | *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

Repeat the procedure 11 times!

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-6	-2	0	5	7	100	4	9	-7	50	4	3

sorted | *unsorted*

Insertion Sort – Algorithm

- Insert the leftmost item of the unsorted list to the proper location of the sorted list

Repeat the procedure 11 times!

index	0	1	2	3	4	5	6	7	8	9	10	11
values	-7	-6	-2	0	3	4	4	5	7	9	50	100

sorted

Insertion Sort – Code

- `def insertion_sort(L: list) -> None:`
- `for i in range(1, len(L)):`
- *# insert L[i] to the proper location of L[:i]*

- `def insertion_sort(L: list) -> None:`
- `for i in range(1, len(L)):`
- `insert(L, i)`

Insertion Sort – Code

- `def insert(L: list, last_idx: int) -> None:`
- `for i in range(last_idx,0,-1):` *# (1) Go backwards*
- `if L[i-1] > L[i]:` *# (2) Check stopping condition*
- `L[i-1], L[i] = L[i], L[i-1]` *# (3) Swap*
- `else:`
- `break`

Insertion Sort – Code

- `def insertion_sort(L: list) -> None:`
- `for i in range(1, len(L)):`
- `for j in range(i,0,-1):` *# (1) Go backwards*
- `if L[j-1] > L[j]:` *# (2) Check stopping condition*
- `L[j-1], L[j] = L[j], L[j-1]` *# (3) Swap*
- `else:`
- `break`

Insertion Sort – Time Complexity

- At i-th iteration, its inner loop (**func insert**) needs to look up $(i+1)/2$ items and swap $i/2$ times on average
 - Look up: $1 + 1.5 + 2 + 2.5 + \dots + (N-1)/2 + N/2$ (When $N = \text{len}(L)$)
 - $= (1 + 2 + 3 + \dots + (N-1) + N)/2 - \frac{1}{2} = N(N+1)/4 - \frac{1}{2}$
 - Swap: $0.5 + 1 + 1.5 + \dots + (N-1)/2$
 - $= (1 + 2 + 3 + \dots + (N-1))/2 = (N-1)N/4$
- **A bit slower** than Selection sort
 - `find_min()` needs to look up the **whole** list
 - `Insert()` needs to look up only **half** on average but also need to swap!
- When a list is almost sorted, insertion sort needs to look up only **kN** items

*Yes, there are better sorting algorithms,
which you will see next time ^0^*

Thanks!