

1-		A logically interrelated collection of shared data which connected over a computer
	netw	ork.
	•	<u>Distributed Database</u>
	•	Distributed DBMS
	•	Data base
2-		
	the di	istribution transparent to users.
	•	Distributed Database
	•	<u>Distributed DBMS</u>
	•	Data base
3-	One	of the characteristics of DDBMS is a collection of logically related shared data.
	•	<u>DDBMS</u>
	•	DBMS
	•	Distributed Database
	•	Reliability
4-	The c	lata is split into a of fragments.
	•	Text
	•	collections
	•	indexs
	•	<u>Number</u>
5-	The c	data at each site is under the control of a
	•	DDB
	•	<u>DBMS</u>
	•	DBS
	•	DDBE
6-	Distr	ibuted Database types can be
	•	Homogenous
	•	Heterogeneous
	•	Both a&b
7-		an intelligent distribution of your data fragments to improve database performance
	and d	lata availability for end-users.
	•	Allocation
	•	Fragmentation
	•	Replication
	•	Reconstructing
8-		"It <u>divides</u> a database into various sub-tables and sub-relations so that data can be
	distri	buted and stored efficiently.
	•	Allocation
	•	<u>Fragmentation</u>
	•	Replication
	•	Reconstructing

- 9- The process of creating and maintaining multiple **copies** of data in different sites to improve database performance and data availability for end-users. Allocation Reconstruction Fragmentation **Replication** 10-A collection of databases scattered across multiple sites over a **network**. **Distributed Database Distributed DBMS** Data base Replication which are applied to relational database system to **partition** the relations 11among network sites. **Data fragmentation** Data allocation Data replication Reconstructing 12-...... in which each fragment is stored at the site with optimal distribution. Data fragmentation **Data allocation** Data replication Reconstructingwhich increases the **availability** and improves the performance of the system. Data fragmentation Data allocation **Data replication** 14-...... dividing a table by assigning each **row** or group of rows of relation to one or more Fragments. **Vertically Fragmentation Horizontal Fragmentation**
 - Mixed or hybrid Fragmentation
- 15- _____refers to the process of decomposing a table by **attributes** to columns.
 - Vertical Fragmentation
 - Horizontal Fragmentation
 - Mixed or hybrid Fragmentation
- 16- The combination of vertical fragmentation of a table followed by further
 - Horizontal fragmentation of some fragments
 - Vertically Fragmentation
 - Horizontal Fragmentation
 - Mixed or hybrid Fragmentation
 - None

- The choice of sites and the degree of replication depend on the
 Performance of the system
 Availability goals of the system
 The types and frequencies of transactions submitted at each site.
 - All of the above
- 18-refers to system live Time, that is system is running efficiently most of the time, the system functions even when failures occur, only delivering reduced performance until the issue is resolved
 - Reliability
 - Fragmentation transparency
 - Availability
 - Modular
- 19- relation can be specified by $\sigma Ci(R)$ operation in the relational algebra.
 - Horizontal fragmentation
 - Vertical fragmentation
 - both
 - None
- 20- Relation can be specified by $\pi Li(R)$ operation in the relational algebra.
 - Horizontal fragmentation
 - Vertical fragmentation
 - both
 - None
- 21- Any table or any fragment of the table can be accessed by the user in such a manner that they are locally stored in the site of the user by using ______.
 - Fragmentation transparency
 - <u>Location transparency</u>
 - Replication transparency
 - None
- 22- Any table can be accessed by the user in such a manner that it is not fragmented.
 - Fragmentation transparency
 - Location transparency
 - Replication transparency
 - None
- 23- Any table can be accessed by the user in such a manner that it is the original table.
 - Fragmentation transparency
 - Location transparency
 - Replication transparency
 - None
- 24- The starting point of bottom-up design is the individual.
 - <u>local conceptual schemas</u>
 - global conceptual schemas

25- T	The source databases are integrated and the integrated database is materialized. These are known as
•	database
•	<u>data warehouses</u>
•	Data fragmentation
26	allows data exchange between applications and perform similar transformation functions
•	LCS
•	<u>EAI</u>
•	ETL
27- T	The global conceptual (or mediated) schema is entirely virtual and not materialized
•	<u>True</u>
•	False
28- T	The global conceptual (or mediated) schema is entirely virtual and not materialized IN
•	Physical integration
•	<u>Logical integration</u>
29-	is the process of identifying that two objects are semantically related
•	Schema mapping
•	Schema matching
30	refers to the transformations between the objects.
•	Schema mapping
•	Schema matching
31-	That determines how to map the elements of each LCS to the other elements of the GCS.
•	Schema mapping
•	Schema matching
32- T	The integration is aided bytools that enable extraction of data from sources, their
trans	sformation to match the GCS.
•	Extract-Transform-Load (ETL)
•	Enterprise Application Integration (EAI)
•	Enterprise Information Integration (EII)
33	in a distributed database refers to the production of union-compatible views for similar
info	rmation expressed dissimilarly in different nodes
•	<u>Data integration</u>
•	Database
•	Data fragmentation
34	
•	Distributed Database
•	Distributed DBMS
•	<u>Distributed processing</u>
oper	rations in parallel, whenever possible, in order to improve performance
•	Parallel DBMS

Distributed processing

Parallel DDBMS

Modular Development Reliability **Lower Communication Cost** Better ResponseUsers access the distributed database via applications, which are classified as those that do not 37require data from other sites **local applications** global applicationsUsers access the distributed database via applications, which are classified as those that do 38require data from other sites local applications global applications The translation step is necessary only if the component databases areand local schemas are 39defined using different data models. **Heterogeneous** Homogenous The Vertical fragmentation should be in such a manner that we can rebuild a table from the fragment 40by taking the natural..... operation. **JOIN** Union Selection projection Each fragment must be assigned to a particular site in the distributed system This process is called Data allocation Data distribution <u>a&b</u> none A fragment containing all the rows of table T is called...... horizontal fragment 42-**Complete Primary Partial** IN Horizontal fragmentationoperation can be performed on the fragments to construct table **Union JOIN** Selection Projection

...........causes no interruptions in the functioning of distributed databases.

- 44- requires evaluation of a large number of query trees each of which produce the required results of a query.
 - <u>Distributed query optimization</u>
 - A distributed system
- 45- the operation is run at the site where the data is stored and not at the client site
 - Operation Shipping
 - Data Shipping
 - Hybrid Shipping
- 46-the data fragments are transferred to the database server, where the operations are executed
 - Operation Shipping
 - Data Shipping
 - Hybrid Shipping
- 47- combination of data and operation shipping
 - Operation Shipping
 - Data Shipping
 - Hybrid Shipping
- 48-data fragments are transferred to the high-speed processors, where the operation runs
 - Operation Shipping
 - Data Shipping
 - Hybrid Shipping
- 49-In query trading algorithm for distributed database systems, the controlling/client site for a distributed query is called
 - The buyer
 - Sellers
- 50-to qualify tuples that are to be joined. This reduces the amount of data transfer which in turn reduces communication cost.
 - <u>Use semi-join operation</u>
 - Perform selection and projection operations
 - Simplify operations on horizontal fragments
 - Merge the common leaves and sub-trees
- 51- Reduces the data flow over communication network
 - Perform selection and projection operations
 - Simplify operations on horizontal fragments
 - Use semi-join operation
 - Merge the common leaves and sub-trees
- 52- Eliminating selection conditions which are not relevant to a particular site.
 - Perform selection and projection operations
 - Simplify operations on horizontal fragments
 - Use semi-join operation
 - Merge the common leaves and sub-trees
- 53- Distributed transparency provides
 - Three dimensions
 - Two dimensions

54- SQL views are of similar fashion, where the user is unaware of the fact that a view of the table and not
the original table is viewed by the user.
• FRAGMENTATION TRANSPARENCY
REPLICATION TRANSPARENCY
• LOCATION TRANSPARENCY
55- Replicated copies facilitate the user in continuing with the queries in case of failure of a site, without
the knowledge of failure, which is known as
Concurrency transparency

- Failure transparency
- 56- The data is updated by the user, it is updated and gets reflected in all the tables of multiple sites, This is hidden from the user, which is known as
 - <u>Concurrency transparency</u>
 - Failure transparency
- 57- _____ database systems are systems with a <u>single</u> logical database located at <u>one</u> site under the control of a single DBMS.
 - Centralized
 - Distributed
 - Both
 - None
- 58- _____database systems are systems with a <u>single or multiple</u> logical databases located at <u>more</u> than one site under the control of a single DDBMS.
 - Centralized
 - <u>Distributed</u>
 - Both
 - None
- 59- DDBMS is a collection of logically ______ shared data.
 - Related
 - Unrelated
 - Far
 - None
- 60- The data is split into several_____.
 - Tracks
 - Parts
 - Patches
 - <u>Fragments</u>
- 61- Implies that a system can be <u>expanded</u> to new locations or units by adding new servers and data to the existing setup and connecting them to the distributed system <u>without</u> interruption.
 - Single Development
 - Reliability
 - Modular Development
 - Lower Communication Cost

62-	With the In a distributed database, the system functions even when failures occur,
01	nly delivering reduced performance until the issue is resolved.
	Single Development
	• Reliability
	Modular Development
	Lower Communication Cost
63-	Each DBMS participates in at least global application.
	• One
	• Two
	• Three
64-	A in a distributed database management system requires the transmission of data between the
co	omputers in a network
	• Query processing
	Distribution strategy
	• Query tree
65-	is a tree data structure representing a relational algebra expression
	 Query processing
	Distribution strategy
	• Query tree
66-	for a query is the ordering of data transmissions and local data processing in a database
sy	ystem
	 Query processing
	• <u>Distribution strategy</u>
	• Query tree
67-	In distributed query processing,operator has been used as an effective operator in reducing
re	lations referenced in the query to reduce the total amount of data transmission.
	• <u>The semi-join</u>
	• Projection
	Natural join
	• Selection
68-	each local DBE performs on the fragments that are stored at the local site, where the local CPU
ar	nd, and the disk input/output (I/O) time are the main drivers.
	• <u>The local level</u>
	• The global level
69-	Global optimization is greatly impacted by
	• The database distribution design
	Distributed database design
	Data distribution
	• None

The tables of the query are represented as 71-Relational operator **Leaf nodes** The root node The result table In a distributed system, the number of alternativesas we apply replication, horizontal fragmentation, and vertical fragmentation to the relations involved in the query **Increases drastically** Decreases drastically The QUERY TREE process continues for all internal nodes is executed and replaced by the result of the query. Relational operator Leaf nodes The root node The result table 89- of the query Is replaced by the result table Relational operator Leaf nodes The root node The nodeLocally storing data reduces communication costs for data manipulation in distributed databases Modular Development Reliability **Lower Communication Cost** Better Response Efficient data distribution in a distributed database system provides a faster response when 76-

Each internal nodes represents a between its child

Relational operator

Leaf nodes

user requests are met locally

Reliability

77-

Better Response

Costly Software

Large Overhead

All of the above

Data Integrity

Modular Development

Lower Communication Cost

The disadvantage of Distributed Database

The root node

The result table

70-

- 78-refers to how the database should be split and allocated to sites so as to optimize certain objective function
 Distributed database design
 Distributed database
 Distributed DBMS
 Centralized database
- 79-which determines how the data should be fragmented
 - Data fragmentation
 - Data allocation
 - Data replication
 - Location transparency
- 80- The small pieces of sub relations or sub tables are called......
 - Fragments
 - Units
 - Tracks
 - Parts
- 81- Replication is useful in improvingof data
 - Integrity
 - Reliability
 - Security
 - Availability
- 82- The most extreme case is replication of the whole database at every site in the distributed system, thus creating a distributed database.
 - Partially replicated
 - <u>Fully replicated</u>
 - Two are true
 - None
- 83- A distributed database is a network of identical database stored on multiple sites
 - Homogeneous
 - Heterogeneous
 - None
- 84- A..... distributed database use different schemas, operating systems, DDBMS and different data models
 - Homogeneous
 - <u>Heterogeneous</u>
 - None
- 85- All database in a collection are linked by a network and communicate with each other
 - Distributed query processing
 - Distributed transaction management
 - Network linking

- 86- database in a collection usually represent a single logical database and they are interconnected
 - Logical integration
 - Physical integration
 - <u>Seamless integration</u>
- 87- A has a number of database servers in the various sites to perform the operations due to a query
 - <u>Distributed system</u>
 - Distributed DBMS
 - Distributed database
- 88- Data Shipping also appropriate in systems where the communication costs are
 - High
 - Low
- 89- The sites where the local queries execute are called
 - Buyer
 - <u>Seller</u>
- 90- The formulates a number of alternatives for choosing sellers and for reconstructing the global results
 - <u>Buyer</u>
 - Seller
- 91- Optimal solution generally involves reduction od solution space so that the cost of query and data transfer.
 - Increase
 - Reduce
- 92- The system built around this system are
 - Shared Disk System
 - Shared nothing disk system
 - <u>Cluster</u>
 - Shared memory system
- 93- The advantage of shared Disk system:
 - Limited scalability
 - Fault tolerance
 - Both
 - None

DEC is: Digital Equipment Corporation

- 94- The DEC is application of
 - Shared memory system
 - Shared disk system
 - Shared nothing disk system
 - NUMA

- 95- Which of the following is application of shared nothing disk system
 - Tera data database machine
 - The grace and Gamma research prototypes
 - DEC
 - <u>A & B</u>
- 96- Which system is caused bandwidth problem
 - Shared memory system
 - Shared disk system
 - Shared nothing disk system
 - NUMA
- 97- which of the following is the advantage if the NUMA
 - Minimize memory bottleneck problem
 - Bandwidth problem
 - Higher cost of Architectures
- 98- In NUMA, if a processor of group (A) access Data from group (B) that mean the processor access data from:
 - The Local memory of Group (B)
 - Remote memory
 - The Local memory
 - <u>A & B</u>
- 99- Database type have the same structured scheme
 - Homogenous
 - Centralized data base
 - Heterogenous
 - Distributed database
- 100- In shared disk memory, If a processor or its memory fails. The other processors can complete the task.

That's called:

- Permanent fault
- Marginal fault
- Fault tolerance
- 101- A is a collection of actions that make consistent transformations of system states while preserving system consistency
 - <u>Transaction</u>
 - Lock
 - Dead Lock
- 102- The relative order of execution of the conflicting operations belonging to unaborted transactions in two histories are the same
 - Conflicting operations
 - Conflict equivalence
 - Result equivalent

- 103- Two incompatible operations (e.g., Read and Write) conflict if they both access the same data item
 - Conflicting operations
 - Conflict equivalence
 - Result equivalent
- 104- No violation of integrity constraints
 - ATOMICITY
 - CONSISTENCY
 - ISOLATION
 - DURABILITY
- 105- Transaction is all or nothing
 - ATOMICITY
 - CONSISTENCY
 - ISOLATION
 - DURABILITY
- 106- Committed updates persist
 - ATOMICITY
 - CONSISTENCY
 - ISOLATION
 - **DURABILITY**
- 107- The problem of synchronizing concurrent transactions such that the consistency of the database is maintained while, at the same time, maximum degree of concurrency is achieved
 - CONCURRENCY CONTROL
 - ISOLATION
 - Inconsistent retrievals
- 108- Transactions indicate their intentions by requesting locks from the scheduler
 - lock manager
 - Transaction manager
 - None
- 109- A Transaction one of whose operations is rejected by a scheduler is restarted by the transaction manager with a new timestamp
 - Basic TO Algorithm
 - Conservative TO Algorithm
 - Multi-version TO Algorithm
- 110-delay each operation until there is an assurance that no operation with a smaller timestamp can arrive at that scheduler. If this condition can be guaranteed, the scheduler will never reject an operation
 - Basic TO Algorithm
 - Conservative TO Algorithm
 - Multi-version TO Algorithm

- 111- methods guarantee that deadlocks cannot occur in the first place. Thus, the transaction manager checks a transaction when it is first initiated and does not permit it to proceed if it may cause a deadlock
 - <u>DEADLOCK PREVENTION</u>
 - DEADLOCK DETECTION
 - DEADLOCK AVOIDANCE
- 112- Make use of transaction timestamps to prioritize transactions and resolve deadlocks by aborting transactions with higher (or lower) priorities
 - DEADLOCK PREVENTION
 - DEADLOCK DETECTION
 - DEADLOCK AVOIDANCE
- 113- Deadlocks that are local to a single site would be detected at that site using the LWFG
 - Centralized Deadlock Detection
 - Hierarchical Deadlock Detection
 - DEADLOCK PREVENTION
- 114- Each lock manager transmits its LWFG to the deadlock detector, which then forms the GWFG and looks for cycles in it
 - <u>Centralized Deadlock Detection</u>
 - Hierarchical Deadlock Detection
 - DEADLOCK PREVENTION
- 115- The effects of some transactions are not reflected on the database
 - Inconsistent retrievals
 - Lost updates
 - Committed Read
- 116- A transaction, if it reads the same data item more than once, should always read the same value
 - <u>Inconsistent retrievals</u>
 - Lost updates
 - committed Read
- 117- TRANSACTIONS PROVIDEexecution in the presence of multiple user accesses
 - Atomic
 - Reliable
 - A & b
 - Correct
- 118- TRANSACTIONS PROVIDE execution in the presence of failures
 - Atomic
 - Reliable
 - <u>A & b</u>
 - Correct
- 119- 2PL schedulers are placed at each site. Each scheduler handles lock requests for data at that site
 - <u>DISTRIBUTED 2PL</u>
 - CENTRALIZED 2PL

- 120- Lock requests are issued to the central scheduler
 - DISTRIBUTED 2PL
 - <u>CENTRALIZED 2PL</u>
- 121- THE PURPOSES OF REPLICATION
 - System availability
 - Scalability
 - Performance
 - All of the above
- 122- allows for a way to support growth with acceptable response times
 - System availability
 - Performance
 - Scalability
 - Application requirements
- 123-remove single points of failure by replicating data, so that data items are accessible from multiple sites
 - System availability
 - Performance
 - Scalability
 - Application requirements
- 124- locate the data closer to their access points, thereby localizing most of the access that contributes to a reduction in response time
 - System availability
 - Performance
 - Scalability
 - Application requirements
- 125-dictated by the applications, which may wish to maintain multiple data copies as part of their operational specifications
 - System availability
 - Performance
 - Scalability
 - Application requirements
- 126-all copies of a data item have the same value at the end of the execution of an update transaction
 - Strong mutual consistency criteria
 - Weak mutual consistency criteria
- 127-if the update activity ceases for some time, the values eventually become identical
 - Strong mutual consistency criteria
 - Weak mutual consistency criteria

- 128-the probability that the system under consideration does not experience any failures in a given time interval
 - Reliability
 - Availability
 - ATOMICITY
- 129- the probability that the system is operational according to its specification at a given point in time
 - Reliability
 - Availability
 - ATOMICITY
- 130-the expected time between subsequent failures in a system with repair
 - MTTR
 - MTBF
 - MTTD
- 131-the expected time to repair a failed system.
 - MTTR
 - MTBF
 - MTTD
- 132- Maintain the atomicity and durability properties of local transactions
 - Local recovery manager
 - LOCAL RELIABILITY PROTOCOLS
 - LOCAL LEVEL
- 133- The fourth failures which is found in distributed database
 - Transaction failure
 - Site failures
 - <u>Communication Line failure</u>
 - Media
- 134- The deviation of a system from the behavior that is described in its specification.
 - Atomicity
 - Availability
 - <u>Failure</u>
- 135- the effects of the transaction on the distributed database is all-or-nothing
 - Atomic commitment
 - Recovery protocols
 - Termination protocols
- 136- Termination protocols for 2PC In Coordinator Timeouts in Wait State is
 - Global-commit
 - Global-abort
 - None

- 137- Dealing with site failure termination protocols for 2PC in Centralized Database Participant Timeouts in The READY state
 - Remain Blocked
 - Abort the transaction following a timeout vote-abort/ignore
 - None
- 138- Dealing with site failure termination protocols for 2PC in Distributed Database Participant Timeouts in if *Pj* is in the INTIAL state
 - It cannot help Pi to terminate the transaction.
 - Reply to Pi with a "vote-abort" message.
 - Send Pi either a "vote commit" or a "vote-abort" message
- 139- Dealing with site failure termination protocols for 2PC in Centralized Database Participant Timeouts in The INITIAL state.
 - Remain Blocked
 - Abort the transaction following a timeout vote-abort/ignore
 - None
- 140- Dealing with site failure termination protocols for 2PC in Distributed Database Participant Timeouts in if *Pj* is in the ABORT or COMMIT state.
 - t cannot help Pi to terminate the transaction.
 - Reply to Pi with a "vote-abort" message.
 - Send Pi either a "vote commit" or a "vote-abort" message
- 141- The coordinator fails while in the it will start the commit process upon recovery
 - INITIAL state
 - WAIT state
 - The COMMIT or ABORT states
- 142- The coordinator fails while in the the coordinator will <u>restart the commit</u> process for this transaction from the beginning by sending the "prepare" message one more time.
 - INITIAL state
 - WAIT state
 - the COMMIT or ABORT states
- 143- The coordinator fails while in the....... In this case, the coordinator will have informed the participants of its decision and terminated the transaction. Thus, upon recovery, it does not need to do anything if all the acknowledgments have been received. Otherwise, the termination protocol is involved.
 - INITIAL state
 - WAIT state
 - The COMMIT or ABORT states
- 144- Dealing with site failure termination protocols for 2PC in Distributed Database Participant Timeouts in if *Pj* is in the Ready state
 - <u>t cannot help Pi to terminate the transaction</u>.
 - Reply to Pi with a "vote-abort" message.
 - Send Pi either a "vote commit" or a "vote-abort" message

- 145- A participant fails in the.....abort the transaction
 - READY state
 - **INITIAL state**
 - ABORT or COMMIT state
- 146- A participant fails while in the the coordinator has been informed of the failed site's affirmative decision about the transaction before the failure.
 - **READY state**
 - INITIAL state
 - ABORT or COMMIT state
- 147- A participant fails while in state. the participant does not need to take any special action.
 - READY state
 - INITIAL state
 - ABORT or COMMIT state
- 148- Which of the following is disadvantages of data fragmentation:
 - Increase the efficiency of the database system
 - Access speed may be very high
 - Maintain the security and privacy
- 149- DEALING WITH SITE FAILURES THREE-PHASE COMMIT PROTOCOL Timeout in the

INITIAL state

- Abort the transaction
- The termination protocol proceeds by electing a new coordinator
- The participant has received the "prepare-to-commit" message and is awaiting the final "global-commit" message from the coordinator
- 150- DEALING WITH SITE FAILURES THREE-PHASE COMMIT PROTOCOL Timeout in the Ready state
 - Abort the transaction
 - The termination protocol proceeds by electing a new coordinator
 - The participant has received the "prepare-to-commit" message and is awaiting the final "global-commit" message from the coordinator
- 151- DEALING WITH SITE FAILURES THREE-PHASE COMMIT PROTOCOL Timeout in the

PRECOMMIT state

- Abort the transaction
- The termination protocol proceeds by electing a new coordinator
- The participant has received the "prepare-to-commit" message and is awaiting the final "global-commit" message from the coordinator
- 152- DEALING WITH SITE FAILURES THREE-PHASE COMMIT PROTOCOL If the new coordinator is in the PRECOMMIT state
 - Globally commit the transaction and send a "global-commit" message
 - All the participants will have moved into the ABORT state as well
- 153- DEALING WITH SITE FAILURES THREE-PHASE COMMIT PROTOCOL If the new coordinator is in the ABORT state
 - Globally commit the transaction and send a "global-commit" message
 - All the participants will have moved into the ABORT state as well

- 154- There is only one 2PL scheduler in the distributed system
 - <u>CENTRALIZED 2PL</u>
 - DISTRIBUTED 2PL
- 155- Transaction may read any of the replicated copies of item x, by obtaining a read lock on one of the copies of x
 - CENTRALIZED 2PL
 - <u>DISTRIBUTED 2PL</u>
- 156- When the databases already exist at a number of sites.
 - Top-down
 - Bottom-up
- 157- LOCAL RELIABILITY PROTOCOLS CHECK POINTING
 - Write a begin checkpoint record into the log.
 - Collect the checkpoint data into the stable storage.
 - Write an end checkpoint record into the log
 - All of the above
- 158-Transactions execute concurrently, but the net effect of the resulting history upon the database is equivalent to some serial history
 - <u>SERIALIZABLE HISTORY</u>
 - TRANSACTION
- 159- Read lock called..... and Write lock called......
 - Shared lock ,exclusive lock
 - Exclusive lock, shared lock
- 160- Finding an optimal or even a good solution to distributed data allocation is a
 - Complex optimization problem
 - Simple optimization problem
 - Clear optimization problem
- 161- A ______ is a DBMS running across multiple processors and disks that is designed to execute operations in parallel whenever possible, in order to improve performance
 - Series DBMS
 - Parallel DBMS
 - Both
 - None
- 162- _____ that a system can be expanded to new locations or units by adding new servers and data to the existing setup and connecting them to the distributed system **without interruption**.
 - Development Modular
 - Replication
 - Reliability
 - Availability

163- means that the system behaves as if there is a <u>single copy</u> of each data item – referred to as single system image

- replication transparency
- System availability
- MUTUAL CONSISTENCY

164-allows a query to see inconsistent data while replicas are being updated but requires that the replicas converge to a one-copy serializable state once the updates are propagated to all of the copies

- Epsilon serializability (ESR)
- Extract Transform load (ETL)

165-the replicas converging to the same value

- Transaction consistency
- Mutual consistency

166-the global execution history be serializable

- Transaction consistency
- Mutual consistency
- **167-** EAGER UPDATE PROPAGATION define <u>read-one/write-all (ROWA) protocols</u>
- 168-applying update on all the replicas at the same time
 - Synchronous propagation
 - Deferred propagation
- 169-the updates are applied to <u>one replica</u> when they are issued, but their application on the other replicas is batched and <u>deferred</u> to the end of the transaction.
 - Synchronous propagation
 - Deferred propagation

170-the transaction does not wait until its updates are applied to all the copies before it commits – it commits as soon as <u>one replica</u> is updated

- LAZY UPDATE PROPAGATION
- EAGER UPDATE PROPAGATION
- Deferred propagation
- 171-have lower response times for update transactions, since an update transaction can commit as soon as it has updated one copy
 - LAZY UPDATE PROPAGATION
 - EAGER UPDATE PROPAGATION
 - Deferred propagation
- 172- Updates are first applied at a master copy and then propagated to other copies which are called
 - Fragments
 - Units
 - Slaves
- 173- Central site can be overloaded and can become a bottleneck
 - DISTRIBUTED TECHNIQUE
 - <u>CENTRALIZED TECHNIQUES</u>
 - LAZY UPDATE

174-Apply the update on the local copy at the site where the update transaction originates, and then the updates are propagated to the other replica sites

- DISTRIBUTED TECHNIQUE
- CENTRALIZED TECHNIQUES
- LAZY UPDATE

175-provide the highest system availability

- **DISTRIBUTED TECHNIQUE**
- CENTRALIZED TECHNIQUES
- LAZY UPDATE

176-different replicas of a data item may be updated at different sites concurrently

- **DISTRIBUTED TECHNIQUE**
- CENTRALIZED TECHNIQUES
- LAZY UPDATE

177- Allow (both read and update) transactions to be <u>submitted</u> at any site and forwarding their operations to either the single master or to the appropriate primary master site

- EAGER CENTRALIZED PROTOCOLS
- LAZY CENTRALIZED PROTOCOLS

178- Query processing in DDBMS is different from query processing in centralized DBMS due to thisof data transfer over the network

- communication cost
- speed
- size

179- The type of optimization is done at....., where communication cost is a prominent factor

- the global level
- the local level

180- can act as the coordinating for both update and read-only transactions

- TM
- RM
- SC

181-the update transactions are not submitted to the master, but to the TM at the site where the application runs (since they don't need to know the master)

- EAGER CENTRALIZED PROTOCOLS
- LAZY CENTRALIZED PROTOCOLS

182- does not sacrifice database consistency, but only allows read-only transactions (queries) to read inconsistent data

- ESR
- EII
- ETL

183-checks a transaction when it is first initiated and does not permit it to proceed if it may cause a deadlock

- transaction manager
- lock manager

	DEADLOCK AVOIDANCE
	DEADLOCK AVOIDANCE
•	DEADLOCK PREVENTION
185-	property of the transaction indicates that either all the operations of the transaction are carried
	at or none of them are carried out.
	• Atomicity
	• Consistency
	 Durability
	Isolation
186-	The phase that the locks is acquired .
100	• Growing
	ShrinkingBoth
197	• None The phase that the locks is released
10/-	The phase that the locks is released .
	• Growing • Shrinking
	• Shrinking • Doth
	• Both
	• None

1- A Distriubted processing in **centralized** database can be accessed over a computer network. True False 2- The DBMS at each site can handle local applications **autonomously**. **True** False **3-** The sites in distributed databases are not linked by a communications network. True **False** 4- The DBMS at each site can not handle local applications autonomously. True **False** 5- A <u>description</u> of the replication of fragments is sometimes (**Replication schema**.) True <u>False</u> 6- Mixed (Hybrid) fragmentation this is achieved by SELECT & PROJECT operations which is represented by $\pi Li(\sigma Ci(R))$ **True** False 7- The replication of the databases is hidden from the user by replication transparency. **True** False 8- All the transparencies are to be maintained in distributed database system. True False

9- Local data storage is not possible in centralized databases.

system, thus creating a fully replicated distributed database.

Each DBMS participates in at least one global application

The most extreme case is replication of the whole database at every site in the distributed

The transmission cost is low when sites are connected through high speed Networks and is quite

The users are not allowed to view the internal details of the distribution

True

False

True

False

False

True

False

True

False

significant in other networks

10-

11-

12-

13-

- Almost all global optimization alternatives ignore the local processing time
 True
 False
- 15- The rows that belong to the vertical fragments are specified by a condition on one or more attributes of the relation
 - True
 - <u>False</u>
- Data shipping is also used in object oriented database as a preferred approach
 - <u>True</u>
 - False
- 17- The algorithm starts with the buyer assigning sub-queries to the seller sites
 - True
 - False
- 18- Data partitioning is not required in shared nothing system
 - True
 - False
- 19- NUMA is hybrid system of shared disk system and nothing shared disk system
 - True
 - Fales
- 20- NUMA Architectures is less than the others
 - True
 - False
- 21- DBMS must make the distribution transparent to users
 - True
 - False
- 22- Transaction (Ti) is assigned a globally unique timestamp ts(Ti)
 - True
 - False
- 23- Timestamps are not assigned to transactions at their initiation but at the beginning of their validation step
 - True
 - False
- A timeout occurs at a destination site when it can get an expected message from a source site within the expected time period
 - True
 - <u>False</u>
- 25- The termination protocols serve the timeouts for both the coordinator and the participant processes
 - <u>True</u>
 - False

- 26- Timeout in the COMMIT or ABORT states sends the global commit or global abort commands to the sites that have not yet responded, and waits for their acknowledgement
 - True
 - False
- 27- The local query optimization re not sufficient for some queries as the data is available locally.
 - True
 - False
- 28- Local applications do not require data from other sites
 - True
 - False
- 29- Global applications require data from other sites
 - <u>True</u>
 - False
- In case of a database failure in a centralized database, the system comes to a complete stop.
 - True
 - False
- 31- Local data storage is **possible** in centralized databases.
 - True
 - False
- When i use homogeneous, if i will integrate more then one sites, Translation step is required.
 - True
 - False
- 33- Just like a local query optimizer, a global query optimizer must evaluate a large number of equivalent query trees
 - True
 - False
- 34- Timestamps are associated with transactions and data items
 - <u>True</u>
 - False
- 35- In horizontal fragmentation, Tuples of the table are placed in one system and the rest are placed in other systems.
 - True
 - False
- 36- In vertical fragmentation, each site may not need all columns of a table.
 - True
 - False
- 37- The specific implementation of the Read and Write interfaces of the transaction monitor differ according to the specific replication protocol
 - True
 - False
- 38- local applications do not require data from other sites
 - True
 - False

- 39- Global applications require data from other sites
 - True
 - False
- 40- In case of a database failure in a centralized database, the system comes to a complete stop.
 - True
 - False
- 41- Local data storage is **possible** in centralized databases.
 - True
 - False
- 42- when i use homogeneous, if i will integrate more then one sites, Translation step is required.
 - True
 - False

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Homogeneous

A homogenous distributed database is a network of **identical databases** stored on multiple sites. The sites have the same operating system, DDBMS, and data structure, making them easily manageable.

Homogenous databases allow users to access data from each of the databases seamlessly. The following diagram shows an example of a homogeneous database:

Heterogeneous

A heterogeneous distributed database uses **different** schemas, operating systems, DDBMS, and different data models.

In the case of a heterogeneous distributed database, a particular site can be completely unaware of other sites causing limited cooperation in processing user requests. The limitation is why translations are required to establish communication between sites.

Concurrency control Definition

The goal of **concurrency control** is to allow multiple transactions to run concurrently (which is great for performance) while ensuring that data access is controlled such that the net effect is the same as if the transactions all ran in some **serial** order. That is, we cannot have the net result be one where transactions read an interim state of data from another transaction.

Concurrency control algorithms

There are two categories of concurrency control: pessimistic and optimistic. **Pessimistic concurrency control** assumes that there is a high likelihood that two or more transactions will try to access the same resources. A transaction will lock access to the resources it needs to keep others from modifying them. **Optimistic concurrency control** assumes that it is not likely that multiple transactions will access the same resource. Optimistic techniques will not lock access to resources. Instead, they will check for conflicts at commit time. If it turns out that multiple transactions did indeed access resources concurrently, one or more of them will have to abort.

Note:

- \circ No-fix/No-flush \Rightarrow redo/undo algorithm
- \circ No-fix/Flush \rightarrow undo/no-redo algorithm
- \circ Fix/No-flush \rightarrow redo/no-undo algorithm
- Fix/Flush → no-undo/no-redo

Distributed database advantages
1- Modular Development
2- Reliability
3- Lower communication cost
4- Better response