

Assignment-2
Cryptography Network Security

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Knapsack algorithm

It is given by Hellman

Asymmetric Key Cryptography

eg: weights = (1, 6, 8, 15 and 24)

In general knapsack, we select weights to achieve a sum. if we want sum = 30

We select 1, 6, 8 and 15

$$\begin{array}{r} \text{Let plaintext} = \begin{array}{ccccc} 1 & 0 & 0 & 1 & 1 \\ \times & \times & \times & \times & \times \\ (*) & 1 & 6 & 8 & 15 & 24 \\ \hline & 1 & +0 & +0 & +15 & +24 \\ \hline & \downarrow & & & & \\ & \Rightarrow 40 & & & & \end{array} \end{array}$$
$$\begin{array}{r} \begin{array}{ccccc} 1 & 1 & 0 & 1 & 0 \\ \times & \times & \times & \times & \times \\ & 1 & 6 & 8 & 15 & 24 \\ \hline & 1 & +6 & +0 & +15 & +0 \\ \hline & \downarrow & & & & \\ & \Rightarrow 22 & & & & \end{array} \end{array}$$

CT = plain text \times corresponding weights

$$\therefore CT = 40 \quad 22$$

* Key Generation: Asymmetric key

1) Public key (Hard knapsack)

2) Private key (easy knapsack)

↓

done first (i.e. we find private key first)

eg: {1, 2, 4, 9, 20, 40}

Weights are always in increasing order

1) First, find private key (Assume)

$$D = \{1, 2, 4, 10, 20, 40\} \rightarrow \text{private key}$$

Select 2 number "n" and "m"

Condition:

$m >$ Sum of all no.s in sequence

$$\text{sum} = 77$$

$$\boxed{m = 110} \quad \text{let}$$

2. $n = \text{select } \phi$ so that it has no common factor with m .
 \therefore let $n=31$

\Rightarrow List out factors of 110, 31 don't have common factor except one (take like that)

now, $(D_i \times n) \bmod m$ \forall elements in n .

$D_i \Rightarrow$ First number in list.

$(1 \times 31) \bmod 110 = 31 \Rightarrow \{1, 2, 4, 10, 20, 40\}$ initial

$(2 \times 31) \bmod 110 = 62 \Rightarrow \{31, 62, 14, 90, 70, 30\}$

$(4 \times 31) \bmod 110 = 14$

\downarrow
public key

$(10 \times 31) \bmod 110 = 90$

$(20 \times 31) \bmod 110 = 70$

$(40 \times 31) \bmod 110 = 30$

* Encryption:

Now Assume PT

let $PT = 100100 \mid 111100 \mid 101110$
 ① ② ③

divide PT into 6-6 parts (no. of elements in sequence = 6)

1st part $\Rightarrow 100100 = 1 \times 31 + 0 \times 62 + 0 \times 14 + 1 \times 90 + 0 \times 70 + 0 \times 30$
 $= 31 + 90 = 121 \Rightarrow CT$

2nd part $\Rightarrow 111100 \Rightarrow 31 + 62 + 14 + 90 + 0 + 0$
 $= 197 \rightarrow CT$

\Rightarrow no. of elements in sequence is 6 elements, so divide into 6 parts PT.

part $\Rightarrow 101110 \Rightarrow 31 + 0 + 14 + 90 + 70 + 0 \Rightarrow 205$
 i.e. multiply with public key with
 PT instead of private key.

$$\therefore CT = [121, 197, 205]$$

* Decryption:

① Calculate $s' = 31^{-1}$
 $31x \bmod 110 = 1$ then we get $x = 7$

\rightarrow until get 1 continue process

$(CT \times x) \bmod m$ from seq $D = \{1, 2, 4, 10, 20, 40\}$
 \downarrow private key

$$(121 \times 7) \bmod 110 = 11 \Rightarrow 100100 \quad (1 + 10 = 11)$$

$$(127 \times 7) \bmod 110 = 17 \Rightarrow 111100 \quad (1 + 2 + 4 + 10 = 17)$$

$$(205 \times 7) \bmod 110 = 35 \Rightarrow 101110 \quad (1 + 4 + 10 + 20 = 35)$$

in encryption we taken public key

in decryption process we taken private key

$\Rightarrow 11$ in sequence add any two no's to get

$$\{1, 2, 4, 10, 20, 40\}$$

\uparrow
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$(0 + 1 = 11) \rightarrow$ here put 1 other case put zero.

$$17 \rightarrow \{1, 2, 4, 10, 20, 40\}$$

$$111100$$

we get same PT back 100100 111100 101110

\Rightarrow knapsack encryption algorithm is the first general public key cryptography algorithm.

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ISHKAA!!

It is developed by Ralph Merkle and
Martin Hellman in 1978.
⇒ we need two different keys
public key & private key
For easy knapsack we will choose
a super increasing knapsack problem.
It is a sequence in which every next
term is greater than the sum of
all preceding terms.