UNIT-1

EXPERIMENT-1

Aim: To study and perform the installation of the Linux (Ubuntu) operating system on a Windows machine using dual-boot or virtualization techniques, ensuring a functional environment for learning and practicing Linux-based software development and system administration.

Description: Linux is a free, open-source operating system based on Unix, created by Linus Torvalds in 1991. It powers everything from servers and supercomputers to smartphones and IoT devices.

Known for its stability, security, and flexibility, Linux is widely used in tech industries. Different distributions (like Ubuntu, Fedora, and Debian) make it adaptable for various needs

File and Directory Commands:

- □ pwd → Print working directory (shows current location).
- \Box ls \rightarrow List files and directories.
- \Box ls -1 \rightarrow Detailed list (permissions, size, owner).
- \square ls -a \rightarrow Show hidden files too.
- \Box cd <directory> \rightarrow Change directory.
- \Box cd .. \rightarrow Go back one level.
- \square mkdir <dirname> \rightarrow Create new directory.
- \square rmdir <dirname> \rightarrow Remove empty directory.
- \square rm -r <dirname> \rightarrow Remove directory and contents.
- \square exit \rightarrow Logout or close terminal.

Steps to Install Linux (Ubuntu) in Windows (via WSL):

- 1. Open PowerShell (Admin) and run:
- 2. wsl –install

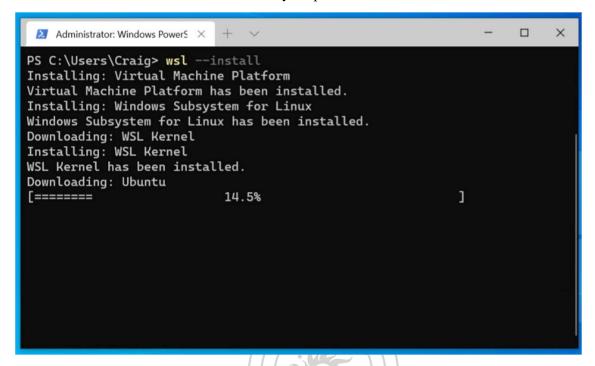
```
Microsoft Windows [Version 10.0.19043.1151]
(c) Microsoft Corporation. All rights reserved.

C: \wsl --install
Installing: Virtual Machine Platform
Virtual Machine Platform has been installed.
Installing: Windows Subsystem for Linux
Windows Subsystem for Linux has been installed.
Downloading: WSL Kernel
Installing: WSL Kernel
WSL Kernel has been installed.
Downloading: Ubuntu
The requested operation is successful. Changes will not be effective until the system is rebooted.

C:\>
```

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3. Wait for installation and select Ubuntu as your preferred distribution.



- 4. After setup, launch Ubuntu from the Start Menu.
- 5. Create a new Linux username and password.
- 6. Test installation by running basic commands in Ubuntu.

Some Basic Commands in Ubuntu:

1. **pwd**

- o Shows the current location in the filesystem.
- Example:
- o pwd
- o /home/student
- 2. man pwd
- o Opens the manual page for the pwd command.
- o Example:
- o man pwd
- 3. mkdir
- o Creates a new directory.
- o Example:
- o mkdir myfolder

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- 4. **ls**
- o Lists files and directories.
- o Example:
- o ls
- o Desktop Documents myfolder
- 5. cat > file
- o Creates a new text file and adds content.
- o Example:
- o cat > notes.txt
- o (type text, press Ctrl+D to save)
- 6. vi file
- o Opens file in *vi editor*.
- o Example:
- o vi notes.txt
- o Inside vi:
- Press $\mathbf{i} \rightarrow \text{insert mode}$.
- Press Esc \rightarrow then type :wq \rightarrow save and quit.



```
Try 'mkdir --help' for more information.
hi@Shannu44-laptop:~/s$ cd
hi@Shannu44-laptop:~$ ls
ni@Shannu44-laptop:~$ cat> x.txt
Welcome Linux^C
hi@Shannu44-laptop:~$ vi x.txt
[5]+ Stopped
                              vi x.txt
ni@Shannu44-laptop:~$ cat> y.txt
Hello
hi@Shannu44-laptop:~$ cp y.txt x.txt
hi@Shannu44-laptop:~$ vi y.txt
[6]+ Stopped
                              vi y.txt
 i@Shannu44-laptop:~$ vi x.txt
 i@Shannu44-laptop:~$ clear
```

```
Shannu44-laptop:~/s$ pwd
/home/hi/s
hi@Shannu44-laptop:~/s$ man pwd
[4]+ Stopped
                              man pwd
hi@Shannu44-laptop:~/s$ mkdir
mkdir: missing operand
Try 'mkdir --help' for more information.
hi@Shannu44-laptop:~/s$ cd
hi@Shannu44-laptop:~$ ls
hi@Shannu44-laptop:~$ cat> x.txt
Welcome Linux^C
hi@Shannu44-laptop:~$ vi x.txt
[5]+ Stopped
                              vi x.txt
hi@Shannu44-laptop:~$
```

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Conclusion: Linux (Ubuntu) was successfully installed on Windows using WSL. Basic commands for file and directory management were executed, demonstrating the ability to work with the Linux

terminal environment.



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EXPERIMENT-2

Aim:To demonstrate the creation of a child process using the fork() system call in Linux. **Description:** The fork() system call is used to create a new process by duplicating the calling process.

- ☐ The original process is the **parent**, and the newly created one is the **child**.
- Both processes execute the same code but have different process IDs.
- ☐ fork() returns:
- o $> 0 \rightarrow$ Parent process (child's PID is returned)
- o $0 \rightarrow \text{Child process}$
- o $< 0 \rightarrow$ Error in creating process

Algorithm:

- 1. Start program execution.
- 2. Call fork() to create a new process.
- 3. If fork() returns 0, it's the child process.
- 4. If fork() returns a positive value, it's the parent process.
- 5. Print process ID for both parent and child.
- 6. End program.

Program:

```
#include <stdio.h>
#include <unistd.h>

#include <unistd.h>

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int main() {
    int pid;
    pid = fork();
    if (pid < 0) {
    printf("Fork failed!\n");
    } else if (pid == 0) {
    printf("This is the child process. PID = %d\n", getpid());
    } else {
    printf("This is the parent process. PID = %d, Child PID = %d\n", getpid(), pid);
    }

return 0;
}</pre>
```

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Output:

This is the parent process. PID = 1851, Child PID = 1852This is the child process. PID = 1852

Conclusion: The fork() system call was successfully demonstrated. A parent process and a child process were created, each with unique process IDs.



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UNIT-2

EXPERIMENT-1A

Aim: To implement the First-Come, First-Served (FCFS) CPU scheduling algorithm.

Description:

- FCFS is the simplest CPU scheduling algorithm.
- Processes are executed in the order they arrive in the ready queue.
- It is non-preemptive.
- Average waiting time is highly affected by the order of process arrival.

Algorithm:

- 1. Input the number of processes and their burst times.
- 2. Sort processes by arrival order (if given).
- 3. Compute waiting time for each process:

$$WT[i] = CT[i] - AT[i] - BT[i]$$

4. Compute turnaround time:

$$TAT[i] = WT[i] + BT[i]$$

5. Display Gantt chart, waiting times, turnaround times, and averages.

Program:

```
#include <stdio.h>
int main() {
                                UNIVERSITY
int n, bt[20], wt[20], tat[20], i;
float avg wt = 0, avg tat = 0;
printf("Enter number of processes: ");
scanf("%d", &n);
printf("Enter Burst Time for each process:\n");
for (i = 0; i < n; i++) {
printf("P%d: ", i+1);
scanf("%d", &bt[i]);
}
wt[0] = 0;
for (i = 1; i < n; i++)
```

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```
wt[i] = wt[i-1] + bt[i-1];
}
for (i = 0; i < n; i++) {
    tat[i] = wt[i] + bt[i];
    avg_wt += wt[i];
avg_tat += tat[i];
}
printf("\nProcess\tBT\tWT\tTAT\n");
for (i = 0; i < n; i++) {
    printf("P%d\t%d\t%d\t%d\n", i+1, bt[i], wt[i], tat[i]);
}
printf("\nAverage Waiting Time = %.2f", avg_wt/n);
    printf("\nAverage Turnaround Time = %.2f\n", avg_tat/n);
return 0;
}</pre>
```

Output:

```
Enter number of processes: 3
Enter Burst Time for each process:
P1: 5
P2: 3
P3: 8
Process BT WT TAT
P1 5
       0
           5
P2 3
       5
           8
P3 8
          16
Average Waiting Time = 4.33
Average Turnaround Time = 9.67
```

Conclusion:

The FCFS scheduling algorithm was successfully implemented and analyzed.

EXPERIMENT-1B

Aim: To implement the Shortest Job First (SJF) CPU scheduling algorithm.

Description:

- □ SJF selects the process with the smallest burst time first.
- ☐ It reduces average waiting time compared to FCFS.
- ☐ Can be preemptive or non-preemptive (here, non-preemptive).

Algorithm:

- 1. Input processes and burst times.
- 2. Sort processes in ascending order of burst time.
- 3. Compute waiting time and turnaround time.
- 4. Display results with averages.

Program:

```
#include <stdio.h>
int main() {
int n, bt[20], wt[20], tat[20], p[20], i, j, temp;
float avg wt = 0, avg tat = 0;
printf("Enter number of processes: ");
scanf("%d", &n);
printf("Enter Burst Time for each process:\n");
for (i = 0; i < n; i++)
                                  UNIVERSITY
printf("P%d: ", i+1);
scanf("%d", &bt[i]);
p[i] = i+1;
for (i = 0; i < n-1; i++)
        for (j = i+1; j < n; j++)
               if (bt[i] > bt[j]) {
                      temp = bt[i]; bt[i] = bt[j]; bt[j] = temp;
                       temp = p[i]; p[i] = p[j]; p[j] = temp;
wt[0] = 0;
for (i = 1; i < n; i++) {
       wt[i] = wt[i-1] + bt[i-1];
```

2 3

3

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```
for (i = 0; i < n; i++) \{ \\ tat[i] = wt[i] + bt[i]; \\ avg\_wt += wt[i]; \\ avg\_tat += tat[i]; \}  printf("\nProcess\tBT\tWT\tTAT\n"); \\ for (i = 0; i < n; i++) \{ \\ printf("P\%d\t\%d\t\%d\t\%d\n", p[i], bt[i], wt[i], tat[i]); \}  printf("\nAverage Waiting Time = \%.2f", avg\_wt/n); \\ printf("\nAverage Turnaround Time = \%.2f\n", avg\_tat/n); \\ return 0; \}
```

Output:

```
Enter number of processes: 3
Enter Burst Time for each process:
P1: 2
P2: 5
P3: 8

Process BT WT TAT
P1 2 0 2
P2 5 2 7
P3 8 7 15

Average Waiting Time = 3.00
Average Turnaround Time = 8.00
```

Conclusion:

The SJF scheduling algorithm was successfully implemented and analyzed.

EXPERIMENT-2A

Aim: To simulate the CPU Scheduling using the **Priority Scheduling algorithm** in C++ and analyze process execution order based on priority.

Description: CPU Scheduling is the process of selecting a process from the ready queue and allocating CPU to it.

Priority Scheduling assigns a priority value to each process.

☐ The CPU is allocated to the process with the **highest priority** (lowest priority number is usually considered highest priority).

☐ If two processes have the same priority, they are scheduled according to their arrival order (or FCFS).

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Algorithm/Steps:

- 1. Start the program.
- 2. Input the number of processes.
- 3. For each process, input:
- Burst Time
- o Priority
- 4. Sort processes based on their priority (highest priority first).
- 5. Calculate:
- o Waiting Time (WT) for each process
- o Turnaround Time (TAT) for each process
- 6. Compute average WT and TAT.
- 7. Display the scheduling order and results.
- 8. End program.

Program:

```
#include <iostream>
#include <algorithm>
using namespace std;
struct Process {
  int pid, bt, priority, wt, tat;
};

// Compare function for sorting based on priority
bool compare(Process a, Process b) {
  return (a.priority < b.priority);
}
int main() {</pre>
```

Operating System

```
int n;
cout << "Enter number of processes: ";</pre>
cin >> n;
Process p[n];
for(int i = 0; i < n; i++) {
        p[i].pid = i + 1;
                cout << "Enter Burst Time and Priority for Process " << p[i].pid << ": ";
                cin >> p[i].bt >> p[i].priority;
// Sort processes by priority
sort(p, p + n, compare);
// Calculate Waiting Time and Turnaround Time
p[0].wt = 0;
for(int i = 1; i < n; i++) {
p[i].wt = p[i-1].wt + p[i-1].bt;
float avgWT = 0, avgTAT = 0;
for(int i = 0; i < n; i++) {
p[i].tat = p[i].wt + p[i].bt;
        avgWT += p[i].wt;
        avgTAT += p[i].tat;
avgWT = n;
avgTAT = n;
cout << "\nProcess\tBT\tPriority\tWT\tTAT\n";</pre>
for(int i = 0; i < n; i++) {
        cout << p[i].pid << "\backslash t" << p[i].bt << "\backslash t" << p[i].priority
                << "\t\t" << p[i].wt << "\t" << p[i].tat << endl;
}
cout << "\nAverage Waiting Time = " << avgWT;</pre>
cout << "\nAverage Turnaround Time = " << avgTAT << endl;</pre>
return 0;
```

Output:

```
Enter number of processes: 3
Enter Burst Time and Priority for Process 1: 10 2
Enter Burst Time and Priority for Process 2: 5 1
Enter Burst Time and Priority for Process 3: 8 3
Process BT Priority
                       WT TAT
2
    5
        1
                0
                    5
    10 2
1
                5
                    15
3
    8
        3
               15 23
Average Waiting Time = 6.66667
Average Turnaround Time = 14.3333
```

Conclusion: The CPU Scheduling using the **Priority Algorithm** was successfully simulated in C++. The program computes **Waiting Time** and **Turnaround Time** for each process and demonstrates how processes are scheduled based on their priority.



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EXPERIMENT-2B

Aim: To simulate CPU Scheduling using the Round Robin (RR) algorithm in C++ and analyze process execution order, waiting time, and turnaround time.

Description:

- Round Robin (RR) is a preemptive CPU scheduling algorithm.
- Each process is assigned a fixed time quantum.
- The CPU executes a process for at most one quantum. If the process is not finished, it is placed at the back of the ready queue.
- This ensures **fairness** since every process gets CPU time in a cyclic order.

Algorithm/Steps:

- 1. Start the program.
- 2. Input the number of processes and their **Burst Times**.
- 3. Input the **Time Quantum**.
- 4. Initialize a queue for scheduling.
- 5. Execute each process for the given quantum:
- If a process finishes before the quantum, mark it complete.
- Otherwise, reduce its burst time and move it to the end of the queue.
- 6. Calculate **Waiting Time (WT)** and **Turnaround Time (TAT)** for each process.

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- 7. Compute average WT and TAT.
- 8. Display process execution details.
- 9. End program.

Program:

```
#include <iostream>
#include <queue>
using namespace std;
struct Process {
int pid, bt, rt, wt, tat;
};
int main() {
int n, tq;
cout << "Enter number of processes: ";
cin >> n;
```

Operating System

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```
Process p[n];
for(int i = 0; i < n; i++) {
p[i].pid = i + 1;
cout << "Enter Burst Time for Process " << p[i].pid << ": ";
cin >> p[i].bt;
p[i].rt = p[i].bt; // remaining time
}
cout << "Enter Time Quantum: ";</pre>
cin >> tq;
int time = 0, completed = 0;
queue<int> q;
for(int i = 0; i < n; i++) q.push(i);
while(!q.empty()) {
int i = q.front();
q.pop();
if(p[i].rt > tq) {
time += tq;
p[i].rt = tq;
q.push(i); // re-queue the process
} else {
time += p[i].rt;
p[i].wt = time - p[i].bt;
p[i].tat = time;
                                  UNIVERSITY
p[i].rt = 0;
       completed++;
} }
float avgWT = 0, avgTAT = 0;
cout << "\nProcess\tBT\tWT\tTAT\n";
for(int i = 0; i < n; i++) {
cout << p[i].pid << "\t" << p[i].bt
       << "\t" << p[i].wt << "\t" << p[i].tat << endl;
avgWT += p[i].wt;
avgTAT += p[i].tat;
cout << "\nAverage Waiting Time = " << avgWT / n;</pre>
cout << "\nAverage Turnaround Time = " << avgTAT / n << endl; return 0;
```

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Output:

```
Enter number of processes: 3
Enter Burst Time for Process 1: 10
Enter Burst Time for Process 2: 5
Enter Burst Time for Process 3: 8
Enter Time Quantum: 3
Process BT WT TAT
   10 13 23
2
    5
       9
           14
       14 22
3
   8
Average Waiting Time = 12
Average Turnaround Time = 19.6667
```

Conclusion: The Round Robin CPU Scheduling algorithm was successfully simulated in C++.



Roll No.: 2 4 B 1 1 2 3 3 Date: A **UNIT-3 EXPERIMENT-1** Aim: To simulate the Banker's Algorithm for deadlock avoidance in C++. **Description:** The Banker's Algorithm is used in operating systems to avoid deadlocks by allocating resources safely. It checks if a system is in a **safe state** before granting resource requests. Algorithm/Steps: 1. Input number of processes and resources. 2. Input Allocation, Maximum, and Available matrices. 3. Compute Need = Max - Allocation. 4. Apply Banker's Algorithm to check if the system is in a safe state. 5. If a safe sequence exists, display it; otherwise, report unsafe state. **Program:** #include <iostream> using namespace std; int main() { int n, m; cout << "Enter number of processes: "; cin >> n; cout << "Enter number of resources: "; cin >> m; int alloc[n][m], max[n][m], avail[m]; cout << "Enter Allocation Matrix:\n";</pre> for(int i=0;i< n;i++) for(int j=0;j< m;j++) cin >> alloc[i][j]; cout << "Enter Maximum Matrix:\n";</pre> for(int i=0; i< n; i++) for(int j=0; j< m; j++) cin >> max[i][j]; cout << "Enter Available Resources:\n";</pre> for(int i=0;i < m;i++) cin >> avail[i]; int need[n][m], $finish[n]=\{0\}$, safeSeq[n], work[m]; for(int i=0;i<m;i++) work[i]=avail[i]; for(int i=0;i< n;i++) for(int j=0;j< m;j++) need[i][j]=max[i][j]-alloc[i][j]; int count=0;

Operating System

```
while(count<n){</pre>
bool found=false;
for(int i=0;i< n;i++){
if(!finish[i]){
int j;
for(j=0;j \le m;j++) if(need[i][j] \ge work[j]) break;
if(j==m){
for(int k=0;k\le m;k++) work[k]+=alloc[i][k];
safeSeq[count++]=i;
finish[i]=1;
found=true;
        if(!found){ cout<<"System is not in a safe state\n"; return 0; }
cout << "System is in a safe state.\nSafe sequence: ";
for(int i=0;i<n;i++) cout<<"P"<<safeSeq[i]<<" "
cout << endl;
return 0;
Output:
Enter number of processes: 5
Enter number of resources: 3
 Enter Allocation Matrix:
2 0 0
3 0 2
```

Conclusion: Banker's Algorithm was successfully implemented to check safe and unsafe states for deadlock avoidance.

2 1 1 0 0 2

3 3 2

Enter Maximum Matrix:

Enter Available Resources:

System is in a safe state. Safe sequence: P1 P3 P4 P0 P2

EXPERIMENT-2:

Aim: To write a C program that illustrates **two processes communicating using shared memory**.

Description:

Shared memory allows **two or more processes to communicate** by accessing a common memory space.

It is one of the fastest IPC mechanisms.

Algorithm/Steps:

- 1. Create a shared memory segment using shmget().
- 2. Attach it to the process using shmat().
- 3. One process writes data into shared memory.
- 4. Another process reads data from shared memory.
- 5. Detach and remove shared memory segment.

Program:

```
#include <stdio.h>
#include <sys/ipc.h>
#include <sys/shm.h>
#include <string.h>
int main() {
key_t key = ftok("shmfile",65);
int shmid = shmget(key,1024,0666|IPC_CREAT);
char *str = (char*) shmat(shmid,(void*)0,0);
printf("Write Data: ");
fgets(str, 1024, stdin);
printf("Data written in memory: %s\n", str);
shmdt(str);
return 0;
}
```

Output:

```
Write Data: LUCKY
Data written in memory: LUCKY
```

Conclusion: Shared memory IPC was successfully demonstrated, showing inter-process data exchange.

Operating System

EXPERIMENT-3:

Aim: To create a thread using the pthreads library in C and run its function.

Description:

- ☐ Threads allow concurrent execution within a process.
- ☐ Pthreads (pthread.h) provide APIs for creating and managing threads.

Algorithm/Steps:

- 1. Include pthread.h.
- 2. Define a function to be executed by a thread.
- 3. Create a thread using pthread create().
- 4. Wait for the thread to finish using pthread join().

Program:

```
#include <stdio.h>
#include <pthread.h>
void* threadFunc(void* arg) {
printf("Hello from thread!\n");
return NULL;
}
int main() {
pthread_t tid;
pthread_create(&tid, NULL, threadFunc, NULL);
pthread_join(tid, NULL);
printf("Thread execution completed.\n");
return 0;
}
```

Output:

```
Hello from thread!
Thread execution completed.
```

Conclusion: Thread creation and execution using the **pthreads library** was successfully demonstrated.

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EXPERIMENT-4

Aim: To illustrate **concurrent execution of threads** using the pthreads library in C. **Description:**

- ☐ Multiple threads can run **concurrently** within the same process.
- Each thread executes independently but shares memory.

Algorithm/Steps:

- 1. Define multiple functions for threads.
- 2. Create multiple threads using pthread create().
- 3. Let each thread execute its function concurrently.
- 4. Use pthread join() to synchronize thread completion.

Program:

```
#include <stdio.h>
#include <pthread.h>
void* task1(void* arg) {
for(int i=0;i<5;i++) printf("Task 1 running\n");
return NULL;
void* task2(void* arg) {
for(int i=0;i<5;i++) printf("Task 2 running\n");
return NULL;
int main() {
                               UNIVERSITY
pthread t t1, t2;
pthread create(&t1, NULL, task1, NULL);
pthread create(&t2, NULL, task2, NULL);
pthread join(t1, NULL);
pthread_join(t2, NULL);
printf("Both threads finished execution.\n");
return 0;
```

Operating System

Output:

```
Task 1 running
Task 2 running
Task 3 running
Task 3 running
Task 4 running
Task 5 running
Task 5 running
Task 6 running
```

Conclusion: Concurrent execution of threads using **pthreads** was successfully simulated, showing independent task execution within a process.



UNIT-4

EXPERIMENT-1

Aim: To simulate memory management using Multiprogramming with a Fixed Number of Tasks (MFT).

Description: MFT divides the memory into fixed partitions. Each process is allocated exactly one partition. Internal fragmentation occurs when a process does not fully use its partition.

Algorithm/Steps:

- 1. Input total memory and partition size.
- 2. Calculate number of partitions = total memory / partition size.
- 3. Input process sizes.

Program:

4. Allocate each process to a partition if it fits.

cout << "Enter total memory size: ";</pre>

5. Show internal fragmentation and unused memory.

#include <iostream> using namespace std; int main() { int ms, ps, nop, mp[10], n;

cin >> ms;
cout << "Enter partition size: ";

```
cin >> ps;

n = ms / ps;

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cout << "Total partitions: " << n << endl;
```

```
} for (int i = 0; i < nop && i < n; i++) { if (mp[i] \le ps)
```

```
cout << "Process " << i+1 << " allocated in partition " << i+1 << " with internal fragmentation: " << ps - mp[i] << endl;
```

```
else  \mbox{cout} << \mbox{"Process"} << i+1 << \mbox{" not allocated (too large)."} << \mbox{endl}; \ \} \\ \mbox{return 0;} \\ \}
```

cin >> mp[i];

Output:

Enter total memory size: 1000
Enter partition size: 200
Total partitions: 5
Enter number of processes: 3
Enter memory required for process 1: 180
Enter memory required for process 2: 20
Enter memory required for process 3: 250
Process 1 allocated in partition 1 with internal fragmentation: 20
Process 2 allocated in partition 2 with internal fragmentation: 180
Process 3 not allocated (too large).

Conclusion: MFT was simulated and internal fragmentation was observed when process sizes were smaller than the partition size.



EXPERIMENT-3

Aim: To simulate the FIFO page replacement algorithm.

Description: FIFO replaces the oldest page in memory when a new page must be loaded and no free frames are available.

Algorithm/Steps:

- 1. Input number of frames.
- 2. Input page reference string.
- 3. Load pages sequentially.
- 4. If page fault occurs and frames are full, replace the oldest page.

Program:

```
#include <iostream>
#include <queue>
#include <vector>
using namespace std;
int main() {
int n, f, page, faults = 0;
cout << "Enter number of frames: ";</pre>
cin >> f;
cout << "Enter number of pages: ";</pre>
cin >> n;
vector<int> pages(n);
cout << "Enter reference string: ";</pre>
for (int i = 0; i < n; i++) cin >> pages[i];
queue<int> q;
vector<int> frame;
for (int i = 0; i < n; i++) {
        page = pages[i];
        bool found = false;
        for (int x: frame) if (x == page) found = true;
        if (!found) {
                if (frame.size() < f) frame.push back(page);
                else {
int victim = q.front(); q.pop();
for (int j = 0; j < \text{frame.size}(); j++)
if (frame[j] == victim) frame[j] = page;
        q.push(page);
        faults++;
```

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```
}
    cout << "Frames: ";
    for (int x : frame) cout << x << " ";
        cout << endl;
}
cout << "Total Page Faults: " << faults << endl;
return 0;
}</pre>
```

Output:

```
Enter number of frames: 3
Enter number of pages: 7
Enter reference string: 1 2 3 4 1 2 5
Frames: 1
Frames: 1 2
Frames: 1 2 3
Frames: 4 2 3
Frames: 4 1 3
Frames: 4 1 2
Frames: 5 1 2
Total Page Faults: 7
```

Conclusion: FIFO page replacement was simulated and the number of page faults was observed.



EXPERIMENT-4

Aim: To simulate the LRU page replacement algorithm.

Description: Least Recently Used (LRU) replaces the page that has not been used for the

longest time.

Algorithm/Steps:

- 1. Input number of frames and reference string.
- 2. Keep track of last used time for each page.
- 3. Replace the page with the longest idle time when a page fault occurs.

Program:

```
#include <iostream>
       #include <vector>
#include <unordered map>
using namespace std;
int main() {
       int n, f, page, faults = 0;
       cout << "Enter number of frames: ";</pre>
       cin >> f;
       cout << "Enter number of pages: ";</pre>
       cin >> n;
       vector<int> pages(n);
       cout << "Enter reference string: ";
       for (int i = 0; i < n; i++) cin >> pages[i];
       vector<int> frame;
       unordered map<int,int> lastUsed;
       for (int i = 0; i < n; i++) {
       page = pages[i];
       bool found = false;
       for (int x : frame) if (x == page) found = true;
       if (!found) {
               if (frame.size() < f) frame.push back(page);
               else {
               int lru index = 0;
               int min last = i;
               for (int j = 0; j < \text{frame.size}(); j++) {
                       if (lastUsed[frame[j]] < min last) {
                              min last = lastUsed[frame[j]];
                              lru index = i;
                       frame[lru index] = page;
```

```
faults++;
}
lastUsed[page] = i;
cout << "Frames: ";
for (int x : frame) cout << x << " ";
cout << endl;
}
cout << "Total Page Faults: " << faults << endl;
return 0;
}
Output:</pre>
```

```
Enter number of frames: 3
Enter number of pages: 7
Enter reference string: 1 2 3 4 1 2 5
Frames: 1
Frames: 1 2
Frames: 1 2 3
Frames: 4 2 3
Frames: 4 1 3
Frames: 4 1 2
Frames: 5 1 2
Total Page Faults: 7
```

Conclusion: LRU page replacement was simulated and the least recently used page was replaced on each fault.

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EXPERIMENT-5: Simulate File Allocation Strategies

Aim:

To simulate file allocation strategies: (a) Sequential, (b) Indexed, (c) Linked.

Description:

- **Sequential:** Files occupy consecutive blocks.
- ☐ **Indexed:** An index block maintains pointers to all file blocks.
- Linked: Each block contains a pointer to the next block.

Algorithm/Steps:

- 1. Input number of blocks and file sizes.
- 2. For Sequential, allocate continuous blocks.
- 3. For Indexed, allocate an index block with pointers.
- 4. For Linked, create chain of blocks.

Program(a):

```
#include <iostream>
#include <vector>
using namespace std;
int main() {
int n, f, start, length;
cout << "Enter total number of disk blocks: ";
cin >> n;
vector\leqint\geq block(n, 0); // 0 = free, 1 = allocated
cout << "Enter number of files: "; UNIVERSITY
cin >> f:
for (int i = 0; i < f; i++) {
cout << \verb"\nEnter starting block of file" << i+1 << ":";
cin >> start;
cout \ll "Enter length of file " \ll i+1 \ll ": ";
cin >> length;
bool allocated = true;
        for (int j = \text{start}; j < \text{start} + \text{length}; j++) {
                if (j \ge n \parallel block[j] == 1) {
                                 allocated = false;
                        break;
        if (allocated) {
```

```
Date:
```

Roll No.: 2 4 B 1 1 A I 2 3 3

```
for (int j = start; j < start + length; j++) \\ block[j] = 1; \\ cout << "File " << i+1 << " allocated blocks: "; \\ for (int j = start; j < start + length; j++) \\ cout << j << " "; \\ cout << endl; \\ } else \{ \\ cout << "File " << i+1 << " cannot be allocated." << endl; \\ } \\  return 0;  \}  Output:
```

```
Enter total number of disk blocks: 10
Enter number of files: 2

Enter starting block of file 1: 2
Enter length of file 1: 3
File 1 allocated blocks: 2 3 4

Enter starting block of file 2: 5
Enter length of file 2: 2
File 2 allocated blocks: 5 6
```

Program(b):

```
#include <iostream> UNIVERSITY
#include <vector>
using namespace std;
int main() {
  int n, f, indexBlock, length;
  cout << "Enter total number of disk blocks: ";
  cin >> n;
  vector<int> block(n, 0);
  cout << "Enter number of files: ";
  cin >> f;
  for (int i = 0; i < f; i++) {</pre>
```

cout << "\nEnter index block for file " << i+1 << ": ";

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```
cin >> indexBlock;
if(block[indexBlock] == 1) {
cout << "Index block already allocated. File cannot be stored.\n";
continue;
block[indexBlock] = 1;
cout << "Enter number of blocks needed for file " << i+1 << ": ";
cin >> length;
vector<int> allocated;
cout << "Enter block numbers: ";</pre>
for (int j = 0; j < length; j++) {
int b;
cin >> b;
if (b \le n \&\& block[b] == 0) {
       block[b] = 1;
              allocated.push back(b);
} else {
              cout << "Block " << b << " already allocated or invalid.\n";
}
       cout << "File " << i+1 << " stored using index block "
                     << indexBlock << " -> ";
       for (int b : allocated) cout << b << " ";
       cout << endl;
return 0;
                                UNIVERSITY
Output:
Enter total number of disk blocks: 20
Enter number of files: 2
Enter index block for file 1: 15
Enter number of blocks needed for file 1: 3
Enter block numbers: 1 2 3
File 1 stored using index block 15 -> 1 2 3
```

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Enter index block for file 2: 15

Index block already allocated. File cannot be stored.

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Program(c):

```
#include <iostream>
#include <vector>
using namespace std;
struct Node {
int block;
Node* next;
};
int main() {
int n, f, length;
cout << "Enter total number of disk blocks: ";</pre>
cin >> n;
vector\leqint\geq block(n, 0);
cout << "Enter number of files: ";</pre>
cin >> f:
for (int i = 0; i < f; i++) {
cout << "\nEnter number of blocks required for file " << i+1 << ": ";
cin >> length;
Node* head = nullptr;
Node* temp = nullptr;
cout << "Enter block numbers: ";</pre>
for (int j = 0; j < length; j++) {
int b:
cin >> b;
if (b < n \&\& block[b] == 0) {
                                              block[b] = 1;
Node* newNode = new Node{b, nullptr};
               if (!head) head = newNode;
               else temp->next = newNode;
               temp = newNode;
} else {
               cout << "Block " << b << " already allocated or invalid.\n";
cout << "File " << i+1 << " allocated blocks (linked): ";
temp = head;
while (temp) {
cout << temp->block << " -> ";
        temp = temp->next;
        cout << "NULL" << endl;</pre>
return 0; }
```

Roll No.: 2 4 B 1 1 A I 2 3 3

Output:

```
Enter total number of disk blocks: 20
Enter number of files: 2

Enter number of blocks required for file 1: 4
Enter block numbers: 10 12 1 8
File 1 allocated blocks (linked): 10 -> 12 -> 1 -> 8 -> NULL

Enter number of blocks required for file 2: 3
Enter block numbers: 1 5 6
Block 1 already allocated or invalid.
File 2 allocated blocks (linked): 5 -> 6 -> NULL
```

Conclusion: We successfully simulated **Sequential, Indexed, and Linked File Allocation** strategies in C++. Each strategy illustrated how files are stored in secondary memory and demonstrated issues like external fragmentation, indexing, and linked block traversal



	Roll No.: 2 4 B 1 1 A I 2 3 3		
	UNIT-5		
AIM	: To compare how Windows and Unix/Linux operating systems implement Secondary		
Stora	age Structures, System Protection, and System Security through case study analysis.		
DES	CRIPTION / THEORY:		
1. Sec	condary Storage Structure:		
	Disk Structure & Attachment:		
o	Windows: NTFS, basic/dynamic disks, SATA/USB/iSCSI attachments.		
o	Unix/Linux: ext4, XFS, ZFS, devices as /dev/sda, mounted with mount.		
	Disk Scheduling:		
o	Windows: SCAN/Elevator algorithms, NCQ.		
o	Linux: CFQ, Deadline, NOOP schedulers.		
	RAID Structure:		
o	Windows: Storage Spaces (software RAID), hardware RAID controllers.		
o	Linux: mdadm for RAID 0,1,5,6,10. Stable Storage: Windows: NTFS journaling, chkdsk, Volume Shadow Copy. Linux: ext4 journaling, fsck, LVM snapshots.		
o			
o			
2. Sys	stem Protection:		
	Goals of Protection: Prevent misuse of CPU, memory, and files; ensure authorized		
acces	s. UNIVERSITY		
	Principles and Domains:		
o	Windows: Access Control Lists (ACLs), Active Directory domains.		
o	Linux: User/Group/Other model, rwx permissions, sudo.		
	Access Matrix & Control:		
o	Windows: ACLs implement access matrix.		
o	Linux: Permission bits, chmod, chown, setfacl.		
	Revocation of Rights: Both OSs allow instant revocation, but require auditing for		
	e sessions.		
active			
	stem Security:		

- Windows: Viruses, ransomware, buffer overflows. Mitigation: Defender, Code
 Signing.
- o Linux: Rootkits, privilege escalation. Mitigation: SELinux, AppArmor.
- ☐ System & Network Threats:
- o Windows: Worms, phishing, DoS. Mitigation: Firewall, BitLocker, auto-updates.
- o Linux: SSH brute force, DDoS. Mitigation: iptables/ufw, fail2ban, Snort.

ALGORITHM / STEPS (Case Study Approach)

- 1. Study secondary storage management in both Windows and Linux.
- 2. Compare protection models: ACLs vs. permission bits.
- 3. Analyze access matrix and revocation implementation.
- 4. Identify common program, system, and network threats in both OSs.
- 5. Compare mitigation strategies (firewalls, encryption, auditing tools).
- 6. Summarize similarities and differences.

COMPARISON:

Feature	Windows	Unix/Linux
File System	NTFS (journaling, EFS)	ext4, XFS, ZFS
Disk Scheduling	SCAN, NCQ	CFQ, Deadline, NOOP
RAID	Storage Spaces, Hardware RAID	mdadm, Hardware RAID
Stable Storage	Journaling, chkdsk, Shadow Copy	Journaling, fsck, LVM snapshots
Protection Model	ACLs, Active Directory	User/Group/Other, rwx, sudo
Access Control	ACLs on files and objects	chmod, chown, setfacl
Program Threats	Viruses, ransomware	Rootkits, privilege escalation
Network Threats	Worms, phishing, DoS	SSH brute force, DDoS
Security Tools	Defender, BitLocker, Firewall	SELinux, iptables, fail2ban, Snort

CONCLUSION: Windows is preferred in enterprise environments requiring integration and centralized management, while Unix/Linux dominates servers and critical systems due to its security, flexibility, and stability. This case study highlights how different OS designs achieve the same fundamental goals of efficiency, protection, and security.

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