



# Dynamic

Video - 131

# Programming



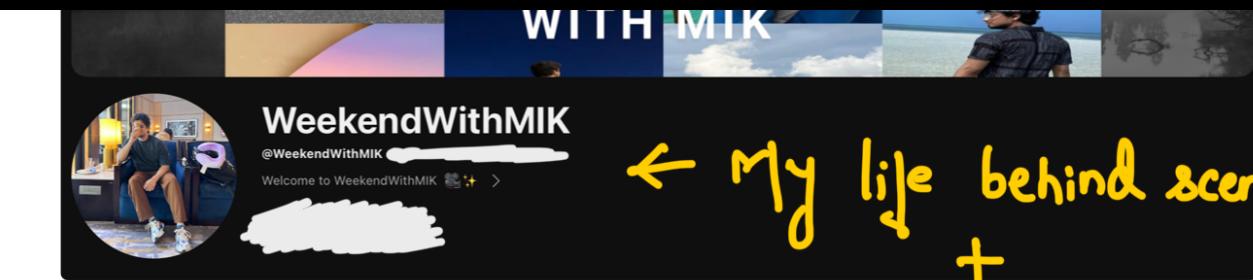
Note :- This playlist is only for explanation of Dns & solutions.

See my "DP Concepts & Dns" playlist for understanding DP from scratch...



- ∞ → codestorywithmik
- X → cswithMIK
- WhatsApp → codestorywithMIK





# Motivation :-

Whatever happens,  
just **Don't Quit.**



**MIK...**

## 3651. Minimum Cost Path with Teleportations

Hard Topics Companies Hint

You are given a  $m \times n$  2D integer array `grid` and an integer `k`. You start at the top-left cell  $(0, 0)$  and your goal is to reach the bottom-right cell  $(m - 1, n - 1)$ .

There are two types of moves available:

- **Normal move**: You can move right or down from your current cell  $(i, j)$ , i.e. you can move to  $(i, j + 1)$  (right) or  $(i + 1, j)$  (down). The cost is the value of the destination cell.
- **Teleportation**: You can teleport from any cell  $(i, j)$ , to any cell  $(x, y)$  such that  $grid[x][y] \leq grid[i][j]$ ; the cost of this move is 0. You may teleport at most  $k$  times.

Return the **minimum** total cost to reach cell  $(m - 1, n - 1)$  from  $(0, 0)$ .

Example:-  $grid =$

0	1	2
1	3	3

$$, k = 2$$

	2	5	4
2	4	3	5

$2 + 5 = 7$   
using  
1 help

Output : 7

### Constraints:

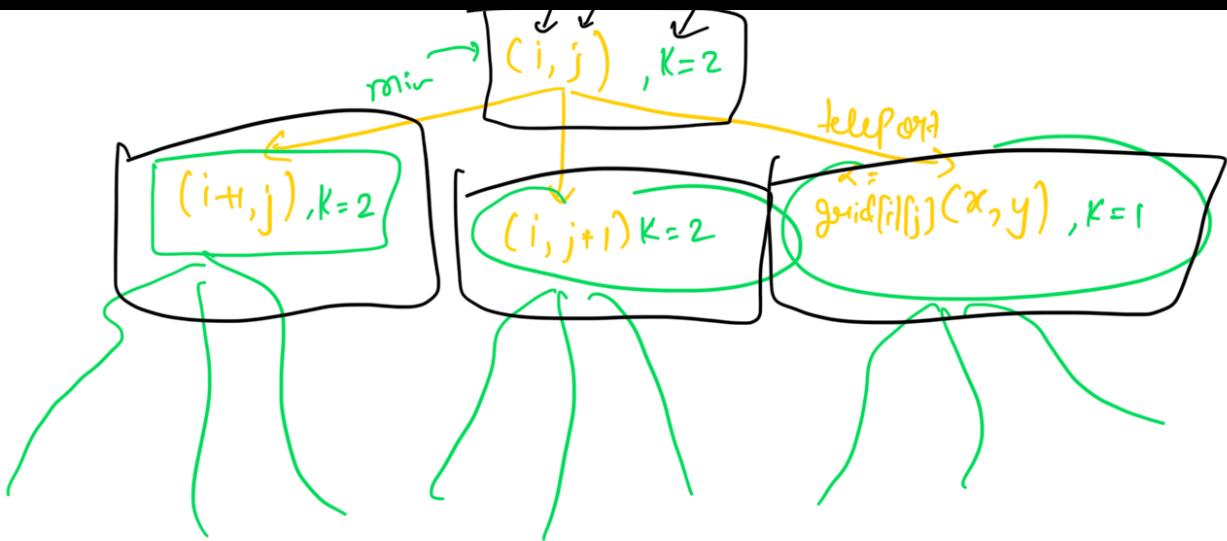
- $2 \leq m, n \leq 80$
- $m == \text{grid.length}$
- $n == \text{grid}[i].length$
- $0 \leq \text{grid}[i][j] \leq 10^4$
- $0 \leq k \leq 10$

Thought Process

0	1	2
1	3	3
2	5	4
3	4	3

$K = 2$

1 1 1



$(m-1)(n-1)$

Memoiz. [ ] [ ] [ ]

3D - DP

Solve(0, 0, 0, K, grid);

Solve(i, j, t, K, grid) {

if ( $i = m-1 \& j = n-1$ ) {

return 0;

}

int result = 1e9;

// Right (i, j) → (i, j+1)      grid[i][j+1];

if ( $j+1 < n$ ) {

int next = Solve(i, j+1, t, K, grid);

result = min(result, grid[i][j+1] + next);

}

// (i, j) → (i+1, j)

if ( $i+1 < m$ ) {

$\gamma(i, j, t) \leftarrow$

```

int next = Solve(i+1, j, t, K, grid);
result = min(result, (grid[i+1][j] + next));
}

// tempo (i,j) → (x,y)
if (t < K) {
    survival = grid[i][j];
    for (int x=0; x<m; x++)
        for (int y=0; y<n; y++) {
            if (x == i && y == j)
                continue;
            if (grid[x][y] <= grid[i][j]) {
                result = min(result,
                    solve(x, y, t+1, K, grid));
            }
        }
}
return result;
}

```

$[m][n][K]$

$$T.C = m \cdot n \cdot K * (m \cdot n)$$

$$= O(m^2 \cdot n^2 \cdot K) \quad \underline{\underline{T \cdot L \cdot E}}$$

# Bottom Up :-

```

int solve(int i, int j, int tPort, int k, vector<vector<int>>& grid) {
    // Reached destination
    if (i == m - 1 && j == n - 1)
        return 0;

    if (t[i][j][tPort] != -1)
        return t[i][j][tPort];

    int result = 1e9; // Large value
    int curVal = grid[i][j];

    // Move Right
    if (j + 1 < n) {
        ↓ ↓ ↓
        int next = solve(i, j + 1, tPort, k, grid);
        result = min(result, grid[i][j + 1] + next);
    }

    // Move Down
    if (i + 1 < m) {
        int next = solve(i + 1, j, tPort, k, grid);
        result = min(result, grid[i + 1][j] + next);
    }

    // Teleport
    if (tPort < k) {
        for (int x = 0; x < m; x++) {
            for (int y = 0; y < n; y++) {
                if ((x != i || y != j) && grid[x][y] <= curVal) {
                    result = min(result, solve(x, y, tPort + 1, k, grid));
                }
            }
        }
    }

    return t[i][j][tPort] = result;
}

```

$$dp = [m+1][n+1][k+1]$$

$dp[i][j][tPort]$  = min cost to  
 reach  $(m-1), (n-1)$   
 from  $(i, j)$  having  
 used  $tPort$

$$(i, j, tPort) =$$

$$dp[x][y][tPort+1]$$

①

for ( $t = 0$  ;  $t \leq K$  ;  $t++$ ) {

$$dp[m-1][n-1][t] = 0;$$

}

②





```

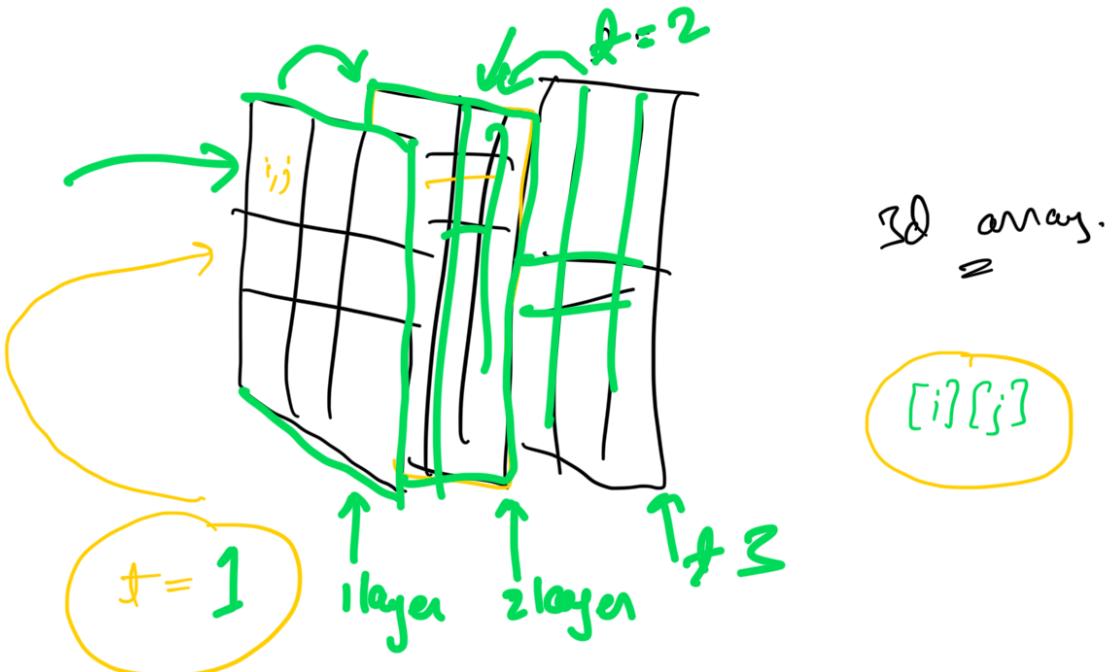
for (tport = K; tport >= 0; tport--) {
    for (int i = m - 1; i >= 0; i--) {
        for (int j = n - 1; j >= 0; j--) {
            if (i == m - 1 && j == n - 1) {
                continue;
            }
            int result = 1e9;
            curr = grid[i][j];
        }
    }
}

```

}

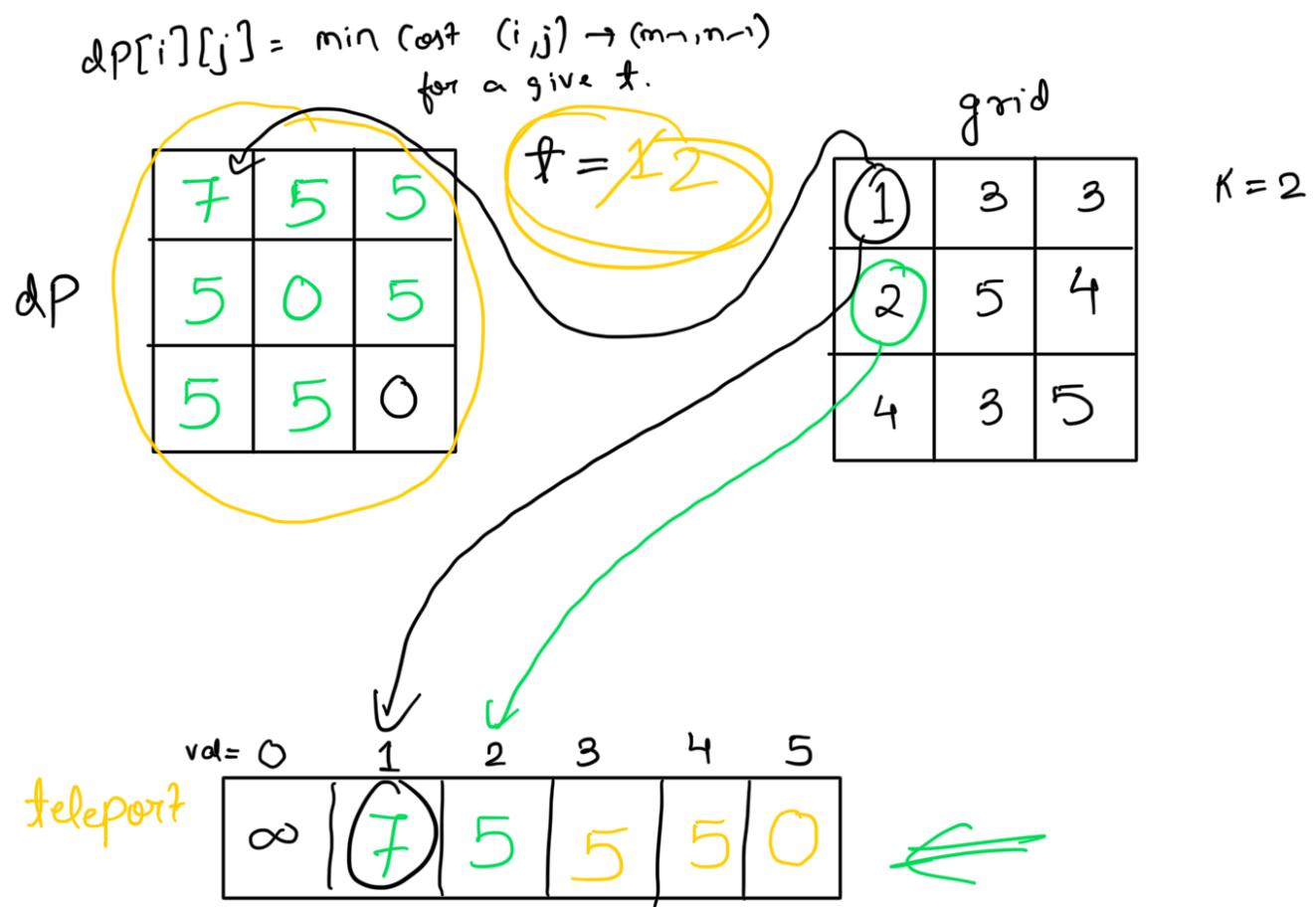
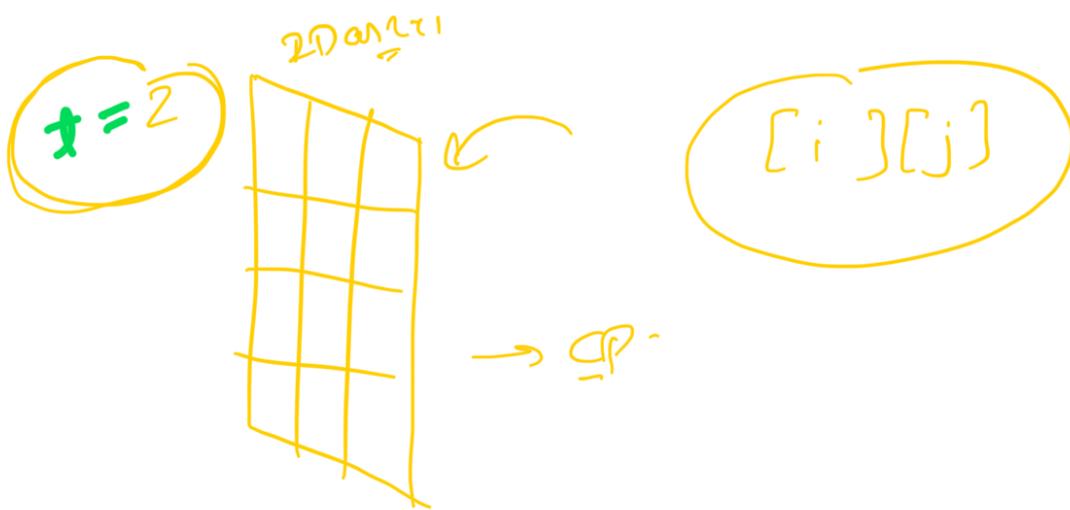
return dp[0][0][0];

# How to further optimize it ???



“Layered D.P.”

It's just a DP technique in which we solve the same problem multiple times, each time allowing one more unit of a limited resource (teleph -)



$t_{\text{teleport}}[\text{val}] =$



for ( $t = 0$ ;  $t \leq k$ ;  $t++$ ) {  $\rightarrow O(k)$

for ( $i = m-1$ ;  $i \geq 0$ ;  $i--$ )  $\rightarrow O(m)$

for ( $j = n-1$ ;  $j \geq 2$ ;  $j--$ ) {  $O(n)$

//  
if ( $i+1 < m$ ) {

$t[i][j] = \min(t[i][j], t[i+1][j] + grid[i+1][j]);$

if ( $j+1 < n$ ) {

$t[i][j] = \min(t[i][j], t[i][j+1] + grid[i][j+1]);$

if ( $t > 0$ ) {

$t[i][j] = \min(t[i][j], teleport[grid[i][j]]);$

}

}

for (int  $i = 0$ ;  $i < m$ ;  $i++$ ) {  $O(n \cdot n)$

for (int  $j = 0$ ;  $j < n$ ;  $j++$ ) {

$telepath[grid[i][j]] = \min(telepath[grid[i][j]], dp[i][j]);$

}

```

    // Prefix-min → Correctly update the DP
    for(int i=1; i<tel.size(); i++) {
        telp[i] = min(telp[i], telp[i-1]);
    }
}

```

return

$t[0][0]$ ;

$$T.C = O(K \cdot (m \cdot n + m \cdot n))$$

~~$$T.C = O(K \cdot m \cdot n)$$~~

What went

wrong?

DP

dp

5	4
4	0

$t = 1$

grid

3	1
10	4

$$K=2 \\ 0+4=4$$

0

teleport = 

i

j