

# # Segment Tree

## Concepts & Qns... #



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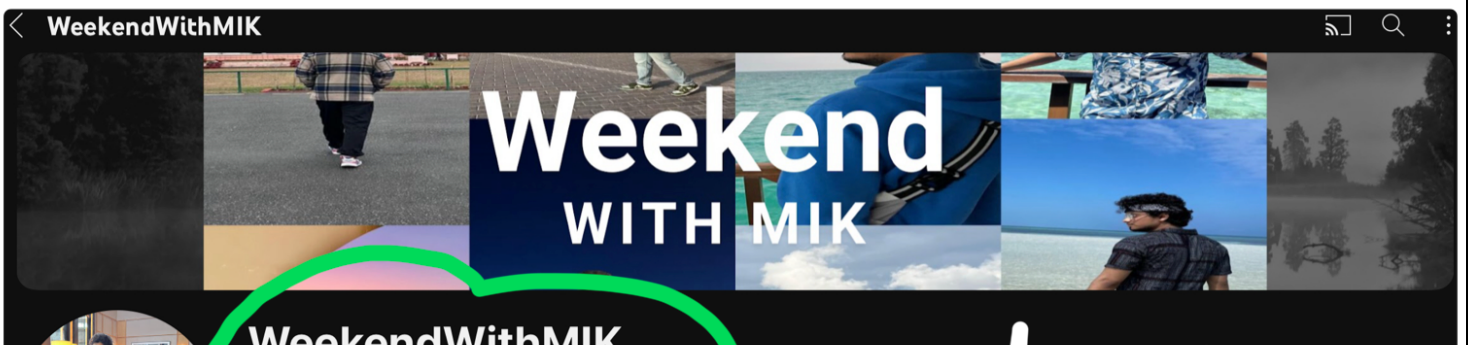
"No more fear of Segment Tree"

video -

15

Leetcode  
- 3721

Hard





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MIK

Try this channel to see  
my "Life behind the scenes + Tech News"

Motivation:-

It's now or never.  
work hard today,  
enjoy the results later...



MIK

## 3721. Longest Balanced Subarray II

Hard

Topics

Companies

Hint

You are given an integer array `nums`.

A **subarray** is called **balanced** if the number of **distinct even numbers** in the subarray is equal to the number of **distinct odd numbers**.

Return the length of the **longest** balanced subarray.

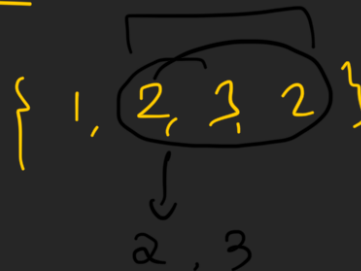
Example 1:

Input: `nums = [2, 5, 4, 3]`

Output: 4

Explanation:

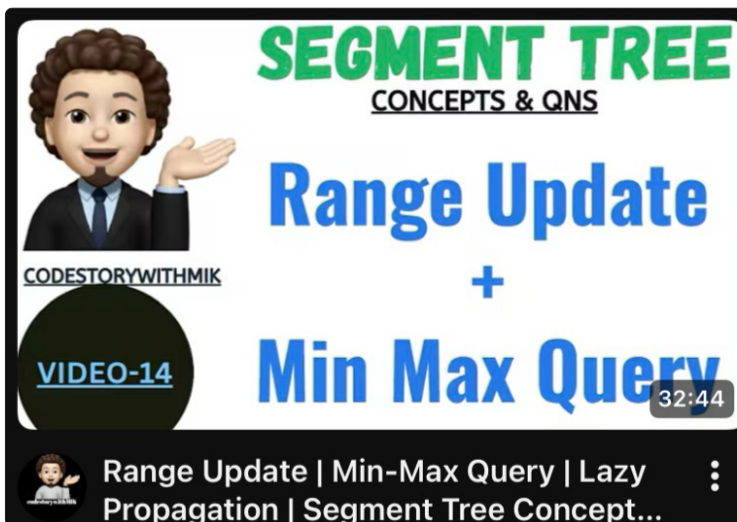
- The longest balanced subarray is `[2, 5, 4, 3]`.
- It has 2 distinct even numbers `[2, 4]` and 2 distinct odd numbers `[5, 3]`. Thus, the answer is 4.



Constraints :  $1 \leq \text{nums.length} \leq 10^5$   
 $1 \leq \text{nums}[i] \leq 10^5$

~~$n^2$~~   
 $O(n)$   
 $n \log n$

# Thought Process



→ Pre-requisite...

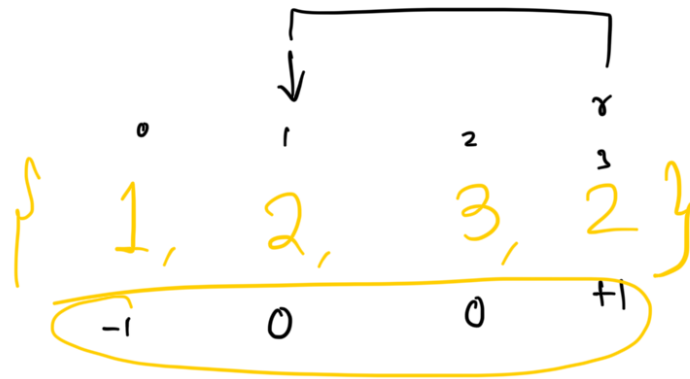
video - 7  
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Brute Force ...





$$\begin{aligned}
 &0 - p + 1 \\
 &= 3 - 0 + 1 \\
 &= 4
 \end{aligned}$$



# Optimising using Segment Tree...

Lazy prop

```

if(prev != -1) { //we have seen this element in past
    //[-1, prev] we are adding (-val) in the range
    for(int l = -1; l <= prev; l++) { //O(n) Range Query Update - Segment tree log(n)
        cumSum[l] += -val;
    }

    //[-1, prev] we are adding val in the range
    for(int l = 0; l <= prev; l++) {
        cumSum[l] += val; //O(n) Range Query Update - Segment tree log(n)
    }

    //We are finding left most 0 in the range [0...r]
    for(int l = 0; l <= r; l++) {
        if(cumSum[l] == 0) {
            maxL = max(maxL, r-l+1); //Range Search - segment in log
            break;
        }
    }
}
    
```

~~XXXX~~  $\tan(\log)$

$\log \sim \log$

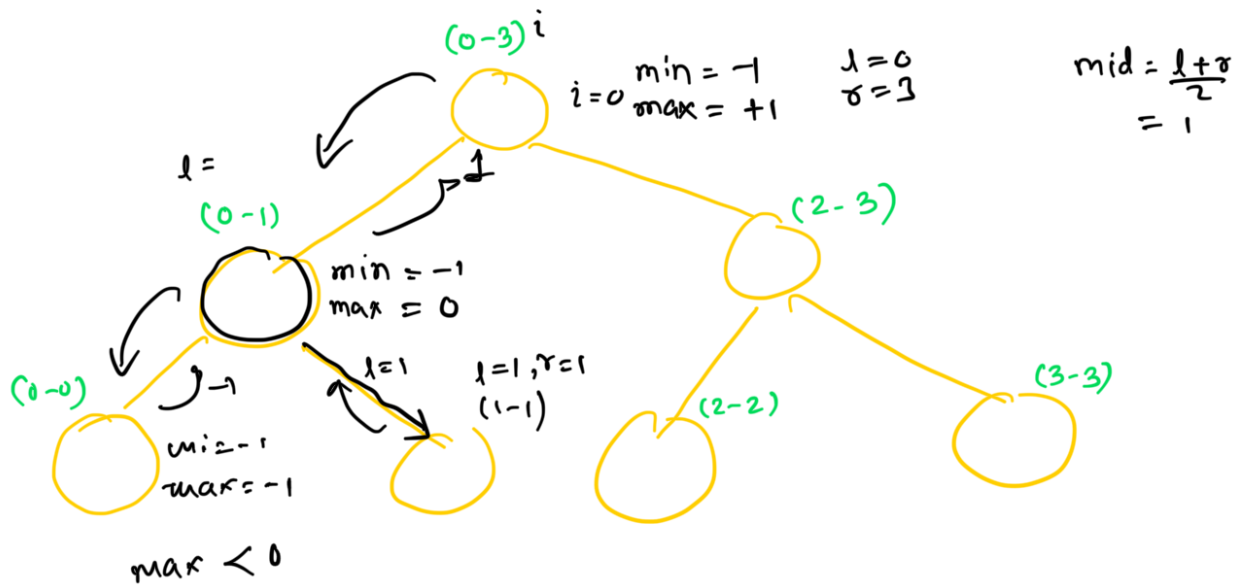
$j = \text{Query}(0-j+1)$



cumSum = [0, 0, 0, 0]



Query ( $i=0, l=0, r=3$ )



Query ( $2i+1, l=0, r=0$ )

$l=0, r=1$

$\text{mid} = \frac{l+r}{2} = 0$