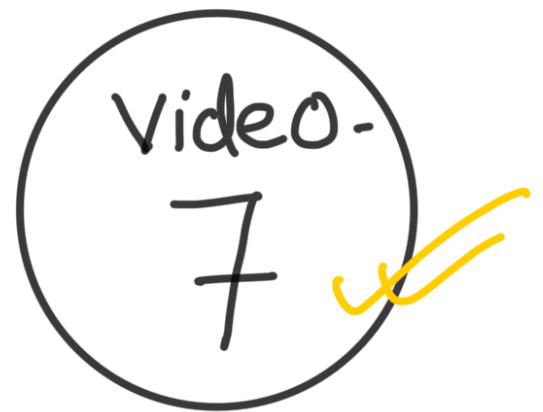


Difference Array

Concepts & Qns



Motivation :-

Use these ~50 days of year 2025 wisely.

Atleast choose any 4 topics which you find difficult. Give each topic 10 days. Relax and enjoy remaining days of this year and prepare yourself }

15
15
15

Strong for the next...

2536. Increment Submatrices by One

Medium

Topics

Companies

Hint

You are given a positive integer n , indicating that we initially have an $n \times n$ 0-indexed integer matrix `mat` filled with zeroes.

You are also given a 2D integer array `query`. For each `query[i] = [row1i, col1i, row2i, col2i]`, you should do the following operation:

- Add 1 to every element in the submatrix with the top left corner $(row1_i, col1_i)$ and the bottom right corner $(row2_i, col2_i)$. That is, add 1 to `mat[x][y]` for all $row1_i \leq x \leq row2_i$ and $col1_i \leq y \leq col2_i$.

Return the matrix `mat` after performing every query.



Example:-

$n = 3$

`queries = [[1, 1, 2, 2], [0, 0, 1, 1]]`

Output:

	0	1	2
0	1	1	0
1	1	2	1
2	0	1	1



Thought Process :-

	0	1	2	3
0	0	0	0	0
1	0	0	0	0
2	0	0	0	0
3	0	0	0	0

2-D

queries = $[(1,1,2,2), (0,0,1,1)]$

Step-1

(l, r)

$diff[l] += 1$

if $(r+1 < n)$ {

$diff[r+1] -= 1$

query = $[(0,2), (1,2), (2,3)]$

diff

0	1	2	3
0	0	0	0

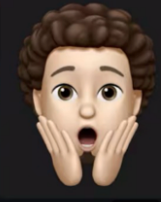
1-D

Step-2 → Find Cumulative Sum of diff.

DIFFERENCE ARRAY TECHNIQUE

CONCEPTS & QNS

VIDEO - 1
INTRODUCTION



26:10

Introduction | What | How | Difference Array Technique: Concepts & Questions - 1 | codestorywithMIK

codestorywithMIK

DAT in 2D-Array

	0	1	2	3
0	1	1	0	0
1	1	2	1	0
2	0	1	1	0
3	0	0	0	0

2-D

queries = $\left[(1, 1, 2, 2), (0, 0, 1, 1) \right]$

$r1 = 0$
 $c1 = 0$

$r2 = 1$
 $c2 = 1$

l

r

for (i = r1 ; i <= r2 ; i++) {

diff[i][c1] += 1;

if (c2+1 < n) {

diff[i][c2+1] -= 1;

}

}

Copy Paste \Rightarrow DAT
(1-D Array)

Store 2

Diff To Code:

```
vector<vector<int>> diff(n, vector<int>(n, 0));
```

```
for (query : queries) {  $\rightarrow O(q * n)$ 
```

```
    r1 = query[0]
```

```
    c1 = query[1]
```

```
    r2 = query[2]
```

```
    c2 = query[3]
```

step-1:-

```
for (int i = r1 ; i <= r2 ; i++) {  $\Leftarrow$   
    diff[i][c1] += 1  
    if (c2+1 < n) {  
        diff[i][c2+1] -= 1  
    }  
}
```

step-2:- Find Cumulative sum:-

$O(n^2)$

```
[ for (i = 0 ; i < n ; i++) {  
    for (j = 1 ; j < n ; j++) {  
        CumSum2  $\Leftarrow$  diff[i][j] += diff[i][j-1];  
    }  
}
```

}

return diff;

$$T.C = O(q \times n + n \times n);$$

$$S.C = O(1).$$