

Segment Tree

Concepts & Qns



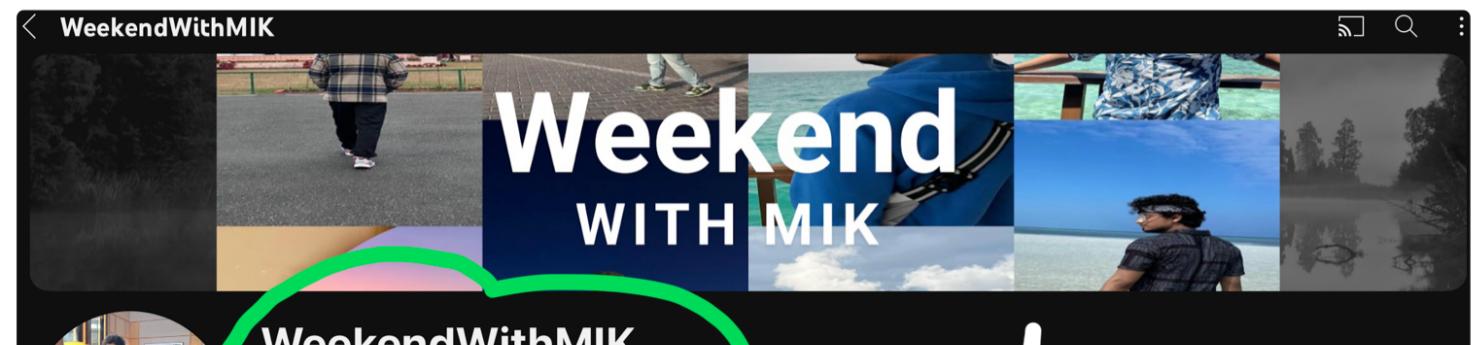
- ∞ → codestorywithMIK
- X → CSwithMIK
- codestorywithMIK

"No more fear of Segment Tree"

video -

15

Leetcode
- 3721
Hard





MIK



Try this channel to see
my "Life behind the Scenes + Tech News"

Motivation:-

It's now or never.
Work hard today,
enjoy the results later...



MIK..

3721. Longest Balanced Subarray II

Hard Topics Companies Hint

You are given an integer array `nums`.

A **subarray** is called **balanced** if the number of **distinct even numbers** in the subarray is equal to the number of **distinct odd numbers**.

Return the length of the **longest** balanced subarray.

Example 1:

Input: `nums = [2, 5, 4, 3]`

Output: 4

Explanation:

- The longest balanced subarray is `[2, 5, 4, 3]`.
- It has 2 distinct even numbers `[2, 4]` and 2 distinct odd numbers `[5, 3]`. Thus,



Constraints :

$1 \leq \text{nums.length} \leq 10^5$

$1 \leq \text{nums}[i] \leq 10^5$

~~n^2~~

$O(n)$

$n \log n$

Thought Process

SEGMENT TREE
CONCEPTS & QNS

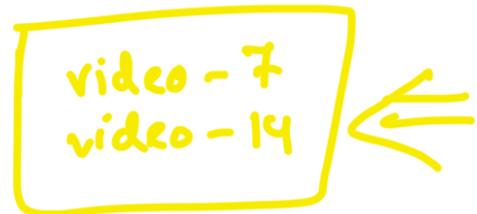
Range Update + **Min Max Query** 32:44

CODESTORYWITHMIK

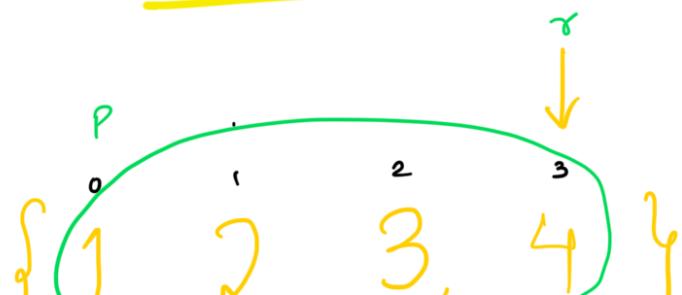
VIDEO-14

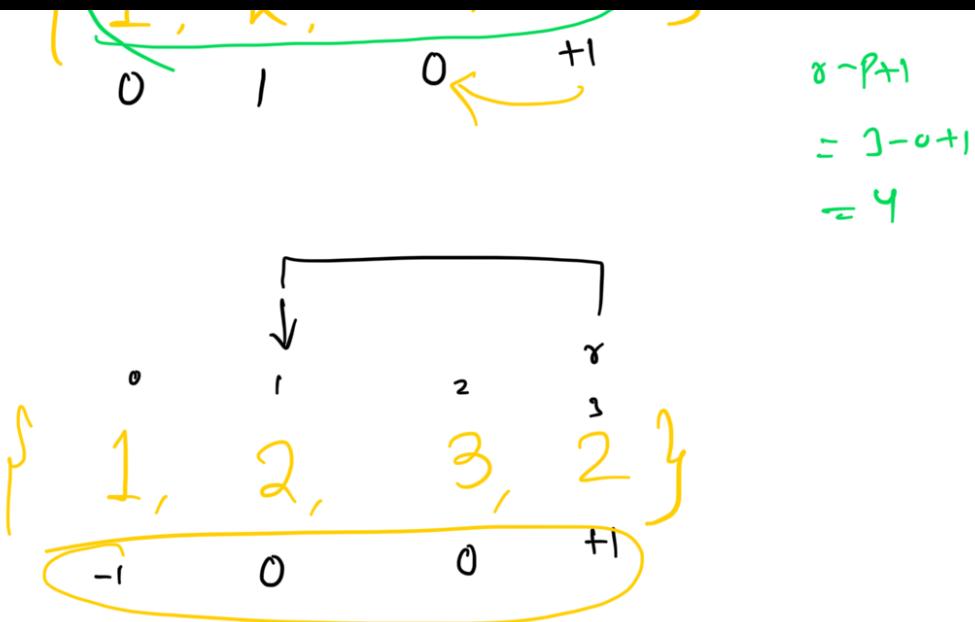
Range Update | Min-Max Query | Lazy Propagation | Segment Tree Concept...

→ Pre - Requisite ...



Brute Force ...





Optimising using
Segment Tree ...

```

if(prev != -1) { //we have seen this element in past
    // [0...prev] we are adding (-val) in the range
    for(int l = 0; l <= prev; l++) { // O(n) Range Query Update - Segment tree log(n)
        cumSum[l] -= val;
    }
}

// [0...r] we are adding val in the range
for(int l = 0; l <= r; l++) {
    cumSum[l] += val; // O(n) Range Query Update - Segment tree log(n)
}

// We are finding left most 0 in the range [0...r]
for(int l = 0; l <= r; l++) {
    if(cumSum[l] == 0) {
        maxL = max(maxL, r-l+1); // Range Search - segment in log
        break;
    }
}

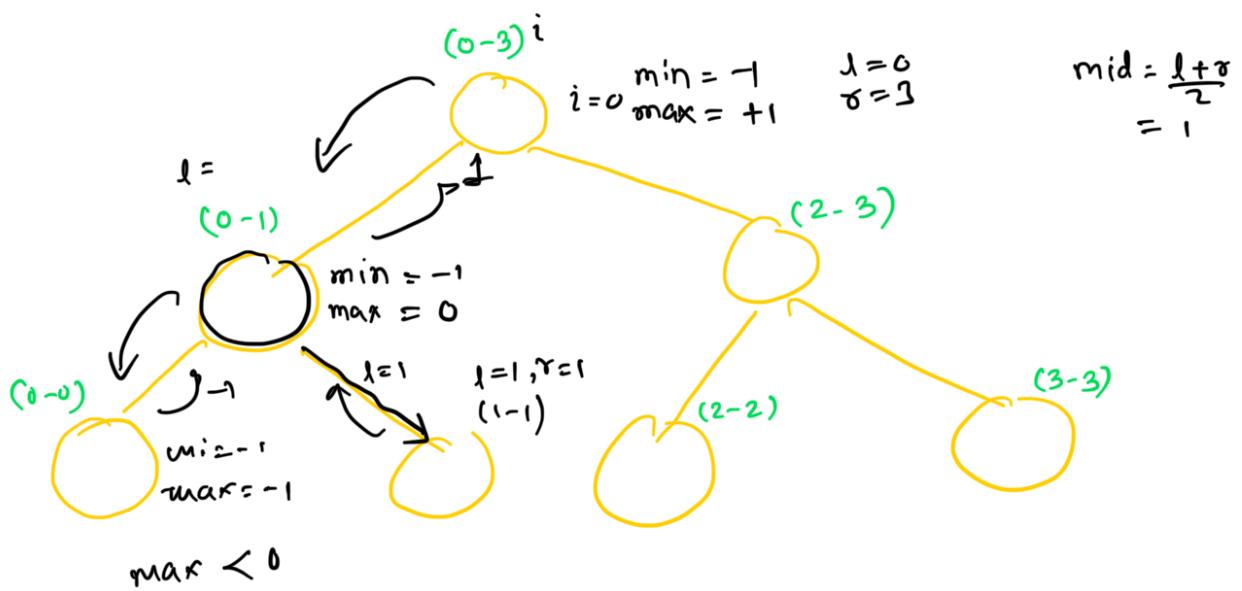
```

$\text{La}^3 \text{ prop}$



$\text{cumSum} = [\dots, \text{min}, \text{max}, \dots]$

Query ($i=0, l=0, r=3$)



Query ($2i+1, l=0, r=0$)

$$l=0, r=1$$

$$\text{mid} = \frac{0+1}{2} = 0$$