

Difference Array

Concepts & Qns



Motivation :-

Use these ~50 days of year 2025

wisely.

At least choose any 4 topics which you find difficult. Give each topic 15 days. Relax and enjoy remaining days 15 of this year and prepare yourself } 15

strong for the next...

2536. Increment Submatrices by One

Medium

Topics

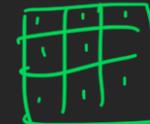
Companies

Hint

You are given a positive integer n , indicating that we initially have an $n \times n$ 0-indexed integer matrix mat filled with zeroes.

You are also given a 2D integer array query . For each $\text{query}[i] = [\text{row1}_i, \text{col1}_i, \text{row2}_i, \text{col2}_i]$, you should do the following operation:

- Add 1 to every element in the submatrix with the **top left corner** $(\text{row1}_i, \text{col1}_i)$ and the **bottom right corner** $(\text{row2}_i, \text{col2}_i)$. That is, add 1 to $\text{mat}[x][y]$ for all $\text{row1}_i \leq x \leq \text{row2}_i$ and $\text{col1}_i \leq y \leq \text{col2}_i$.



Return the matrix mat after performing every query.

Example:- $n = 3$. $\text{queries} = [[1, 1, 2, 2], [0, 0, 1, 1]]$

Output:

	1	1	0
	1	2	1
0	0	1	1

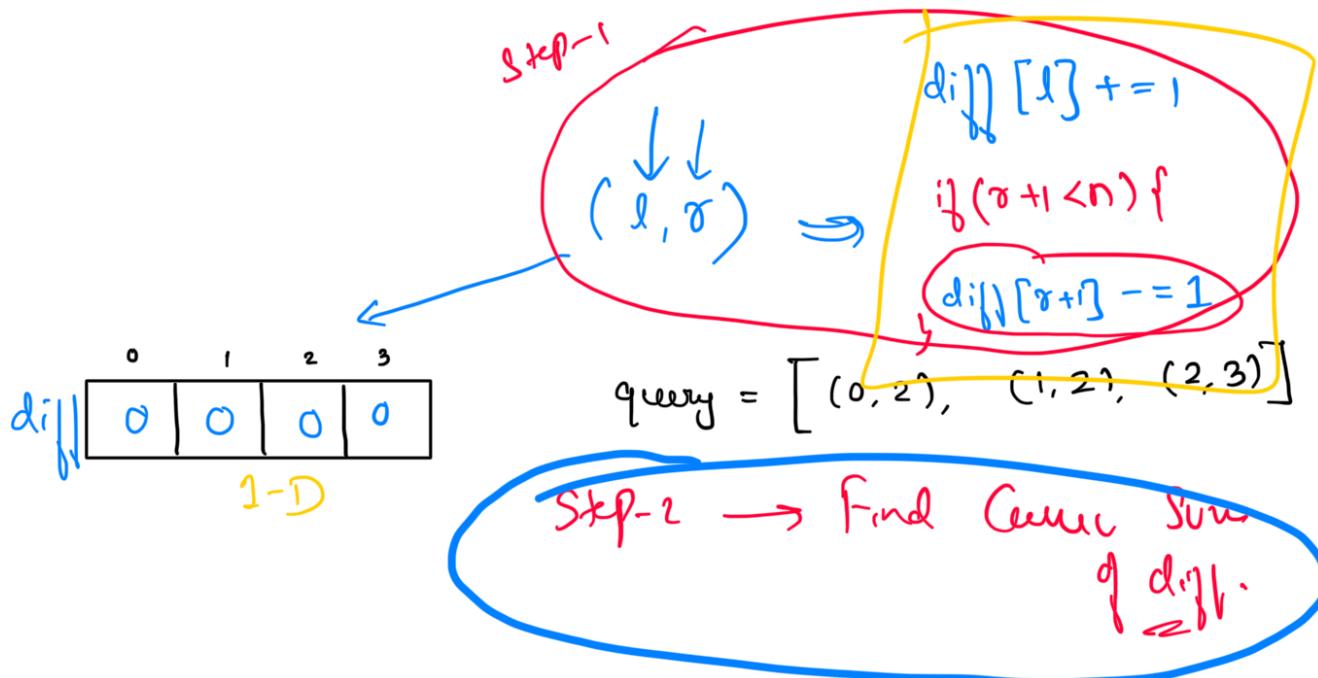


Thought Process :-

	0	1	2	3
0	0	0	0	0
1	0	0	0	0
2	0	0	0	0
3	0	0	0	0

2-D

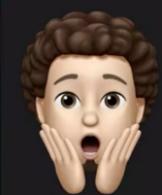
$$\text{queries} = [(1, 1, 2, 2), (0, 0, 1, 1)]$$



DIFFERENCE ARRAY TECHNIQUE

CONCEPTS & QNS

VIDEO - 1
INTRODUCTION



Introduction | What | How | Difference Array Technique: Concepts & Questions - 1 | codestorywithMIK

codestorywithMIK

DAT in 2D-Array

	0	1	2	3
0	1	1	0	0
1	1	2	1	0
2	0	1	1	0
3	0	0	0	0

2-D

$$\text{queries} = \left[(1, 1, 2, 2), (0, 0, 1, 1) \right]$$

$$\begin{aligned} r_1 &= 0 \\ c_1 &= 0 \end{aligned}$$

$$\begin{aligned} r_2 &= 1 \\ c_2 &= 1 \end{aligned}$$

for (i = r1 ; i <= r2 ; i++) {

 dij[i][c1] += 1;

 if (c2+1 < n) {

 dij[i][c2+1] -= 1;

}

Copy Paste \Rightarrow DAT
 (1-D Array)

Chas

? , , ,

Q10g To Code:

vector<vector<int>> diff(n, vector<int>(n, 0));

for (query : queries) { $\rightarrow O(q * n)$

$$\begin{aligned} r_1 &= \text{query}[0] \\ c_1 &= \text{query}[1] \\ r_2 &= \text{query}[2] \\ c_2 &= \text{query}[3] \end{aligned}$$

Step - 1 :-
for (int i = r₁; i <= r₂; i++) { \Leftarrow
 diff[i][c₁] += 1
 if (c₂+1 < n) {
 diff[i][c₂+1] -= 1
 }
}

Step - 2 :- Find Cumulative sum:-

$O(n^2)$

[for (i = 0; i < n; i++) {
 [for (j = 1; j < n; j++) {
 CumSum \Leftarrow diff[i][j] += diff[i][j-1];
 }
}

}

return ~~diff~~;

$$T.C = O(q*n + n*n);$$

$$S.C = O(1).$$