# Towards Practical Debugging of Wireless Sensor Networks

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### What is a Wireless Sensor Network

A wireless sensor network (WSN) is a collection of computing devices called motes, they have:

- a short range wireless radio
- an array of sensors such as light, heat and humidity
- a simple low powered CPU
- a battery with limited power supply

Motes communicate with each other to form a WSN. WSNs perform data gathering tasks such as environment monitoring.

## The Problem of Debugging Distributed Systems

- Multiple tasks running simultaneously leads to non-deterministic interactions
- Traditional debugging tools are unsuited
- Timing and synchronisation issues

## Complications to the Problem

- Motes are energy constrained
- Sending messages is the most expensive task
- Receiving messages is the next most expensive task [Shnayder et al., 2004]
- Motes have low computing power and a small memory
- WSNs deployed in hard to reach areas physical access after deployment is difficult [Herbert et al., 2007]

### Related Work

- Global Predicate Detection [Garg and Waldecker, 1996]
- H-SEND [Herbert et al., 2007]
- NodeMD [Krunic et al., 2007]
- TinyOS [Levis et al., 2005], Contiki

## Project Aims

- Develop tools to aid in debugging distributed programs running on WSNs.
- Implement libraries that check predicates, with a focus on correctly evaluating these predicates.
- Investigate if there are places in the network where evaluation is more efficient.
- Visualise some of the state of the network, as part of a tool to inform system users what state the network is in.

## Project Management

- Matthew Group Leader
- Amit Technical Manager
- Everyone was involved with development and research
- Bitbucket.org Git repository

# Implemented Libraries

## Implemented Libraries — Containers

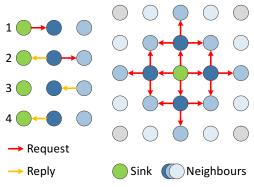
- Linked List
  - Ours: Standard linked list
  - Contiki's: Intrusive linked list
- Array List
- Unique Array
- Map
- Unit tests implemented for each

#### Benefits:

- Abstraction
- Reduced code duplication
- Simplified memory management

## Implemented Libraries — N-Hop Request

- Used by predicate evaluation
- Floods request N hops away from sink
- Asks for mote's current state
- Returned along the chain created by the flooding stage



## Implemented Libraries — N–Hop Flood

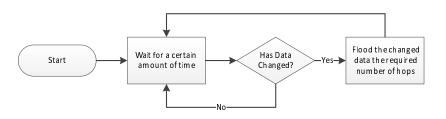
- Floods a given packet N hops
- Surprised that Contiki did not provide this as a library
  - Had to implement ourselves using TTLs in packet headers

## Implemented Libraries — Event Update

- Used by predicate evaluation
- Periodically checks if node's data has changed
- If it has, floods the new data the required number of hops

### Depends On:

N-Hop Flood



## Implemented Libraries — Multi-Packet Unicast

- Contiki packet size: 128 bytes
- This is too small for some of our data
- Alternative APIs Contiki implements are convoluted
  - Targeted towards sending file chunks
- We split packet up, send pieces and then reassemble

## Implemented Libraries — Tree Aggregation

- Leaf node generates data, forwards to parent
- Parent waits for a period, aggregating received data
- The node adds its own data to the aggregate
- The node then forwards the message to its parent
- Sepeat until the aggregated message reaches the base station

Again surprised that Contiki didn't have an implementation

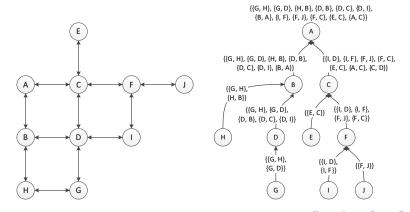
### Depends On:

Multi-Packet



## Implemented Libraries — Neighbour Detection

- To debug a WSN we need to know the network topology
- Uses Tree Aggregation to send neighbours to sink
- Contiki library does neighbour detection



## Predicate Evaluation

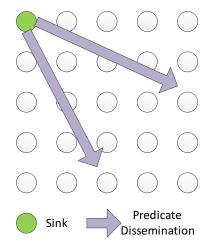
### Predicate Evaluation

- Disseminate predicate to network
- ② Evaluate predicate
- Return response to base station

#### Consider:

- Where to evaluate: Local vs. Global
- When to propagate mote data: Periodic vs. Event
- How to respond to a failure or success

## Predicate Evaluation — Dissemination



• Targeted to individual node or entire network



### Predicate Evaluation — Libraries I

	Periodic	Event
Local	PELP  • Evaluated in–network  • Data is requested when needed to evaluate predicate  • Previous round's data is forgotten after a round completes	PELE  • Evaluated in–network  • Data is sent by data sources, when it changes  • Data is never forgotten, simply updated

### Predicate Evaluation — Libraries II

	Periodic	Event
	PEGP	PEGE
Global	<ul> <li>Similar to PELP, except data is aggregated to the base station</li> <li>Predicate evaluation at base station</li> </ul>	<ul> <li>Similar to PELE, except data is aggregated to the base station</li> <li>Predicate evaluation at base station</li> </ul>

### Predicate Evaluation — Libraries Used

	Periodic	Event
Local	PELP  • N-Hop Request	PELE
Global	PEGP  • Neighbour Detect  • Tree Aggregation	PEGE      Neighbour Detect     Tree Aggregation

## Predicate Evaluation — Response

#### Failure

- Only predicate failures are reported to the base station
- Cannot say much about the network state, either:
  - we have not been informed of a failure
  - · we have been informed
- We chose this one due to the reduced traffic

#### Failure and Success

- Both failures and successes are reported
- Supports detecting the network is in the following states:
  - Unknown
  - Failed
  - Succeeded
  - Failed and then later succeeded
- Uses more energy

Note: Global PE may as well take advantage of "Failure and Success" messages, as the target of them is the node the predicates are evaluated on



### Predicate Evaluation — Evaluation

- Implemented a virtual machine in C to evaluate predicates on the nodes
  - Optimised for low memory environment
  - Opcodes for high-level operations to reduce program size
- Implemented a DSL with a JavaCC parser and a Java compiler and assembler
  - Functional language
  - Expects a boolean output

### Predicate Evaluation — DSL

```
\forall n \in \mathsf{Nodes}
                \forall n' \in \text{Neighbours}(n, 2).
                     slot(n) \neq slot(n')
[all]
function 1 as slot returning int in
    using Neighbours(2) as twohopn in
         @(x : twohopn ~
              slot(x) != slot(this)
```

### Predicate Evaluation — DSL

```
\forall n \in \mathsf{Nodes}
        \forall n' \in \mathsf{Neighbours}(n,1) \cup \{n\}.
              \forall n'' \in \mathsf{Neighbours}(n,1) \cup \{n\}
                   addr(n') \neq addr(n'')
                          \implies slot(n') \neq slot(n'')
[all]
function 0 as addr returning int in
function 1 as slot returning int in
     using Neighbours(1) as ohn in
          @(a : ohn ~
               @(b : ohn ~ addr(a) != addr(b)
                    => slot(a) != slot(b))
               & slot(a) != slot(this)
                              Towards Practical Debugging of Wireless Sensor Networks
```

### Predicate Evaluation — DSL

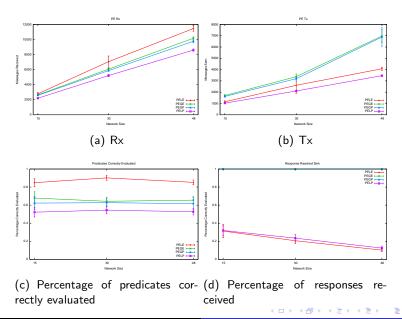
```
//TARGETING all
                                    //slot(a[*1])
                   INEQ
//FUNC 0 AS addr JZ end2
                                    end2: VIFAFC 1 255 1
                                    //slot(this)
//FUNC 1 AS slot //addr(a[*1])
//STORING 1 IN ohn VIFAFC 1 255 0
                                    THISC 1
IVAR 1
                   //addr(b[*2])
                                    INEQ
IPUSH 1
                   VIFAFC 2 255 0
                                    AND
TPUSH 0
                   INEQ
                                    AND
TSTORE 1
                   //slot(a[*1])
                                    VIINC 1
start1: ALEN 255
                VTFAFC 1 255 1
                                    JMP start1
INEQ
                   //slot(b[*2])
                                    end1: HALT
J7. end1
                   VIFAFC 2 255 1 //VD 1 = 255
TVAR. 2
                   INEQ
IPUSH 1
                   IMPLIES
IPUSH 0
                   AND
ISTORE 2
                   VIINC 2
start2: ALEN 255
                   JMP start2
```

## Results

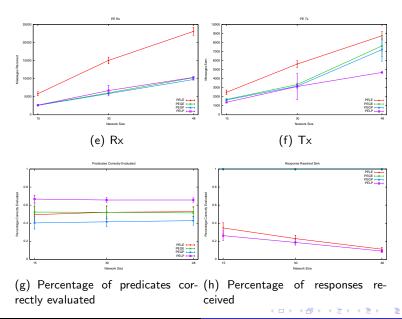
## Results Methodology

- Run and measure energy usage of TDMA algorithm
- Measure energy cost of predicate evaluation algorithm
  - Checking for slot collisions
  - Vary predicate distance (1–hop and 2–hop)
  - Vary predicate evaluation algorithm
- Network was laid out as a grid
- N, S, E, W communication possible
- 5 minutes setup time for PE, start TDMA
- 35 minutes total runtime

# Results when period=4.0 minutes using a 1-hop predicate



# Results when period=4.0 minutes using a 2-hop predicate



## Demo

### Visualisation Tool and Network Interface

#### Features:

- Creating and compiling predicates to monitor
- Deploying predicates to the WSN
- Recording history of evaluation results
- Use of serialdump-linux to interface with sink mote

#### Views:

- Predicate view
- Network view

## Conclusions

## Summary

### Developed:

- Libraries for use in Contiki (Container and Network)
- Predicate Evaluation Libraries (Global and Local)
- GUI tool to interface with network

#### Found:

- In-network, event-based evaluation suitable for "small" predicates
- In-network, periodic-based evaluation suitable for "large" predicate where evaluation is concerned
- Global evaluation more suitable for "large" predicates where energy is concerned



### Future Work

- Improve memory management
- Improve C containers developed
- Stateful predicates
- Handle mote mobility
- Improve failure response message deliver ratio

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### The End

Any Questions?