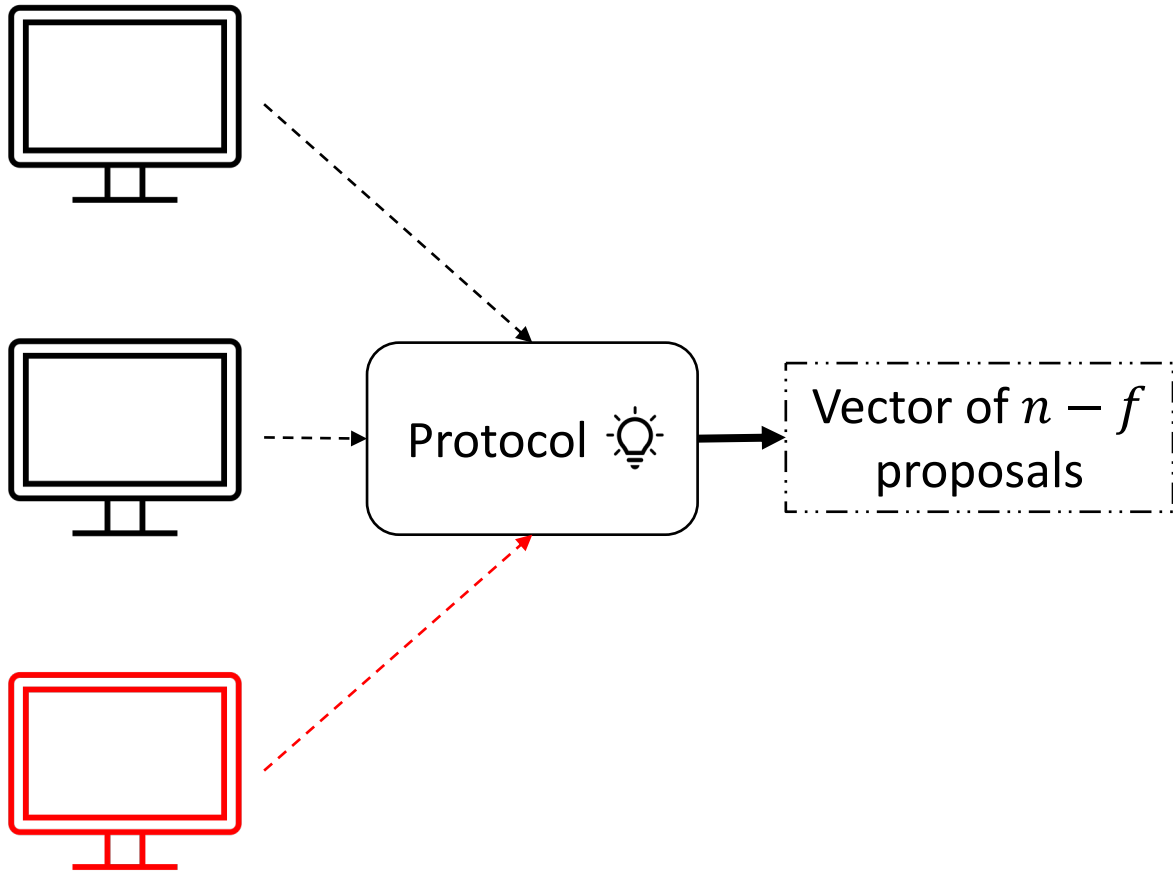


Improving the Complexity of Byzantine Vector Consensus

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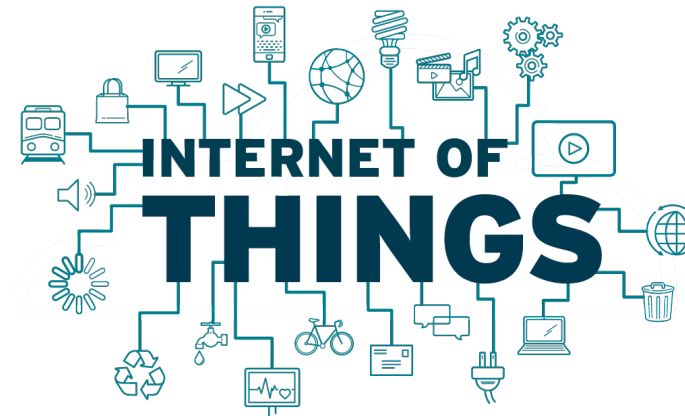
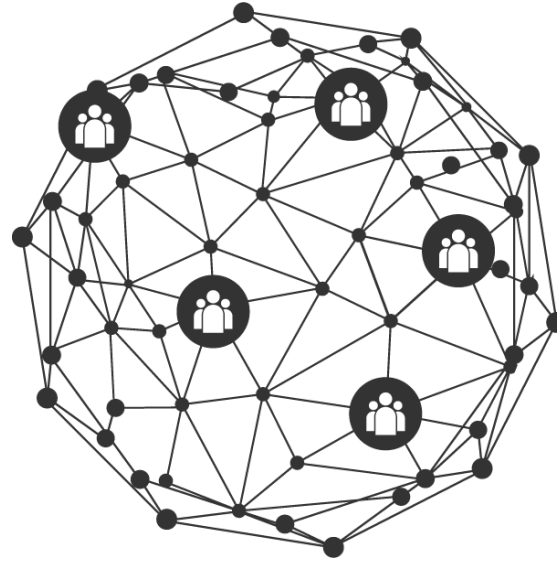
What is Byzantine Vector Consensus?



- System: n processes
 - f Byzantine
 - $n = 3f + 1$
- Decide a **vector of $(n - f)$ distinct proposals**
- Partial Synchrony: Delays are initially unbounded, until ***GST***
 - Delays bounded by δ
 - Clocks do not drift

Why do we care?

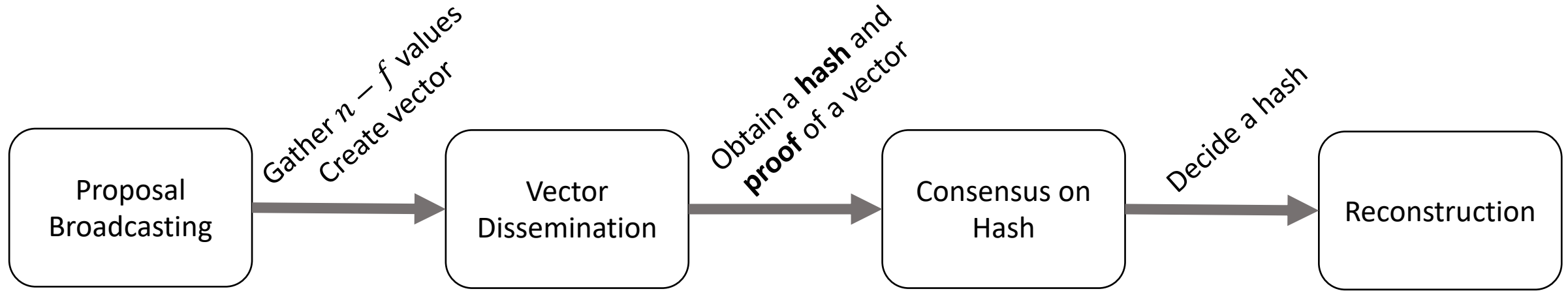
- Interesting problem for distributed computing
- Real-world applications
 - Blockchain
 - Decentralized Computing
 - IoT



Existing solutions cannot work

- Byzantine Consensus is $O(n^2)$: QUAD
 - For vector consensus: $O(n^3)$ communication complexity, $O(f)$ latency
 - Message size is $O(n)$
- Alternative: Work with hashes instead of vectors
 - $O(n^2 \log(n))$ communication complexity, $O(n^f)$ latency
- Can we do better?
 - We achieved **$O(n^2 \sqrt{n})$ communication complexity** and **$O(n \sqrt{n})$ latency!**

Structure of Byzantine Vector Consensus



Communication Complexity

$O(n^2)$

???

$O(n^2)$

$O(n^2 \log(n))$

Latency

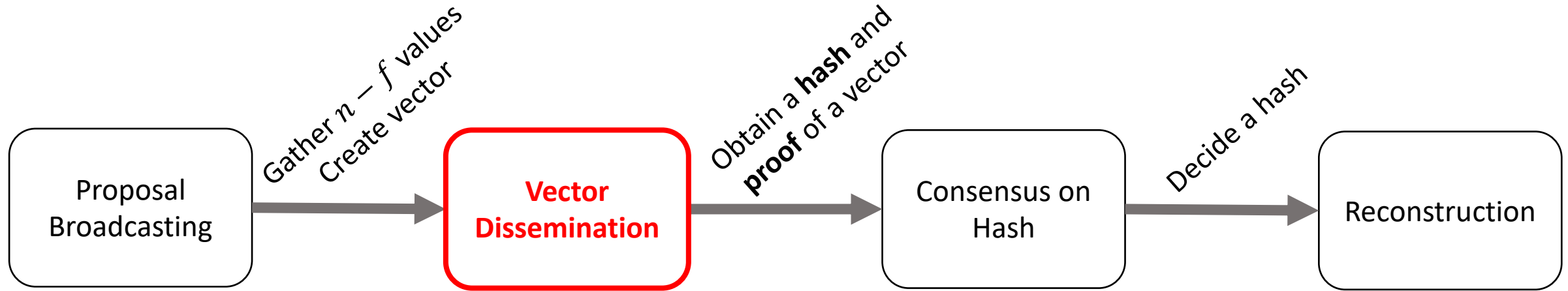
$O(1)$

???

$O(f)$

$O(1)$

Structure of Byzantine Vector Consensus



Communication Complexity

$O(n^2)$

$O(n^2\sqrt{n})$

$O(n^2)$

$O(n^2 \log(n))$

Latency

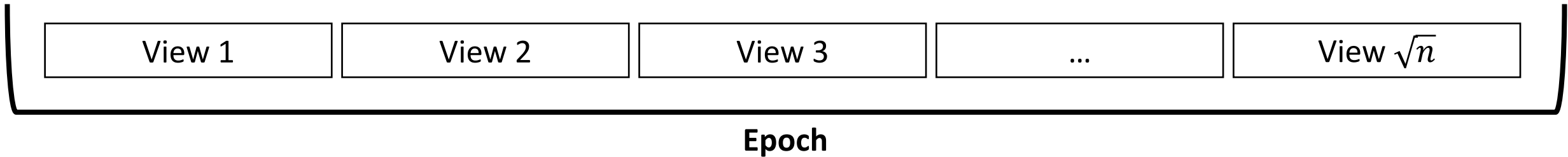
$O(1)$

$O(n\sqrt{n})$

$O(f)$

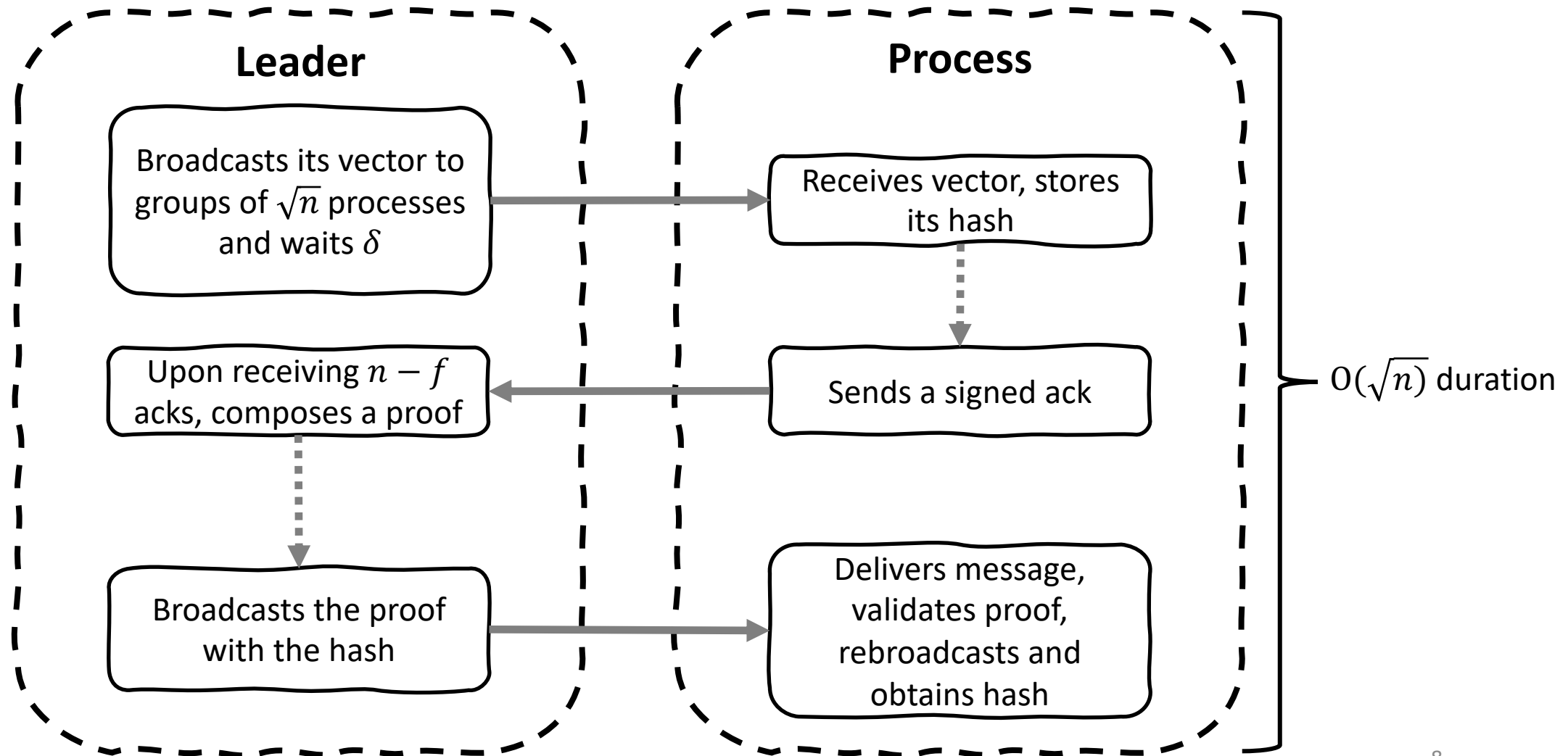
$O(1)$

Leader-Based Vector Dissemination



- Each view has a leader process that tries to disseminate their vector (**View Core**)
- Hash obtained when all processes overlap in a view with correct leader
 - Processes advance over views until they synchronize (**View Synchronizer**)

View Core

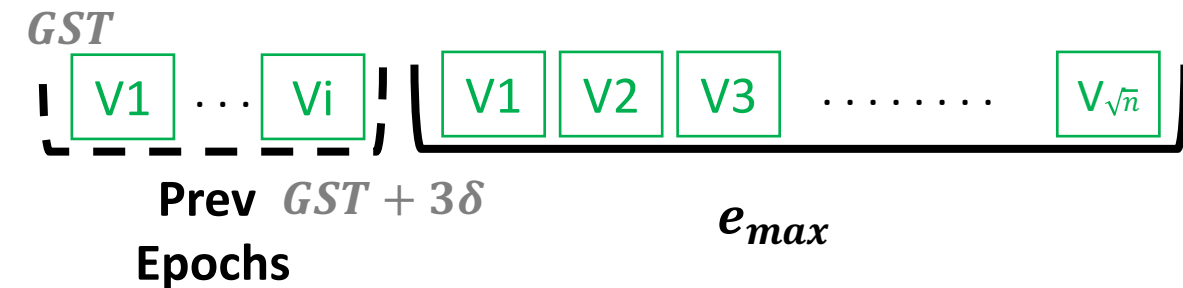


View Synchronizer

- RareSync: Processes overlap for every view of each epoch entered after GST
 - \sqrt{n} views
 - $f + 1$ views (worst case) to reach a correct leader
- Processes communicate only at start/end of epoch
 - Local clocks to advance inside epochs
 - Wait δ before entering new epoch
- **Synchronization time:** Processes synchronized, correct leader: Obtain hash!

Communication Complexity

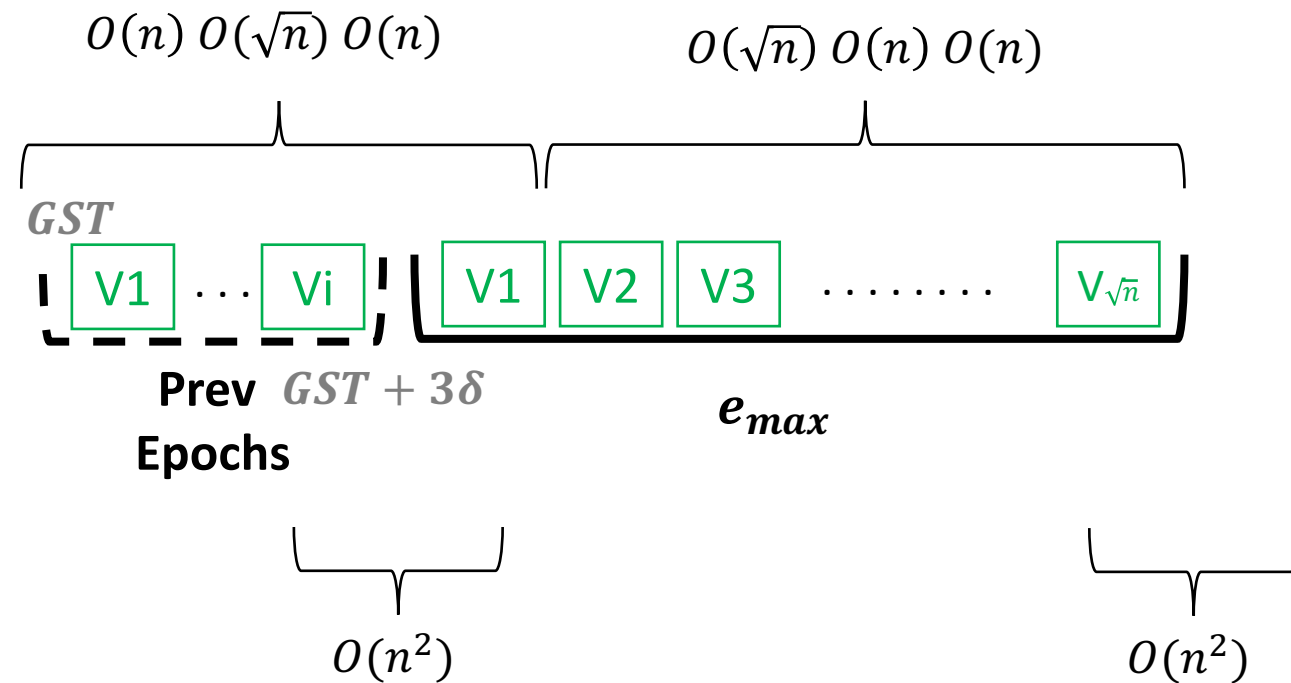
- e_{max} : Epoch of most advanced process at GST .



- Worst case:
 - At GST , all correct processes are leaders
 - Most advanced process in e_{max} : All correct processes will enter e_{max} by $GST + 3\delta$
 - All leaders correct but processes misaligned

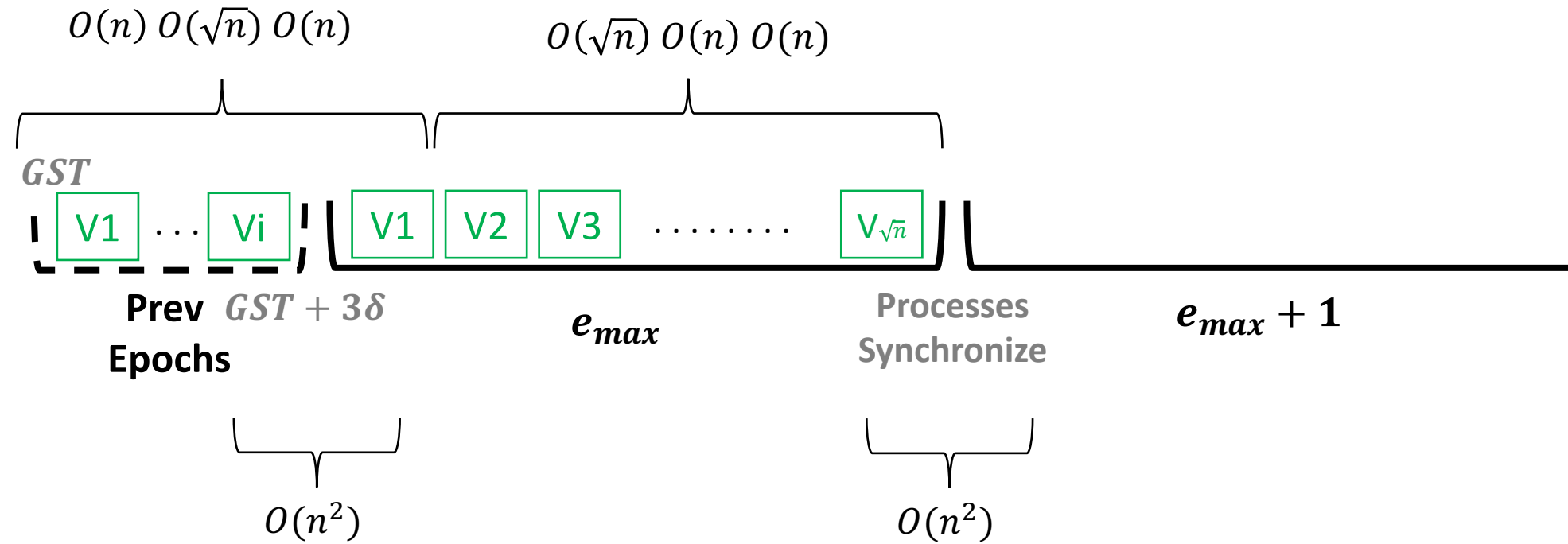
Communication Complexity

- e_{max} : Epoch of most advanced process at GST .

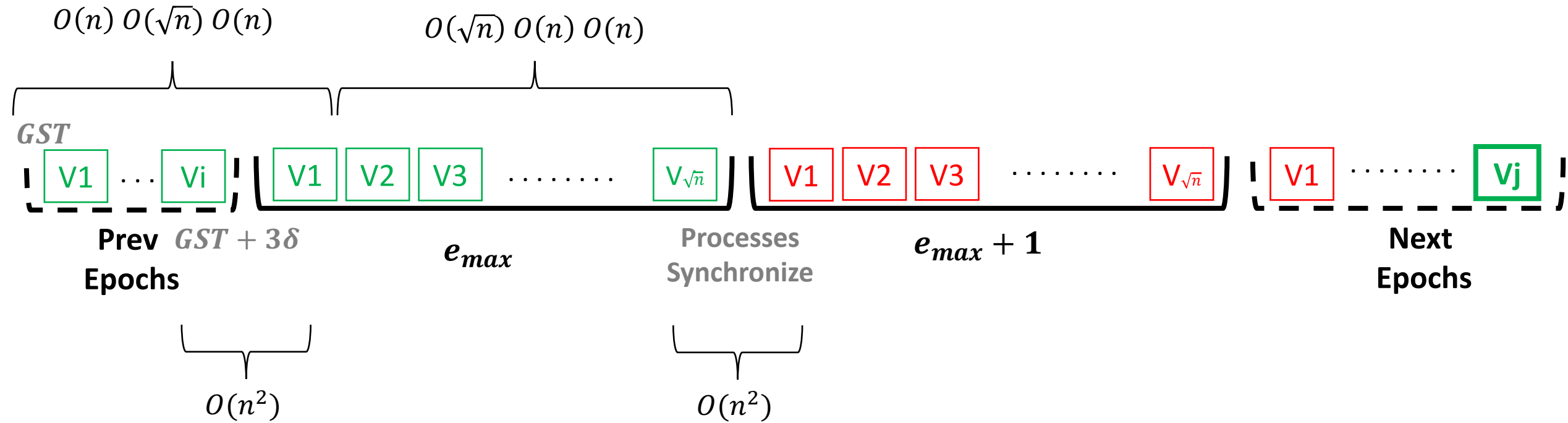


- Worst case:
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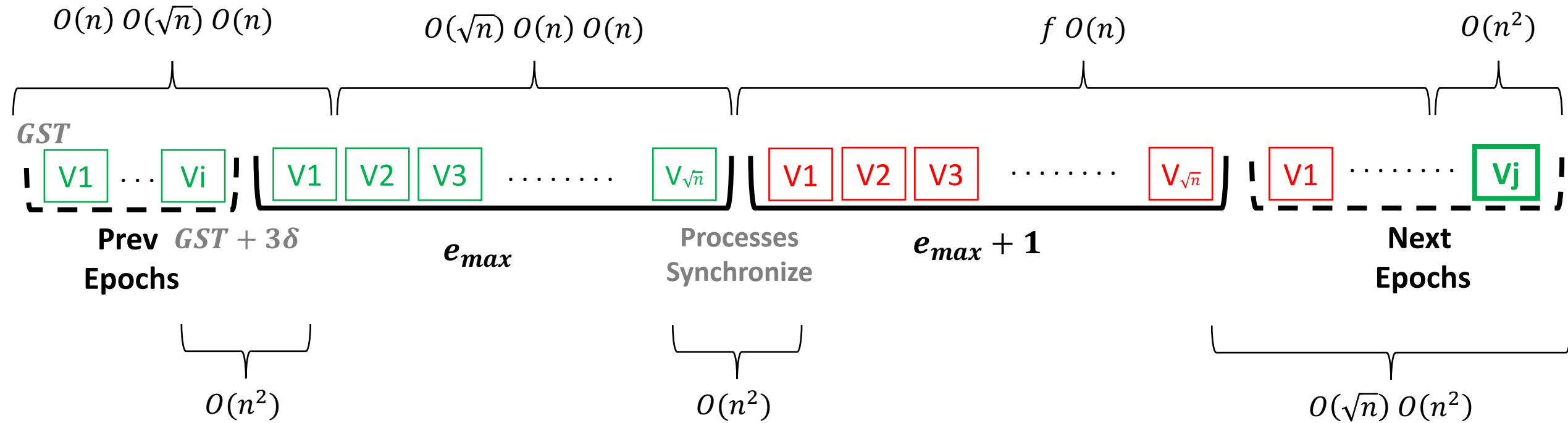
Communication Complexity



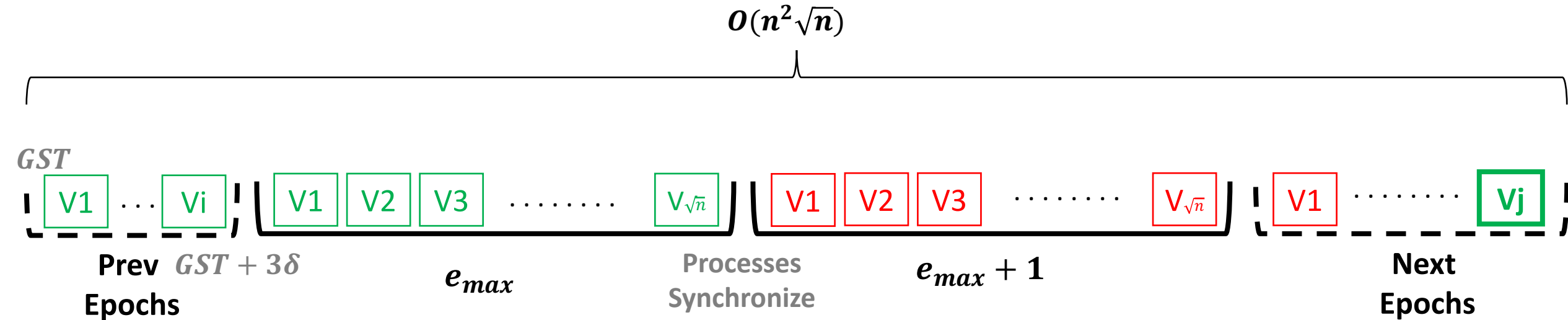
Communication Complexity



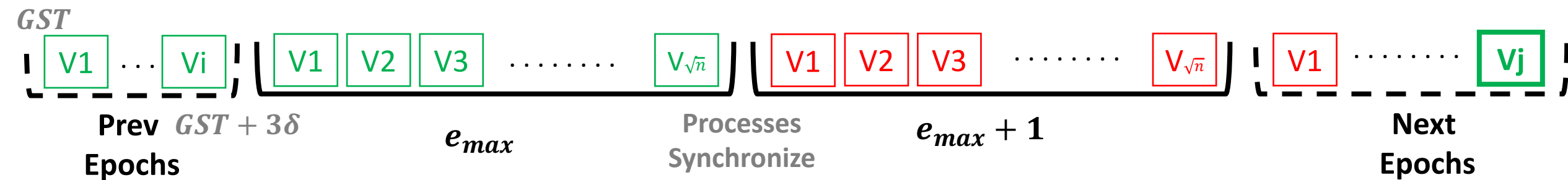
Communication Complexity



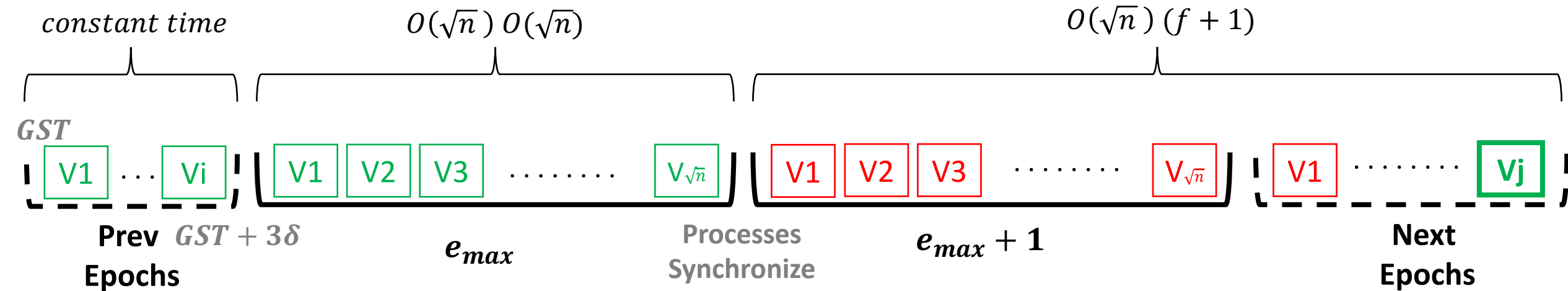
Communication Complexity



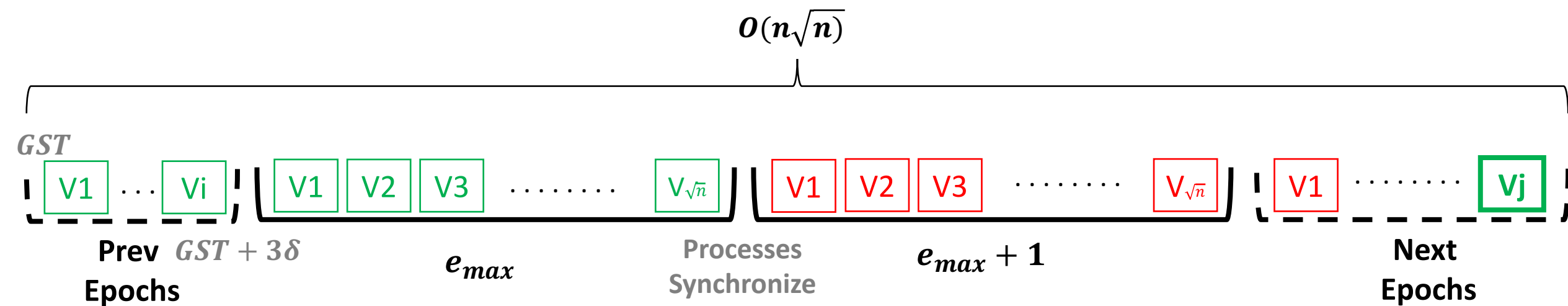
Latency



Latency

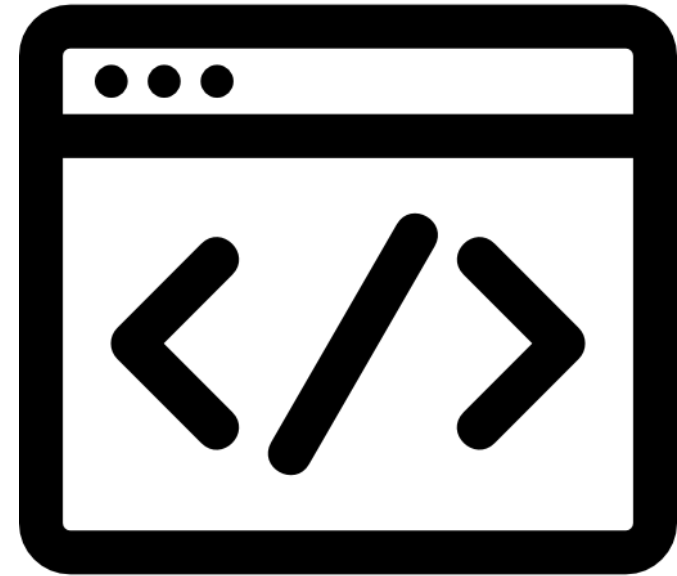


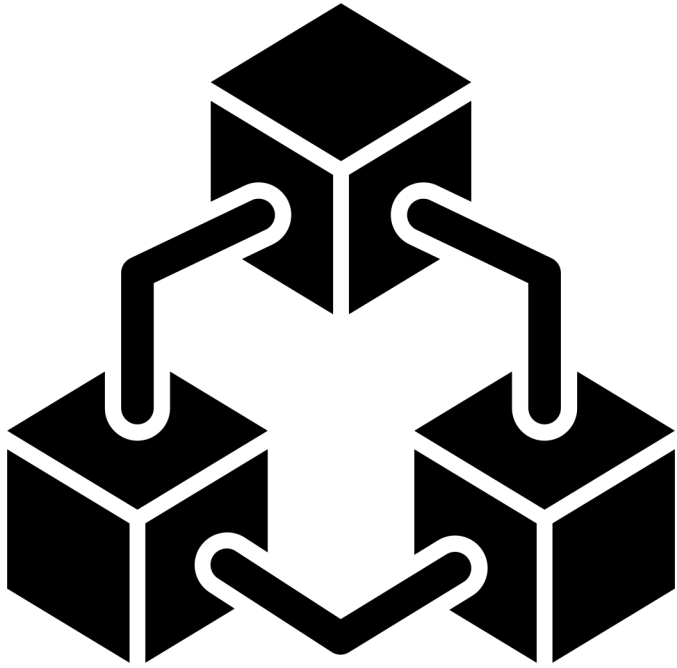
Latency



Summary

- Byzantine Vector Consensus is a difficult problem
 - With many real-world applications!
- Previous solutions could not work
 - Either cubic communication complexity or exponential latency
- **Leader-Based Vector Dissemination**
 - Achieves $O(n^2\sqrt{n})$ communication complexity and $O(n\sqrt{n})$ latency!





Thank you!



Questions?