

Trees

| | |
|---------------|-----------|
| ☰ Chapter No. | 14 |
| ▼ Status | Completed |

- A **tree** is a non-linear, hierarchical data structure

▼ Tree Terminology

▼ Node

- A node is represented by an oval and each node stores a value

▼ Edge

- An edge is represented by an arrow and each edge connects two nodes to show the relationship between them

▼ Root

- The root is the only node that has no incoming edge

▼ Parent, Child & Sibling

- A parent node is connected to its child node through an edge
- The children nodes of the same parent node are sibling nodes

▼ Subtree (Branch)

- A subtree or branch is a set of nodes and edges comprising a parent node and all the descendant nodes of that parent node

▼ Leaf

- A leaf node is a node that has no children

▼ Path

- A path is an ordered list of nodes that are connected by edges

▼ Level

- The level of a node refers to its depth in a tree
- By definition, the level of the root node is 0

▼ Height

- The height of the tree refers to how deep the entire tree is

▼ Size

- The size of the tree refers to the total number of nodes present

▼ Binary Tree

- A **binary tree** is a tree where each node has at most two children
- The two children of a node are commonly referred to as the **left child** and the **right child** of the node

▼ Tree Traversal

▼ Pre-Order Traversal

- **Root → Left Subtree → Right Subtree**

▼ In-Order Traversal

- **Left Subtree → Root → Right Subtree**

▼ Post-Order Traversal

- **Left Subtree → Right Subtree → Root**

▼ Binary Search Tree

- A **binary search tree** is a binary tree in which all values smaller than a node are stored in that node's left subtree and all values larger than a node are stored in that node's right subtree

▼ Common Binary Search Tree Operations

▼ create

- Create a new binary search tree with a root node of a given value without any children

▼ edit

- Change the value of a node in the binary search tree

▼ search

- Search a given value in the binary search tree
- ▼ insert
 - Insert a given value into the binary search tree
- ▼ delete
 - Delete a node of a given value in the binary search tree and possibly rearrange the tree
 - If a node to be deleted has no children (leaf node), there is no need to shift any of the remaining nodes
 - If the node to be deleted has one child, the child node will take the deleted node's place
 - ▼ If the node to be deleted has two children, there are two possibilities
 - The largest node in the left subtree takes the deleted node's place
 - The smallest node in the right subtree takes the deleted node's place
- In-order traversal of a binary search tree returns the values sorted in increasing order
- ▼ Rebalancing Binary Search Trees
 - A binary search tree is unbalanced when one side has many more element than the other side
 - This makes binary search of the binary search tree inefficient because the time complexity of binary search becomes almost $O(n)$ instead of $O(\lg n)$
 - In this case, the binary search tree needs to be re-created or rebalanced with the same items to make it balanced
 - This distributes the items more evenly between the left and right subtrees, ensuring that the time complexity of binary search is $O(\lg n)$
- ▼ Advantages of Binary Search Trees

- [Binary search](#) can be used to search for items instead of less efficient searching algorithms, such as linear search
- Insertion and deletion of items involves the [updating of pointers](#), similar to linked lists, [instead of rearranging elements](#)