

The Network Layer

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- Network Layer Overview
 - Transport segment from sending to receiving host
 - * Sender: encapsulates segments into packets, passes to link layer
 - * Receiver: extracts segments from packets and delivers segments to transport layer protocol
- Network Layer Functions
 - Forwarding: move packets from router's input link to appropriate router's output link
 - Routing: determine route taken by packets from source to destination
 - * Routing Algorithms
 - Analogy: Taking a Trip
 - * Forwarding: process of getting through single intersection
 - * Routing: process of planning trip from source to destination
- Data Plane
 - Local, per-router function
 - Determines how packet arriving on router input port is forwarded to router output port
- Control Plane
 - Network-wide logic
 - Determines how packet is routed among routers along end-end path from source host to destination host
 - Two control-plane approaches

- * Traditional routing algorithms: implemented in routers
- * Software-Defined Networking (SDN): implemented in (remote) servers
- Traditional Control Plane Algorithms
 - Individual routing algorithm components in each and every router interact in the control plane
- SDN Control Plane
 - Remote controller interacts with local Control Agents (CAs) to compute, install forwarding tables in routers
- Network Layer Service Model
 - A network layer service model defines the characteristics of end-to-end transport of packets between sending and receiving hosts
 - Examples of possible services (this is only a partial list, there are countless variants):
 - * Guaranteed delivery
 - * Guaranteed delivery with bounded delay
 - * In-order packet delivery
 - * Guaranteed minimum transmission rate
 - * Security
 - Services provided by the network layer: two main options
 1. Connection-oriented service
 - * A path from source all the way to destination must be established before any data packets can be sent
 - This connection is called a Virtual Circuit (VC)
 - The network is called a virtual-circuit network
 - Each VC requires router table space and reservation of resources
 - * Designed to provide some quality of service (QoS) (*i.e.* maximum delay guarantees, minimum losses, minimum throughput guarantees, etc.)
 - * Example: Asynchronous Transfer Mode (ATM) → popular in the 90s early 200, being replaced by all-IP architectures
 2. Connectionless service
 - * Best-effort service
 - * Packets are injected into the network individually and routed independently of each other
 - * No advance setup is needed
 - * No error or flow service functionalities provided

- The transport layer might do something end-to-end
 - The link layer might do something at the link level
 - * For example, IP (internet protocol)
- Reflections on Best-Effort Service
 - Simplicity of mechanism has allowed Internet to be widely deployed and adopted
 - Sufficient provisioning of capacity allows performance of real-time applications (*e.g.* interactive voice, video) to be “good enough” for “most of the time”
 - Replicated, application-layer distributed services (data centers, content distribution networks) connecting close to clients’ networks, allow services to be provided from multiple locations
 - Congestion control at the transport layer of “elastic” services helps
- Input Ports
 - Decentralized Switching:
 - * Using header field values, lookup output port using forwarding table in input port memory (“match plus action”)
 - Destination-based forwarding: forward based only on destination IP address (traditional)
 - Generalized forwarding: forward based on any set of header field values
 - Input port queueing: if packets arrive faster than forwarding rate into switch fabric
- Input Port Queueing
 - If switch fabric slower than input ports combined → queueing may occur at input queues
 - * Queueing delay and loss due to input buffer overflow
 - Head-of-the-Line (HOL) blocking: queued packet at front of queue prevents others in queue from moving forward
- Output Ports
 - Buffering required when packets arrive from fabric faster than link transmission rate
 - Drop policy: which packets to drop if no free buffers?
 - Scheduling discipline chooses among queued packets for next transmission
 - * FCFS (First Come, First Served), priority, ...
- The Internet Protocol

- The glue that holds the whole Internet together (data plane)
 - * Designed with internetworking in mind
- Provides a best-effort (no guarantee) way to transport IP packets (aka datagrams) from source to destination
 - * Without regard to whether these machines are on the same network or whether there are other networks between them
- There are two versions of IP in use today
 - * IPv4 (IP version 4)
 - The first “major version” of IP and currently the dominant protocol of the Internet
 - * IPv6
- IP Fragmentation
 - Network links have MTU (maximum transmission unit)
 - * MTU: largest possible payload in link-level frame → maximum IP packet size
 - * Different link types, different MTUs
 - Problem: IP packet larger than MTU of output link
 - * Solution: Fragmentation?
 - Typically, IPv6 does not allow fragmentation
 - Typically, TCP does not allow fragmentation
- IP Alternative to Fragmentation
 - If fragmentation is not allowed → “path MTU discovery”
 - Path MTU Discovery
 - * Each IPv4 packet is sent with its header bits set to indicate that fragmentation is not allowed to be performed (flag DF=1)
 - * Added start-up delay
 - * The transport layer can learn about the MTU to adapt the Maximum Segment Size (MSS)
- IP Addressing: Introduction
 - IPv4 Address: 32-bit identifier associated with each host or router interface
 - Interface: connection between host/router and physical link
 - * Routers typically have multiple interfaces
 - * Host typically has one or two interfaces (e.g, wired Ethernet, wireless 802.11)
- Subnets

- Device interfaces that can physically reach each other without passing through an intervening router
- IP Addresses have structure:
 - * Network portion (aka subnet portion): high order bits
 - Devices in same subnet have common network portion
 - * Host portion: remaining low order bits
- IP Addressing in Subnets: CIDR
 - CIDR: Classless Inter Domain Routing (pronounced “cider”)
 - * Network portion (aka prefix) of address of arbitrary length
 - * Address format (by convention): A.B.C.D.X, where X is the number of bits in the network portion of the address
 - Network address (subnet address): network portion and 0s in the host portion/x
 - Subnet mask: binary mask of 1s in the subnet portion and 0s in the host portion → X
 - * The subnet mask can be ANDed with an IP address to obtain the network address
 - Recipe for identifying subnets
 - * Detach each interface from its host or router, creating “islands” of isolated networks
 - * Each isolated network is a subnet
- Longest Prefix Matching
 - When looking for forwarding table entry for a given destination address, use longest address prefix that matches destination address.
- Forwarding in Access Networks
 - Forwarding tables in routers of an access network have an entry for their subnets
 - When a datagram reaches a router in an access network, it looks at the destination address of the datagram, and checks which subnet inside the network it belongs to. How?
 - * AND the destination address with the mask for each subnet entry in the table
 - * Check to see if the result is the prefix in the entry
- Forwarding in the Network Core
 - Routers in ISPs and backbones in the middle of the internet must know which way to go to get to every network and no simple default will work

- * This can make for a very large table
 - Routers must perform a lookup in this table for every datagram they forward
- Hierarchical Addressing: Route Aggregation
 - Hierarchical addressing allows efficient advertisement for routing information
- How Are IP Addresses Assigned?
 - Hard-coded by system administrator → fixed IP address
 - DHCP: Dynamic Host Configuration Protocol
 - * Can renew its lease on address in use
 - * Allows reuse of addresses (only hold address while connected/on)
 - * Support for mobile users who join/leave network
- DHCP: More than IP Addresses
 - DHCP can return more than just allocated IP addresses on a subnet:
 - * Address of first-hop router for client
 - * Name and IP address of local DNS server
 - * Subnet mask (indicating network versus host portion of address)
- NAT: Network Address Translation
 - All devices in local network have 32-bit IP addresses in a “private” IP address space that can only be used in local networks
 - “Private” IP address space corresponds to prefixes: 10.0.0.0/8, 172.16.0.0/12 and 192.168.0.0/16
 - * Defined by IANA (Internet Assigned Numbers Authority) → department of ICANN
 - Advantages:
 - * Private IP addresses can be reused in different private networks
 - * Just one IP address needed from provider ISP for all devices
 - Implementation: NAT router (or NAT box)
 - Outgoing datagrams: replace (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
 - * Remote clients/servers will respond using (NAT IP address, new port #) as destination address
 - Store in NAT translation table every (source IP address, port #) to (NAT IP address, new port) translation pair

- Incoming datagrams: replace (NAT IP address, new port #) in destination fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table
- NAT has been controversial:
 - * Routers “should” only process up to layer 3
 - * Address “shortage” should be solved by IPv6
 - * Violates end-to-end argument (port number manipulation by network-layer device)
 - * NAT traversal: what if client wants to connect to server behind NAT?
- But NAT is here to stay:
 - * Extensively used in home and institutional networks, 4G/5G cellular networks
- IPv6: Motivation
 - Initial motivation: 32-bit IPv4 address space would be completely allocated
 - Additional motivation:
 - * Speed processing/forwarding: 40-byte fixed length header
 - Extension headers: optional headers can be added after the fixed IPv6 header
 - * Enable different network-layer treatment of “flows”
- Transition from IPv4 to IPv6
 - Not all routers can be upgraded simultaneously
 - * No “flag delays”
 - * How will network operate with mixed IPv4 and IPv6 routers?
 - Tunneling: Packet within a packet
 - * IPv6 datagram carried as payload in IPv4 datagram among IPv4 routers (“packet within a packet”)
 - * Tunneling used extensively in other contexts (4G/5G)
- Flow Table Abstraction
 - Flow: defined by header fields
 - Generalized Forwarding
 - * Match: pattern values in packet header fields
 - * Actions: for matched packet: drop, forward, modify matched packet or send matched packet to controller
 - * Priority: disambiguate overlapping packets
 - * Counters: # bytes and packets

- OpenFlow Abstraction

- Match and Action: abstraction unifies different kinds of devices
- Router:
 - * Match: longest destination IP prefix
 - * Action: forward out a link
- Firewall:
 - * Match: IP addresses and TCP/UDP port numbers
 - * Action: permit or deny
- Switch
 - * Match: destination MAC address (link layer address)
 - * Action: Forward or flood
- NAT
 - * Match: IP address and port
 - * Action: rewrite address and port

- Middlebox

- Any intermediary box performing functions apart from normal, standard functions of an IP router
- Initially: proprietary (closed) hardware solutions
- Move towards “whitebox” hardware implementing open API
 - * Move away from proprietary hardware solutions
 - * Programmable local actions via match and action
 - * Move towards innovation/differentiation in software
- SDN: (logically) centralized control and configuration management often in private/public cloud
 - * Network Function Virtualization (NFV): programmable services over white box networking, computation, and storage
 - Allows for programmable network devices