Automata CW2

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Q1

a) L(G) over $\Sigma = \{0,1\} \ Y0 ->$

1Y0 | Y1

Y1-> 0 Y0 | 1 Y1 | Y2

 $Y2 -> 0 Y2 |1 Y1| \epsilon$

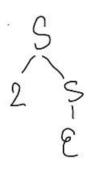
b) i) For 2:

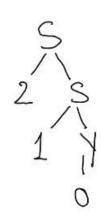
For 210

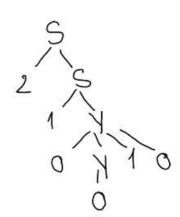
For 210010

S-> 2S->2
$$\epsilon$$
 -> 2

S-> 2S-> 21Y -> 210Y10 -> 210010







ii) The grammar is not in Chomsky normal form because the rules of this automaton do not respect this rule: CNF grammar must not contain it's start variable in the right hand side of any grammar rule.

Thus, S-> 2S is not respecting the rule and as a consequence the grammar loses the Chomsky properly.

iii)

Since S appears in a RHS rule, we first apply the 'start' routine:

$$SO \rightarrow S$$
 , $SO = start variable$

 $S \rightarrow 1Y \mid 2S \mid \epsilon$

Y -> S 1 | 0Y10 | 0

There are several rules on the RHS with length > 2, so we apply the BIN routine:

S0 -> S

```
S \rightarrow 1 Y \mid 2 S \mid \epsilon
```

X -> Y X1

X1 -> 10

We apply the DEL routine to remove the S -> ϵ rule:

 $SO \rightarrow S \mid \epsilon$

S->1Y|2S|2

Y -> S 1 | 0 X | 0 | 1

X -> Y X1

X1-> 10

Next we apply the UNIT routine to remove the unit rules: S

$$SO \rightarrow 1 Y | 2 S | 2 | \epsilon$$

X -> Y X1

X1-> 10

Then we apply TERM to remove terminals that appear in binary rules: U1-> 1, U0-> 0, U2-> 2

SO -> U1 Y | U2 S | 2 | ϵ

S -> U1 Y | U2 S | 2

Y -> S U1 | U0 X | 0 | 1

X -> Y X1

X1-> U1 U0

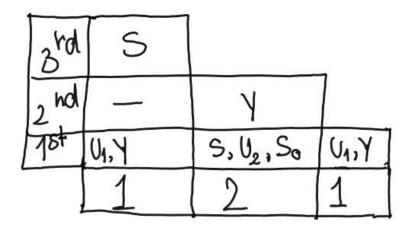
U1-> 1

U0-> 0

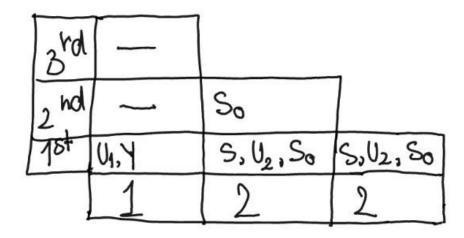
U2-> 2

IV) G: 121, 122

For 121: S0-> U1 Y ->U1 S U1->U1 2 U1-> 121



For 122: S0 -> U1 Y , AND THE REST DOESN'T WORK because Y doesn't lead to any word containing '2'.



c) i)

 $\Sigma = \{0,1,2,3\}$

YES	ΝΟ (Σ*)		
132031	132033		
31202013	ε		
13120131	12032031		

ii) for 132031

S->1S1->13S31->13Y31->132031

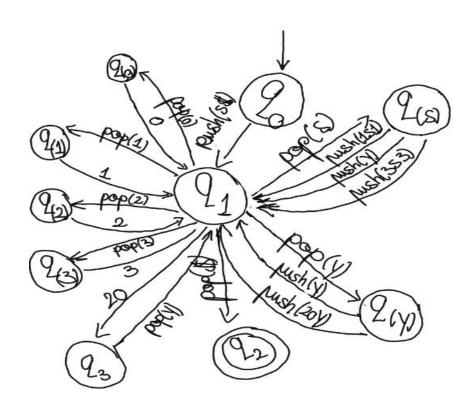
For 31202013

S->3S3->31S13->31Y13->3120Y13

for 13120131

iii) L(G) is the language of all words (of even length) following this rule: the second half is the first half in a reverse order (it's like looking at the first half in a mirror) but the only problem would be that the 2's will be transformed into 0's and 2's will be replaced by 0's.

iv)



The language of the PDA is: $A=(\{1,2,3\}, \{\$,1,2,3\}, \{q0,q1,q2,q(1),q(2),q(3),q(s)\}, \delta, q0,q2\}$

With δ =({q0, push(S\$),q1}, {q1, pop(1), q(1)}, {q(1), 1, q1}, {q1, pop(2), q(2)}, {q(2), 2, q1}, {q1, pop(3), q(3)}, {q(3), 3, q1}, {q1, pop(\$), q2}, {q1, pop(\$), q(\$)}, {q(\$), push(1\$\$1), q1}, {q(\$), push(3\$\$3), q1}, {q(\$), push(\$), q1}, , {q(\$), \$\epsilon\$, q1}, {q1, pop(\$), q(\$)}, {q(\$), push(20), q1}, {q(\$), push(20\$), q1}, {q(\$), \$\epsilon\$, q1}, {q1, pop(\$), q2}).

(v) For 132031:

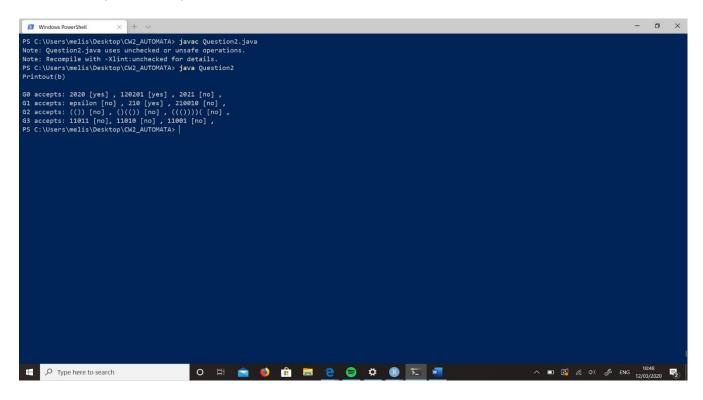
$$\begin{array}{l} (Q_{3},3) \xrightarrow{\text{push}(S\$)} (Q_{10}S\$) \xrightarrow{\text{pop}(S)} (Q_{10}S\$) \xrightarrow{\text{push}(1S1)} (Q_{11},1S1\$) \xrightarrow{\text{pop}(1)} (Q_{11},S1\$) \xrightarrow{1} (Q_{11},S1\$) \xrightarrow{\text{pop}(S)} (Q_{10},S1\$) \xrightarrow{\text{pop}(S)} (Q_{10},S1\$) \xrightarrow{\text{pop}(S)} (Q_{10},S1\$) \xrightarrow{\text{pop}(S)} (Q_{10},S1\$) \xrightarrow{\text{pop}(S)} (Q_{11},S1\$) \xrightarrow$$

For 13120131:

```
Q2
(a) (i)
public static CFG generateCFG1() {
    String[] alphabet = new String[]{ "0", "1", "2" };
    String[] vars = new String[]{ "S", "Y" };
                                                Rule[]
R = new Rule[] {
                        new Rule("S",new
String[]{"1","Y"}),
                        new Rule("S",new
String[]{"2","S"}),
                        new Rule("S",new
String[]{""}),
                    new Rule("Y",new
String[]{"S","1"}),
                         new Rule("Y",new
String[]{"0","0","Y","1","0"}),
                                    new Rule("Y",new
String[]{"0"})
    };
    return new CFG(alphabet,vars,R,"S");
  }
(ii)
public static CFG generateCFG2() {
    String[] alphabet = new String[]{ "(", ")" };
    String[] vars = new String[]{ "S" };
    Rule[] R = new Rule[] {
       new Rule("S",new String[]{""),
new Rule("S",new String[]{"(","S",")"}),
new Rule("Y",new String[]{"S","S"})
    };
```

```
return new CFG(alphabet,vars,R,"S");
  }
(iii)
 public static CFG generateCFG3() {
    String[] alphabet = new String[]{ "0", "1", "2", "3" };
    String[] vars = new String[]{ "S", "Y" };
Rule[] R = new Rule[] {
                               new
Rule("S",new String[]{"3","X"}),
new Rule("S",new String[]{"1","X"}),
new Rule("X",new String[]{""}),
new Rule("X",new String[]{"3","S"}),
new Rule("X",new String[]{"1","Y"}),
new Rule("Y",new String[]{"0","X"}),
new Rule("Y",new String[]{"3","X"}),
new Rule("Y",new String[]{""}),
    };
    return new CFG(alphabet,vars,R,"S");
  }
```

(b) The snapshot of the printout of runThem is:



```
The code for this implementation is:
public static void runThem(){
    CFG G0=generateCFG0();
    CFG G1=generateCFG1();
    CFG G2=generateCFG2();
    CFG G3=generateCFG3();
    // first CFG: G0
    String [] alphabet1=new String []{"2","0", "2","0"};
    String [] alphabet2=new String []{"1","2","0", "2","0","1"};
    String [] alphabet3=new String[]{"2","0", "2","1"};
     System.out.println("");
     System.out.print("G0"+" accepts: "+ "2020 [");
if( isAccepted(G0,alphabet1) == true)
      {
         System.out.print("yes");
      } else
         System.out.print("no");
       System.out.print("] ,");
       System.out.print(" 120201 [" );
if( isAccepted(G0,alphabet2) == true)
         System.out.print("yes");
      } else
         System.out.print("no");
       System.out.print("] ,");
```

```
System.out.print(" 2021 [" );
if( isAccepted(G0,alphabet3) == true)
         System.out.print("yes");
      } else
         System.out.print("no");
      }
       System.out.print("] ,");
    // second CFG: G1
    String [] alphabet11=new String []{ };
    String [] alphabet22=new String []{"2","1","0"};
    String [] alphabet33=new String[]{"2","1","0", "0","1","0"};
    System.out.println("");
     System.out.print("G1"+" accepts: "+ "epsilon [");
      if( isAccepted(G1,alphabet11) == true)
      {
         System.out.print("yes");
      } else
        System.out.print("no");
      }
       System.out.print("] ,");
       System.out.print(" 210 [" );
if( isAccepted(G1,alphabet22) == true)
         System.out.print("yes");
      } else
```

```
System.out.print("no");
      }
       System.out.print("] ,");
       System.out.print(" 210010 [" );
if( isAccepted(G1,alphabet33) == true)
      {
         System.out.print("yes");
      } else
         System.out.print("no");
       System.out.print("] ,");
       // 3rd CFG :G2
   System.out.println("");
     System.out.print("G2"+" accepts: "+ "(()) [");
      if( isAccepted(G2,new String[]{"(","(", ")",")"}) == true)
      {
         System.out.print("yes");
      } else
         System.out.print("no");
      }
       System.out.print("] ,");
       System.out.print(" ()(()) [" );
      if( isAccepted(G2,new String[]{"(",")","(", "(",")",")"}) == true)
         System.out.print("yes");
      } else
      {
         System.out.print("no");
```

```
}
       System.out.print("] ,");
       System.out.print(" ((())))( [" );
      if( isAccepted(G2,new String[]{"(","(", "(",")", ")",")",")","("}) == true)
      {
         System.out.print("yes");
      } else
         System.out.print("no");
       System.out.print("] ,");
       //4th CFG: G3
   String [] alphabet1111=new String[]{"1","1","0","1","1"};
   String [] alphabet2222=new String[]{"1","1","0","1","1"};
   String [] alphabet3333=new String[]{"1","1","0", "0","1"};
    System.out.println("");
     System.out.print("G3"+" accepts: "+ "11011 [");
if( isAccepted(G1,alphabet1111) == true)
      {
         System.out.print("yes");
      } else
         System.out.print("no");
       System.out.print("],");
       System.out.print(" 11010 [" );
if( isAccepted(G1,alphabet2222) == true)
         System.out.print("yes");
```

```
} else
{
    System.out.print("no");
}
System.out.print("],");

System.out.print(" 11001 [" );
if( isAccepted(G1,alphabet3333) == true)
{
    System.out.print("yes");
} else
{
    System.out.print("no");
}
System.out.print("],");
}
```

(c) I implemented my code following these rules:

If its rules have one of these forms then the CFG is not in CFN and the method returns false immediately after it finds one of these rules applied:

- **O** a rule $X \rightarrow \varepsilon$ where X is not the start variable.
- **O** a rule $X \rightarrow w$ where w has length greater than 1 and contains some x from the alphabet.
- **O** a rule $X \rightarrow Y$ where Y is a variable
- **O** a rule $X \rightarrow w$ where w has length greater than 2
- **O** a rule $X \rightarrow (var) S (var)$ where S is the start variable

Otherwise, it returns true at the end of the program which means the CFG is in CFN.

The code of my implementation is:

```
return false;
           }
       for(int i = 0; i < G.rules[j].rhs.length; i++)</pre>
      {
       //a rule X \rightarrow ... S ... where S is the start variable
       String s = G.rules[j].rhs[i];
           if(s.equals(G.startVar)) {
              return false;
           }
          if(G.rules[j].rhs.length > 2) {
          //a rule X \rightarrow w where w has length greater than 2
              return false;
          }
          if(G.rules[j].rhs.length > 1 && G.checkLetter(s)) {
           // a rule X \rightarrow w where w has length greater than 1 and contains some x from \Sigma.
                     return false;
          }
           if(G.rules[j].rhs.length == 1 && G.checkVar(s)) {
          // a rule X \rightarrow Y where Y is a variable
              return false;
             }
           }
       }
     return true;
The printout of my code can be found in the photos above:
```

if(!G.rules[j].var.equals(G.startVar) && G.rules[j].rhs.length==0) {

