

TREES AND GRAPHS

DAVIDE YI XIAN HU

EMAIL: DAVIDEYI.HU@POLIMI.IT



POLITECNICO
MILANO 1863



INTRO PRESENTATION

DAVIDE YI XIAN HU



PH.D. STUDENT INFORMATION TECHNOLOGY



MAIL: DAVIDEYI.HU@POLIMI.IT



RESEARCH INTERESTS:



DEEP LEARNING TESTING



COMPUTER VISION APPLICATIONS



POLITECNICO
MILANO 1863




□ TREES

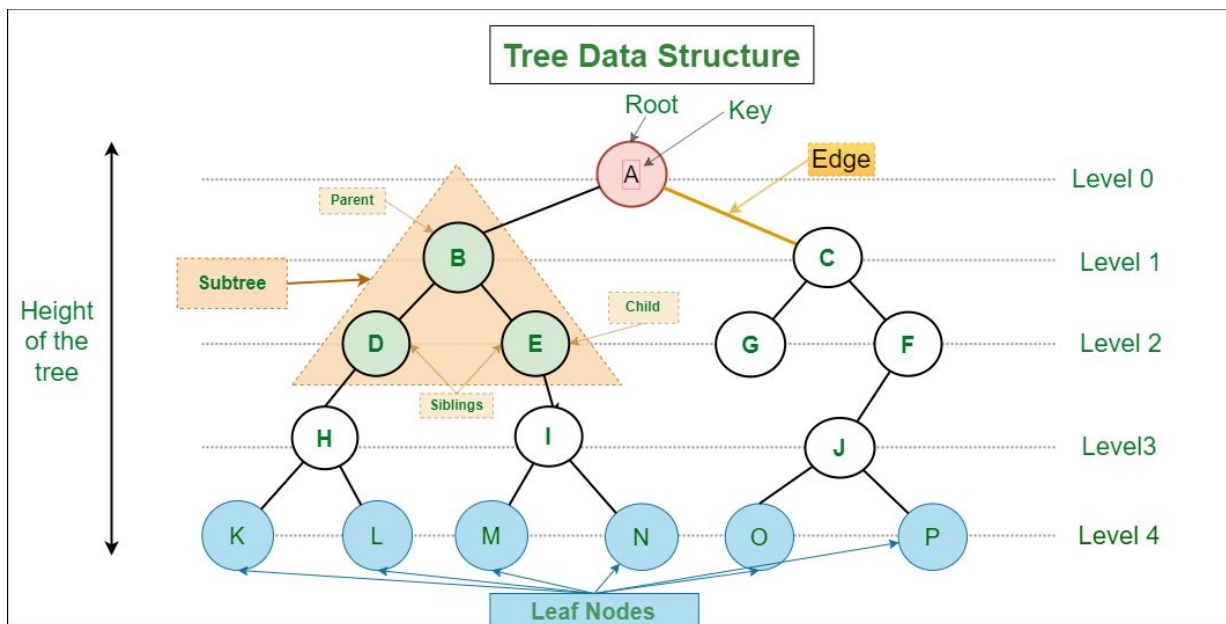


TREE DEFINITION

TREE DATA STRUCTURE

- 👉 A **TREE**  IS A **HIERARCHICAL** DATA STRUCTURE WITH A SET OF NODES
- 👉 EACH NODE CAN HAVE **MANY CHILDREN**
- 👉 EACH NODE HAVE **EXACTLY ONE PARENT**
(EXCEPT FOR THE ROOT NODE)
- 👉 **No CYCLES/LOOPS**

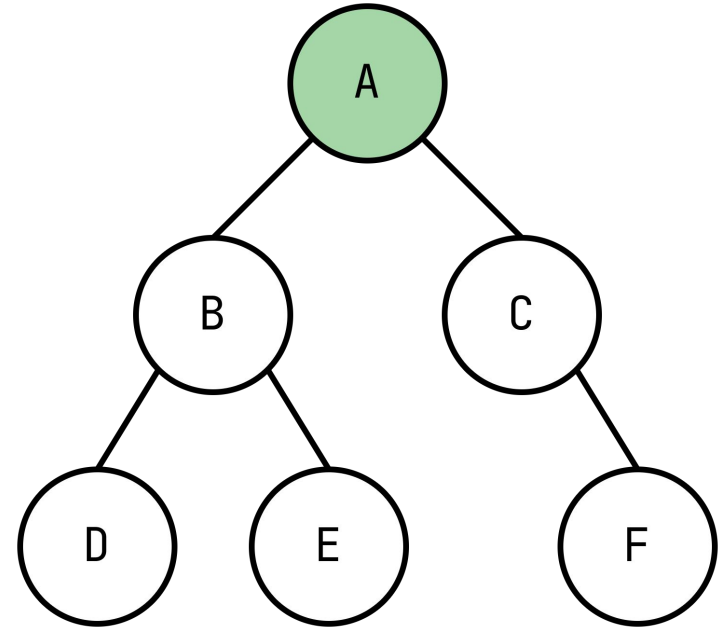
EXAMPLE OF TREE



BINARY TREE

👉 ONE OF THE **MOST COMMONLY**
USED TYPE OF TREES

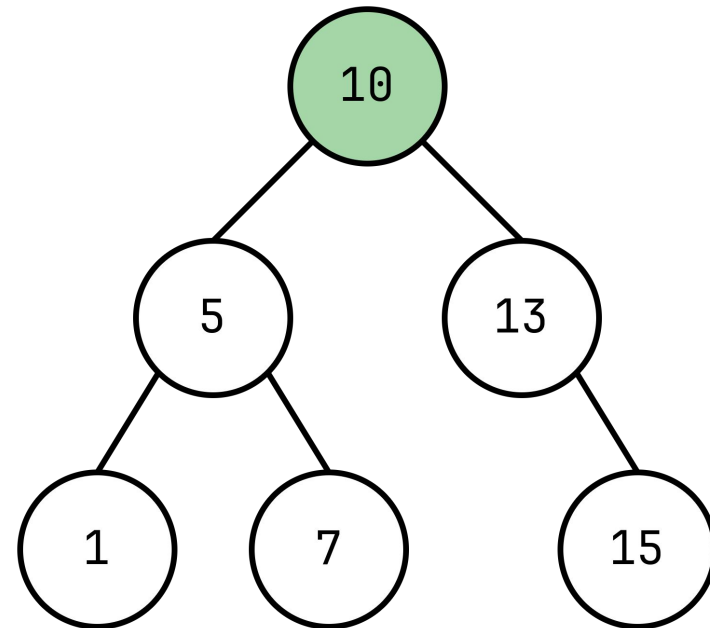
👉 EACH NODE CAN HAVE
AT MOST **TWO CHILDREN**
(K-ARY TREE WITH $K=2$)



BINARY SEARCH TREE

THE **KEY** OF A NODE IS:

- ➡ GREATER THAN THE KEYS IN THE LEFT SUBTREE.
- ➡ LESS THAN THE KEYS IN THE RIGHT SUBTREE.



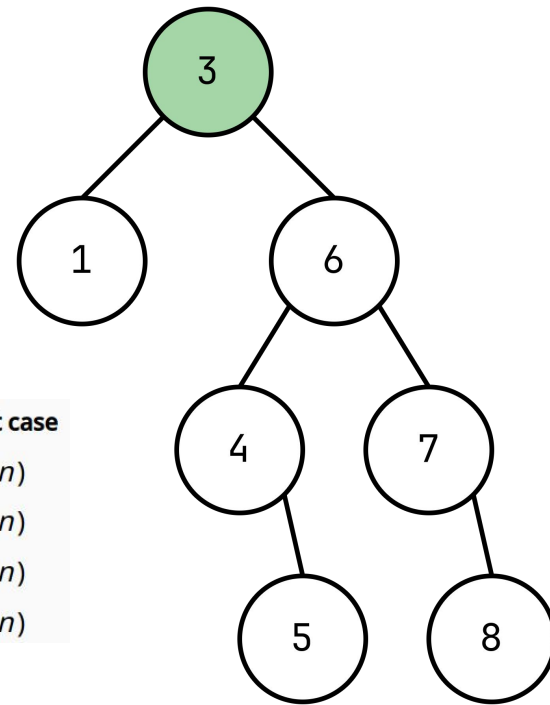
BALANCED TREES

BINARY SEARCH TREES
ARE **NOT BALANCED**

BALANCED TREE:

THE **HEIGHT** OF THE
LEFT AND THE RIGHT
SUBTREE FOR EACH NODE
IS EITHER ZERO OR ONE

Algorithm	Average	Worst case
Space	$O(n)$	$O(n)$
Search	$O(\log n)$	$O(n)$
Insert	$O(\log n)$	$O(n)$
Delete	$O(\log n)$	$O(n)$



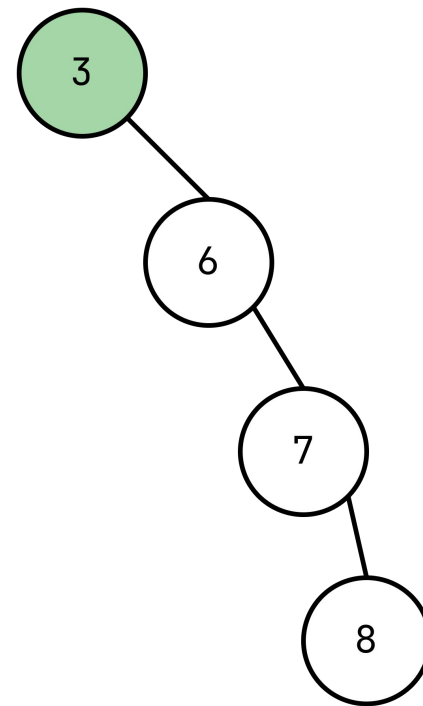
BALANCED TREES

BINARY SEARCH TREES
ARE **NOT BALANCED**

BALANCED TREE:




THE **HEIGHT** OF THE
LEFT AND THE RIGHT
SUBTREE FOR EACH NODE
IS EITHER ZERO OR ONE

Algorithm	Average	Worst case
Space	$O(n)$	$O(n)$
Search	$O(\log n)$	$O(n)$
Insert	$O(\log n)$	$O(n)$
Delete	$O(\log n)$	$O(n)$

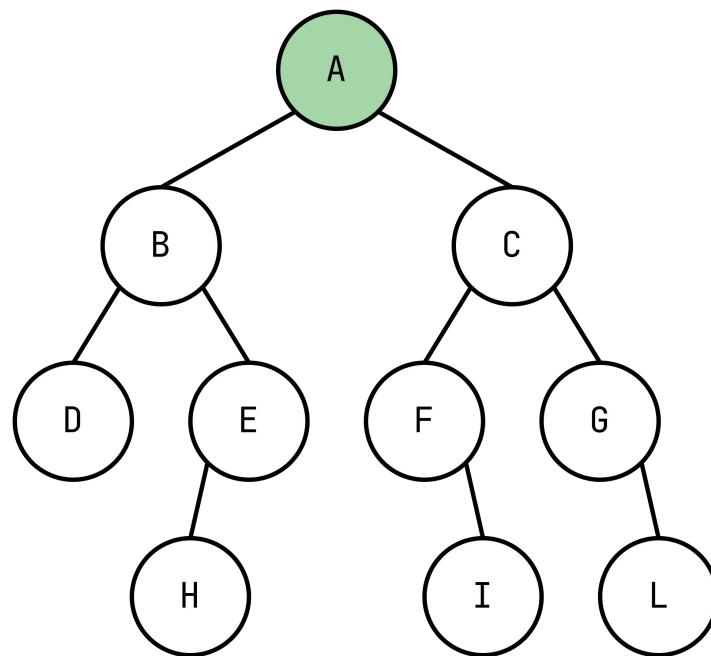


BALANCED SEARCH TREES

EXAMPLES OF BALANCED TREES:

-  AVL (BINARY) TREE
-  RED-BLACK (BINARY) TREE
-  B-TREE (K-ARY)

SEARCH, INSERTION AND DELETION
COST **$O(\log N)$**






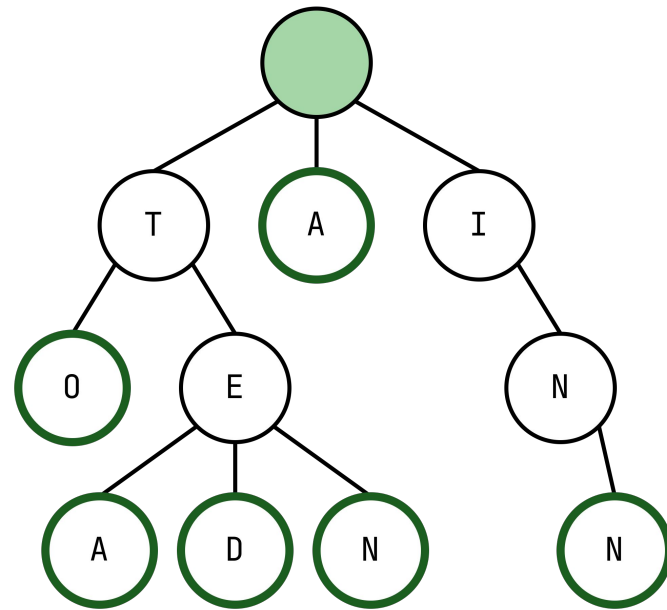
TRIE

KNOWN ALSO AS **DIGITAL TREE** OF **PREFIX TREE**

STRING-INDEXED LOOKUP DATA STRUCTURE

APPLICATIONS:

-  AUTOCOMPLETE
-  FULL-TEXT SEARCH
-  BLOCKCHAIN (MERKLE-PATRICIA)






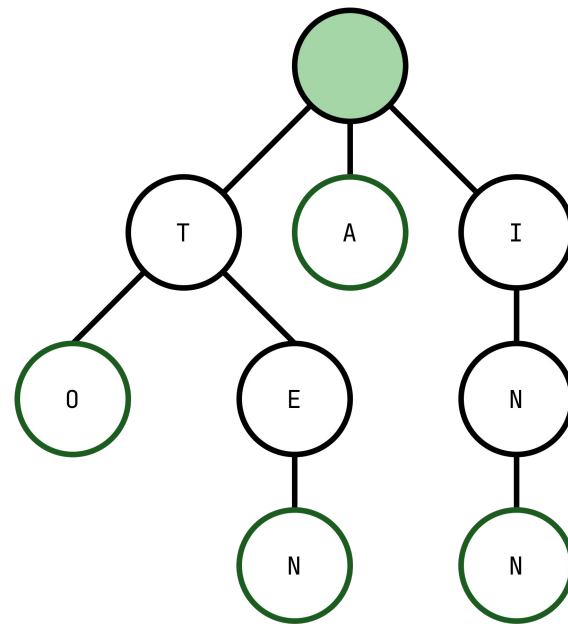
TRIE

KNOWN ALSO AS **DIGITAL TREE** OF **PREFIX TREE**

STRING-INDEXED LOOKUP DATA STRUCTURE

APPLICATIONS:

-  AUTOCOMPLETE
-  FULL-TEXT SEARCH
-  BLOCKCHAIN (MERKLE-PATRICIA)






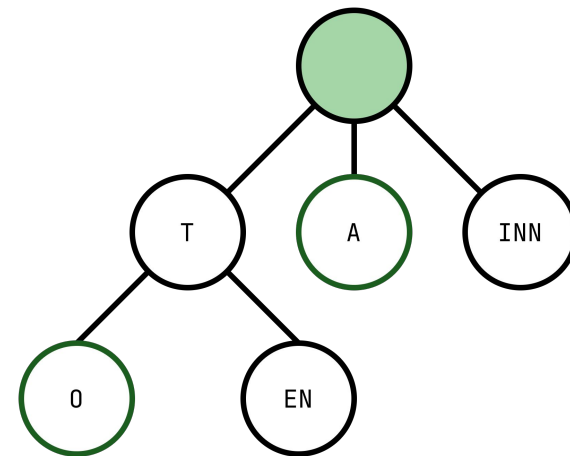
TRIE

KNOWN ALSO AS **DIGITAL TREE** OF **PREFIX TREE**

STRING-INDEXED LOOKUP DATA STRUCTURE

APPLICATIONS:

-  AUTOCOMPLETE
-  FULL-TEXT SEARCH
-  BLOCKCHAIN (MERKLE-PATRICIA)



COMPRESSED TRIE





POLITECNICO
MILANO 1863



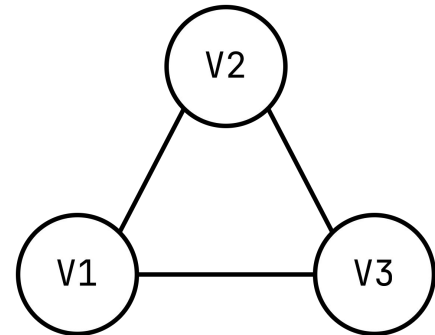
□ GRAPHS

GRAPHS

DEFINITION

GRAPH DATA STRUCTURE

- 👉 A **GRAPH** 🍇 IS A DATA STRUCTURE WITH A SET OF **VERTICES** AND **EDGES**
- 👉 GRAPHS CAN BE **DIRECTED (DiGRAPH)** OR **UNDIRECTED (GRAPH)**
- 👉 BOTH VERTICES AND EDGES CAN HOLD VALUES
- 👉 TREES ARE A PARTICULAR TYPE OF GRAPH



GRAPH IMPLEMENTATION

GRAPHS CAN BE IMPLEMENTED AS:



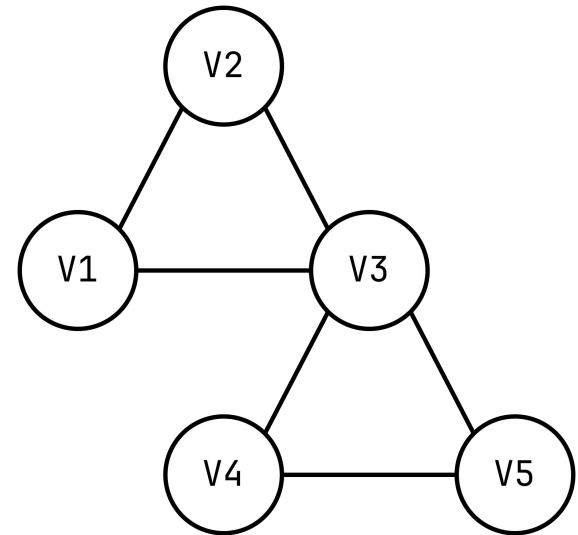
ADJACENCY LIST



ADJACENCY MATRIX



INCIDENCE MATRIX

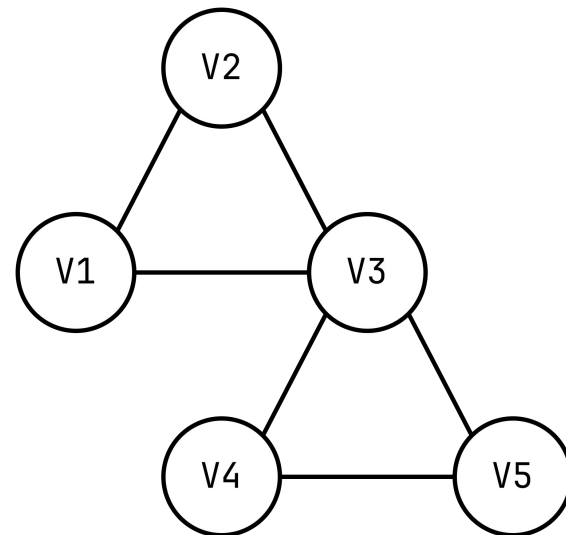


GRAPH IMPLEMENTATION

ADJACENCY LIST

LIST OF VERTICES AND
EACH VERTEX HAS A LIST OF
ADJACENT VERTICES

```
[ v1 [v2, v3],  
  v2 [v1, v3],  
  v3 [v1, v2, v3, v4],  
  v4 [v3, v5]  
  v5 [v3, v4]  
]
```



GRAPH IMPLEMENTATION

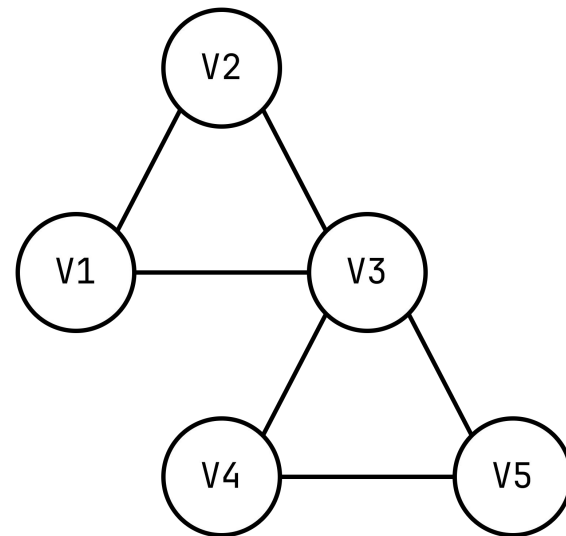
ADJACENCY MATRIX

TWO-DIMENSIONAL MATRIX

ROWS REPRESENTS **SOURCE** VERTICES

COLUMNS REPRESENTS **DESTINATION** VERTICES

	V1	V2	V3	V4	V5
V1	-	X	X	0	0
V2	X	-	X	0	0
V3	X	X	-	X	X
V4	0	0	X	-	X
V5	0	0	X	X	-





GRAPH IMPLEMENTATION

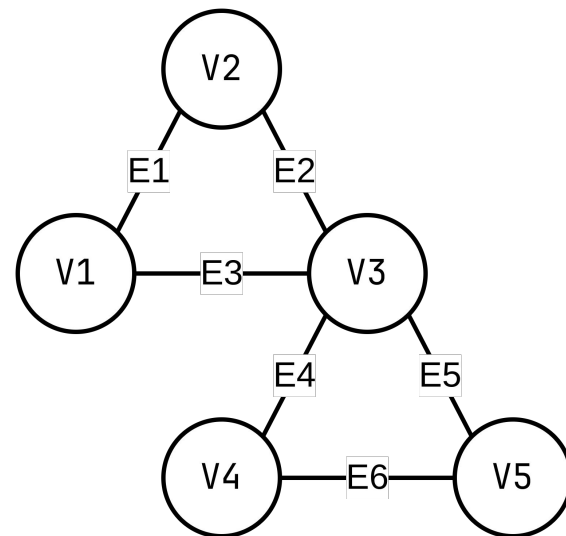
INCIDENCE MATRIX

TWO-DIMENSIONAL MATRIX

ROWS REPRESENTS **VERTICES**

COLUMNS REPRESENTS **EDGES**

	E1	E2	E3	E4	E5	E6
V1	X	0	X	0	0	0
V2	X	X	0	0	0	0
V3	0	X	X	X	0	X
V4	0	0	0	X	X	0
V5	0	0	0	0	X	X



GRAPH IMPLEMENTATION

	Adjacency list	Adjacency matrix	Incidence matrix
Store graph	$O(V + E)$	$O(V ^2)$	$O(V \cdot E)$
Add vertex	$O(1)$	$O(V ^2)$	$O(V \cdot E)$
Add edge	$O(1)$	$O(1)$	$O(V \cdot E)$
Remove vertex	$O(E)$	$O(V ^2)$	$O(V \cdot E)$
Remove edge	$O(V)$	$O(1)$	$O(V \cdot E)$
Are vertices x and y adjacent (assuming that their storage positions are known)?	$O(V)$	$O(1)$	$O(E)$
Remarks	Slow to remove vertices and edges, because it needs to find all vertices or edges	Slow to add or remove vertices, because matrix must be resized/copied	Slow to add or remove vertices and edges, because matrix must be resized/copied

GRAPH COMPRESSED REPRESENTATION

SPARSE MATRIX

IS THE MOST EFFICIENT WAY TO REPRESENT
A MATRIX WITH LOT OF ZEROS

FAST CONSTRUCTION:

LIL - LIST OF LISTS (ADJACENCY LIST)

DOK - DICTIONARY OF KEYS

COO - COORDINATE LIST

(FAST DATA ACCESS):

CSR - COMPRESSED SPARSE ROW

CSC - COMPRESSED SPARSE COLUMN

GRAPH COMPRESSED REPRESENTATION

LIL

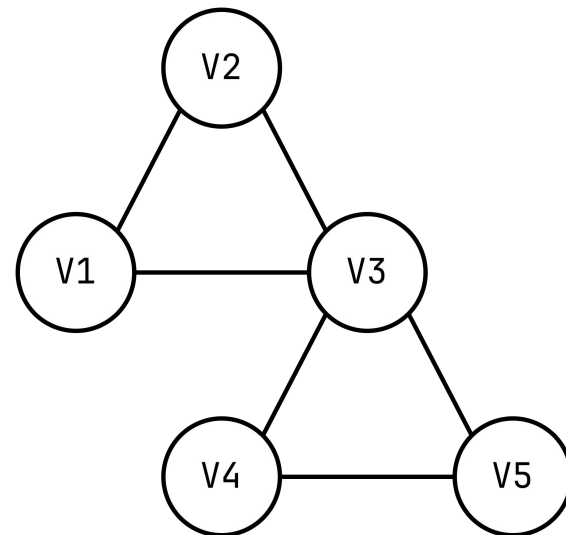
```
[ v1 [v2, v3],  
  v2 [v1, v3],  
  v3 [v1, v2, v3, v4],  
  v4 [v3, v5]  
  v5 [v3, v4]  
]
```

DOK

```
{ (v1, v2),  
  (v1, v3),  
  (v2, v3),  
  (v3, v4),  
  (v3, v5),  
  (v4, v5)}
```

COO

```
[ (v1, v2),  
  (v1, v3),  
  (v2, v3),  
  (v3, v4),  
  (v3, v5),  
  (v4, v5)]
```





POLITECNICO
MILANO 1863

SEARCH ALGORITHMS



SEARCH ALGORITHMS

SEARCH ALGORITHMS

- 👉 THEY ARE USED TO TRAVERSE OR EXPLORE A GRAPH OR TREE IN ORDER TO FIND A SPECIFIC ELEMENT OR PATH.
- 👉 THE MOST POPULAR ALGORITHMS ARE:
DFS (DEPTH FIRST SEARCH) AND
BFS (BREADTH FIRST SEARCH).



DFS

DEPTH FIRST SEARCH

DEPTH FIRST SEARCH

- 👉 USES A STACK DATA STRUCTURE.
- 👉 VISITS NODES DEPTH-FIRST (I.E., GOES AS DEEP AS POSSIBLE ALONG EACH BRANCH BEFORE **BACKTRACKING**).
- 👉 GOOD FOR TRAVERSING AN ENTIRE TREE OR GRAPH.
- 👉 CAN GET STUCK IN AN INFINITE LOOP IF THE GRAPH HAS CYCLES.

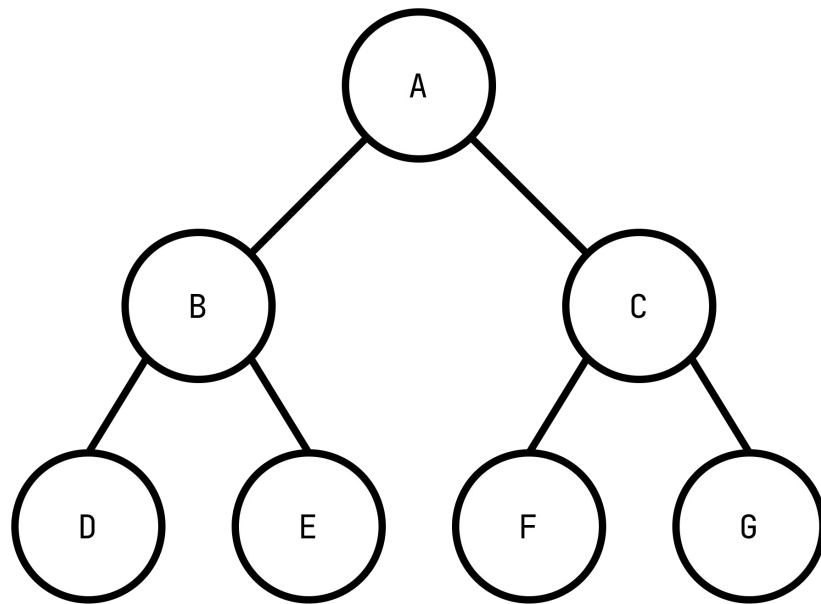
DEPTH FIRST SEARCH

EXAMPLE:

FIND A PATH FROM A TO F

VISITED: []

STACK: []



DEPTH FIRST SEARCH

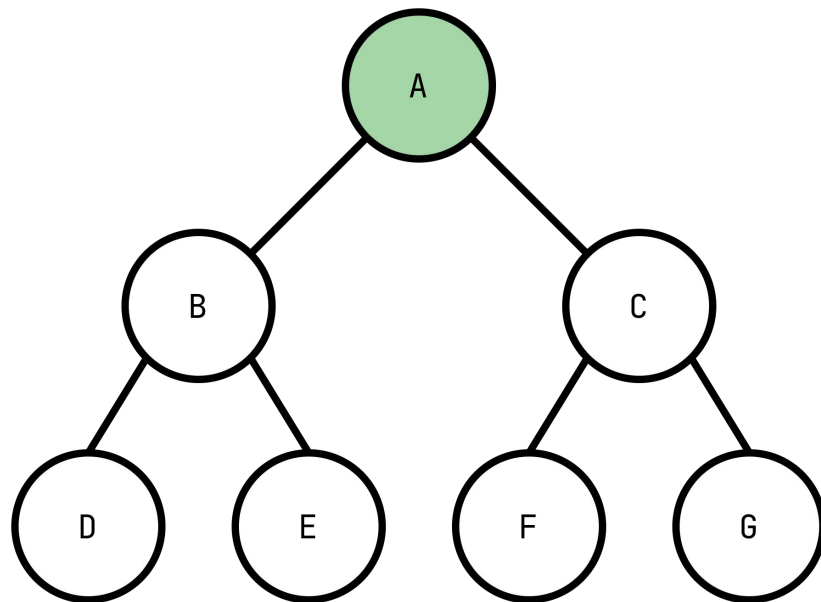
EXAMPLE:

FIND A PATH FROM A TO F

VISIT NODE A

VISITED: [A]

STACK: [B, C]



DEPTH FIRST SEARCH

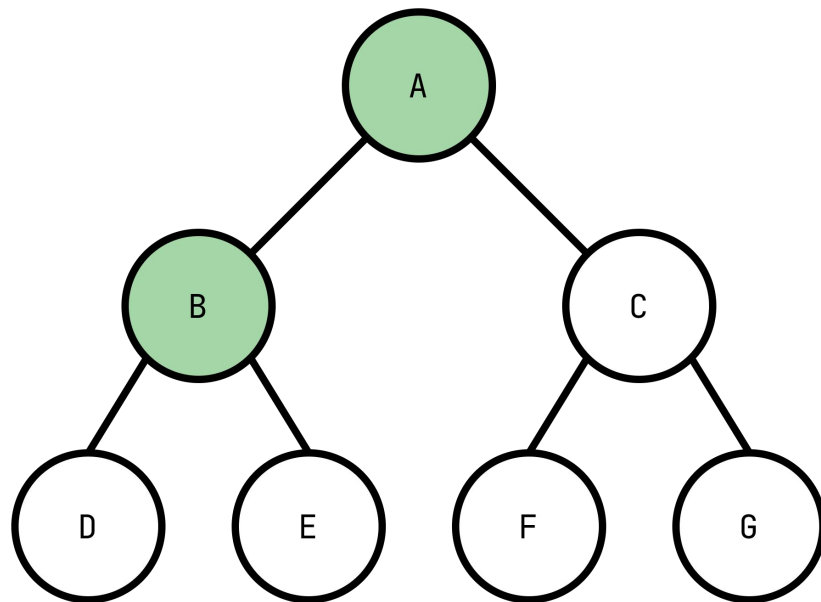
EXAMPLE:

FIND A PATH FROM A TO F

VISIT NODE B

VISITED: [A, B]

STACK: [D, E, C]



DEPTH FIRST SEARCH

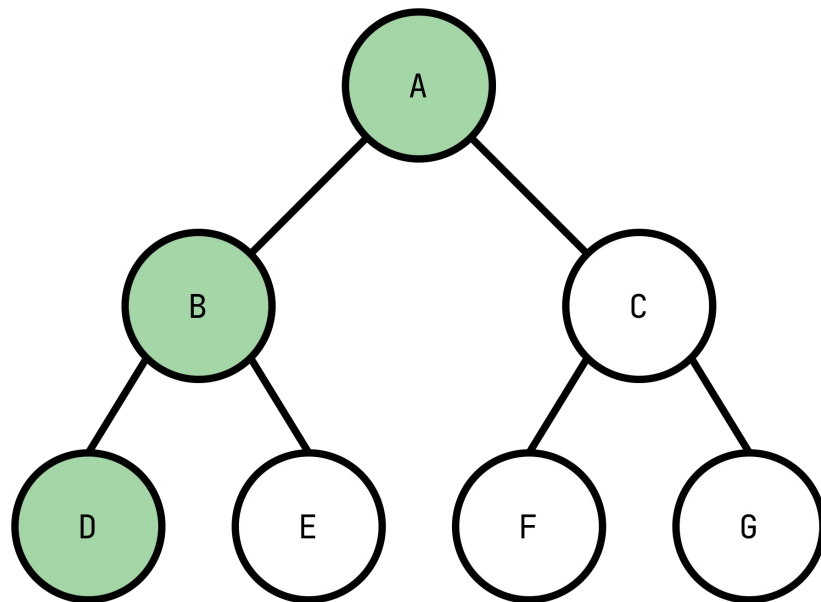
EXAMPLE:

FIND A PATH FROM A TO F

VISIT NODE D

VISITED: [A, B, D]

STACK: [E, C]



DEPTH FIRST SEARCH

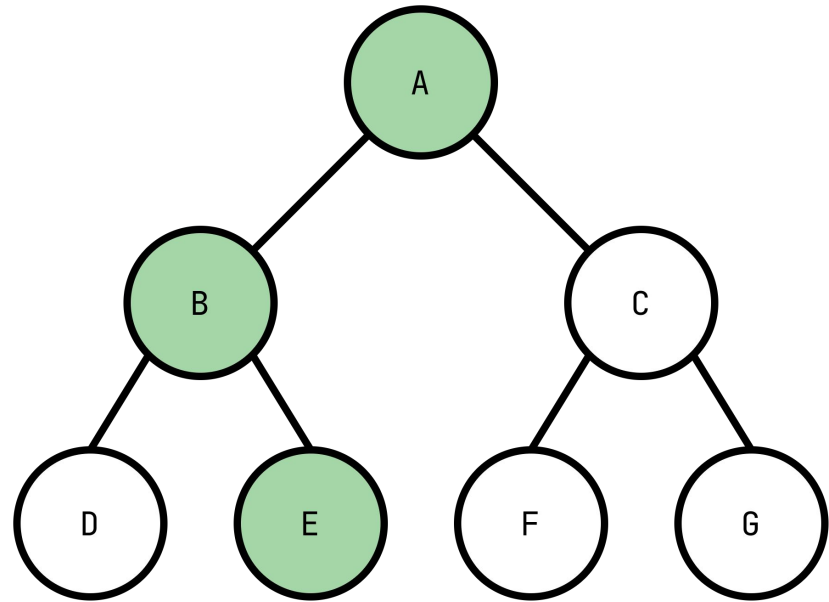
EXAMPLE:

FIND A PATH FROM A TO F

BACKTRACKING
AND VISIT NODE E

VISITED: [A, B, D, E]

STACK: [C]



DEPTH FIRST SEARCH

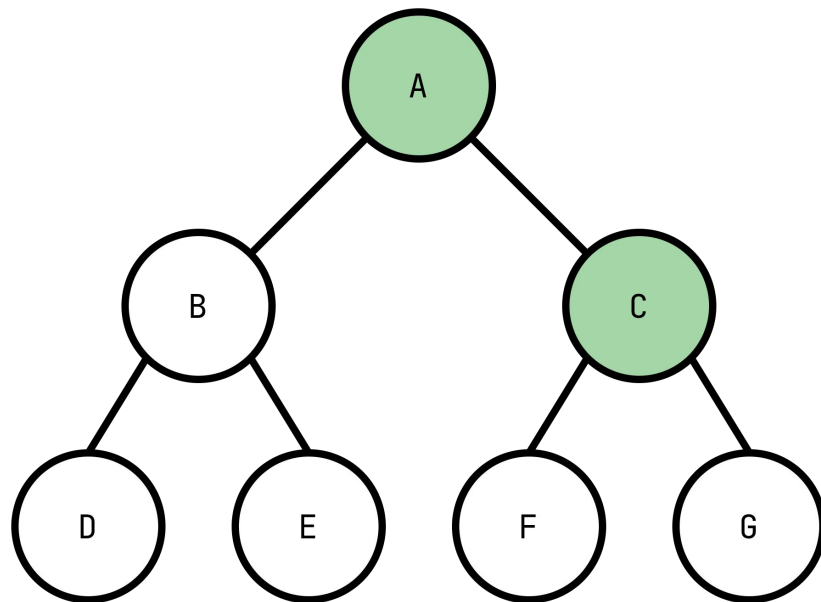
EXAMPLE:

FIND A PATH FROM A TO F

BACKTRACKING
AND VISIT NODE C

VISITED: [A, B, D, E, C]

STACK: [F, G]



DEPTH FIRST SEARCH

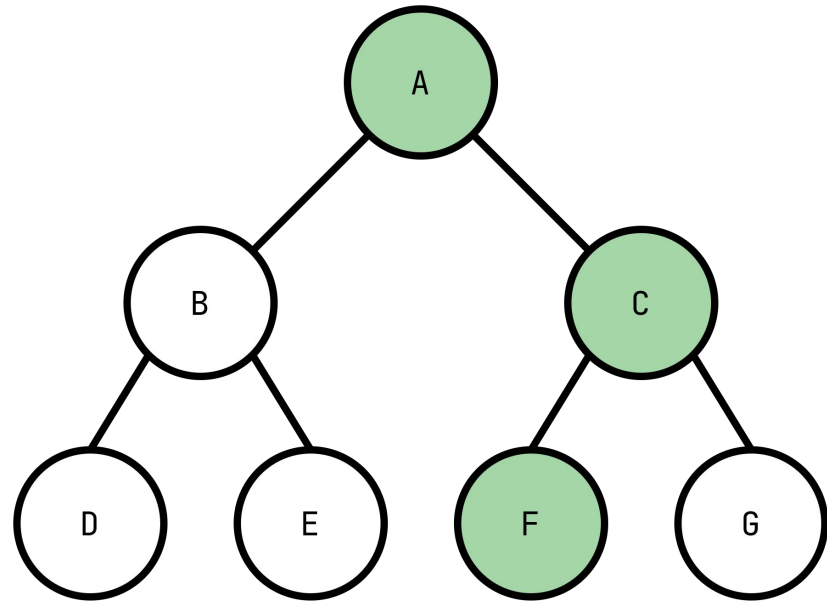
EXAMPLE:

FIND A PATH FROM A TO F

VISIT NODE F

VISITED: [A, B, D, E, C, F]

STACK: [C]





BFS

BREADTH FIRST SEARCH

BREADTH FIRST SEARCH

- 👉 USES A QUEUE DATA STRUCTURE.
- 👉 VISITS NODES BREADTH-FIRST (I.E., EXPLORES ALL THE NEIGHBORS OF A NODE BEFORE MOVING ON TO THE NEXT LEVEL).
- 👉 GOOD FOR FINDING THE SHORTEST PATH BETWEEN TWO NODES.
- 👉 MAY USE MORE MEMORY THAN DFS IF THE GRAPH IS LARGE.
- 👉 MAY BE SLOWER THAN DFS FOR TRAVERSING THE ENTIRE TREE OR GRAPH.

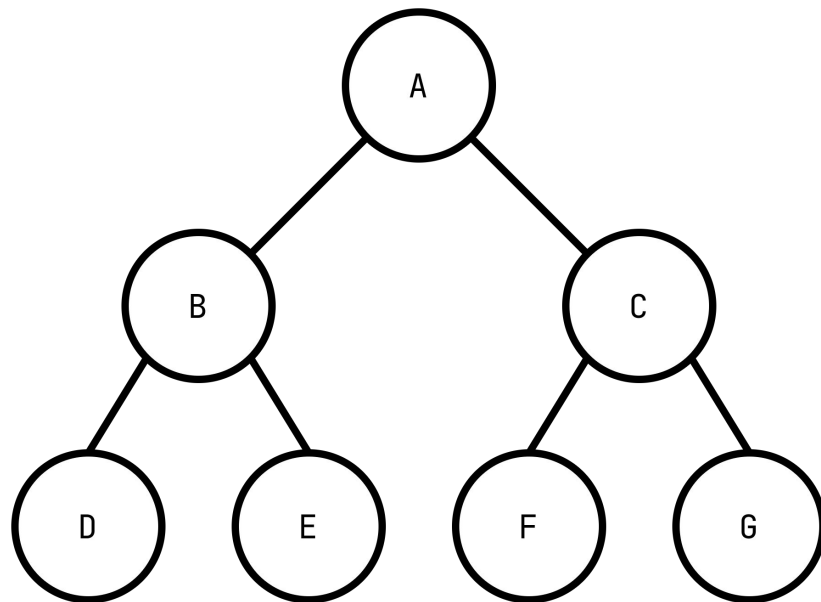
BREADTH FIRST SEARCH

EXAMPLE:

FIND A PATH FROM A TO F

VISITED: []

QUEUE: []



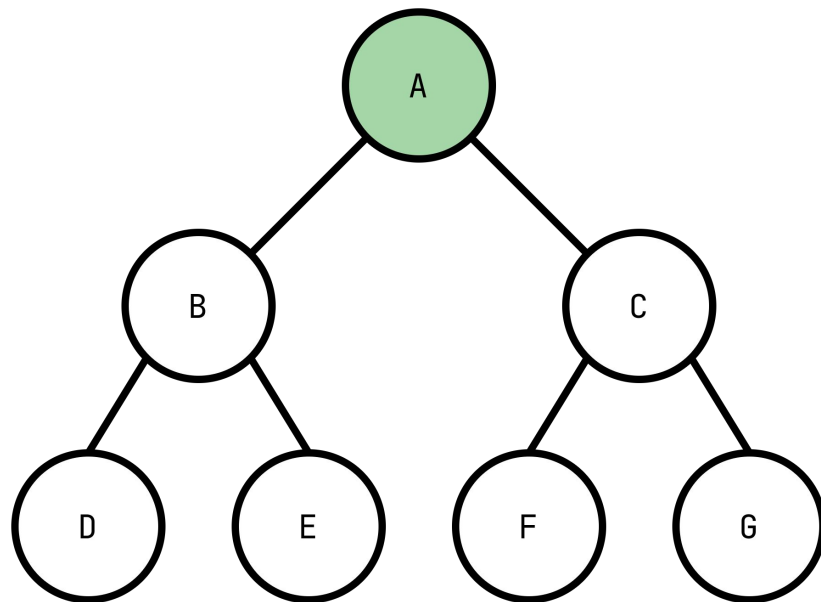
BREADTH FIRST SEARCH

EXAMPLE:

FIND A PATH FROM A TO F

VISITED: [A]

QUEUE: [B, C]



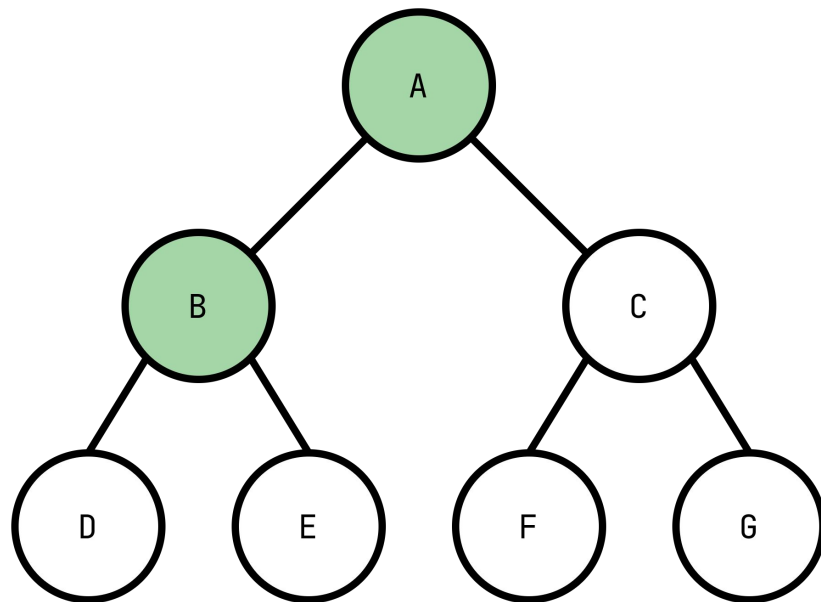
BREADTH FIRST SEARCH

EXAMPLE:

FIND A PATH FROM A TO F

VISITED: [A, B]

QUEUE: [C, D, E]



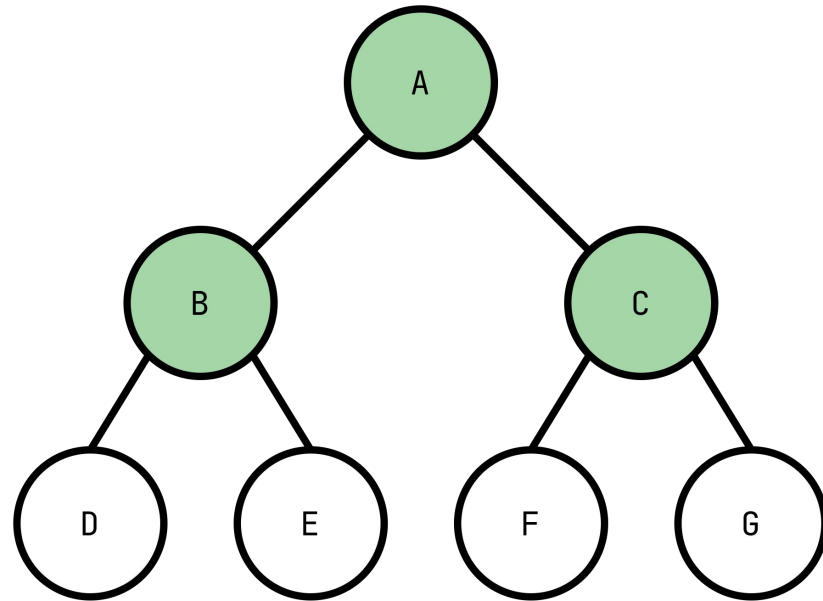
BREADTH FIRST SEARCH

EXAMPLE:

FIND A PATH FROM A TO F

VISITED: [A, B, C]

QUEUE: [D, E, F, G]



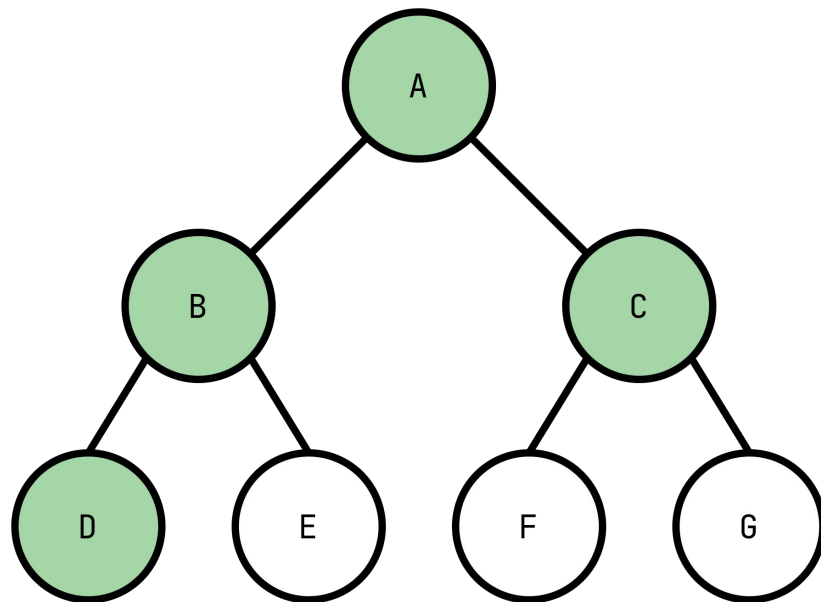
BREADTH FIRST SEARCH

EXAMPLE:

FIND A PATH FROM A TO F

VISITED: [A, B, C, D]

QUEUE: [E, F, G]



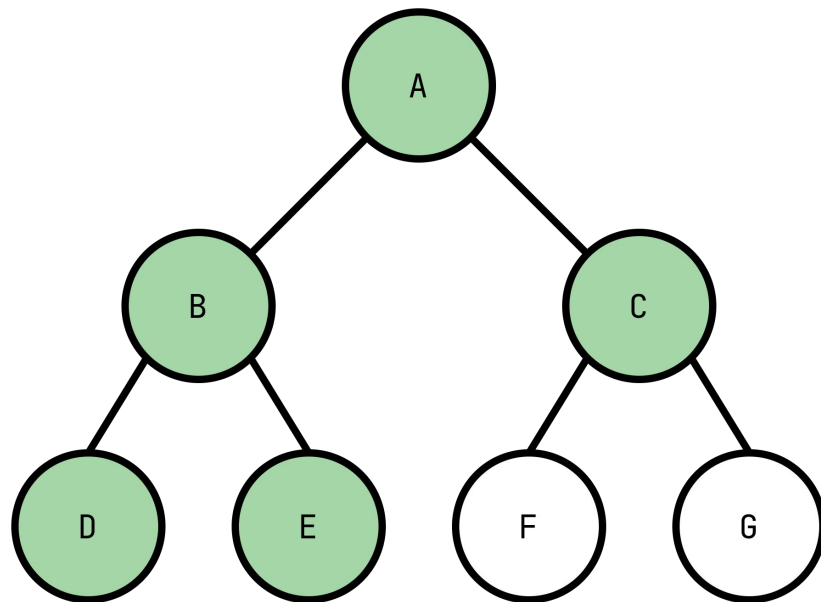
BREADTH FIRST SEARCH

EXAMPLE:

FIND A PATH FROM A TO F

VISITED: [A, B, C, D, E]

QUEUE: [F, G]



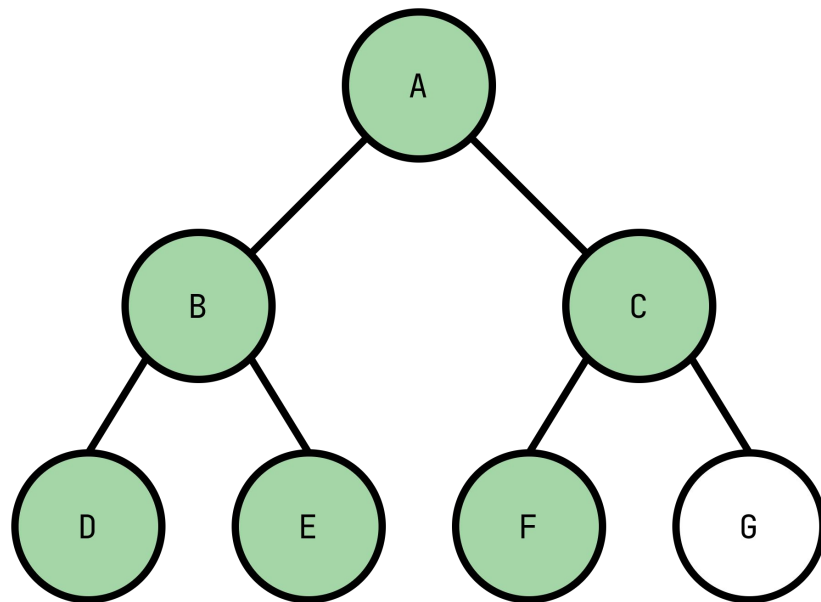
BREADTH FIRST SEARCH

EXAMPLE:

FIND A PATH FROM A TO F

VISITED: [A, B, C, D, E, F]

QUEUE: [G]





SEARCH ALGORITHMS

COMPLEXITY AND OPTIMALITY



WORST CASE COMPLEXITIES

OF BOTH ALGORITHMS ARE THE SAME.



TIME COMPLEXITY: $O(|V| + |E|)$



SPACE COMPLEXITY: $O(|V|)$



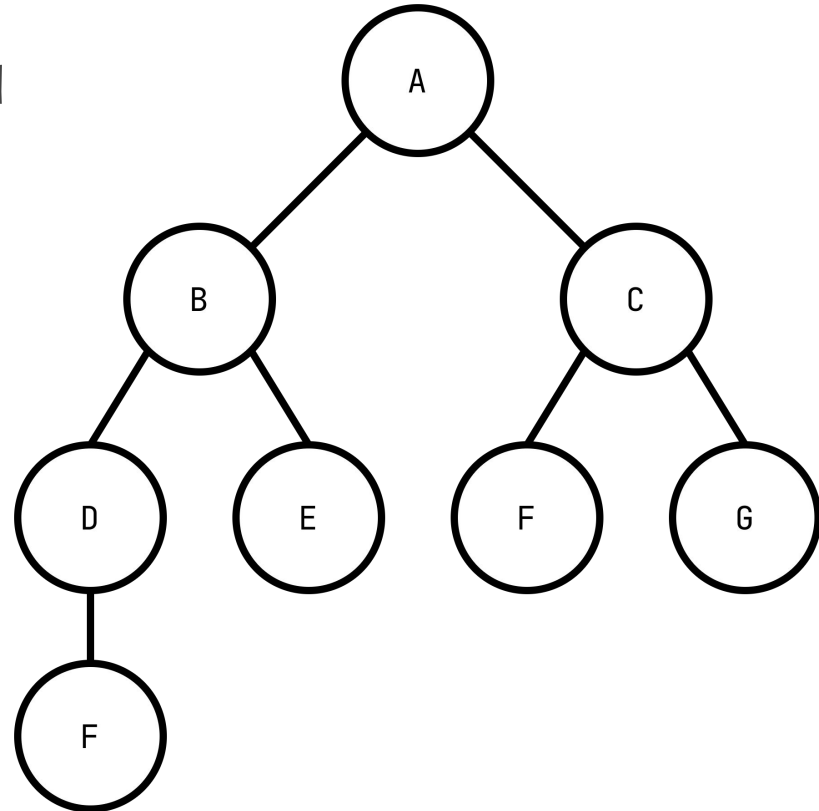
BFS IS OPTIMAL IF THE PATH COST IS A
NON-DECREASING FUNCTION OF $D(\text{DEPTH})$.

NORMALLY, BFS IS APPLIED WHEN ALL THE ACTIONS
HAVE THE SAME COST.

OPTIMALITY COMPARISON

➡ ASSUME THERE ARE
TWO PATHS TO REACH F

➡ ASSUME EACH STEP
HAVE THE SAME COST
(FOR SIMPLICITY 1)



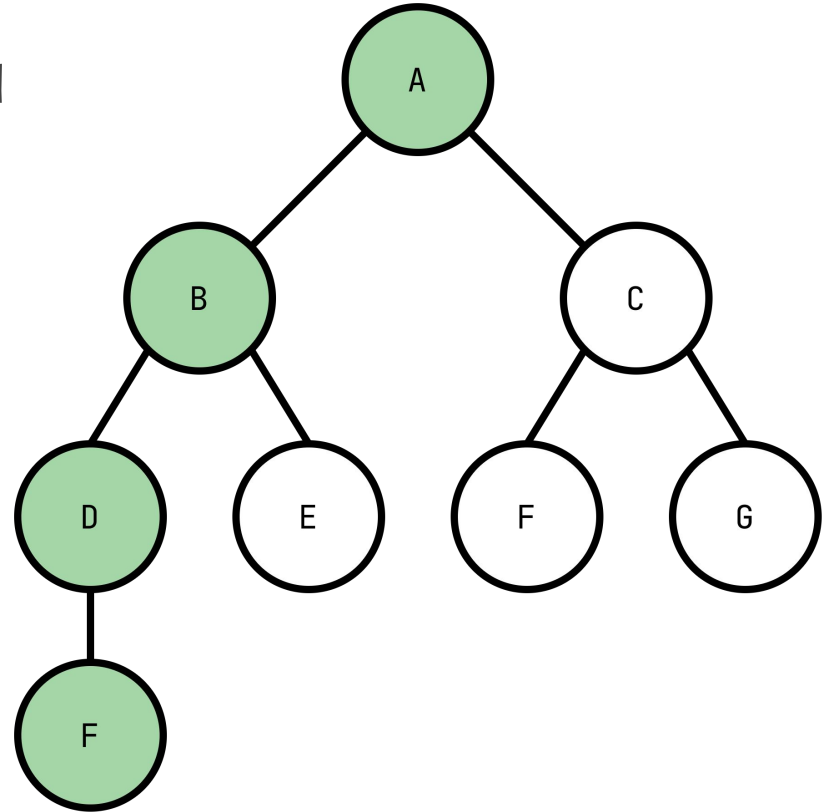
OPTIMALITY COMPARISON

DEPTH FIRST SEARCH

PATH FOUND:

[A → B → D → F]

PATH COST: 4



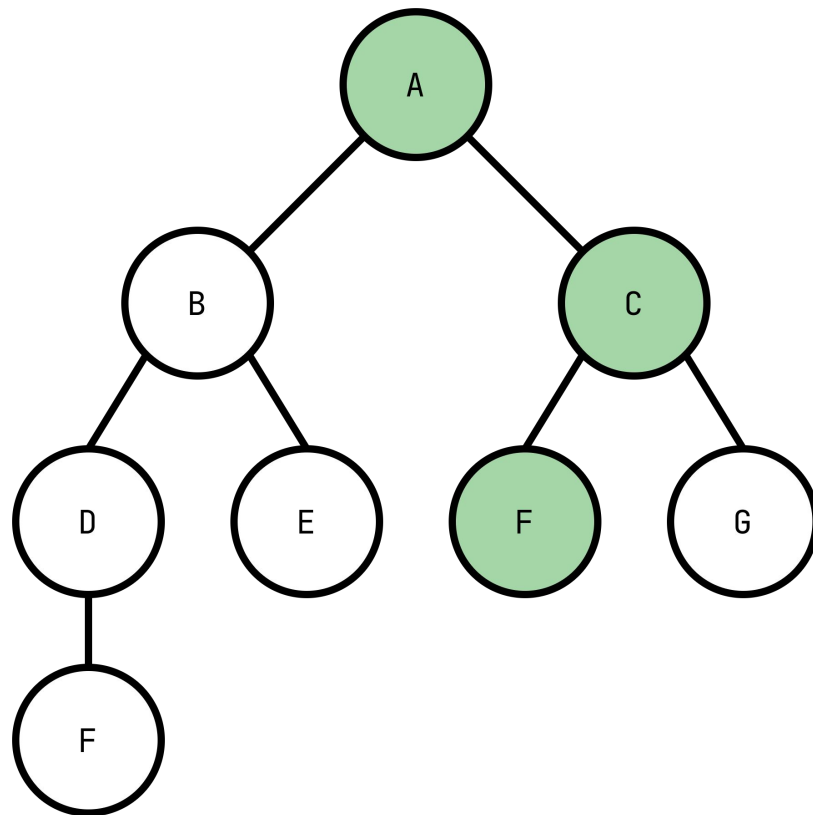
OPTIMALITY COMPARISON

BREADTH FIRST SEARCH

PATH FOUND:

[A → C → F]

PATH COST: 3





POLITECNICO
MILANO 1863

LEETCODE

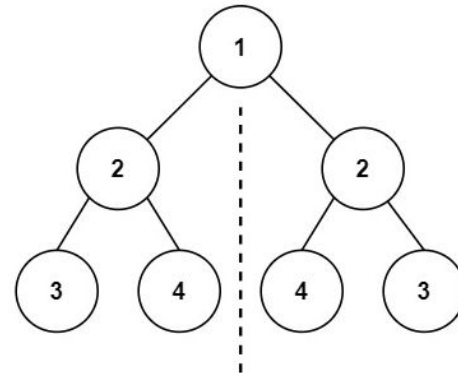


LEETCODE PROBLEM 1

101 SYMMETRIC TREE

[HTTPS://LEETCODE.COM/PROBLEMS/SYMMETRIC-TREE/](https://leetcode.com/problems/symmetric-tree/)

GIVEN THE ROOT OF A BINARY TREE, CHECK WHETHER IT IS A MIRROR OF ITSELF (I.E., SYMMETRIC AROUND ITS CENTER).





LEETCODE

PROBLEM 2

98 VALIDATE BINARY TREE

[HTTPS://LEETCODE.COM/PROBLEMS/VALIDATE-BINARY-SEARCH-TREE/](https://leetcode.com/problems/validate-binary-search-tree/)

GIVEN THE ROOT OF A BINARY TREE, DETERMINE IF IT IS A VALID BINARY SEARCH TREE (BST). A VALID BST IS DEFINED AS FOLLOWS:

- 👉 THE LEFT SUBTREE OF A NODE CONTAINS ONLY NODES WITH KEYS LESS THAN THE NODE'S KEY.
- 👉 THE RIGHT SUBTREE OF A NODE CONTAINS ONLY NODES WITH KEYS GREATER THAN THE NODE'S KEY.
- 👉 BOTH THE LEFT AND RIGHT SUBTREES MUST ALSO BE BINARY SEARCH TREES.



LEETCODE

PROBLEM 3

207 COURSE SCHEDULE

[HTTPS://LEETCODE.COM/PROBLEMS/VALIDATE-BINARY-SEARCH-TREE/](https://leetcode.com/problems/validate-binary-search-tree/)

THERE ARE A TOTAL OF `NUMCOURSES` COURSES YOU HAVE TO TAKE, LABELED FROM 0 TO `NUMCOURSES - 1`. YOU ARE GIVEN AN ARRAY `PREREQUISITES` WHERE `PREREQUISITES[i] = [ai, bi]` INDICATES THAT YOU MUST TAKE COURSE `bi` FIRST IF YOU WANT TO TAKE COURSE `ai`.

FOR EXAMPLE, THE PAIR `[0, 1]`, INDICATES THAT TO TAKE COURSE 0 YOU HAVE TO FIRST TAKE COURSE 1.

RETURN TRUE IF YOU CAN FINISH ALL COURSES. OTHERWISE, RETURN FALSE.