# μC/OS-II Reference Manual

This chapter provides a reference to  $\mu C/OS$ -II services. Each of the user-accessible kernel services is presented in alphabetical order. The following information is provided for each of the services:

- · A brief description
- The function prototype
- The filename of the source code
- The #define constant needed to enable the code for the service
- A description of the arguments passed to the function
- A description of the returned value(s)
- Specific notes and warnings on using the service
- One or two examples of how to use the function

# OS\_ENTER\_CRITICAL() OS\_EXIT\_CRITICAL()

Chapter	File	Called from	Code enabled by
3	OS_CPU.H	Task or ISR	N/A

OS\_ENTER\_CRITICAL() and OS\_EXIT\_CRITICAL() are macros used to disable and enable, respectively, the processor's interrupts.

### Arguments

none

### **Returned Values**

none

### **Notes/Warnings**

- 1. These macros must be used in pairs.
- 2. If OS\_CRITICAL\_METHOD is set to 3, your code is assumed to have allocated local storage for a variable of type OS CPU SR, which is called cpu sr, as follows

```
#if OS_CRITICAL_METHOD == 3 /* Allocate storage for CPU status reg. */
    OS_CPU_SR cpu_sr;
#endif
```

# OSEventNameGet()

Chapter	File	Called from	Code enabled by	
New in V2.60	OS_CORE.C	Task	OS_EVENT_NAME_EN	

OSEventNameGet () allows you to obtain the name that you assigned to a semaphore, a mutex, a mailbox or a message queue. This function is typically used by a debugger to allow associating a name to a resource.

### **Arguments**

is a pointer to the event control block. pevent can point either to a semaphore, a mutex, a pevent mailbox or a queue. Where this function is concerned, the actual type is irrelevant. This pointer is returned to your application when the semaphore, mutex, mailbox or queue is created (see OSSemCreate(), OSMutexCreate(), OSMboxCreate() and OSQCreate()). is a pointer to a pointer to the name of the semaphore, mutex, mailbox or queue. pname a pointer to an error code and can be any of the following: perr If pname now points to the name of the semaphore, mutex, OS ERR NONE mailbox or queue. You are not pointing to either a semaphore, mutex, mailbox or OS ERR EVENT TYPE message queue. You passed a NULL pointer for pevent. OS ERR PEVENT NULL

You tried calling this function from an ISR.

### **Returned Values**

The size of the ASCII string pointed to by pname or 0 if an error is encountered.

OS ERR NAME GET ISR

# **Notes/Warnings**

1. The semaphore, mutex, mailbox or message queue must be created before you can use this function and obtain the name of the resource.

# OSEventNameSet()

Chapter	File	Called from	Code enabled by	
New in V2.60	OS_CORE.C	Task	OS_EVENT_NAME_EN	

OSEventNameSet () allows you to assign a name to a semaphore, a mutex, a mailbox or a message queue. This function is typically used by a debugger to allow associating a name to a resource.

### Arguments

pevent is a pointer to the event control block that you want to name. pevent can point either to a

semaphore, a mutex, a mailbox or a queue. Where this function is concerned, the actual type is irrelevant. This pointer is returned to your application when the semaphore, mutex, mailbox or queue is created (see OSSemCreate(), OSMutexCreate(), OSMboxCreate() and

OSQCreate()).

pname is a pointer to the name of the semaphore, mutex, mailbox or queue.

perr a pointer to an error code and can be any of the following:

OS ERR NONE If the call was successfull.

OS\_ERR\_EVENT\_TYPE You are not pointing to either a semaphore, mutex, mailbox or

message queue.

OS\_ERR\_PEVENT\_NULL You passed a NULL pointer for pevent.

OS ERR NAME SET ISR You called this function from an ISR.

# **Returned Values**

none

### **Notes/Warnings**

1. The semaphore, mutex, mailbox or message queue must be created before you can use this function and set the name of the resource or the event.

```
OS_EVENT *PrinterSem;

void Task (void *p_arg)
{
    INT8U err;

    (void)p_arg;
    for (;;) {
        OSEventNameSet(PrinterSem, "Printer #1", &err);
        .
        .
        .
    }
}
```

# OSEventPendMulti()

```
INT16U OSEventPendMulti(OS EVENT **pevents pend,
                         OS EVENT **pevents rdy,
                         void
                                   **pmsgs rdy,
                         INT16U
                                     timeout,
                         INT8U
                                    *perr);
```

Chapter	File	Called from	Code enabled by
	OS_CORE.C	Task only	OS_EVENT_MULTI_EN

OSEventPendMulti() is used when a task expects to wait on multiple events. If multiple events are ready when OSEventPendMulti() is called, then all available events and messages, if any, are returned as ready to the caller. If no events are ready, OSEventPendMulti() suspends the current task until either an event is ready or a user-specified timeout expires. If an event becomes ready and multiple tasks are waiting for the event, μC/OS-II resumes the highest priority task waiting to run.

A pended task that has been suspended with OSTaskSuspend() can still receive a message from a multipended mailbox or message queue or obtain a multi-pended semaphore. However, the task remains suspended until it is resumed by calling OSTaskResume ().

### Arguments

pevents pend is a pointer to a null-terminated array of OS EVENT pointers. These event pointers are returned to your application when the mailboxes, message queues, and semaphores are created [see OSMboxCreate(), OSQCreate(), and OSSemCreate()].

pevents rdy

is a pointer to an array to return the available OS EVENT pointers. The size of the array must be greater than or equal to the size of the pevents pend array, including the terminating NULL.

pmsgs rdy

is a pointer to an array to return messages from any multi-pended mailbox or message queue events. The size of the array must be greater than or equal to the size of the pevents pend array, excluding the terminating NULL. Since NULL messages are valid messages, this array cannot be NULL-terminated. Instead, every available message is returned in the pmsqs rdy array at the same index as the ready mailbox or message queue event is returned in the pevents rdy array. All other pmsgs rdy array indices are filled with NULL messages.

timeout

allows the task to resume execution if no multi-pended event is ready within the specified number of clock ticks. A timeout value of 0 indicates that the task wants to wait forever for any of the multi-pended events. The maximum timeout is 65,535 clock ticks. The timeout value is not synchronized with the clock tick. The timeout count begins decrementing on the next clock tick, which could potentially occur immediately.

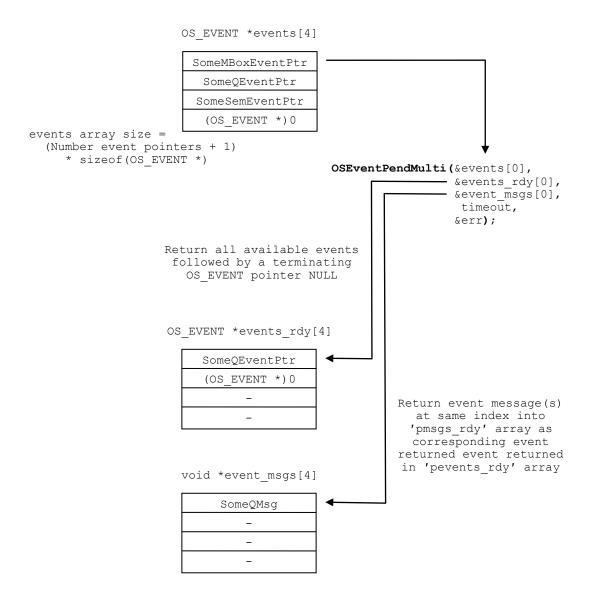
perr	is a pointer to a variable that holds an error code. OSEventPendMulti() sets *perr to one of the following:			
	OS_ERR_NONE	if any of the multi-pended events are ready; check the pevents_rdy array for which events are available.		
	OS_ERR_TIMEOUT	if no multi-pended event is ready within the specified timeout.		
	OS_ERR_PEND_ABORT	indicates that a multi-pended event was aborted; check the pevents_rdy array for which events were aborted.		
	OS_ERR_EVENT_TYPE	if $pevents\_pend$ is not pointing to an array of valid mailbox, message queue, or semaphore events.		
	OS_ERR_PEND_LOCKED	if you called this function when the scheduler is locked.		
	OS_ERR_PEND_ISR	if you call this function from an ISR and $\mu C/OS\text{-}II$ suspends it. In general, you should not call <code>OSEventPendMulti()</code> from an ISR, but $\mu C/OS\text{-}II$ checks for this situation anyway.		
	OS_ERR_PEVENT_NULL	if pevents_pend, pevents_rdy, or pmsgs_rdy is a NULL pointer.		

### **Returned Value**

OSEventPendMulti() returns the number of multi-pended events that are ready or have been aborted, and \*perr is set to OS\_ERR\_NONE or OS\_ERR\_PEND\_ABORT, respectively. If no multi-pended event is ready within the specified timeout period or because of any error, then the pevents\_rdy and pmsgs\_rdy array are returned as NULL pointers, and \*perr is set to OS\_ERR\_TIMEOUT or to the respective error.

# **Notes/Warnings**

- 1. Mailbox, message queue, or semaphore events must be created before they are used.
- 2. You should not call OSEventPendMulti() from an ISR.
- 3. You cannot multi-pend on event flags and mutexes.



```
void EventTask(void *p arg)
   OS_EVENT *events[4];
   OS_EVENT *events_rdy[4];
   void    *event_msgs[4];
   INT16U timeout;
   INT8U err;
   (void)p_arg;
   for (;;) {
       events[0] = (OS_EVENT *)SomeMBoxEventPtr;
       events[1] = (OS EVENT *)SomeQEventPtr;
       events[2] = (OS_EVENT *)SomeSemEventPtr;
       events[3] = (OS_EVENT *)0;
       events nbr rdy = OSEventsPendMulti(&events[0]
                                      &events_rdy[0],
                                      &event_msgs[0],
                                       timeout,
                                      &err);
       if (err == OS_ERR_NONE) {
          } else {
          . /* Code for events not ready within timeout */
       }
     }
```

# OSFlagAccept()

OS\_FLAGS OSFlagAccept(OS\_FLAG\_GRP \*pgrp,
OS\_FLAGS flags,
INT8U wait\_type,
INT8U \*perr);

Chapter	File	Called from	Code enabled by
9	OS_FLAG.C	Task and ISR	OS_FLAG_EN && OS_FLAG_ACCEPT_EN

OSFlagAccept () allows you to check the status of a combination of bits to be either set or cleared in an event flag group. Your application can check for **any** bit to be set/cleared or **all** bits to be set/cleared. This function behaves exactly as OSFlagPend() does, except that the caller does NOT block if the desired event flags are not present.

# **Arguments**

is a pointer to the event flag group. This pointer is returned to your application when the event

flag group is created [see  ${\tt OSFlagCreate}$  ()].

flags is a bit pattern indicating which bit(s) (i.e., flags) you wish to check. The bits you want are

specified by setting the corresponding bits in flags.

wait\_type specifies whether you want all bits to be set/cleared or any of the bits to be set/cleared. You

can specify the following arguments:

 ${\tt OS\_FLAG\_WAIT\_CLR\_ALL} \qquad You check \ \mbox{all bits in flags to be clear} \ (0)$ 

OS FLAG WAIT CLR ANY You check any bit in flags to be clear (0)

OS\_FLAG\_WAIT\_SET\_ALL You check all bits in flags to be set (1)

OS FLAG WAIT SET ANY You check any bit in flags to be set (1)

You can add OS\_FLAG\_CONSUME if you want the event flag(s) to be consumed by the call. For example, to wait for **any** flag in a group and then clear the flags that are present, set wait type to

OS FLAG WAIT SET ANY + OS FLAG CONSUME

perr a pointer to an error code and can be any of the following:

OS\_ERR\_NONE No error

OS ERR EVENT TYPE You are not pointing to an event flag group

OS ERR FLAG WAIT TYPE You didn't specify a proper wait type argument.

OS\_ERR\_FLAG\_INVALID\_PGRP You passed a NULL pointer instead of the event flag

handle.

 ${\tt OS\_ERR\_FLAG\_NOT\_RDY} \qquad \qquad {\tt The \ desired \ flags \ for \ which \ you \ are \ waiting \ are \ not}$ 

available.

#### **Returned Values**

The flag(s) that cause the task to be ready or, 0 if either none of the flags are ready or an error occurred.

### **Notes/Warnings**

- 1. The event flag group must be created before it is used.
- 2. This function does **not** block if the desired flags are not present.

#### **IMPORTANT**

The return value of OSFlagAccept() is different as of V2.70. In previous versions, OSFlagAccept() returned the current state of the flags and now, it returns the flag(s) that are ready, if any.

```
#define ENGINE OIL PRES OK
                             0x01
#define ENGINE_OIL_TEMP_OK
                              0x02
#define ENGINE_START
                             0x04
OS FLAG GRP *EngineStatus;
void Task (void *p arg)
   INT8U err;
   OS FLAGS value;
    (void)p_arg;
    for (;;) {
       value = OSFlagAccept(EngineStatus,
                            ENGINE OIL PRES OK + ENGINE OIL TEMP OK,
                            OS_FLAG_WAIT_SET_ALL,
                            &err);
       switch (err) {
           case OS_ERR_NONE:
                 /* Desired flags are available */
                break;
            case OS_ERR_FLAG_NOT_RDY:
                 /* The desired flags are NOT available */
                break;
```

# OSFlagCreate()

```
OS FLAG GRP *OSFlagCreate(OS_FLAGS flags,
                          INT8U
                                 *perr);
```

Chapter	File	Called from	Code enabled by
9	OS_FLAG.C	Task or startup code	OS_FLAG_EN

OSFlagCreate() is used to create and initialize an event flag group.

# **Arguments**

contains the initial value to store in the event flag group. flags

is a pointer to a variable that is used to hold an error code. The error code can be one of the perr following:

OS\_ERR\_NONE if the call is successful and the event flag group has been

created.

if you attempt to create an event flag group from an ISR. OS ERR CREATE ISR

if no more event flag groups are available. You need to OS ERR FLAG GRP DEPLETED

increase the value of OS MAX FLAGS in OS CFG.H.

### **Returned Values**

A pointer to the event flag group if a free event flag group is available. If no event flag group is available, OSFlagCreate() returns a NULL pointer.

# **Notes/Warnings**

1. Event flag groups must be created by this function before they can be used by the other services.

### **Example**

```
OS FLAG GRP *EngineStatus;
void main (void)
   INT8U err;
               /* Initialize µC/OS-II */
   OSInit();
                 /* Create a flag group containing the engine's status */
   EngineStatus = OSFlagCreate(0x00, &err);
   OSStart(); /* Start Multitasking */
```

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# OSFlagDel()

Chapter	File	Called from	Code enabled by
9	OS_FLAG.	Task	OS_FLAG_EN and OS_FLAG_DEL_EN

OSFlagDel() is used to delete an event flag group. This function is dangerous to use because multiple tasks could be relying on the presence of the event flag group. You should always use this function with great care. Generally speaking, before you delete an event flag group, you must first delete all the tasks that access the event flag group.

### **Arguments**

7 XI Suments			
pgrp	is a pointer to the event flag group. This pointer is returned to your application when the event flag group is created [see OSFlagCreate()].		
opt	specifies whether you want to delete the event flag group only if there are no pending tasks (OS_DEL_NO_PEND) or whether you always want to delete the event flag group regardless of whether tasks are pending or not (OS DEL ALWAYS). In this case, all pending task are readied.		
perr	is a pointer to a variable that is used to hold an error code. The error code can be one of following:		
	OS_ERR_NONE	if the call is successful and the event flag group has been deleted.	
	OS_ERR_DEL_ISR	if you attempt to delete an event flag group from an ISR.	
	OS_ERR_FLAG_INVALID_PGRP	if you pass a NULL pointer in pgrp.	
	OS_ERR_EVENT_TYPE	if pgrp is not pointing to an event flag group.	
	OS_ERR_INVALID_OPT	if you do not specify one of the two options mentioned in the opt argument.	
	OS_ERR_TASK_WAITING	if one or more task are waiting on the event flag group and you specify OS_DEL_NO_PEND.	

### **Returned Values**

A NULL pointer if the event flag group is deleted or pgrp if the event flag group is not deleted. In the latter case, you need to examine the error code to determine the reason for the error.

# **Notes/Warnings**

- 1. You should use this call with care because other tasks might expect the presence of the event flag group.
- 2. This call can potentially disable interrupts for a long time. The interrupt-disable time is directly proportional to the number of tasks waiting on the event flag group.

# **Example**

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# OSFlagNameGet()

Chapter	File	Called from	Code enabled by	
New in V2.60	OS_FLAG.C	Task or ISR	OS_FLAG_NAME_EN	

OSFlagNameGet() allows you to obtain the name that you assigned to an event flag group. This function is typically used by a debugger to allow associating a name to a resource.

# **Arguments**

pgrp is a pointer to the event flag group.

pname is a pointer to a pointer to the name of the event flag group.

perr a pointer to an error code and can be any of the following:

OS ERR NONE If the call was successfull.

OS\_ERR\_EVENT\_TYPE You are not pointing to a flag group.

OS\_ERR\_PNAME\_NULL You passed a NULL pointer for pname.

OS\_ERR\_INVALID\_PGRP You passed a NULL pointer for pgrp.

### **Returned Values**

The size of the ASCII string pointed to by pname or 0 if an error is encountered.

# **Notes/Warnings**

1. The event flag group must be created before you can use this function and obtain the name of the resource.

# OSFlagNameSet()

Chapter	File	Called from	Code enabled by	
New in V2.60	OS_FLAG.C	Task	OS_FLAG_NAME_EN	

OSFlagNameSet() allows you to assign a name to an event flag group. This function is typically used by a debugger to allow associating a name to a resource.

# Arguments

pgrp is a pointer to the event flag group that you want to name. This pointer is returned to your

application when the event flag group is created (see OSFlagCreate()).

pname is a pointer to an ASCII string that contains the name of the event flag group.

perr a pointer to an error code and can be any of the following:

OS ERR NONE If the call was successfull.

OS ERR EVENT TYPE You are not pointing to an event flag group.

OS ERR PNAME NULL You passed a NULL pointer for pname.

OS\_ERR\_INVALID\_PGRP You passed a NULL pointer for pgrp.

OS\_ERR\_NAME\_SET\_ISR You called this function from an ISR.

### **Returned Values**

none

### **Notes/Warnings**

1. The event flag group must be created before you can use this function to set the name of the resource.

```
OS_FLAG_GRP *EngineStatus;

void Task (void *p_arg)
{
    INT8U err;

    (void)p_arg;
    for (;;) {
        OSFlagNameSet(EngineStatus, "Engine Status Flags", &err);
        .
        .
        .
     }
}
```

# OSFlagPend()

```
OS FLAGS OSFlagPend(OS FLAG GRP *pgrp,
                     OS FLAGS
                                   flags,
                     INT8U
                                   wait type,
                     INT32U
                                   timeout,
                     INT8U
                                  *perr);
```

Chapter	File	Called from	Code enabled by
9	OS_FLAG.C	Task only	OS_FLAG_EN

OSFlagPend() is used to have a task wait for a combination of conditions (i.e., events or bits) to be set (or cleared) in an event flag group. You application can wait for any condition to be set or cleared or for all conditions to be set or cleared. If the events that the calling task desires are not available, then the calling task is blocked until the desired conditions are satisfied or the specified timeout expires.

# Arguments

is a pointer to the event flag group. This pointer is returned to your application when the event pgrp

flag group is created [see OSFlagCreate()].

is a bit pattern indicating which bit(s) (i.e., flags) you wish to check. The bits you want are flags

specified by setting the corresponding bits in flags.

specifies whether you want all bits to be set/cleared or any of the bits to be set/cleared. You wait type

can specify the following arguments:

OS\_FLAG\_WAIT\_CLR ALL You check all bits in flags to be clear (0)

OS FLAG WAIT CLR ANY You check any bit in flags to be clear (0)

You check all bits in flags to be set (1) OS FLAG WAIT SET ALL

OS FLAG WAIT SET ANY You check any bit in flags to be set (1)

You can also specify whether the flags are consumed by adding OS FLAG CONSUME to the wait type. For example, to wait for any flag in a group and then clear the flags that satisfy the condition, set wait type to

OS FLAG WAIT SET ANY + OS FLAG CONSUME

timeout

allows the task to resume execution if the desired flag(s) is(are) not received from the event flag group within the specified number of clock ticks. A timeout value of 0 indicates that the task wants to wait forever for the flag(s). The timeout value is not synchronized with the clock tick. The timeout count begins decrementing on the next clock tick, which could potentially occur immediately.

is a pointer to an error code and can be: perr

> No error. OS ERR NONE

You try to call OSFlagPend from an ISR, which is not OS ERR PEND ISR

allowed.

You pass a NULL pointer instead of the event flag handle. OS ERR FLAG INVALID PGRP

OS ERR EVENT TYPE You are not pointing to an event flag group.

The flags are not available within the specified amount of OS ERR TIMEOUT

time.

You don't specify a proper wait\_type argument. OS ERR FLAG WAIT TYPE

# **Returned Values**

The flag(s) that cause the task to be ready or, 0 if either none of the flags are ready or an error occurred.

# **Notes/Warnings**

1. The event flag group must be created before it's used.

### **IMPORTANT**

The return value of OSFlagPend() is different as of V2.70. In previous versions, OSFlagPend() returned the current state of the flags and now, it returns the flag(s) that are ready, if any.

```
#define ENGINE_OIL_PRES_OK 0x01
#define ENGINE_OIL_TEMP_OK 0x02
#define ENGINE START
                       0x04
OS_FLAG_GRP *EngineStatus;
void Task (void *p_arg)
   INT8U err;
   OS FLAGS value;
    (void)p_arg;
    for (;;) {
       value = OSFlagPend(EngineStatus,
                          ENGINE_OIL_PRES_OK + ENGINE_OIL_TEMP_OK,
                          OS_FLAG_WAIT_SET_ALL + OS_FLAG_CONSUME,
                          10,
                          &err);
       switch (err) {
           case OS_ERR_NONE:
                /* Desired flags are available
                                                                 */
                break;
           case OS_ERR_TIMEOUT:
                /* The desired flags were NOT available before .. */
                /* .. 10 ticks occurred
                                                                 */
                break;
```

# OSFlagPendGetFlagsRdy()

OS\_FLAGS OSFlagPendGetFlagsRdy(void)

Chapter	File	Called from	Code enabled by
Added in V2.60 OS_FLAG.C		Task only	OS_FLAG_EN

 ${\tt OSFlagPendGetFlagsRdy()}\ is\ used\ to\ obtain\ the\ flags\ that\ caused\ the\ current\ task\ to\ become\ ready\ to\ run.\ In\ other\ words,\ this\ function\ allows\ you\ to\ know\ "Who\ done\ It!"$ 

# **Arguments**

None

### **Returned Value**

The value of the flags that caused the current task to become ready to run.

# **Notes/Warnings**

1. The event flag group must be created before it's used.

```
#define ENGINE_OIL_PRES_OK 0x01
#define ENGINE_OIL_TEMP_OK 0x02
#define ENGINE START
                             0x04
OS_FLAG_GRP *EngineStatus;
void Task (void *p_arg)
   INT8U
            err;
    OS_FLAGS value;
    (void)p_arg;
    for (;;) {
        value = OSFlagPend(EngineStatus,
                           ENGINE_OIL_PRES_OK + ENGINE_OIL_TEMP_OK,
                          OS_FLAG_WAIT_SET_ALL + OS_FLAG_CONSUME,
                          10,
                          &err);
        switch (err) {
            case OS ERR NONE:
                /* Find out who made task ready
                                                                  */
                flags = OSFlagPendGetFlagsRdy();
                break;
            case OS_ERR_TIMEOUT:
                 /* The desired flags were NOT available before .. */
                /* .. 10 ticks occurred
                                                                  */
                break;
```

# OSFlagPost()

```
OS_FLAGS OSFlagPost(OS_FLAG_GRP *pgrp,
OS_FLAGS flags,
INT8U opt,
INT8U *perr);
```

Chapter	File	Called from	Code enabled by
9	OS_FLAG.C	Task or ISR	OS_FLAG_EN

You set or clear event flag bits by calling OSFlagPost(). The bits set or cleared are specified in a *bit mask*. OSFlagPost() readies each task that has its desired bits satisfied by this call. You can set or clear bits that are already set or cleared.

# Arguments

is a pointer to the event flag group. This pointer is returned to your application when the event

flag group is created [see  ${\tt OSFlagCreate}$  ()].

specifies which bits you want set or cleared. If opt is OS\_FLAG\_SET, each bit that is set in

flags sets the corresponding bit in the event flag group. For example to set bits 0, 4, and 5, you set flags to 0x31 (note, bit 0 is the least significant bit). If opt is OS\_FLAG\_CLR, each bit that is set in flags will **clears** the corresponding bit in the event flag group. For example to clear bits

0, 4,and 5,you specify flags as 0x31 (note, bit 0 is the least significant bit).

opt indicates whether the flags are set ( $OS_FLAG_SET$ ) or cleared ( $OS_FLAG_CLR$ ).

perr is a pointer to an error code and can be:

OS ERR NONE The call is successful.

OS ERR FLAG INVALID PGRP You pass a NULL pointer.

OS ERR EVENT TYPE You are not pointing to an event flag group.

OS ERR FLAG INVALID OPT You specify an invalid option.

#### **Returned Value**

The new value of the event flags.

### **Notes/Warnings**

- 1. Event flag groups must be created before they are used.
- 2. The execution time of this function depends on the number of tasks waiting on the event flag group. However, the execution time is deterministic.
- 3. The amount of time interrupts are **disabled** also depends on the number of tasks waiting on the event flag group.

# OSFlagQuery()

Chapter	File	Called from	Code enabled by
9	OS_FLAG.C	Task or ISR	OS_FLAG_EN && OS_FLAG_QUERY_EN

OSFlagQuery() is used to obtain the current value of the event flags in a group. At this time, this function does **not** return the list of tasks waiting for the event flag group.

# Arguments

is a pointer to the event flag group. This pointer is returned to your application when the event

flag group is created [see OSFlagCreate()].

perr is a pointer to an error code and can be:

OS\_ERR\_NONE The call is successful.

 $\hbox{\tt OS\_ERR\_FLAG\_INVALID\_PGRP} \quad You\ pass\ a\ \hbox{\tt NULL}\ pointer.$ 

OS\_ERR\_EVENT\_TYPE You are not pointing to an event flag groups.

#### **Returned Value**

The state of the flags in the event flag group.

# **Notes/Warnings**

- 1. The event flag group to query must be created.
- 2. You can call this function from an ISR.

# OSInit()

void OSInit(void);

Chapter	File	Called from	Code enabled by
3	OS_CORE.C	Startup code only	N/A

 $\label{eq:osinit} \textit{OSInit()} \ \ initializes \ \mu C/OS-II \ and \ must \ be \ called \ prior \ to \ calling \ \textit{OSStart()}, \ which \ actually \ starts \ multitasking.$ 

# Arguments

none

# **Returned Values**

none

# **Notes/Warnings**

1. OSInit() must be called before OSStart().

```
void main (void)
{
          .
          .
          OSInit();     /* Initialize \( \mu C/OS-II */ \)
          .
          OSStart();     /* Start Multitasking */
}
```

# OSIntEnter()

void OSIntEnter(void);

Chapter	File	Called from	Code enabled by
3	OS_CORE.C	ISR only	N/A

OSIntEnter() notifies  $\mu$ C/OS-II that an ISR is being processed, which allows  $\mu$ C/OS-II to keep track of interrupt nesting. OSIntEnter() is used in conjunction with OSIntExit().

# **Arguments**

none

### **Returned Values**

none

### **Notes/Warnings**

- 1. This function must not be called by task-level code.
- 2. You can increment the interrupt-nesting counter (OSIntNesting) directly in your ISR to avoid the overhead of the function call/return. It's safe to increment OSIntNesting in your ISR because interrupts are assumed to be disabled when OSIntNesting needs to be incremented.
- 3. You are allowed to nest interrupts up to 255 levels deep.

# Example 1

### (Intel 80x86, real mode, large model)

Use OSIntEnter() for backward compatibility with  $\mu C/OS$ .

```
ISRx PROC
            FAR
     PUSHA
                                    ; Save interrupted task's context
     PUSH
            ES
     PUSH
            DS
     CALL
           FAR PTR OSIntEnter ; Notify \mu C/OS-II of start of ISR
     POP
            DS
                                    ; Restore processor registers
     POP
            ES
     POPA
     IRET
                                    ; Return from interrupt
ISRx ENDP
```

# Example 2 (Intel 80x86, real mode, large model)

```
ISRx
        PROC
               FAR
        PUSHA
                                       ; Save interrupted task's context
        PUSH
               ES
        PUSH
               DS
        MOV
               AX, SEG(_OSIntNesting) ; Reload DS
        MOV
               DS, AX
        INC
               BYTE PTR _OSIntNesting ; Notify \mu\text{C/OS-II} of start of ISR
        POP
               DS
                                     ; Restore processor registers
        POP
               ES
        POPA
        IRET
                                       ; Return from interrupt
ISRx
        ENDP
```

# OSIntExit()

void OSIntExit(void);

Chapter	File	Called from	Code enabled by
3	OS_CORE.C	ISR only	N/A

OSIntExit() notifies  $\mu$ C/OS-II that an ISR is complete, which allows  $\mu$ C/OS-II to keep track of interrupt nesting. OSIntExit() is used in conjunction with OSIntEnter(). When the last nested interrupt completes, OSIntExit() determines if a higher priority task is ready to run, in which case, the interrupt returns to the higher priority task instead of the interrupted task.

# **Arguments**

none

#### **Returned Value**

none

### **Notes/Warnings**

1. This function must not be called by task-level code. Also, if you decided to increment OSIntNesting, you still need to call OSIntExit().

# **Example**

(Intel 80x86, real mode, large model)

```
ISRx
        PROC
                 FAR
        PUSHA
                                      ; Save processor registers
        PUSH
                 ES
        PUSH
                 DS
                 FAR PTR _OSIntExit ; Notify \mu\text{C/OS-II} of end of ISR
        CALL
        POP
                 DS
                                      ; Restore processor registers
        POP
                 ES
        POPA
        IRET
                                      ; Return to interrupted task
ISRx
        ENDP
```

# OSMboxAccept()

void \*OSMboxAccept(OS EVENT \*pevent);

Chapter	File	Called from	Code enabled by
10	OS_MBOX.C	Task or ISR	OS_MBOX_EN && OS_MBOX_ACCEPT_EN

OSMboxAccept() allows you to see if a message is available from the desired mailbox. Unlike OSMboxPend(), OSMboxAccept() does not suspend the calling task if a message is not available. In other words, OSMboxAccept() is non-blocking. If a message is available, the message is returned to your application, and the content of the mailbox is cleared. This call is typically used by ISRs because an ISR is not allowed to wait for a message at a mailbox.

# Arguments

pevent

is a pointer to the mailbox from which the message is received. This pointer is returned to your application when the mailbox is created [see OSMboxCreate()].

#### **Returned Value**

A pointer to the message if one is available; NULL if the mailbox does not contain a message.

### **Notes/Warnings**

1. Mailboxes must be created before they are used.

# OSMboxCreate()

OS EVENT \*OSMboxCreate(void \*pmsg);

Chapter	File	Called from	Code enabled by
10	OS_MBOX.C	Task or startup code	OS_MBOX_EN

OSMboxCreate() creates and initializes a mailbox. A mailbox allows tasks or ISRs to send a pointer-sized variable (message) to one or more tasks.

# **Arguments**

pmsq

is used to initialize the contents of the mailbox. The mailbox is empty when pmsg is a NULL pointer. The mailbox initially contains a message when pmsg is non-NULL.

### **Returned Value**

A pointer to the event control block allocated to the mailbox. If no event control block is available, OSMboxCreate() returns a NULL pointer.

# **Notes/Warnings**

1. Mailboxes must be created before they are used.

# OSMboxDel()

Chapter	File	Called from	Code enabled by
10	OS_MBOX.C	Task	OS_MBOX_EN and
			OS_MBOX_DEL_EN

OSMboxDel() is used to delete a message mailbox. This function is dangerous to use because multiple tasks could attempt to access a deleted mailbox. You should always use this function with great care. Generally speaking, before you delete a mailbox, you must first delete all the tasks that can access the mailbox.

### Arguments

is a pointer to the mailbox. This pointer is returned to your application when the mailbox is created [see OSMboxCreate()].

specifies whether you want to delete the mailbox only if there are no pending tasks (OS\_DEL\_NO\_PEND) or whether you always want to delete the mailbox regardless of whether tasks are pending or not (OS\_DEL\_ALWAYS). In this case, all pending task are readied.

perr is a pointer to a variable that is used to hold an error code. The error code can be one of the following:

OS ERR NONE if the call is successful and the mailbox has been deleted.

OS ERR DEL ISR if you attempt to delete the mailbox from an ISR.

OS\_ERR\_INVALID\_OPT if you don't specify one of the two options mentioned in the

opt argument.

OS ERR TASK WAITING One or more tasks is waiting on the mailbox.

OS\_ERR\_EVENT\_TYPE if pevent is not pointing to a mailbox.

OS ERR PEVENT NULL if pevent is a NULL pointer.

#### **Returned Value**

A NULL pointer if the mailbox is deleted or pevent if the mailbox is not deleted. In the latter case, you need to examine the error code to determine the reason.

# **Notes/Warnings**

- 1. You should use this call with care because other tasks might expect the presence of the mailbox.
- 2. Interrupts are disabled when pended tasks are readied, which means that interrupt latency depends on the number of tasks that are waiting on the mailbox.
- 3. OSMboxAccept () callers do not know that the mailbox has been deleted.

# OSMboxPend()

Chapter	File	Called from	Code enabled by
10	OS_MBOX.C	Task only	OS_MBOX_EN

OSMboxPend() is used when a task expects to receive a message. The message is sent to the task either by an ISR or by another task. The message received is a pointer-sized variable, and its use is application specific. If a message is present in the mailbox when OSMboxPend() is called, the message is retrieved, the mailbox is emptied, and the retrieved message is returned to the caller. If no message is present in the mailbox, OSMboxPend() suspends the current task until either a message is received or a user-specified timeout expires. If a message is sent to the mailbox and multiple tasks are waiting for the message,  $\mu$ C/OS-II resumes the highest priority task waiting to run. A pended task that has been suspended with OSTaskSuspend() can receive a message. However, the task remains suspended until it is resumed by calling OSTaskResume().

### Arguments

pevent is a pointer to the mailbox from which the message is received. This pointer is returned to your

application when the mailbox is created [see OSMboxCreate()].

allows the task to resume execution if a message is not received from the mailbox within the specified number of clock ticks. A timeout value of 0 indicates that the task wants to wait

forever for the message. The timeout value is not synchronized with the clock tick. The timeout count begins decrementing on the next clock tick, which could potentially occur immediately.

perr is a pointer to a variable that holds an error code. OSMboxPend() sets \*perr to one of the

following:

OS\_ERR\_NONE if a message is received.

OS ERR TIMEOUT if a message is not received within the specified timeout period.

OS ERR PEND ABORT indicates that the pend was aborted by another task or ISR by

calling OSMboxPendAbort().

OS\_ERR\_EVENT\_TYPE if pevent is not pointing to a mailbox.

OS ERR PEND LOCKED if you called this function when the scheduler is locked.

OS ERR PEND ISR if you call this function from an ISR and μC/OS-II suspends it.

In general, you should not call OSMboxPend() from an ISR,

but µC/OS-II checks for this situation anyway.

OS\_ERR\_PEVENT\_NULL if pevent is a NULL pointer.

#### **Returned Value**

OSMboxPend() returns the message sent by either a task or an ISR, and \*perr is set to OS\_ERR\_NONE. If a message is not received within the specified timeout period, the returned message is a NULL pointer, and \*perr is set to OS\_ERR\_TIMEOUT.

### **Notes/Warnings**

- 1. Mailboxes must be created before they are used.
- 2. You should not call OSMboxPend() from an ISR.

```
OS_EVENT *CommMbox;
void CommTask(void *p_arg)
    INT8U err;
     void *pmsg;
     (void)p_arg;
     for (;;) {
        pmsg = OSMboxPend(CommMbox, 10, &err);
        if (err == OS_ERR_NONE) {
                                                              */
            . /* Code for received message
        } else {
            . /* Code for message not received within timeout */
```

# OSMboxPendAbort()

New Function	File	Called from	Code enabled by
V2.84	OS_MBOX.C	Task only	OS_MBOX_EN
			& &
			OS_MBOX_PEND_ABORT_EN

OSMboxPendAbort() aborts & readies any tasks currently waiting on a mailbox. This function should be used to fault-abort the wait on the mailbox, rather than to normally signal the mailbox via OSMboxPost() or OSMboxPostOpt().

#### **Arguments**

pevent is a pointer to the mailbox for which pend(s) need to be aborted. This pointer is returned to your

application when the mailbox is created [see OSMboxCreate()].

opt determines what type of abort is performed.

OS\_PEND\_OPT\_NONE Aborts the pend of only the highest priority task waiting on the

mailbox.

OS PEND OPT BROADCAST Aborts the pend of all the tasks waiting on the mailbox.

perr is a pointer to a variable that holds an error code. OSMboxPendAbort() sets \*perr to one of

the following:

OS ERR NONE if no tasks were waiting on the mailbox. In this case, the return

value is also 0.

OS ERR PEND ABORT at least one task waiting on the mailbox was readied and

informed of the aborted wait. Check the return value for the

number of tasks whose wait on the mailbox was aborted.

OS ERR EVENT TYPE if pevent is not pointing to a mailbox.

OS ERR PEVENT NULL if pevent is a NULL pointer.

#### **Returned Value**

OSMboxPendAbort () returns the number of tasks made ready to run by this function. Zero indicates that no tasks were pending on the mailbox and thus this function had no effect.

#### **Notes/Warnings**

1. Mailboxes must be created before they are used.

```
OS_EVENT *CommMbox;
void CommTask(void *p_arg)
     INT8U err;
     INT8U nbr_tasks;
     (void)p_arg;
      for (;;) {
        nbr_tasks = OSMboxPendAbort(CommMbox, OS_PEND_OPT_BROADCAST, &err);
        if (err == OS_ERR_NONE) {
             . /* No tasks were waiting on the mailbox
                                                                       */
        } else {
             . /* All pends of tasks waiting on mailbox were aborted \dots */
             . /* ... 'nbr_tasks' indicates how many were made ready. */
         }
```

# OSMboxPost()

Chapter	File	Called from	Code enabled by
10	OS_MBOX.C	Task or ISR	OS_MBOX_EN &&
			OS_MBOX_POST_EN

OSMboxPost() sends a message to a task through a mailbox. A message is a pointer-sized variable and, its use is application specific. If a message is already in the mailbox, an error code is returned indicating that the mailbox is full. OSMboxPost() then immediately returns to its caller, and the message is not placed in the mailbox. If any task is waiting for a message at the mailbox, the highest priority task waiting receives the message. If the task waiting for the message has a higher priority than the task sending the message, the higher priority task is resumed, and the task sending the message is suspended. In other words, a context switch occurs.

#### **Arguments**

pevent is a pointer to the mailbox into which the message is deposited. This pointer is returned to your

application when the mailbox is created [see OSMboxCreate()].

pmsg is the actual message sent to the task. pmsg is a pointer-sized variable and is application

specific. You must never post a NULL pointer because this pointer indicates that the mailbox is

empty.

### **Returned Value**

OSMboxPost () returns one of these error codes:

OS ERR NONE if the message is deposited in the mailbox.

OS ERR MBOX FULL if the mailbox already contains a message.

OS ERR EVENT TYPE if pevent is not pointing to a mailbox.

OS ERR PEVENT NULL if pevent is a NULL pointer.

OS ERR POST NULL PTR if you are attempting to post a NULL pointer. By convention a

NULL pointer is not supposed to point to anything.

### Notes/Warnings

- 1. Mailboxes must be created before they are used.
- 2. You must never post a NULL pointer because this pointer indicates that the mailbox is empty.

# OSMboxPostOpt()

INT8U OSMboxPostOpt(OS EVENT \*pevent, void \*pmsg, INT8U opt);

Chapter	File	Called from	Code enabled by
10	OS_MBOX.C	Task or ISR	OS_MBOX_EN and
			OS_MBOX_POST_OPT_EN

OSMboxPostOpt() works just like OSMboxPost() except that it allows you to post a message to multiple tasks. In other words, OSMboxPostOpt () allows the message posted to be broadcast to all tasks waiting on the mailbox. OSMboxPostOpt() can actually replace OSMboxPost() because it can emulate OSMboxPost().

OSMboxPostOpt() is used to send a message to a task through a mailbox. A message is a pointer-sized variable, and its use is application specific. If a message is already in the mailbox, an error code is returned indicating that the mailbox is full. OSMboxPostOpt () then immediately returns to its caller, and the message is not placed in the mailbox. If any task is waiting for a message at the mailbox, OSMboxPostOpt() allows you either to post the message to the highest priority task waiting at the mailbox (opt set to OS POST OPT NONE) or to all tasks waiting at the mailbox (opt is set to OS POST OPT BROADCAST). In either case, scheduling occurs and, if any of the tasks that receives the message have a higher priority than the task that is posting the message, then the higher priority task is resumed, and the sending task is suspended. In other words, a context switch occurs.

## **Arguments**

is a pointer to the mailbox. This pointer is returned to your application when the mailbox is pevent

created [see OSMboxCreate()].

is the actual message sent to the task(s). pmsq is a pointer-sized variable and is application pmsq

specific. You must never post a NULL pointer because this pointer indicates that the mailbox is

empty.

specifies whether you want to send the message to the highest priority task waiting at the opt

mailbox (when opt is set to OS POST OPT NONE) or to all tasks waiting at the mailbox (when

opt is set to OS POST OPT BROADCAST).

When set to OS POST OPT NO SCHED, the scheduler will not be called to see if a higher

priority task has been made ready to run.

Note that options are additive and thus, you can specify:

OS POST OPT BROADCAST | OS POST OPT NO SCHED

#### **Returned Value**

Returns one of the following error codes:

if the call is successful and the message has been sent. OS ERR NONE

OS ERR MBOX FULL if the mailbox already contains a message. You can only send

one message at a time to a mailbox, and thus the message must

be consumed before you are allowed to send another one.

if pevent is not pointing to a mailbox. OS ERR EVENT TYPE

OS ERR PEVENT NULL if pevent is a NULL pointer.

if you are attempting to post a NULL pointer. By convention, a OS ERR POST NULL PTR

NULL pointer is not supposed to point to anything.

## **Notes/Warnings**

- 1. Mailboxes must be created before they are used.
- 2. You must **never** post a NULL pointer to a mailbox because this pointer indicates that the mailbox is empty.
- 3. If you need to use this function and want to reduce code space, you can disable code generation of OSMboxPost() because OSMboxPostOpt() can emulate OSMboxPost().
- 4. The execution time of OSMboxPostOpt() depends on the number of tasks waiting on the mailbox if you set opt to OS\_POST\_OPT\_BROADCAST.

## **Example**

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# OSMboxQuery()

Chapter	File	Called from	Code enabled by	
10	OS_MBOX.C	Task or ISR	OS_MBOX_EN && OS_MBOX_QUERY_EN	

OSMboxQuery() obtains information about a message mailbox. Your application must allocate an OS\_MBOX\_DATA data structure, which is used to receive data from the event control block of the message mailbox. OSMboxQuery() allows you to determine whether any tasks are waiting for a message at the mailbox and how many tasks are waiting (by counting the number of 1s in the .OSEventTbl[] field). You can also examine the current contents of the mailbox. Note that the size of .OSEventTbl[] is established by the #define constant OS EVENT TBL SIZE (see uCOS II.H).

### **Arguments**

pevent is a pointer to the mailbox. This pointer is returned to your application when the mailbox is created [see OSMboxCreate()].

P mbox data is a pointer to a data structure of type OS MBOX DATA, which contains the following fields:

#### **Returned Value**

OSMboxQuery() returns one of these error codes:

```
OS_ERR_NONE if the call is successful.

OS_ERR_PEVENT_NULL if pevent is a NULL pointer.

OS_ERR_EVENT_TYPE if you don't pass a pointer to a message mailbox.

OS_ERR_PNAME_NULL You passed a NULL pointer for p mbox data.
```

#### **Notes/Warnings**

1. Message mailboxes must be created before they are used.

# OSMemCreate()

```
OS_MEM *OSMemCreate(void *addr,
INT32U nblks,
INT32U blksize,
INT8U *perr);
```

Chapter	File	Called from	Code enabled by
12	OS_MEM.C	Task or startup code	OS_MEM_EN

OSMemCreate() creates and initializes a memory partition. A memory partition contains a user-specified number of fixed-size memory blocks. Your application can obtain one of these memory blocks and, when done, release the block back to the partition.

# Arguments

Arguments				
addr	is the address of the start of a memory area that is used to create fixed-size memory blocks. Memory partitions can be created either using static arrays or $malloc()$ during startup. Note that the partition MUST align on a pointer boundary. Thus, if a pointer is 16 bits wide then the partition must start on a memory location with an address that ends with 0, 2, 4, 6, 8, etc. If a pointer is 32 bits wide then the partition must start on a memory location with and address that ends with 0, 4, 8 of C.			
nblks	contains the number of memory blocks available from the specified partition. You must specify at least two memory blocks per partition.			
blksize	specifies the size (in bytes) of each memory block within a partition. A memory block must be large enough to hold at least a pointer. Also, the size of a memory block must be a multiple of the size of a pointer. In other words, if a pointer is 32 bits wide then the block size must be 4, 8, 12, 16, 20, etc. bytes (i.e. a multiple of 4 bytes).			
perr	is a pointer to a variable that holds an error code. OSMemCreate() sets *perr to:			
	OS_ERR_NONE	if the memory partition is created successfully		
	OS_ERR_MEM_INVALID_ADDR if you are specifying an invalid address (i.e., addr is a pointer) or your partition is not properly aligned.			
	OS_ERR_MEM_INVALID_PART	if a free memory partition is not available		
	OS_ERR_MEM_INVALID_BLKS	if you don't specify at least two memory blocks per partition		
OS_ERR_MEM_INVALID_SIZE if you don't specify a block size that can contain at pointer variable and if it's not a multiple of a point variable.				

### **Returned Value**

OSMemCreate() returns a pointer to the created memory-partition control block if one is available. If no memory-partition control block is available, OSMemCreate() returns a NULL pointer.

### **Notes/Warnings**

1. Memory partitions must be created before they are used.

# OSMemGet()

Chapter	File	Called from	Code enabled by
12	OS_MEM.C	Task or ISR	OS_MEM_EN

OSMemGet obtains a memory block from a memory partition. It is assumed that your application knows the size of each memory block obtained. Also, your application must return the memory block [using OSMemPut()] when it no longer needs it. You can call OSMemGet() more than once until all memory blocks are allocated.

#### **Arguments**

pmem is a pointer to the memory-partition control block that is returned to your application from the

OSMemCreate() call.

perr is a pointer to a variable that holds an error code. OSMemGet() sets \*perr to one of the

following:

OS ERR NONE if a memory block is available and returned to your application.

OS\_ERR\_MEM\_NO\_FREE\_BLKS if the memory partition doesn't contain any more memory

blocks to allocate.

OS\_ERR\_MEM\_INVALID\_PMEM if pmem is a NULL pointer.

#### **Returned Value**

OSMemGet() returns a pointer to the allocated memory block if one is available. If no memory block is available from the memory partition, OSMemGet() returns a NULL pointer.

#### **Notes/Warnings**

1. Memory partitions must be created before they are used.

# OSMemNameGet()

Chapter	File	Called from	Code enabled by
New in V2.60	OS_MEM.C	Task	OS_MEM_NAME_EN

OSMemNameGet () allows you to obtain the name that you assigned to a memory partition. This function is typically used by a debugger to allow associating a name to a resource.

### Arguments

pmem is a pointer to the memory partition.

pname is a pointer to a pointer to an ASCII string that contains the name of the memory partition.

perr a pointer to an error code and can be any of the following:

OS ERR NONE If the call was successfull.

OS\_ERR\_INVALID\_PMEM You passed a NULL pointer for pmem.

OS\_ERR\_PNAME\_NULL You passed a NULL pointer for pname.

OS\_ERR\_NAME\_GET\_ISR You called this function from an ISR.

#### **Returned Values**

The size of the ASCII string pointed to by pname or 0 if an error is encountered.

## **Notes/Warnings**

1. The memory partition must be created before you can use this function and obtain the name of the resource.

```
OS_MEM *CommMem;
INT8U *CommMemName;

void Task (void *pdata)
{
    INT8U err;
    INT8U size;

    pdata = pdata;
    for (;;) {
        size = OSMemNameGet(CommMem, &CommMemName, &err);
        .
        .
     }
}
```

# OSMemNameSet()

Chapter	File	Called from	Code enabled by
New in V2.60	OS_MEM.C	Task	OS_MEM_NAME_EN

OSMemNameSet () allows you to assign a name to a memory partition. This function is typically used by a debugger to allow associating a name to a resource.

### Arguments

pmem is a pointer to the memory partition that you want to name. This pointer is returned to your

application when the memory partition is created (see OSMemCreate()).

pname is a pointer to an ASCII string that contains the name for the memory partition.

perr a pointer to an error code and can be any of the following:

OS\_ERR\_NONE If the call was successfull.

 $\hbox{\tt OS\_ERR\_MEM\_INVALID\_PMEM}\qquad You\ passed\ a\ \hbox{\tt NULL}\ pointer\ for\ pmem.$ 

 ${\tt OS\_ERR\_PNAME\_NULL} \qquad \qquad {\tt You~passed~a~NULL~pointer~for~pname}.$ 

 $\verb|OS_ERR_MEM_NAME_TOO_LONG| If the name is not able to fit in the specified storage.$ 

OS\_ERR\_NAME\_SET\_ISR You called this function from an ISR.

### **Returned Values**

none

### **Notes/Warnings**

1. The memory partition must be created before you can use this function to set the name of the resource.

```
OS_MEM *CommMem;

void Task (void *p_arg)
{
   INT8U err;

   (void)p_arg;
   for (;;) {
        OSMemNameSet(CommMem, "Comm. Buffer", &err);
        .
        .
        .
    }
}
```

# OSMemPut()

Chapter	File	Called from	Code enabled by
12	OS_MEM.C	Task or ISR	OS_MEM_EN

OSMemPut() returns a memory block to a memory partition. It is assumed that you return the memory block to the appropriate memory partition.

### Arguments

pmem is a pointer to the memory-partition control block that is returned to your application from the

OSMemCreate() call.

pblk is a pointer to the memory block to be returned to the memory partition.

#### **Returned Value**

OSMemPut() returns one of the following error codes:

OS\_ERR\_NONE if the memory block was returned to the memory partition.

OS\_ERR\_MEM\_FULL if the memory partition can not accept more memory blocks.

This code is surely an indication that something is wrong because you are returning more memory blocks than you

obtained using OSMemGet().

OS\_ERR\_MEM\_INVALID\_PMEM if pmem is a NULL pointer.

OS ERR MEM INVALID PBLK if pblk is a NULL pointer.

### **Notes/Warnings**

- 1. Memory partitions must be created before they are used.
- 2. You must return a memory block to the proper memory partition.

# OSMemQuery()

Chapter	File	Called from	Code enabled by
12	OS_MEM.C	Task or ISR	OS_MEM_EN && OS_MEM_QUERY_EN

OSMemQuery() obtains information about a memory partition. Basically, this function returns the same information found in the OS\_MEM data structure but in a new data structure called OS\_MEM\_DATA. OS\_MEM\_DATA also contains an additional field that indicates the number of memory blocks in use.

### Arguments

pmem is a pointer to the memory-partition control block that is returned to your application from the OSMemCreate() call.

p mem data is a pointer to a data structure of type OS MEM DATA, which contains the following fields

```
void *OSAddr; /* Points to beginning address of the memory partition */
void *OSFreeList; /* Points to beginning of the free list of memory blocks */
INT32U OSBlkSize; /* Size (in bytes) of each memory block */
INT32U OSNBlks; /* Total number of blocks in the partition */
INT32U OSNFree; /* Number of memory blocks free */
INT32U OSNUsed; /* Number of memory blocks used */
```

#### **Returned Value**

OSMemQuery() returns one of the following error codes:

```
OS_ERR_NONE if *p_mem_data was filled successfully.

OS_ERR_MEM_INVALID_PMEM if pmem is a NULL pointer.

OS_ERR_MEM_INVALID_PDATA if pdata is a NULL pointer.
```

### **Notes/Warnings**

1. Memory partitions must be created before they are used.

# OSMutexAccept()

Chapter	File	Called from	Code enabled by
8	OS MUTEX.C	Task	OS MUTEX EN

OSMutexAccept() allows you to check to see if a resource is available. Unlike OSMutexPend(), OSMutexAccept() does not suspend the calling task if the resource is not available. In other words, OSMutexAccept() is non-blocking.

#### **Arguments**

pevent is a pointer to the mutex that guards the resource. This pointer is returned to your application

when the mutex is created [see OSMutexCreate()].

perr is a pointer to a variable used to hold an error code. OSMutexAccept() sets \*perr to one of

the following:

OS\_ERR\_NONE if the call is successful.

OS\_ERR\_EVENT\_TYPE if pevent is not pointing to a mutex.

OS\_ERR\_PEVENT\_NULL if pevent is a NULL pointer.

 ${\tt OS\_ERR\_PEND\_ISR} \qquad \qquad \text{if you call } {\tt OSMutexAccept()} \ \ \text{from an } ISR.$ 

OS ERR PIP LOWER If the priority of the task that owns the Mutex is HIGHER (i.e.

a lower number) than the PIP. This error indicates that you did not set the PIP higher (lower number) than ALL the tasks that compete for the Mutex. Unfortunately, this is something that could not be detected when the Mutex is created because we

don't know what tasks will be using the Mutex.

### **Returned Value**

If the mutex is available, OSMutexAccept() returns OS\_TRUE. If the mutex is owned by another task, OSMutexAccept() returns OS FALSE.

#### **Notes/Warnings**

- 1. Mutexes must be created before they are used.
- 2. This function **must not** be called by an ISR.
- 3. If you acquire the mutex through OSMutexAccept(), you must call OSMutexPost() to release the mutex when you are done with the resource.

# OSMutexCreate()

Chapter	File	Called from	Code enabled by
8	OS_MUTEX.C	Task or startup code	OS_MUTEX_EN

OSMutexCreate() is used to create and initialize a mutex. A mutex is used to gain exclusive access to a resource.

#### Arguments

prio is the priority inheritance priority (PIP) that is used when a high priority task attempts to acquire

the mutex that is owned by a low priority task. In this case, the priority of the low priority task

is *raised* to the PIP until the resource is released.

perr is a pointer to a variable that is used to hold an error code. The error code can be one of the

following:

OS ERR NONE if the call is successful and the mutex has been created.

OS\_ERR\_CREATE\_ISR if you attempt to create a mutex from an ISR.

OS ERR PRIO EXIST if a task at the specified priority inheritance priority already

exists.

OS ERR PEVENT NULL if no more OS EVENT structures are available.

OS ERR PRIO INVALID if you specify a priority with a higher number than

OS LOWEST PRIO.

#### **Returned Value**

A pointer to the event control block allocated to the mutex. If no event control block is available, OSMutexCreate() returns a NULL pointer.

#### **Notes/Warnings**

- 1. Mutexes must be created before they are used.
- 2. You **must** make sure that prio has a higher priority than **any** of the tasks that use the mutex to access the resource. For example, if three tasks of priority 20, 25, and 30 are going to use the mutex, then prio must be a number **lower** than 20. In addition, there **must not** already be a task created at the specified priority.

# OSMutexDel()

```
OS_EVENT *OSMutexDel(OS_EVENT *pevent, INT8U opt, INT8U *perr);
```

Chapter	File	Called from	Code enabled by
8	OS_MUTEX.C	Task	OS_MUTEX_EN and
			OS_MUTEX_DEL_EN

OSMutexDel() is used to delete a mutex. This function is dangerous to use because multiple tasks could attempt to access a deleted mutex. You should always use this function with great care. Generally speaking, before you delete a mutex, you must first delete all the tasks that can access the mutex.

### Arguments

pevent	is a pointer to the mutex. This pointer is returned to your application when the mutex is created [see OSMutexCreate()].		
opt	specifies whether you want to delete the mutex only if there are no pending tasks (OS_DEL_NO_PEND) or whether you always want to delete the mutex regardless of whether tasks are pending or not (OS_DEL_ALWAYS). In this case, all pending task are readied.		
perr	is a pointer to a variable that i following:	s used to hold an error code. The error code can be one of the	
	OS_ERR_NONE	if the call is successful and the mutex has been deleted.	
	OS_ERR_DEL_ISR	if you attempt to delete a mutex from an ISR.	
	OS_ERR_INVALID_OPT	if you don't specify one of the two options mentioned in the opt argument.	
	OS_ERR_TASK_WAITING	if one or more task are waiting on the mutex and you specify $OS\_DEL\_NO\_PEND$ .	
	OS_ERR_EVENT_TYPE	if pevent is not pointing to a mutex.	
	OS_ERR_PEVENT_NULL	if pevent is a NULL pointer.	

#### **Returned Value**

A NULL pointer if the mutex is deleted or pevent if the mutex is not deleted. In the latter case, you need to examine the error code to determine the reason.

#### **Notes/Warnings**

1. You should use this call with care because other tasks might expect the presence of the mutex.

# OSMutexPend()

Chapter	File	Called from	Code enabled by
8	OS_MUTEX.C	Task only	OS_MUTEX_EN

OSMutexPend() is used when a task desires to get exclusive access to a resource. If a task calls OSMutexPend() and the mutex is available, then OSMutexPend() gives the mutex to the caller and returns to its caller. Note that nothing is actually given to the caller except for the fact that if perr is set to OS\_ERR\_NONE, the caller can assume that it owns the mutex. However, if the mutex is already owned by another task, OSMutexPend() places the calling task in the wait list for the mutex. The task thus waits until the task that owns the mutex releases the mutex and thus the resource or until the specified timeout expires. If the mutex is signaled before the timeout expires,  $\mu$ C/OS-II resumes the highest priority task that is waiting for the mutex. Note that if the mutex is owned by a lower priority task, then OSMutexPend() raises the priority of the task that owns the mutex to the PIP, as specified when you created the mutex [see OSMutexCreate()].

## Arguments

pevent is a pointer to the mutex. This pointer is returned to your application when the mutex is created

[see OSMutexCreate()].

timeout is used to allow the task to resume execution if the mutex is not signaled (i.e., posted to) within

the specified number of clock ticks. A timeout value of 0 indicates that the task desires to wait forever for the mutex. The timeout value is not synchronized with the clock tick. The timeout count starts being decremented on the next clock tick, which could potentially occur

immediately.

perr is a pointer to a variable that is used to hold an error code. OSMutexPend() sets \*perr to one

of the following:

OS ERR NONE if the call is successful and the mutex is available.

OS\_ERR\_TIMEOUT if the mutex is not available within the specified timeout.

OS ERR EVENT TYPE if you don't pass a pointer to a mutex to OSMutexPend().

OS ERR PEVENT NULL if pevent is a NULL pointer.

OS ERR PEND LOCKED if you called this function when the scheduler is locked

OS ERR PEND ABORT if OSMutexPend() was aborted by another task

OS ERR PEND ISR if you attempt to acquire the mutex from an ISR.

OS ERR PIP LOWER If the priority of the task that owns the Mutex is HIGHER (i.e.

a lower number) than the PIP. This error indicates that you did not set the PIP higher (lower number) than ALL the tasks that compete for the Mutex. Unfortunately, this is something that could not be detected when the Mutex is created because we

don't know what tasks will be using the Mutex.

#### **Returned Value**

none

## **Notes/Warnings**

- 1. Mutexes must be created before they are used.
- 2. You should **not** suspend the task that owns the mutex, have the mutex owner wait on any other  $\mu$ C/OS-II objects (i.e., semaphore, mailbox, or queue), and delay the task that owns the mutex. In other words, your code should hurry up and release the resource as quickly as possible.

## **Example**

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# OSMutexPost()

INT8U OSMutexPost(OS EVENT \*pevent);

Chapter	File	Called from	Code enabled by
8	OS_MUTEX.C	Task	OS_MUTEX_EN

A mutex is signaled (i.e., released) by calling OSMutexPost(). You call this function only if you acquire the mutex by first calling either OSMutexAccept() or OSMutexPend(). If the priority of the task that owns the mutex has been raised when a higher priority task attempts to acquire the mutex, the original task priority of the task is restored. If one or more tasks are waiting for the mutex, the mutex is given to the highest priority task waiting on the mutex. The scheduler is then called to determine if the awakened task is now the highest priority task ready to run, and if so, a context switch is done to run the readied task. If no task is waiting for the mutex, the mutex value is simply set to available (0xff).

### **Arguments**

pevent

is a pointer to the mutex. This pointer is returned to your application when the mutex is created [see OSMutexCreate()].

#### **Returned Value**

OSMutexPost() returns one of these error codes:

OS_ERR_NONE	if the call is successful and the mutex is released.
OS_ERR_EVENT_TYPE	if you don't pass a pointer to a mutex to ${\tt OSMutexPost}$ () .
OS_ERR_PEVENT_NULL	if pevent is a NULL pointer.
OS_ERR_POST_ISR	if you attempt to call ${\tt OSMutexPost}$ () from an ISR.
OS_ERR_NOT_MUTEX_OWNER	if the task posting (i.e., signaling the mutex) doesn't actually own the mutex.
OS_ERR_PIP_LOWER	If the priority of the new task that owns the Mutex is HIGHER (i.e. a lower number) than the PIP. This error indicates that you did not set the PIP higher (lower number) than ALL the tasks that compete for the Mutex. Unfortunately, this is something that could not be detected when the Mutex is created because we don't know what tasks will be using the Mutex.

### **Notes/Warnings**

- 1. Mutexes must be created before they are used.
- 2. You cannot call this function from an ISR.

```
OS_EVENT *DispMutex;
void TaskX (void *p_arg)
   INT8U err;
    (void)p_arg;
   for (;;) {
       err = OSMutexPost(DispMutex);
       switch (err) {
          case OS_ERR_NONE: /* Mutex signaled */
              break;
          case OS_ERR_EVENT_TYPE:
              break;
          case OS_ERR_PEVENT_NULL:
              break;
          case OS_ERR_POST_ISR:
               break;
```

# OSMutexQuery()

Chapter	File	Called from	Code enabled by
8	OS_MUTEX.C	Task	OS_MUTEX_EN && OS_MUTEX_QUERY_EN

OSMutexQuery() is used to obtain run-time information about a mutex. Your application must allocate an OS\_MUTEX\_DATA data structure that is used to receive data from the event control block of the mutex. OSMutexQuery() allows you to determine whether any task is waiting on the mutex, how many tasks are waiting (by counting the number of 1s) in the .OSEventTbl[] field, obtain the PIP, and determine whether the mutex is available (OS\_TRUE) or not (OS\_FALSE). Note that the size of .OSEventTbl[] is established by the #define constant OS\_EVENT\_TBL\_SIZE (see uCOS\_II.H).

#### Arguments

pevent is a pointer to the mutex. This pointer is returned to your application when the mutex is created [see OSMutexCreate()].

p\_mutex\_data is a pointer to a data structure of type OS\_MUTEX\_DATA, which contains the following fields

```
INT8U
                      /* The PIP of the mutex
       OSMutexPIP;
INT8U OSOwnerPrio; /* The priority of the mutex owner
                                                                     */
BOOLEAN OSValue;
                      /* The current mutex value
                                                                     */
                      /* OS TRUE means available
                                                                     * /
                      /* OS FALSE means unavailable
#if OS LOWEST PRIO <= 63
       OSEventGrp;
                      /* Copy of the mutex wait list
INT8U
       OSEventTbl[OS EVENT TBL SIZE];
INT8U
#else
                      /* Copy of the mutex wait list
INT16U OSEventGrp;
INT16U OSEventTbl[OS EVENT TBL SIZE];
#endif
```

#### **Returned Value**

OSMutexQuery() returns one of these error codes:

```
OS_ERR_NONE if the call is successful.

OS_ERR_EVENT_TYPE if you don't pass a pointer to a mutex to OSMutexQuery().

OS_ERR_PEVENT_NULL if pevent is a NULL pointer.

OS_ERR_PDATA_NULL if p_mutex_data is a NULL pointer.

OS_ERR_QUERY_ISR if you attempt to call OSMutexQuery() from an ISR.
```

### **Notes/Warnings**

- 1. Mutexes must be created before they are used.
- You cannot call this function from an ISR.

In this example, we check the contents of the mutex to determine the highest priority task that is waiting for it.

```
OS_EVENT *DispMutex;
void Task (void *p_arg)
   OS_MUTEX_DATA mutex_data;
   INT8U
               err;
   INT8U
               highest; /* Highest priority task waiting on mutex
   INT8U
               х;
   INT8U
          у;
   (void)p_arg;
   for (;;) {
       err = OSMutexQuery(DispMutex, &mutex_data);
       if (err == OS_ERR_NONE) {
           /* Examine Mutex data */
           }
       }
```

# OSQAccept()

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task or ISR	OS_Q_EN

OSQAccept() checks to see if a message is available in the desired message queue. Unlike OSQPend(), OSQAccept() does not suspend the calling task if a message is not available. In other words, OSQAccept() is non-blocking. If a message is available, it is extracted from the queue and returned to your application. This call is typically used by ISRs because an ISR is not allowed to wait for messages at a queue.

### Arguments

pevent is a pointer to the message queue from which the message is received. This pointer is returned to

your application when the message queue is created [see OSQCreate()].

perr is a pointer to a variable that is used to hold an error code. OSQAccept () sets \*perr to one of

the following:

OS ERR NONE if the call is successful and a message is available in the desired

message queue.

OS ERR EVENT TYPE if you don't pass a pointer to a queue to OSQAccept ().

OS\_ERR\_PEVENT\_NULL if pevent is a NULL pointer.

OS ERR Q EMPTY if the queue doesn't contain any messages.

#### **Returned Value**

A pointer to the message if one is available; NULL if the message queue does not contain a message or the message received is a NULL pointer. If a message was available in the queue, it will be removed before OSQAccept() returns.

#### **Notes/Warnings**

- 1. Message queues must be created before they are used.
- 2. The API (Application Programming Interface) has changed for this function in V2.60 because you can now post NULL pointers to queues. Specifically, the perr argument has been added to the call.

# OSQCreate()

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task or startup code	OS_Q_EN

OSQCreate() creates a message queue. A message queue allows tasks or ISRs to send pointer-sized variables (messages) to one or more tasks. The meaning of the messages sent are application specific.

### Arguments

start is the base address of the message storage area. A message storage area is declared as an array

of pointers to voids.

size is the size (in number of entries) of the message storage area.

#### **Returned Value**

OSQCreate() returns a pointer to the event control block allocated to the queue. If no event control block or no queue control block is available, OSQCreate() returns a NULL pointer.

## **Notes/Warnings**

1. Queues must be created before they are used.

# OSQDel()

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task	OS_Q_EN and OS_Q_DEL_EN

OSQDel() is used to delete a message queue. This function is dangerous to use because multiple tasks could attempt to access a deleted queue. You should always use this function with great care. Generally speaking, before you delete a queue, you must first delete all the tasks that can access the queue.

is a pointer to the queue. This pointer is returned to your application when the queue is created

if pevent is not pointing to a queue.

if pevent is a NULL pointer.

### **Arguments**

pevent

	[see OSQCreate()].		
opt	(OS_DEL_NO_PEND) or whether	to delete the queue only if there are no pending tasks or you always want to delete the queue regardless of whether tasks LWAYS). In this case, all pending task are readied.	
perr	is a pointer to a variable that is used to hold an error code. The error code can be one of the following:		
	OS_ERR_NONE	if the call is successful and the queue has been deleted.	
	OS_ERR_DEL_ISR	if you attempt to delete the queue from an ISR.	
	OS_ERR_INVALID_OPT	if you don't specify one of the two options mentioned in the opt argument.	
	OS_ERR_TASK_WAITING	if one or more tasks are waiting for messages at the message	

OS ERR EVENT TYPE

OS\_ERR\_PEVENT\_NULL

A NULL pointer if the queue is deleted or pevent if the queue is not deleted. In the latter case, you need to examine the error code to determine the reason.

queue.

### **Notes/Warnings**

**Returned Value** 

- 1. You should use this call with care because other tasks might expect the presence of the queue.
- 2. Interrupts are disabled when pended tasks are readied, which means that interrupt latency depends on the number of tasks that are waiting on the queue.

# OSQFlush()

INT8U \*OSQFlush(OS EVENT \*pevent);

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task or ISR	OS_Q_EN && OS_Q_FLUSH_EN

OSQFlush () empties the contents of the message queue and eliminates all the messages sent to the queue. This function takes the same amount of time to execute regardless of whether tasks are waiting on the queue (and thus no messages are present) or the queue contains one or more messages.

### **Arguments**

pevent

is a pointer to the message queue. This pointer is returned to your application when the message queue is created [see OSQCreate()].

### **Returned Value**

OSQFlush() returns one of the following codes:

```
OS_ERR_NONE if the message queue is flushed.

OS_ERR_EVENT_TYPE if you attempt to flush an object other than a message queue.

OS_ERR_PEVENT_NULL if pevent is a NULL pointer.
```

### **Notes/Warnings**

- 1. Queues must be created before they are used.
- 2. You should use this function with great care because, when to flush the queue, you LOOSE the references to what the queue entries are pointing to and thus, you could cause 'memory leaks'. In other words, the data you are pointing to that's being referenced by the queue entries should, most likely, need to be de-allocated (i.e. freed). To flush a queue that contains entries, you should instead repeateadly use OSQAccept ().

## OSQPend()

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task only	OS_Q_EN

OSQPend() is used when a task wants to receive messages from a queue. The messages are sent to the task either by an ISR or by another task. The messages received are pointer-sized variables, and their use is application specific. If at least one message is present at the queue when OSQPend() is called, the message is retrieved and returned to the caller. If no message is present at the queue, OSQPend() suspends the current task until either a message is received or a user-specified timeout expires. If a message is sent to the queue and multiple tasks are waiting for such a message, then  $\mu C/OS-II$  resumes the highest priority task that is waiting. A pended task that has been suspended with OSTaskSuspend() can receive a message. However, the task remains suspended until it is resumed by calling OSTaskResume().

### **Arguments**

pevent is a pointer to the queue from which the messages are received. This pointer is returned to your

application when the queue is created [see OSQCreate()].

allows the task to resume execution if a message is not received from the mailbox within the specified number of clock ticks. A timeout value of 0 indicates that the task wants to wait forever for the message. The timeout value is not synchronized with the clock tick. The timeout

count starts decrementing on the next clock tick, which could potentially occur immediately.

is a pointer to a variable used to hold an error code. OSQPend() sets \*perr to one of the

following:

OS\_ERR\_NONE if a message is received.

OS ERR TIMEOUT if a message is not received within the specified timeout.

OS ERR EVENT TYPE if pevent is not pointing to a message queue.

OS\_ERR\_PEVENT\_NULL if pevent is a NULL pointer.

OS ERR PEND ABORT if OSQPend() was aborted by another task who called

OSQPendAbort().

os err pend isr if you call this function from an ISR and  $\mu$ C/OS-II has to

suspend it. In general, you should not call OSQPend() from an

ISR. µC/OS-II checks for this situation anyway.

OS ERR PEND LOCKED if you called this function when the scheduler is locked.

### **Returned Value**

OSQPend() returns a message sent by either a task or an ISR, and \*perr is set to OS\_ERR\_NONE. If a timeout occurs, OSQPend() returns a NULL pointer and sets \*perr to OS\_ERR\_TIMEOUT.

### **Notes/Warnings**

- 1. Queues must be created before they are used.
- 2. You should not call OSQPend() from an ISR.
- 3. OSQPend() was changed in V2.60 to allow it to receive NULL pointer messages.

```
OS_EVENT *CommQ;
void CommTask(void *p_arg)
    INT8U err;
    void *pmsg;
     (void)p_arg;
     for (;;) {
       pmsg = OSQPend(CommQ, 100, &err);
        if (err == OS_ERR_NONE) {
         . /* Message received within 100 ticks! */
        } else {
                  /* Message not received, must have timed out */
```

## OSQPendAbort()

New Function	File	Called from	Code enabled by
V2.84	os_Q.c	Task only	OS_Q_EN &&
			OS_Q_PEND_ABORT_EN

OSQPendAbort() aborts & readies any tasks currently waiting on a queue. This function should be used to fault-abort the wait on the queue, rather than to normally signal the queue via OSQPost(), OSQPostFront() or OSQPostOpt().

### **Arguments**

pevent is a pointer to the queue for which pend(s) need to be aborted. This pointer is returned to your

application when the queue is created [see OSQCreate()].

opt determines what type of abort is performed.

OS\_PEND\_OPT\_NONE Aborts the pend of only the highest priority task waiting on the

queue.

OS PEND OPT BROADCAST Aborts the pend of all the tasks waiting on the queue.

perr is a pointer to a variable that holds an error code. OSQPendAbort () sets \*perr to one of the

following:

OS ERR NONE if no tasks were waiting on the queue. In this case, the return

value is also 0.

OS ERR PEND ABORT at least one task waiting on the queue was readied and informed

of the aborted wait. Check the return value for the number of

tasks whose wait on the queue was aborted.

OS\_ERR\_EVENT\_TYPE if pevent is not pointing to a queue.

OS ERR PEVENT NULL if pevent is a NULL pointer.

### **Returned Value**

OSQPendAbort () returns the number of tasks made ready to run by this function. Zero indicates that no tasks were pending on the queue and thus this function had no effect.

### Notes/Warnings

1. Queues must be created before they are used.

```
OS_EVENT *CommQ;
void CommTask(void *p_arg)
     INT8U err;
     INT8U nbr_tasks;
     (void)p_arg;
      for (;;) {
        nbr_tasks = OSQPendAbort(CommQ, OS_PEND_OPT_BROADCAST, &err);
        if (err == OS_ERR_NONE) {
            . /* No tasks were waiting on the queue
                                                                   */
        } else {
             . /* All pends of tasks waiting on queue were aborted \dots */
            . /* ... 'nbr_tasks' indicates how many were made ready. */
        }
```

## OSQPost()

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task or ISR	OS_Q_EN && OS_Q_POST_EN

OSQPOST () sends a message to a task through a queue. A message is a pointer-sized variable, and its use is application specific. If the message queue is full, an error code is returned to the caller. In this case, OSQPOST () immediately returns to its caller, and the message is not placed in the queue. If any task is waiting for a message at the queue, the highest priority task receives the message. If the task waiting for the message has a higher priority than the task sending the message, the higher priority task resumes, and the task sending the message is suspended; that is, a context switch occurs. Message queues are first-in first-out (FIFO), which means that the first message sent is the first message received.

### Arguments

pevent is a pointer to the queue into which the message is deposited. This pointer is returned to your

application when the queue is created [see OSQCreate()].

pmsg is the actual message sent to the task. pmsg is a pointer-sized variable and is application

specific. As of V2.60, you are allowed to post a NULL pointer.

### **Returned Value**

OSQPost () returns one of these error codes:

OS\_ERR\_NONE if the message is deposited in the queue.

OS ERR Q FULL if the queue is already full.

OS\_ERR\_EVENT\_TYPE if pevent is not pointing to a message queue.

OS ERR PEVENT NULL if pevent is a NULL pointer.

### **Notes/Warnings**

- 1. Queues must be created before they are used.
- 2. As of V2.60, you are now allowed to post a NULL pointer. It is up to you're application to check the perr variable accordingly.

```
OS_EVENT *CommQ;
INT8U CommRxBuf[100];
void CommTaskRx (void *p_arg)
    INT8U err;
      (void)p_arg;
     for (;;) {
        err = OSQPost(CommQ, (void *)&CommRxBuf[0]);
        switch (err) {
            case OS_ERR_NONE:
                /* Message was deposited into queue */
                break;
            case OS_ERR_Q_FULL:
               /* Queue is full
                                                      */
                Break;
```

## OSQPostFront()

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task or ISR	OS_Q_EN && OS_Q_POST_FRONT_EN

OSQPOSTFront() sends a message to a task through a queue. OSQPOSTFront() behaves very much like OSQPOSTFront(), except that the message is inserted at the front of the queue. This means that OSQPOSTFront() makes the message queue behave like a last-in first-out (LIFO) queue instead of a first-in first-out (FIFO) queue. The message is a pointer-sized variable, and its use is application specific. If the message queue is full, an error code is returned to the caller. OSQPOSTFront() immediately returns to its caller, and the message is not placed in the queue. If any tasks are waiting for a message at the queue, the highest priority task receives the message. If the task waiting for the message has a higher priority than the task sending the message, the higher priority task is resumed, and the task sending the message is suspended; that is, a context switch occurs.

### Arguments

pevent is a pointer to the queue into which the message is deposited. This pointer is returned to your

application when the queue is created [see OSQCreate()].

pmsg is the actual message sent to the task. pmsg is a pointer-sized variable and is application

specific. As of V2.60, you are allowed to post a NULL pointer.

#### **Returned Value**

OSQPostFront() returns one of these error codes:

OS\_ERR\_NONE if the message is deposited in the queue.

OS\_ERR\_Q\_FULL if the queue is already full.

OS\_ERR\_EVENT\_TYPE if pevent is not pointing to a message queue.

OS ERR PEVENT NULL if pevent is a NULL pointer.

### **Notes/Warnings**

- 1. Queues must be created before they are used.
- 2. As of V2.60, you are now allowed to post a NULL pointer. It is up to you're application to check the perr variable accordingly.

```
OS_EVENT *CommQ;
INT8U CommRxBuf[100];
void CommTaskRx (void *p_arg)
    INT8U err;
      (void)p_arg;
     for (;;) {
        err = OSQPostFront(CommQ, (void *)&CommRxBuf[0]);
        switch (err) {
            case OS_ERR_NONE:
                /* Message was deposited into queue */
                break;
            case OS_ERR_Q_FULL:
               /* Queue is full
                                                      */
                break;
```

# OSQPostOpt()

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task or ISR	OS_Q_EN && OS_Q_POST_OPT_EN

OSQPOSTOPT() is used to send a message to a task through a queue. A message is a pointer-sized variable, and its use is application specific. If the message queue is full, an error code is returned indicating that the queue is full. OSQPOSTOPT() then immediately returns to its caller, and the message is not placed in the queue. If any task is waiting for a message at the queue, OSQPOSTOPT() allows you to either post the message to the highest priority task waiting at the queue (opt set to OS\_POST\_OPT\_NONE) or to all tasks waiting at the queue (opt is set to OS\_POST\_OPT\_BROADCAST). In either case, scheduling occurs, and, if any of the tasks that receive the message have a higher priority than the task that is posting the message, then the higher priority task is resumed, and the sending task is suspended. In other words, a context switch occurs.

OSQPostOpt() emulates both OSQPost() and OSQPostFront() and also allows you to post a message to **multiple** tasks. In other words, it allows the message posted to be broadcast to **all** tasks waiting on the queue. OSQPostOpt() can actually replace OSQPost() and OSQPostFront() because you specify the mode of operation via an option argument, opt. Doing this allows you to reduce the amount of code space needed by  $\mu C/OS-II$ .

### **Arguments**

is a pointer to the queue. This pointer is returned to your application when the queue is created pevent [see OSQCreate()]. is the actual message sent to the task(s). pmsg is a pointer-sized variable, and what pmsg points pmsg to is application specific. As of V2.60, you are now allowed to post a NULL pointer. determines the type of POST performed: opt POST to a single waiting task [identical to OSQPost()]. OS POST OPT NONE OS POST OPT BROADCAST POST to all tasks waiting on the queue. OS POST OPT FRONT POST as LIFO [simulates OSQPostFront()]. Do not call the scheduler after the post. OS POST OPT NO SCHED Below is a list of some of the possible combination of these flags: OS\_POST\_OPT\_NONE is identical to OSQPost() is identical to OSQPostFront() OS POST OPT FRONT is identical to OSQPost() but broadcasts pmsq to all waiting OS POST OPT BROADCAST tasks OS POST OPT FRONT + OS POST OPT BROADCAST is identical to OSQPostFront() except that broadcasts pmsg to all waiting tasks. OS POST OPT FRONT + OS POST OPT BROADCAST + OS POST OPT NO SCHED is identical to OSQPostFront() except that broadcasts pmsg

to all waiting tasks and the scheduler will not be called

### **Returned Value**

OSQPostOpt returns one of the following error codes:

```
OS_ERR_NONE if the call is successful and the message has been sent.

OS_ERR_Q_FULL if the queue can no longer accept messages because it is full.

OS_ERR_EVENT_TYPE if pevent is not pointing to a mailbox.

OS_ERR_PEVENT_NULL if pevent is a NULL pointer.
```

### **Notes/Warnings**

- 1. Queues must be created before they are used.
- 2. If you need to use this function and want to reduce code space, you can disable code generation of OSQPost() (set OS\_Q\_POST\_EN to 0 in OS\_CFG.H) and OSQPOSTFront() (set OS\_Q\_POST\_FRONT\_EN to 0 in OS\_CFG.H) because OSQPOSTOPt() can emulate these two functions.
- 3. The execution time of OSQPostOpt() depends on the number of tasks waiting on the queue if you set opt to OS\_POST\_OPT\_BROADCAST.

# OSQQuery()

Chapter	File	Called from	Code enabled by
11	os_Q.c	Task or ISR	OS_Q_EN && OS_QUERY_EN

OSQQuery() obtains information about a message queue. Your application must allocate an OS\_Q\_DATA data structure used to receive data from the event control block of the message queue. OSQQuery() allows you to determine whether any tasks are waiting for messages at the queue, how many tasks are waiting (by counting the number of 1s in the .OSEventTbl[] field), how many messages are in the queue, and what the message queue size is. OSQQuery() also obtains the next message that is returned if the queue is not empty. Note that the size of .OSEventTbl[] is established by the #define constant OS EVENT TBL SIZE (see uCOS II.H).

### **Arguments**

is a pointer to the message queue. This pointer is returned to your application when the queue is created [see OSQCreate()].

p\_q\_data is a pointer to a data structure of type OS\_Q\_DATA, which contains the following fields

### **Returned Value**

 ${\tt OSQQuery} \ \hbox{()} \ \ returns \ one \ of these \ error \ codes:$ 

```
OS_ERR_NONE if the call is successful.

OS_ERR_EVENT_TYPE if you don't pass a pointer to a message queue.

OS_ERR_PEVENT_NULL if pevent is a NULL pointer.

OS_ERR_PDATA_NULL if p_q_data is a NULL pointer.
```

### **Notes/Warnings**

1. Message queues must be created before they are used.

# OSSchedLock()

void OSSchedLock(void);

Chapter	File	Called from	Code enabled by
3	OS_CORE.C	Task or ISR	OS_SCHED_LOCK_EN

OSSchedLock() prevents task rescheduling until its counterpart, OSSchedUnlock(), is called. The task that calls OSSchedLock() keeps control of the CPU even though other higher priority tasks are ready to run. However, interrupts are still recognized and serviced (assuming interrupts are enabled). OSSchedLock() and OSSchedUnlock() must be used in pairs.  $\mu$ C/OS-II allows OSSchedLock() to be nested up to 255 levels deep. Scheduling is enabled when an equal number of OSSchedUnlock() calls have been made.

### Arguments

none

#### **Returned Value**

none

### **Notes/Warnings**

1. After calling OSSchedLock(), your application must not make system calls that suspend execution of the current task; that is, your application cannot call OSTimeDly(), OSTimeDlyHMSM(), OSFlagPend(), OSSemPend(), OSMutexPend(), OSMboxPend(), or OSQPend(). Because the scheduler is locked out, no other task is allowed to run, and your system will lock up.

# OSSchedUnlock()

void OSSchedUnlock(void);

Chapter	File	Called from	Code enabled by
3	OS_CORE.C	Task or ISR	OS_SCHED_LOCK_EN

OSSchedUnlock() re-enables task scheduling whenever it is paired with OSSchedLock().

### **Arguments**

none

### **Returned Value**

none

### **Notes/Warnings**

1. After calling OSSchedLock(), your application must not make system calls that suspend execution of the current task; that is, your application cannot call OSTimeDly(), OSTimeDlyHMSM(), OSFlagPend(), OSSemPend(), OSMutexPend(), OSMboxPend(), or OSQPend(). Because the scheduler is locked out, no other task is allowed to run, and your system will lock up.

# OSSemAccept()

INT16U OSSemAccept(OS EVENT \*pevent);

Chapter	File	Called from	Code enabled by
7	OS_SEM.C	Task or ISR	OS_SEM_EN && OS_SEM_ACCEPT_EN

OSSemAccept() checks to see if a resource is available or an event has occurred. Unlike OSSemPend(), OSSemAccept() does not suspend the calling task if the resource is not available. In other words, OSSemAccept() is non-blocking. Use OSSemAccept() from an ISR to obtain the semaphore.

### Arguments

pevent

is a pointer to the semaphore that guards the resource. This pointer is returned to your application when the semaphore is created [see OSSemCreate()].

### Returned Value

When OSSemAccept() is called and the semaphore value is greater than 0, the semaphore value is decremented, and the value of the semaphore before the decrement is returned to your application. If the semaphore value is 0 when OSSemAccept() is called, the resource is not available, and 0 is returned to your application.

### **Notes/Warnings**

1. Semaphores must be created before they are used.

# OSSemCreate()

OS EVENT \*OSSemCreate(INT16U value);

Chapter	File	Called from	Code enabled by
7	OS_SEM.C	Task or startup code	OS_SEM_EN

OSSemCreate() creates and initializes a semaphore. A semaphore

- allows a task to synchronize with either an ISR or a task (you initialize the semaphore to 0),
- gains exclusive access to a resource (you initialize the semaphore to a value greater than 0), and
- signals the occurrence of an event (you initialize the semaphore to 0).

### **Arguments**

value

is the initial value of the semaphore and can be between 0 and 65,535. A value of 0 indicates that a resource is not available or an event has not occurred.

### **Returned Value**

OSSemCreate() returns a pointer to the event control block allocated to the semaphore. If no event control block is available, OSSemCreate() returns a NULL pointer.

### **Notes/Warnings**

1. Semaphores must be created before they are used.

# OSSemDel()

Chapter	File	Called from	Code enabled by
7	OS_SEM.C	Task	OS_SEM_EN and OS_SEM_DEL_EN

OSSemDel() is used to delete a semaphore. This function is dangerous to use because multiple tasks could attempt to access a deleted semaphore. You should always use this function with great care. Generally speaking, before you delete a semaphore, you must first delete all the tasks that can access the semaphore.

### **Arguments**

pevent	is a pointer to the semaphore. This pointer is returned to your application when the semaphore is created [see OSSemCreate()].		
opt	specifies whether you want to delete the semaphore only if there are no pending tasks (OS_DEL_NO_PEND) or whether you always want to delete the semaphore regardless of whether tasks are pending or not (OS_DEL_ALWAYS). In this case, all pending task are readied.		
perr	is a pointer to a variable that i following:	s used to hold an error code. The error code can be one of the	
	OS_ERR_NONE	if the call is successful and the semaphore has been deleted.	
	OS_ERR_DEL_ISR	if you attempt to delete the semaphore from an ISR.	
	OS_ERR_INVALID_OPT	if you don't specify one of the two options mentioned in the opt argument.	
	OS_ERR_TASK_WAITING	if one or more tasks are waiting on the semaphore.	
	if pevent is not pointing to a semaphore.		
	OS_ERR_PEVENT_NULL	if pevent is a NULL pointer.	

### **Returned Value**

A NULL pointer if the semaphore is deleted or pevent if the semaphore is not deleted. In the latter case, you need to examine the error code to determine the reason.

### Notes/Warnings

- 1. You should use this call with care because other tasks might expect the presence of the semaphore.
- 2. Interrupts are disabled when pended tasks are readied, which means that interrupt latency depends on the number of tasks that are waiting on the semaphore.

## OSSemPend()

Chapter	File	Called from	Code enabled by
7	OS_SEM.C	Task only	OS_SEM_EN

OSSemPend() is used when a task wants exclusive access to a resource, needs to synchronize its activities with an ISR or a task, or is waiting until an event occurs. If a task calls OSSemPend() and the value of the semaphore is greater than 0, OSSemPend() decrements the semaphore and returns to its caller. However, if the value of the semaphore is 0, OSSemPend() places the calling task in the waiting list for the semaphore. The task waits until a task or an ISR signals the semaphore or the specified timeout expires. If the semaphore is signaled before the timeout expires,  $\mu$ C/OS-II resumes the highest priority task waiting for the semaphore. A pended task that has been suspended with OSTaskSuspend() can obtain the semaphore. However, the task remains suspended until it is resumed by calling OSTaskResume().

### Arguments

pevent is a pointer to the semaphore. This pointer is returned to your application when the semaphore is

created [see OSSemCreate()].

timeout allows the task to resume execution if a message is not received from the mailbox within the

specified number of clock ticks. A timeout value of 0 indicates that the task waits forever for the message. The timeout value is not synchronized with the clock tick. The timeout count begins

decrementing on the next clock tick, which could potentially occur immediately.

perr is a pointer to a variable used to hold an error code. OSSemPend() sets \*perr to one of the

following:

OS\_ERR\_NONE if the semaphore is available.

OS ERR TIMEOUT if the semaphore is not signaled within the specified timeout.

OS ERR EVENT TYPE if pevent is not pointing to a semaphore.

OS\_ERR\_PEND\_ISR if you called this function from an ISR and  $\mu C/OS$ -II has to

suspend it. You should not call OSSemPend() from an ISR.

 $\mu$ C/OS-II checks for this situation.

OS ERR PEND LOCKED if you called this function when the scheduler is locked.

OS\_ERR\_PEVENT\_NULL if pevent is a NULL pointer.

### **Returned Value**

none

### **Notes/Warnings**

1. Semaphores must be created before they are used.

# OSSemPendAbort()

New Function	File	Called from	Code enabled by
V2.84	OS_SEM.C	Task only	OS_SEM_EN &&
			OS_SEM_PEND_ABORT_EN

OSSemPendAbort () aborts & readies any tasks currently waiting on a semaphore. This function should be used to fault-abort the wait on the semaphore, rather than to normally signal the semaphore via OSSemPost ().

### **Arguments**

pevent is a pointer to the semaphore for which pend(s) need to be aborted. This pointer is returned to your application when the semaphore is created [see OSSemCreate()].

your application when the semaphore is created [see OSSemCreate ()

opt determines what type of abort is performed.

OS\_PEND\_OPT\_NONE Aborts the pend of only the highest priority task waiting on the

semaphore.

OS\_PEND\_OPT\_BROADCAST Aborts the pend of all the tasks waiting on the semaphore.

perr is a pointer to a variable that holds an error code. OSSemPendAbort () sets \*perr to one of the

following:

OS ERR NONE if no tasks were waiting on the semaphore. In this case, the

return value is also 0.

OS ERR PEND ABORT at least one task waiting on the semaphore was readied and

informed of the aborted wait. Check the return value for the

number of tasks whose wait on the semaphore was aborted.

OS\_ERR\_EVENT\_TYPE if pevent is not pointing to a semaphore.

OS ERR PEVENT NULL if pevent is a NULL pointer.

### **Returned Value**

OSSemPendAbort () returns the number of tasks made ready to run by this function. Zero indicates that no tasks were pending on the semaphore and thus this function had no effect.

### **Notes/Warnings**

1. Semaphores must be created before they are used.

```
OS_EVENT *CommSem;
void CommTask(void *p_arg)
     INT8U err;
     INT8U nbr_tasks;
      (void)p_arg;
      for (;;) {
        nbr_tasks = OSSemPendAbort(CommSem, OS_PEND_OPT_BROADCAST, &err);
        if (err == OS_ERR_NONE) {
            . /* No tasks were waiting on the semaphore
                                                                        */
        } else {
             . /* All pends of tasks waiting on semaphore were aborted \dots */
             . /* ... 'nbr_tasks' indicates how many were made ready. */
         }
```

## OSSemPost()

INT8U OSSemPost(OS EVENT \*pevent);

Chapter	File	Called from	Code enabled by
7	OS_SEM.C	Task or ISR	OS_SEM_EN

A semaphore is signaled by calling OSSemPost(). If the semaphore value is 0 or more, it is incremented, and OSSemPost() returns to its caller. If tasks are waiting for the semaphore to be signaled, OSSemPost() removes the highest priority task pending for the semaphore from the waiting list and makes this task ready to run. The scheduler is then called to determine if the awakened task is now the highest priority task ready to run.

### Arguments

pevent

is a pointer to the semaphore. This pointer is returned to your application when the semaphore is created [see OSSemCreate()].

### **Returned Value**

OSSemPost () returns one of these error codes:

OS ERR NONE if the semaphore is signaled successfully.

OS\_ERR\_SEM\_OVF if the semaphore count overflows.

OS ERR EVENT TYPE if pevent is not pointing to a semaphore.

OS\_ERR\_PEVENT\_NULL if pevent is a NULL pointer.

### **Notes/Warnings**

1. Semaphores must be created before they are used.

```
OS_EVENT *DispSem;
void TaskX (void *p_arg)
{
     INT8U err;
      (void)p_arg;
     for (;;) {
        err = OSSemPost(DispSem);
        switch (err) {
            case OS_ERR_NONE:
                 /* Semaphore signaled */
                break;
            case OS_ERR_SEM_OVF:
                /* Semaphore has overflowed */
                 break;
```

# OSSemQuery()

Chapter	File	Called from	Code enabled by
7	OS_SEM.C	Task or ISR	OS_SEM_EN && OS_SEM_QUERY_EN

OSSemQuery() obtains information about a semaphore. Your application must allocate an OS\_SEM\_DATA data structure used to receive data from the event control block of the semaphore. OSSemQuery() allows you to determine whether any tasks are waiting on the semaphore and how many tasks are waiting (by counting the number of 1s in the .OSEventTb1[] field) and obtains the semaphore count. Note that the size of .OSEventTb1[] is established by the #define constant OS EVENT TBL SIZE (see uCOS II.H).

### Arguments

pevent is a pointer to the semaphore. This pointer is returned to your application when the semaphore is created [see OSSemCreate()].

P\_sem\_data is a pointer to a data structure of type OS\_SEM\_DATA, which contains the following fields

### **Returned Value**

OSSemQuery() returns one of these error codes:

```
OS_ERR_NONE if the call is successful.

OS_ERR_EVENT_TYPE if you don't pass a pointer to a semaphore.

OS_ERR_PEVENT_NULL if pevent is is a NULL pointer.

OS_ERR_PDATA_NULL if p sem_data is is a NULL pointer.
```

### **Notes/Warnings**

1. Semaphores must be created before they are used.

In this example, the contents of the semaphore is checked to determine the highest priority task waiting at the time the function call was made.

```
OS_EVENT *DispSem;
void Task (void *p_arg)
      OS_SEM_DATA sem_data;
      INT8U
                  err;
                  highest; /* Highest priority task waiting on sem. */
      INT8U
      INT8U
      INT8U
                  у;
      (void)p_arg;
      for (;;) {
         err = OSSemQuery(DispSem, &sem_data);
         if (err == OS_ERR_NONE) {
            /* Examine sem_data */
```

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## OSSemSet()

Chapter	File	Called from	Code enabled by
7	OS_SEM.C	Task or ISR	OS_SEM_EN && OS_SEM_SET_EN

OSSemSet() is used to change the current value of the semaphore count. This function would normally be used when a semaphore is used as a signaling mechanism. OSSemSet() can then be used to reset the count to any value. If the semaphore count is already 0 then, the count is only changed if there are no tasks waiting on the semaphore.

### **Arguments**

pevent is a pointer to the semaphore that is used as a signaling mechanism. This pointer is returned to

your application when the semaphore is created [see OSSemCreate()].

is the desired count that you want the semaphore set to.

perr is a pointer to a variable used to hold an error code. OSSemSet() sets \*perr to one of the

following:

OS ERR NONE if the count was changed or, not changed because there was one

or more tasks waiting on the semaphore.

OS\_ERR\_EVENT\_TYPE if pevent is not pointing to a semaphore.

OS ERR PEVENT NULL if pevent is a NULL pointer.

OS ERR TASK WAITING if tasks are waiting on the semaphore.

### **Returned Value**

None

### **Notes/Warnings**

1. You should **NOT** use this function if the semaphore is used to protect a shared resource.

# OSStart()

void OSStart(void);

Chapter	File	Called from	Code enabled by
3	OS_CORE.C	Startup code only	N/A

 ${\tt OSStart()} \ \, \text{starts multitasking under } \mu C/OS\text{-}II. \ \, \text{This function is typically called from your startup code but after you call } {\tt OSInit()}.$ 

### **Arguments**

none

### **Returned Value**

none

### **Notes/Warnings**

1. OSInit() must be called prior to calling OSStart(). OSStart() should only be called once by your application code. If you do call OSStart() more than once, it does not do anything on the second and subsequent calls.

# OSStatInit()

void OSStatInit(void);

Chapter	File	Called from	Code enabled by
3	OS_CORE.C	Startup code only	OS_TASK_STAT_EN &&
			OS_TASK_CREATE_EXT_EN

OSStatInit() determines the maximum value that a 32-bit counter can reach when no other task is executing. This function must be called when only one task is created in your application and when multitasking has started; that is, this function must be called from the first and, only, task created.

### **Arguments**

none

### **Returned Value**

none

### **Notes/Warnings**

none

# OSTaskChangePrio()

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task only	OS_TASK_CHANGE_PRIO_EN

OSTaskChangePrio() changes the priority of a task.

### **Arguments**

oldprio is the priority number of the task to change.

newprio is the new task's priority.

### **Returned Value**

OSTaskChangePrio() returns one of the following error codes:

OS\_ERR\_NONE if the task's priority is changed.

OS\_ERR\_PRIO\_INVALID if either the old priority or the new priority is equal to or exceeds OS\_LOWEST\_PRIO.

OS\_ERR\_PRIO\_EXIST if newprio already exists.

OS\_ERR\_PRIO if no task with the specified old priority exists (i.e., the task specified by oldprio does not exist).

OS\_ERR\_TASK\_NOT\_EXITS if the task is assigned to a Mutex PIP.

### **Notes/Warnings**

1. The desired priority must not already have been assigned; otherwise, an error code is returned. Also, OSTaskChangePrio() verifies that the task to change exists.

# OSTaskCreate()

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task or startup code	OS_TASK_CREATE_EN

OSTaskCreate () creates a task so it can be managed by  $\mu$ C/OS-II. Tasks can be created either prior to the start of multitasking or by a running task. A task cannot be created by an ISR. A task must be written as an infinite loop, as shown below, and must not return.

 ${\tt OSTaskCreate()} \ \ is \ used \ \ for \ backward \ \ compatibility \ with \ } \mu C/OS \ \ and \ \ when \ \ the \ \ added \ \ features \ \ of \ \ OSTaskCreateExt() \ \ are \ not \ needed.$ 

Depending on how the stack frame is built, your task has interrupts either enabled or disabled. You need to check with the processor-specific code for details.

```
void Task (void *p arg)
                          /* Do something with 'pdata'
                          /* Task body, always an infinite loop.
                                                                         */
      for (;;) {
                                                                         */
         /* Must call one of the following services:
         /*
                                                                         */
               OSMboxPend()
               OSFlagPend()
                                                                         */
         /*
               OSMutexPend()
                                                                         */
         /*
                                                                         */
               OSQPend()
               OSSemPend()
                                                                         */
         /*
               OSTimeDly()
                                                                         */
                                                                         */
               OSTimeDlyHMSM()
               OSTaskSuspend()
                                    (Suspend self)
                                                                         */
         /*
                                                                         */
               OSTaskDel()
                                    (Delete self)
```

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### Arguments

is a pointer to the task's code. task

is a pointer to an optional data area used to pass parameters to the task when it is created. Where pdata

the task is concerned, it thinks it is invoked and passes the argument pdata. pdata can be used to pass arguments to the task created. For example, you can create a generic task that handles an asynchronous serial port. pdata can be used to pass this task information about the serial

port it has to manage: the port address, the baud rate, the number of bits, the parity, and more.

is a pointer to the task's top-of-stack. The stack is used to store local variables, function ptos parameters, return addresses, and CPU registers during an interrupt. The size of the stack is determined by the task's requirements and the anticipated interrupt nesting. Determining the size

of the stack involves knowing how many bytes are required for storage of local variables for the task itself and all nested functions, as well as requirements for interrupts (accounting for nesting). If the configuration constant OS STK GROWTH is set to 1, the stack is assumed to grow downward (i.e., from high to low memory). ptos thus needs to point to the highest valid

memory location on the stack. If OS STK GROWTH is set to 0, the stack is assumed to grow in

the opposite direction (i.e., from low to high memory).

is the task priority. A unique priority number must be assigned to each task, and the lower the prio

number, the higher the priority (i.e., the task importance).

### **Returned Value**

OSTaskCreate() returns one of the following error codes:

OS ERR NONE if the function is successful.

OS\_ERR PRIO EXIST if the requested priority already exists.

OS ERR PRIO INVALID if prio is higher than OS LOWEST PRIO.

if  $\mu C/OS$ -II doesn't have any more OS\_TCBs to assign. OS ERR NO MORE TCB

OS ERR TASK CREATE ISR if you attempted to create the task from an ISR.

### Notes/Warnings

- The stack for the task must be declared with the OS STK type.
- A task must always invoke one of the services provided by μC/OS-II to wait for time to expire, suspend the task, or wait for an event to occur (wait on a mailbox, queue, or semaphore). This allows other tasks to gain control of the CPU.
- 3. You should not use task priorities 0, 1, 2, 3, OS LOWEST PRIO-3, OS LOWEST PRIO-2, OS LOWEST PRIO-1, and OS LOWEST PRIO because they are reserved for use by  $\mu$ C/OS-II.

This example shows that the argument that Taskl() receives is not used, so the pointer pdata is set to NULL. Note that I assume the stack grows from high to low memory because I pass the address of the highest valid memory location of the stack Tasklstk[]. If the stack grows in the opposite direction for the processor you are using, pass Tasklstk[] as the task's top-of-stack.

Assigning pdata to itself is used to prevent compilers from issuing a warning about the fact that pdata is not being used. In other words, if I had not added this line, some compilers would have complained about 'WARNING - variable pdata not used.'

```
OS_STK Task1Stk[1024];
void main (void)
      INT8U err;
      OSInit();
                              /* Initialize μC/OS-II
      OSTaskCreate (Task1,
                    (void *)0,
                   &Task1Stk[1023],
                   25);
                               /* Start Multitasking
      OSStart();
                                                                    */
}
void Task1 (void *p arg)
      (void)p_arg;
                                /* Prevent compiler warning
                                                                   */
      for (;;) {
                                 /* Task code
      }
```

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You can create a generic task that can be instantiated more than once. For example, a task that handles a serial port could be passed the address of a data structure that characterizes the specific port (i.e., port address and baud rate). Note that each task has its own stack space and its own (different) priority. In this example, I arbitrarily decided that COM1 is the most important port of the two.

```
*Comm1Stk[1024];
OS STK
                               /* Data structure containing COMM port
COMM DATA Comm1Data;
                               /* Specific data for channel 1
                                                                        */
OS STK
          *Comm2Stk[1024];
COMM DATA Comm2Data;
                              /* Data structure containing COMM port */
                               /* Specific data for channel 2
                                                                        */
void main (void)
  INT8U err;
  OSInit();
                              /* Initialize μC/OS-II
                                                                        */
                               /* Create task to manage COM1
  OSTaskCreate(CommTask,
                (void *) &CommlData,
                &Comm1Stk[1023],
                25);
                               /* Create task to manage COM2
  OSTaskCreate(CommTask,
                (void *) &Comm2Data,
                &Comm2Stk[1023],
                26);
  OSStart();
                              /* Start Multitasking
void CommTask (void *p_arg) /* Generic communication task
      for (;;) {
                               /* Task code
```

## OSTaskCreateExt()

```
(*task) (void *pd),
INT8U OSTaskCreateExt(void
                       void
                               *pdata,
                       OS STK *ptos,
                       INT8U
                                prio,
                       INT16U
                                id,
                       OS STK *pbos,
                       INT32U
                               stk size,
                       void
                               *pext,
                       INT16U
                                opt);
```

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task or startup code	OS_TASK_CREATE_EXT_EN

OSTaskCreate(), except that it allows you to specify additional information about your task to  $\mu$ C/OS-II. Tasks can be created either prior to the start of multitasking or by a running task. A task cannot be created by an ISR. A task must be written as an infinite loop, as shown below, and must not return. Depending on how the stack frame is built, your task has interrupts either enabled or disabled. You need to check with the processor-specific code for details. Note that the first four arguments are exactly the same as the ones for OSTaskCreate(). This was done to simplify the migration to this new and more powerful function. It is highly recommended that you use OSTaskCreateExt() instead of the older OSTaskCreate() function because it's much more flexible.

```
void Task (void *p arg)
                           /* Do something with 'pdata'
                                                                             */
      for (;;) {
                            /* Task body, always an infinite loop.
                                                                             */
         /* Must call one of the following services:
                                                                             */
          /*
                OSMboxPend()
                                                                             */
          /*
                OSFlagPend()
                                                                             */
                OSMutexPend()
                                                                             */
          /*
                OSQPend()
                                                                             */
          /*
                OSSemPend()
                                                                             */
          /*
                OSTimeDly()
                                                                             * /
          /*
                OSTimeDlyHMSM()
                                                                             */
          /*
                OSTaskSuspend()
                                     (Suspend self)
                                                                             */
                OSTaskDel()
                                     (Delete self)
                                                                             */
```

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#### **Arguments**

task is a pointer to the task's code.

is a pointer to an optional data area, which is used to pass parameters to the task when it is pdata

> created. Where the task is concerned, it thinks it is invoked and passes the argument pdata. pdata can be used to pass arguments to the task created. For example, you can create a generic task that handles an asynchronous serial port. pdata can be used to pass this task information about the serial port it has to manage: the port address, the baud rate, the number of bits, the

parity, and more.

is a pointer to the task's top-of-stack. The stack is used to store local variables, function parameters, return addresses, and CPU registers during an interrupt.

> The size of this stack is determined by the task's requirements and the anticipated interrupt nesting. Determining the size of the stack involves knowing how many bytes are required for storage of local variables for the task itself and all nested functions, as well as requirements for interrupts (accounting for nesting).

> If the configuration constant OS STK GROWTH is set to 1, the stack is assumed to grow downward (i.e., from high to low memory). ptos thus needs to point to the highest valid memory location on the stack. If OS STK GROWTH is set to 0, the stack is assumed to grow in the opposite direction (i.e., from low to high memory).

> is the task priority. A unique priority number must be assigned to each task: the lower the number, the higher the priority (i.e., the importance) of the task.

> is the task's ID number. At this time, the ID is not currently used in any other function and has simply been added in OSTaskCreateExt() for future expansion. You should set id to the same value as the task's priority.

> is a pointer to the task's bottom-of-stack. If the configuration constant OS STK GROWTH is set to 1, the stack is assumed to grow downward (i.e., from high to low memory); thus, pbos must point to the lowest valid stack location. If OS STK GROWTH is set to 0, the stack is assumed to grow in the opposite direction (i.e., from low to high memory); thus, pbos must point to the highest valid stack location. pbos is used by the stack-checking function OSTaskStkChk().

> specifies the size of the task's stack in number of elements. If OS STK is set to INT8U, then stk size corresponds to the number of bytes available on the stack. If OS STK is set to INT16U, then stk size contains the number of 16-bit entries available on the stack. Finally, if OS STK is set to INT32U, then stk size contains the number of 32-bit entries available on the stack.

> is a pointer to a user-supplied memory location (typically a data structure) used as a TCB extension. For example, this user memory can hold the contents of floating-point registers during a context switch, the time each task takes to execute, the number of times the task is switched in, and so on.

> contains task-specific options. The lower 8 bits are reserved by  $\mu$ C/OS-II, but you can use the upper 8 bits for application-specific options. Each option consists of one or more bits. The option is selected when the bit(s) is set. The current version of μC/OS-II supports the following options:

> specifies that there are no options. OS TASK OPT NONE specifies whether stack checking is allowed for the task. OS TASK OPT STK CHK specifies whether the stack needs to be cleared. OS TASK OPT STK CLR specifies whether floating-point registers are saved. This option OS TASK OPT SAVE FP is only valid if your processor has floating-point hardware and the processor-specific code saves the floating-point registers.

Refer to ucos II. H for other options.

ptos

prio

id

pbos

stk size

pext

opt

#### **Returned Value**

OSTaskCreateExt() returns one of the following error codes:

### **Notes/Warnings**

- 1. The stack must be declared with the OS STK type.
- A task must always invoke one of the services provided by μC/OS-II to wait for time to expire, suspend the
  task, or wait an event to occur (wait on a mailbox, queue, or semaphore). This allows other tasks to gain
  control of the CPU.
- 3. You should not use task priorities 0, 1, 2, 3, OS\_LOWEST\_PRIO-3, OS\_LOWEST\_PRIO-2, OS\_LOWEST\_PRIO-1, and OS\_LOWEST\_PRIO because they are reserved for use by  $\mu C/OS$ -II.

- E1(1) The task control block is extended using a user-defined data structure called OS\_TASK\_USER\_DATA, which in this case contains the name of the task as well as other fields.
- E1(2) The task name is initialized with the standard library function strcpy().
- E1(4) Note that stack checking has been enabled for this task, so you are allowed to call OSTaskStkChk().
- E1(3) Also, assume here that the stack grows downward on the processor used (i.e., OS\_STK\_GROWTH is set to 1; TOS stands for top-of-stack and BOS stands for bottom-of-stack).

```
typedef struct {
                            /* User defined data structure */
                                                              (1)
   char OSTaskName[20];
  INT16U OSTaskCtr;
  INT16U OSTaskExecTime;
   INT32U OSTaskTotExecTime;
} OS_TASK_USER_DATA;
OS STK TaskStk[1024];
TASK_USER_DATA TaskUserData;
void main (void)
    INT8U err;
    OSInit();
                                   /* Initialize μC/OS-II */
    err = OSTaskCreateExt(Task,
           (void *)0,
            &TaskStk[1023], /* Stack grows down (TOS) */ (3)
            10,
           10,
                                   /* Stack grows down (BOS) */ (3)
           &TaskStk[0],
           1024,
           (void *)&TaskUserData,
                                   /* TCB Extension */
                                   /* Stack checking enabled */ (4)
           OS_TASK_OPT_STK_CHK);
    OSStart();
                                   /* Start Multitasking */
}
void Task(void *p arg)
                                   /* Avoid compiler warning */
    (void)p_arg;
    for (;;) {
                                    /* Task code
                                                          */
```

We now create a task, but this time on a processor for which the stack grows upward. The Intel MCS-51 is an example of such a processor. In this case, OS\_STK\_GROWTH is set to 0.

E2(2) Note that stack checking has been enabled for this task so you are allowed to call OSTaskStkChk() (TOS stands for top-of-stack and BOS stands for bottom-of-stack).

```
OS_STK *TaskStk[1024];
void main (void)
    INT8U err;
    OSInit();
                                 /* Initialize μC/OS-II */
    err = OSTaskCreateExt(Task,
           (void *)0,
            &TaskStk[0], /* Stack grows up (TOS) */ (1)
            10,
            10,
            &TaskStk[1023], /* Stack grows up (BOS) */ (1)
           1024,
           (void *)0,
            OS_TASK_OPT_STK_CHK); /* Stack checking enabled */ (2)
                                    /* Start Multitasking */
    OSStart();
void Task (void *p_arg)
                                   /* Avoid compiler warning */
    (void)p_arg;
    for (;;) {
                                                           */
                                     /* Task code
    }
```

# OSTaskDel()

INT8U OSTaskDel(INT8U prio);

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task only	OS_TASK_DEL_EN

OSTaskDel() deletes a task by specifying the priority number of the task to delete. The calling task can be deleted by specifying its own priority number or OS\_PRIO\_SELF (if the task doesn't know its own priority number). The deleted task is returned to the dormant state. The deleted task can be re-created by calling either OSTaskCreate() or OSTaskCreateExt() to make the task active again.

## Arguments

prio

is the priority number of the task to delete. You can delete the calling task by passing OS\_PRIO\_SELF, in which case the next highest priority task is executed.

#### **Returned Value**

OSTaskDel() returns one of the following error codes:

OS_ERR_NONE	if the task doesn't delete itself.
OS_ERR_TASK_IDLE	if you try to delete the idle task, which is of course is not allowed.
OS_ERR_TASK_DEL	if the task to delete does not exist.
OS_ERR_PRIO_INVALID	if you specify a task priority higher than OS_LOWEST_PRIO.
OS_ERR_TASK_DEL_ISR	if you try to delete a task from an ISR.
OS_ERR_TASK_DEL	if the task is assigned to a Mutex.
OS_ERR_TASK_NOT_EXIST	if the task is assigned to a Mutex PIP.

## **Notes/Warnings**

- 1. OSTaskDel() verifies that you are not attempting to delete the  $\mu C/OS$ -II idle task.
- 2. You must be careful when you delete a task that owns resources. Instead, consider using OSTaskDelReq() as a safer approach.

# OSTaskDelReq()

INT8U OSTaskDelReq(INT8U prio);

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task only	OS_TASK_DEL_EN

OSTaskDelReq() requests that a task delete itself. Basically, use OSTaskDelReq() when you need to delete a task that can potentially own resources (e.g., the task might own a semaphore). In this case, you don't want to delete the task until the resource is released. The requesting task calls OSTaskDelReq() to indicate that the task needs to be deleted. Deletion of the task is, however, deferred to the task being deleted. In other words, the task is actually deleted when it regains control of the CPU. For example, suppose Task 10 needs to be deleted. The task wanting to delete this task (example Task 5) calls OSTaskDelReq(10). When Task 10 executes, it calls OSTaskDelReq(OS\_PRIO\_SELF) and monitors the return value. If the return value is OS\_ERR\_TASK\_DEL\_REQ, then Task 10 is asked to delete itself. At this point, Task 10 calls OSTaskDel(OS\_PRIO\_SELF). Task 5 knows whether Task 10 has been deleted by calling OSTaskDelReq(10) and checking the return code. If the return code is OS\_ERR\_TASK\_NOT\_EXIST, then Task 5 knows that Task 10 has been deleted. Task 5 might have to check periodically until OS\_ERR\_TASK\_NOT\_EXIST is returned.

## **Arguments**

prio

is the task's priority number of the task to delete. If you specify OS\_PRIO\_SELF, you are asking whether another task wants the current task to be deleted.

#### **Returned Value**

OSTaskDelReq() returns one of the following error codes:

OS_ERR_NONE	if the task deletion has been registered.	
OS_ERR_TASK_NOT_EXIST	if the task does not exist. The requesting task can monitor this return code to see if the task is actually deleted.	
OS_ERR_TASK_IDLE	if you ask to delete the idle task (which is obviously not allowed).	
OS_ERR_PRIO_INVALID	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	
OS_ERR_TASK_DEL	if the task is assigned to a Mutex.	
OS_ERR_TASK_DEL_REQ	if a task (possibly another task) requests that the running task be deleted.	

### Notes/Warnings

1. OSTaskDelReq() verifies that you are not attempting to delete the  $\mu$ C/OS-II idle task.

```
void TaskThatDeletes (void *p_arg) /* My priority is 5
                                                                   */
    INT8U err;
     for (;;) {
       err = OSTaskDelReq(10); /* Request task #10 to delete itself */
       if (err == OS_ERR_NONE) {
         while (err != OS_ERR_TASK_NOT_EXIST) {
            err = OSTaskDelReq(10);
            OSTimeDly(1); /* Wait for task to be deleted */
          }
                                /* Task #10 has been deleted */
        }
    }
void TaskToBeDeleted (void *p_arg) /* My priority is 10
                                                                  */
 (void)p_arg;
  for (;;) {
    OSTimeDly(1);
    if (OSTaskDelReq(OS_PRIO_SELF) == OS_ERR_TASK_DEL_REQ) {
       /* Release any owned resources;
                                                                    */
       /* De-allocate any dynamic memory;
                                                                    */
       OSTaskDel(OS_PRIO_SELF);
```

## OSTaskNameGet()

Chapter	File	Called from	Code enabled by
New in V2.60	OS_TASK.C	Task	OS_TASK_NAME_EN

OSTaskNameGet () allows you to obtain the name that you assigned to a task. This function is typically used by a debugger to allow associating a name to a task.

#### **Arguments**

prio is the priority of the task from which you would like to obtain the name from. If you specify OS PRIO SELF, you would obtain the name of the current task. is a pointer to a pointer to an ASCII string that point to the name of the task. pname a pointer to an error code and can be any of the following: perr If the call is successfull. OS ERR NONE The task you specified was not created or has been deleted. OS ERR TASK NOT EXIST OS\_ERR\_PRIO INVALID If you specified an invalid priority - a priority higher than the idle task (OS LOWEST PRIO) or you didn't specify OS PRIO SELF. If you passed a NULL pointer for pname. OS ERR PNAME NULL OS ERR NAME GET ISR You called this function from an ISR.

#### **Returned Values**

The size of the ASCII string pointed to by pname or 0 if an error is encountered.

### **Notes/Warnings**

1. The task must be created before you can use this function and obtain the name of the task.

```
INT8U *EngineTaskName;

void Task (void *p_arg)
{
    INT8U     err;
    INT8U     size;

    (void)p_arg;
    for (;;) {
        size = OSTaskNameGet(OS_PRIO_SELF, &EngineTaskName, &err);
        .
        .
        .
     }
}
```

# OSTaskNameSet()

Chapter	File	Called from	Code enabled by
New in V2.60	OS_TASK.C	Task	OS_TASK_NAME_EN

OSTaskNameSet() allows you to assign a name to a task. This function is typically used by a debugger to allow associating a name to a task.

### Arguments

prio is the priority of the task that you want to name. If you specify OS\_PRIO\_SELF, you would set

the name of the current task.

pname is a pointer to an ASCII string that contains the name of the task.

perr a pointer to an error code and can be any of the following:

OS\_ERR\_NONE If the call was successfull.

OS ERR TASK NOT EXIST The task you specified was not created or has been

deleted.

OS ERR PRIO INVALID If you specified an invalid priority - a priority higher than

the idle task (OS\_LOWEST\_PRIO) or you didn't specify

OS\_PRIO\_SELF.

OS ERR PNAME NULL If you passed a NULL pointer for pname.

OS\_ERR\_NAME\_SET\_ISR You called this function from an ISR.

#### **Returned Values**

None.

#### Notes/Warnings

1. The task must be created before you can use this function to set the name of the task.

```
void Task (void *p_arg)
{
    INT8U err;

    (void)p_arg;
    for (;;) {
        OSTaskNameSet(OS_PRIO_SELF, "Engine Task", &err);
        .
        .
        .
     }
}
```

# OSTaskResume()

INT8U OSTaskResume(INT8U prio);

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task only	OS_TASK_SUSPEND_EN

 ${\tt OSTaskResume()} \quad resumes \quad a \quad task \quad suspended \quad through \quad the \quad {\tt OSTaskSuspend()} \quad function. \quad In \quad fact, \\ {\tt OSTaskResume()} \quad is \quad the \quad only \quad function \quad that \quad can \quad unsuspend \quad a \quad suspended \quad task.$ 

#### Arguments

prio specifies the priority of the task to resume.

#### **Returned Value**

OSTaskResume () returns one of the these error codes:

```
OS_ERR_NONE if the call is successful.

OS_ERR_TASK_RESUME_PRIO if the task you are attempting to resume does not exist.

OS_ERR_TASK_NOT_SUSPENDED if the task to resume has not been suspended.

OS_ERR_PRIO_INVALID if prio is higher or equal to OS_LOWEST_PRIO.

OS_ERR_TASK_NOT_EXIST if the task is assigned to a Mutex PIP.
```

### **Notes/Warnings**

none

# OSTaskStkChk()

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task code	OS_TASK_CREATE_EXT

OSTaskStkChk() determines a task's stack statistics. Specifically, it computes the amount of free stack space, as well as the amount of stack space used by the specified task. This function requires that the task be created with OSTaskCreateExt() and that you specify OS TASK OPT STK CHK in the opt argument.

Stack sizing is done by walking from the bottom of the stack and counting the number of 0 entries on the stack until a nonzero value is found. Of course, this assumes that the stack is cleared when the task is created. For that purpose, you need to set OS\_TASK\_OPT\_STK\_CLR to 1 as an option when you create the task. You could set OS\_TASK\_OPT\_STK\_CLR to 0 if your startup code clears all RAM and you never delete your tasks. This reduces the execution time of OSTaskCreateExt().

#### **Arguments**

prio

is the priority of the task about which you want to obtain stack information. You can check the stack of the calling task by passing OS PRIO SELF.

P stk data is a pointer to a variable of type OS STK DATA, which contains the following fields:

```
INT32U OSFree; /* Number of bytes free on the stack */
INT32U OSUsed; /* Number of bytes used on the stack */
```

#### **Returned Value**

OSTaskStkChk() returns one of the these error codes:

OS_ERR_NONE	if you specify valid arguments and the call is successful.
OS_ERR_PRIO_INVALID	if you specify a task priority higher than OS_LOWEST_PRIO or you don't specify OS_PRIO_SELF.
OS_ERR_TASK_NOT_EXIST	if the specified task does not exist.
OS_ERR_TASK_OPT_ERR	if you do not specify $OS_TASK_OPT_STK_CHK$ when the task was created by $OSTaskCreateExt()$ or if you create the task by using $OSTaskCreate()$ .
OS_ERR_PDATA_NULL	if p_stk_data is a NULL pointer.

#### **Notes/Warnings**

- 1. Execution time of this function depends on the size of the task's stack and is thus nondeterministic.
- 2. Your application can determine the total task stack space (in number of bytes) by adding the two fields .OSFree and .OSUsed of the OS STK DATA data structure.
- 3. Technically, this function can be called by an ISR, but because of the possibly long execution time, it is not advisable.

# OSTaskSuspend()

INT8U OSTaskSuspend(INT8U prio);

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task only	OS_TASK_SUSPEND_EN

OSTaskSuspend() suspends (or blocks) execution of a task unconditionally. The calling task can be suspended by specifying its own priority number or OS\_PRIO\_SELF if the task doesn't know its own priority number. In this case, another task needs to resume the suspended task. If the current task is suspended, rescheduling occurs, and  $\mu$ C/OS-II runs the next highest priority task ready to run. The only way to resume a suspended task is to call OSTaskResume().

Task suspension is additive, which means that if the task being suspended is delayed until n ticks expire, the task is resumed only when both the time expires and the suspension is removed. Also, if the suspended task is waiting for a semaphore and the semaphore is signaled, the task is removed from the semaphore-wait list (if it is the highest priority task waiting for the semaphore), but execution is not resumed until the suspension is removed.

#### Arguments

prio

specifies the priority of the task to suspend. You can suspend the calling task by passing OS\_PRIO\_SELF, in which case, the next highest priority task is executed.

#### **Returned Value**

OSTaskSuspend() returns one of the these error codes:

OS_ERR_NONE	if the call is successful.
OS_ERR_TASK_SUSPEND_IDLE	if you attempt to suspend the _C/OS-II idle task, which is not allowed.
OS_ERR_PRIO_INVALID	if you specify a priority higher than the maximum allowed (i.e., you specify a priority of OS_LOWEST_PRIO or more) or you don't specify OS_PRIO_SELF.
OS_ERR_TASK_SUSPEND_PRIO	if the task you are attempting to suspend does not exist.
OS ERR TASK NOT EXITS	if the task is assigned to a Mutex PIP.

#### **Notes/Warnings**

- 1. OSTaskSuspend() and OSTaskResume() must be used in pairs.
- 2. A suspended task can only be resumed by OSTaskResume().

# OSTaskQuery()

Chapter	File	Called from	Code enabled by
4	OS_TASK.C	Task or ISR	OS_TASK_QUERY_EN

OSTaskQuery() obtains information about a task. Your application must allocate an OS\_TCB data structure to receive a snapshot of the desired task's control block. Your copy contains *every* field in the OS\_TCB structure. You should be careful when accessing the contents of the OS\_TCB structure, especially OSTCBNext and OSTCBPrev, because they point to the next and previous OS\_TCBs in the chain of created tasks, respectively. You could use this function to provide a debugger kernel awareness.

#### **Arguments**

prio is the priority of the task from which you wish to obtain data. You can obtain information about the calling task by specifying OS PRIO SELF.

p\_task\_data is a pointer to a structure of type OS\_TCB, which contains a copy of the task's control block.

#### **Returned Value**

OSTaskQuery() returns one of these error codes:

OS\_ERR\_NONE if the call is successful.

OS\_ERR\_PRIO\_INVALID if you specify a priority higher than OS\_LOWEST\_PRIO.

OS\_ERR\_PRIO if you try to obtain information from an invalid task.

OS\_ERR\_TASK\_NOT\_EXIST if the task is assigned to a Mutex PIP.

OS\_ERR\_PDATA\_NULL if p\_task\_data is a NULL pointer.

#### **Notes/Warnings**

- 1. The fields in the task control block depend on the following configuration options (see OS CFG.H):
  - OS\_TASK\_CREATE\_EN
  - OS\_Q\_EN
  - OS FLAG EN
  - OS MBOX EN
  - OS SEM EN
  - OS\_TASK\_DEL\_EN

# OSTimeDly()

void OSTimeDly(INT32U ticks);

Chapter	File	Called from	Code enabled by
5	OS_TIME.C	Task only	N/A

OSTimeDly() allows a task to delay itself for an integral number of clock ticks. Rescheduling always occurs when the number of clock ticks is greater than zero. Valid delays range from one to  $2^{32}$ -1 ticks. A delay of 0 means that the task is not delayed, and OSTimeDly() returns immediately to the caller. The actual delay time depends on the tick rate (see OS TICKS PER SEC in the configuration file OS CFG.H).

### **Arguments**

ticks

is the number of clock ticks to delay the current task.

#### **Returned Value**

none

### **Notes/Warnings**

- 1. Note that calling this function with a value of 0 results in no delay, and the function returns immediately to the caller.
- 2. To ensure that a task delays for the specified number of ticks, you should consider using a delay value that is one tick higher. For example, to delay a task for at least 10 ticks, you should specify a value of 11.

## **Example**

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# OSTimeDlyHMSM()

Chapter	File	Called from	Code enabled by
5	OS_TIME.C	Task only	N/A

OSTIMEDlyHMSM() allows a task to delay itself for a user-specified amount of time specified in hours, minutes, seconds, and milliseconds. This format is more convenient and natural than ticks. Rescheduling always occurs when at least one of the parameters is nonzero.

### Arguments

ms

hours is the number of hours the task is delayed. The valid range of values is 0 to 255.

minutes is the number of minutes the task is delayed. The valid range of values is 0 to 59.

seconds is the number of seconds the task is delayed. The valid range of values is 0 to 59.

is the number of milliseconds the task is delayed. The valid range of values is 0 to 999. Note that the resolution of this argument is in multiples of the tick rate. For instance, if the tick rate is set to 100Hz, a delay of 4ms results in no delay. The delay is rounded to the nearest tick. Thus, a delay of 15ms actually results in a delay of 20ms.

#### **Returned Value**

OSTimeDlyHMSM() returns one of the these error codes:

OS\_ERR\_NONE if you specify valid arguments and the call is successful.

OS\_ERR\_TIME\_INVALID\_MINUTES if the minutes argument is greater than 59.

OS\_ERR\_TIME\_INVALID\_MS if the seconds argument is greater than 59.

OS\_ERR\_TIME\_INVALID\_MS if the milliseconds argument is greater than 999.

OS\_ERR\_TIME\_ZERO\_DLY if all four arguments are 0.

OS\_ERR\_TIME\_DLY\_ISR if you called this function from an ISR.

#### **Notes/Warnings**

1. Note that OSTimeDlyHMSM(0,0,0,0) (i.e., hours, minutes, seconds, milliseconds) results in no delay, and the function returns to the caller.

# OSTimeDlyResume()

INT8U OSTimeDlyResume(INT8U prio);

Chapter	File	Called from	Code enabled by
5	OS_TIME.C	Task only	N/A

 ${\tt OSTimeDlyResume()} \ \ \textbf{resumes} \ \ \textbf{a} \ \ \textbf{task} \ \ \textbf{that} \ \ \textbf{has} \ \ \textbf{been} \ \ \textbf{delayed} \ \ \textbf{through} \ \ \textbf{a} \ \ \textbf{call} \ \ \textbf{to} \ \ \textbf{either} \ \ \texttt{OSTimeDlyHMSM()}.$ 

### Arguments

prio specifies the priority of the task to resume.

## **Returned Value**

OSTimeDlyResume() returns one of the these error codes:

```
OS_ERR_NONE if the call is successful.

OS_ERR_PRIO_INVALID if you specify a task priority greater than OS_LOWEST_PRIO.

OS_ERR_TIME_NOT_DLY if the task is not waiting for time to expire.

OS_ERR_TASK_NOT_EXIST if the task has not been created or has been assigned to a Mutex PIP.
```

### **Notes/Warnings**

1. Note that you must not call this function to resume a task that is waiting for an event with timeout. This situation makes the task look like a timeout occurred (unless you desire this effect).

# OSTimeGet()

INT32U OSTimeGet(void);

Chapter	File	Called from	Code enabled by
5	OS_TIME.C	Task or ISR	N/A

OSTimeGet() obtains the current value of the system clock. The system clock is a 32-bit counter that counts the number of clock ticks since power was applied or from a value set by OSTimeSet().

## **Arguments**

none

#### **Returned Value**

The current system clock value (in number of ticks).

## **Notes/Warnings**

none

## **Example**

539

# OSTimeSet()

void OSTimeSet(INT32U ticks);

Chapter	File	Called from	Code enabled by
5	OS_TIME.C	Task or ISR	N/A

OSTimeSet() sets the system clock. The system clock is a 32-bit counter that counts the number of clock ticks since power was applied or since the system clock was last set.

## **Arguments**

ticks is the desired value for the system clock, in ticks.

#### **Returned Value**

none

## **Notes/Warnings**

none

# OSTimeTick()

void OSTimeTick(void);

Chapter	File	Called from	Code enabled by
5	OS_TIME.C	Task or ISR	N/A

OSTimeTick() processes a clock tick.  $\mu$ C/OS-II checks all tasks to see if they are either waiting for time to expire [because they called OSTimeDly() or OSTimeDlyHMSM()] or waiting for events to occur until they timeout.

#### **Arguments**

none

#### **Returned Value**

none

## **Notes/Warnings**

1. The execution time of OSTimeTick() is directly proportional to the number of tasks created in an application. OSTimeTick() can be called by either an ISR or a task. If called by a task, the task priority should be very high (i.e., have a low priority number) because this function is responsible for updating delays and timeouts.

#### (Intel 80x86, real mode, large model)

```
OSTickISR PROC FAR
            PUSHA
                                            ; Save processor context
            PUSH ES
            PUSH DS
            MOV
                  AX, SEG(_OSIntNesting) ; Reload DS
                  DS, AX
            MOV
                  BYTE PTR DS: OSIntNesting ; Notify \mu C/OS-II of ISR
             INC
             CMP
                  BYTE PTR DS:_OSIntNesting, 1 ; if (OSIntNesting == 1)
            JNE SHORT _OSTickISR1
            MOV AX, SEG(_OSTCBCur)
                                     ; Reload DS
            MOV
                  DS, AX
            LES BX, DWORD PTR DS:_OSTCBCur ; OSTCBCur->OSTCBStkPtr = SS:SP
            MOV ES:[BX+2], SS
             MOV ES:[BX+0], SP
             CALL FAR PTR _OSTimeTick ; Process clock tick
                                            ; User Code to clear interrupt
            CALL FAR PTR _OSIntExit
                                           ; Notify \mu\text{C/OS-II} of end of ISR
                                            ; Restore processor registers
             POP DS
             POP ES
             POPA
            IRET
                                            ; Return to interrupted task
OSTickISR
            ENDP
```

# OSTmrCreate()

Chapter	File	Called from	Code enabled by
New in V2.83	OS_TMR.C	Task	OS_TMR_EN

OSTMTCreate() allows you to create a timer. The timer can be configured to run continuously (opt set to OS\_TMR\_OPT\_PERIODIC) or only once (opt set to OS\_TMR\_OPT\_ONE\_SHOT). When the timer counts down to 0 (from the value specified in period), an optional 'callback' function can be executed. The callback can be used to signal a task that the timer expired or, perform any other function. However, it's recommended that you keep the callback function as short as possible.

You MUST call OSTmrStart() to actually start the timer. If you configured the timer for one shot mode and the timer expired, you need to call OSTmrStart() to retrigger the timer or OSTmrDel() to delete the timer if you don't plan on retriggering it and or not use the timer anymore. Note that you can use the callback function to delete the timer if you use the one shot mode.

#### **Arguments**

dly specifies an initial delay used by the timer (see drawing below).

In ONE-SHOT mode, this is the time of the one-shot.

If in PERIODIC mode, this is the initial delay before the timer enters periodic mode.

The units of this time depends on how often you call <code>OSTmrSignal()</code>. In other words, if <code>OSTmrSignal()</code> is called every 1/10 of a second (i.e. <code>OS\_TMR\_CFG\_TICKS\_PER\_SEC</code> set to 10) then, <code>dly</code> specifies the number of 1/10 of a second before the delay expires. Note that the timer is **NOT** started when it is created.

period

specifies the amount of time it will take before the timer expires. You should set the 'period' to 0 when you use one-shot mode. The units of this time depends on how often you call OSTmrSignal(). In other words, if OSTmrSignal() is called every 1/10 of a second (i.e. OS\_TMR\_CFG\_TICKS\_PER\_SEC set to 10) then, period specifies the number of 1/10 of a second before the timer times out.

```
opt OS TMR OPT PERIODIC:
```

specifies whether you want to timer to automatically reload itself.

OS TMR OPT ONE SHOT:

specifies to stop the timer when it times out.

Note that you MUST select one of these two options.

callback

specifies the address of a function (optional) that you want to execute when the timer expires or, is terminated before it expires (i.e. by calling OSTmrStop ()). The callback function must be declared as follows:

#### void MyCallback (void \*ptmr, void \*callback\_arg);

When the timer expires, this function will be called and passed the timer 'handle' of the expiring timer as well as the argument specified by callback arg.

You should note that you don't have to specify a callback and, in this case, simply pass a NULL pointer.

callback\_arg Is the argument passed to the callback function when the timer expires or is terminated. callback arg can be a NULL pointer if the callback function doesn't require arguments.

pname

Is a pointer to an ASCII string that allows you to give a name to your timer. You can retrieve this name by calling OSTmrNameGet().

perr

a pointer to an error code and can be any of the following:

OS ERR NONE

If the timer was created successfully.

OS ERR TMR INVALID DLY

You specified a delay of 0 when in ONE SHOT mode.

OS ERR TMR INVALID PERIOD

You specified a period of 0 when in PERIODIC mode.

OS ERR TMR INVALID OPT

If you did not specify either OS TMR OPT PERIODIC or OS TMR OPT ONE SHOT.

OS ERR TMR ISR

If you called this function from an ISR, which you are not allowed to do.

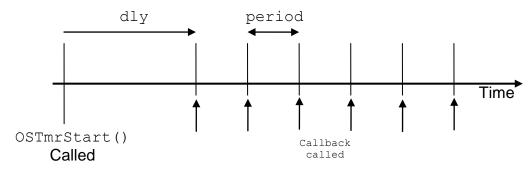
OS ERR TMR NON AVAIL

You get this error when you cannot start a timer because all timer elements (i.e. objects) have already been allocated.

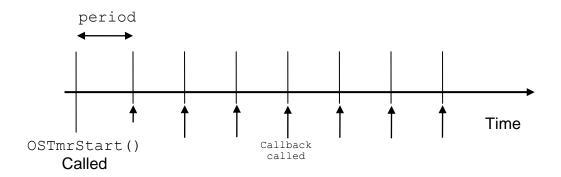
OS ERR TMR NAME TOO LONG

The name you are giving to the timer is too long and must be less than OS TMR CFG NAME SIZE.

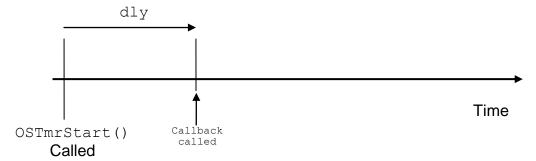
## PERIODIC MODE (see 'opt') - dly > 0



## **PERIODIC MODE** (see 'opt') – dly == 0



## ONE-SHOT MODE (see 'opt') - dly MUST be non-zero



#### **Returned Values**

A pointer to an OS\_TMR object that you MUST use to reference the timer that you just created. A NULL pointer is returned if the timer was not created because of errors (see returned error codes).

#### **Notes/Warnings**

- 1. You should examine the return value to make sure what you get from this function is valid.
- 2. You MUST NOT call this function from an ISR.
- 3. Note that the timer is **NOT** started when it is created. To start the timer, you **MUST** call OSTmrStart().

```
OS TMR *CloseDoorTmr;
void Task (void *p_arg)
{
    INT8U
              err;
    (void)p arg;
    for (;;) {
        CloseDoorTmr = OSTmrCreate( 10,
                                    OS_TMR_OPT_PERIODIC,
                                    DoorCloseFnct,
                                    (void *)0,
                                    "Door Close",
                                    &err);
        if (err == OS ERR NONE) {
            /* Timer was created but NOT started */
    }
```

# OSTmrDel()

BOOLEAN OSTmrDel(OS\_TMR \*ptmr, INT8U \*perr);

Chapter	File	Called from	Code enabled by
New in V2.83	OS_TMR.C	Task	OS_TMR_EN

OSTmrDel() allows you to delete a timer. If a timer was running, it will be stopped and then deleted. If the timer has already timed out and is thus stopped, it will simply be deleted.

It is up to you to delete unused timers. If you delete a timer you MUST NOT reference it anymore.

### Arguments

ptmr is a pointer to the timer that you want to delete. This pointer is returned to you when the timer

is created (see OSTmrCreate()).

perr a pointer to an error code and can be any of the following:

OS ERR NONE If the timer was deleted successfully.

OS ERR TMR INVALID If you passed a NULL pointer for the ptmr argument.

OS ERR TMR INVALID TYPE 'ptmr' is not pointing to a timer.

OS ERR TMR ISR You called this function from an ISR which is NOT allowed.

OS\_ERR\_TMR\_INACTIVE ptmr is pointing to an inactive timer. In other words, you

would get this error if you are pointing to a timer that has been

deleted or was not created.

#### **Returned Values**

OS\_TRUE if the timer was deleted OS FALSE if an error occurred.

#### **Notes/Warnings**

- 1. You should examine the return value to make sure what you get from this function is valid.
- 2. You MUST NOT call this function from an ISR.
- 3. If you delete a timer you **MUST NOT** reference it anymore.

# OSTmrNameGet()

Chapter	File	Called from	Code enabled by
New in V2.81	OS_TMR.C	Task	OS_TMR_EN && OS_TMR_CFG_NAME_EN

OSTmrNameGet() allows you to retrieve the name associated with the specified timer. OSTmrNameGet() places the name of the timer in an array of characters which must be as big as OS\_TMR\_CFG\_NAME\_SIZE (see OS CFG.H).

### Arguments

ptmr is a pointer to the timer that you are inquiring about. This pointer is returned to you when the

timer is created (see OSTmrCreate()).

pdest is a pointer to a pointer to the name of the timer.

perr a pointer to an error code and can be any of the following:

OS\_ERR\_NONE If the name of the task was copied to the array pointed to by

pname.

OS\_ERR\_TMR\_INVALID\_DEST You specified a NULL pointer for pdest.

OS\_ERR\_TMR\_INVALID If you passed a NULL pointer for the ptmr argument.

OS ERR TMR INVALID TYPE 'ptmr' is not pointing to a timer.

OS ERR NAME GET ISR You called this function from an ISR which is NOT allowed.

OS\_ERR\_TMR\_INACTIVE ptmr is pointing to an inactive timer. In other words, you

would get this error if you are pointing to a timer that has been

deleted or was not created.

#### **Returned Values**

The length of the timer name (in number of characters).

- 1. You should examine the return value of this function.
- 2. You MUST NOT call this function from an ISR.

# OSTmrRemainGet()

Chapter	File	Called from	Code enabled by	
New in V2.81	OS_TMR.C	Task	OS_TMR_EN	

OSTMTRemainGet () allows you to obtain the time remaining (before it times out) of the specified timer. The value returned depends on the rate (in Hz) at which the timer task is signaled (see OS\_TMR\_CFG\_TICKS\_PER\_SEC in OS\_CFG.H). In other words, if OS\_TMR\_CFG\_TICKS\_PER\_SEC is set to 10 then the value returned is the number of 1/10 of a second before the timer times out. If the timer has timed out, the value returned will be 0.

### Arguments

ptmr is a pointer to the timer that you are inquiring about. This pointer is returned to you when the

timer is created (see OSTmrCreate()).

perr a pointer to an error code and can be any of the following:

OS\_ERR\_NONE If the function returned the time remaining for the timer.

OS\_ERR\_TMR\_INVALID If you passed a NULL pointer for the ptmr argument.

OS\_ERR\_TMR\_INVALID\_TYPE 'ptmr' is not pointing to a timer.

OS ERR TMR ISR You called this function from an ISR which is NOT allowed.

OS\_ERR\_TMR\_INACTIVE ptmr is pointing to an inactive timer. In other words, you

would get this error if you are pointing to a timer that has been deleted or was not created.

**Returned Values** 

The time remaining for the timer. The value returned depends on the rate (in Hz) at which the timer task is signaled (see OS\_TMR\_CFG\_TICKS\_PER\_SEC in OS\_CFG.H). In other words, if OS\_TMR\_CFG\_TICKS\_PER\_SEC is set to 10 then the value returned is the number of 1/10 of a second before the timer times out. If you specified an invalid timer, the returned value will be 0. If the timer has already expired then the returned value will be 0.

- 1. You should examine the return value to make sure what you get from this function is valid.
- 2. You MUST NOT call this function from an ISR.

# OSTmrSignal()

INT8U OSTmrSignal(void);

Chapter	File	Called from	Code enabled by	
New in V2.81	OS_TMR.C	Task or ISR	OS_TMR_EN	

OSTmrSignal() is called either by a task or an ISR to indicate that it's time to update the timers. Typically, OSTmrSignal() would be called by OSTimeTickHook() at a multiple of the tick rate. In other words, if OS\_TICKS\_PER\_SEC is set to 1000 in OS\_CFG.H then you should call OSTmrSignal() every  $10^{th}$  or  $100^{th}$  tick interrupt (100 Hz or 10 Hz, respectively). You should typically call OSTmrSignal() every 1/10 of a second. The higher the timer rate, of course, the more overhead timer management will impose on your system. Generally, we recommend 10 Hz (1/10 of a second).

You 'could' call <code>OSTmrSignal()</code> from the  $\mu C/OS$ -II tick ISR hook function (see example below). If the tick rate occurs at 1000 Hz then you can simply call <code>OSTmrSignal()</code> every  $100^{th}$  tick. Of course, you would have to implement a simple counter to do this.

## **Arguments**

None.

#### **Returned Values**

OSTmrSignal() uses semaphores to implement the signaling mechanism. Because of that, OSTmrSignal() can return one of the following errors. However, it's very unlikely you will get anything else but OS ERR NONE.

OS_ERR_NONE	The call was successful and the timer task was signaled.
OS ERR SEM OVF	If OSTmrSignal() was called more often than OSTmr Task() can handle the
	timers. This would indicate that your system is heavily loaded.
OS_ERR_EVENT_TYPE	Unlikely you would get this error because the semaphore used for signaling is
	created by μC/OS-II.
OS_ERR_PEVENT_NULL	Again, unlikely you would ever get this error because the semaphore used for
	signaling is created by μC/OS-II.

## **Notes/Warnings**

None.

```
#if OS_TMR_EN > 0
static INT16U OSTmrTickCtr = 0;
#endif

void OSTimeTickHook (void)
{
#if OS_TMR_EN > 0
    OSTmrTickCtr++;
    if (OSTmrTickCtr >= (OS_TICKS_PER_SEC / OS_TMR_CFG_TICKS_PER_SEC)) {
        OSTmrTickCtr = 0;
        OSTmrSignal();
    }
#endif
}
```

# OSTmrStart()

Chapter	File	Called from	Code enabled by	
New in V2.81	OS_TMR.C	Task	OS_TMR_EN	

OSTmrStart () allows you to start (or restart) the countdown process of a timer. The timer to start MUST have previously been created.

### **Arguments**

ptmr is a pointer to the timer that you want to start (or restart). This pointer is returned to you when

the timer is created (see OSTmrCreate()).

perr a pointer to an error code and can be any of the following:

OS ERR NONE If the timer was started.

OS\_ERR\_TMR\_INVALID If you passed a NULL pointer for the ptmr argument.

OS ERR TMR INVALID TYPE 'ptmr' is not pointing to a timer.

OS ERR TMR ISR You called this function from an ISR which is NOT allowed.

OS\_ERR\_TMR\_INACTIVE ptmr is pointing to an inactive timer. In other words, you

would get this error if you are pointing to a timer that has been

deleted or was not created.

### **Returned Values**

OS\_TRUE if the timer was started OS\_FALSE if an error occurred.

- 1. You should examine the return value to make sure what you get from this function is valid.
- 2. You MUST NOT call this function from an ISR.
- 3. The timer to start **MUST** have previously been created.

# OSTmrStateGet()

Chapter	File	Called from	Code enabled by	
New in V2.83	OS_TMR.C	Task	OS TMR EN	

OSTmrStateGet() allows you to obtain the current state of a timer. A timer can be in one of 4 states:

OS\_TMR\_STATE\_UNUSED
OS\_TMR\_STATE\_STOPPED
The timer has not been created
The timer has been created but has not been started or has been stopped.

OS\_TMR\_STATE\_COMPLETED
OS\_TMR\_STATE\_RUNNING
The timer is in ONE-SHOT mode and has completed its delay.
The timer is currently running

#### **Arguments**

ptmr is a pointer to the timer that you are inquiring about. This pointer is returned to you when the

timer is created (see OSTmrCreate()).

perr a pointer to an error code and can be any of the following:

OS\_ERR\_NONE If the function returned the time remaining for the timer.

OS\_ERR\_TMR\_INVALID If you passed a NULL pointer for the ptmr argument.

OS\_ERR\_TMR\_INVALID\_TYPE 'ptmr' is not pointing to a timer.

OS ERR TMR ISR You called this function from an ISR which is NOT allowed.

OS\_ERR\_TMR\_INACTIVE ptmr is pointing to an inactive timer. In other words, you

would get this error if you are pointing to a timer that has been

deleted or was not created.

#### **Returned Values**

The state of the timer (see description).

- 1. You should examine the return value to make sure what you get from this function is valid.
- 2. You MUST NOT call this function from an ISR.

# OSTmrStop()

Chapter	File	Called from	Code enabled by	
New in V2.81	OS_TMR.C	Task	OS_TMR_EN	

OSTMrStop() allows you to stop (i.e. abort) a timer. You can execute the callback function of the timer when it's stopped and pass this callback function a different argument than what was specified when the timer was started. This allows your callback function to *know* that the timer was stopped because the callback argument can be made to indicate this (this, of course, is application specific). If the timer is already stopped, the callback function is not called.

#### **Arguments**

ptmr

Is a pointer to the timer you want to stop. This 'handle' was returned to your application when you called OSTmrStart() and uniquely identifies the timer.

opt

specifies whether you want the timer to:

- 1) OS\_TMR\_OPT\_NONE:
  - Do NOT call the callback function.
- 2) OS\_TMR\_OPT\_CALLBACK: Call the callback function and pass it the callback argument specified when you started the timer (see OSTmrCreate()).
- 3) OS\_TMR\_OPT\_CALLBACK\_ARG:
  Call the callback function BUT pass it the callback argument specified in the
  OSTmrStop() function INSTEAD of the one defined in OSTmrCreate().

callback\_arg If you set opt to OS\_TMR\_OPT\_CALLBACK\_ARG then this is the argument passed to the callback function when it's executed.

perr a pointer to an error code and can be any of the following:

OS ERR NONE If the timer was started. If you passed a NULL pointer for the ptmr argument. OS ERR TMR INVALID OS ERR TMR INVALID TYPE 'ptmr' is not pointing to a timer. You called this function from an ISR which is NOT OS ERR TMR ISR allowed. OS\_ERR\_TMR\_INVALID\_OPT You specified an invalid option for 'opt'. OS ERR TMR STOPPED The timer was already stopped. However, this is NOT considered an actual error since it's OK to attempt to stop a timer that is already stopped. OS ERR TMR INACTIVE ptmr is pointing to an inactive timer. In other words, you

been deleted or was not created.

would get this error if you are pointing to a timer that has

OS\_ERR\_TMR\_NO\_CALLBACK If you wanted the callback to be called but no callback has been specified for this timer.

#### **Returned Values**

```
OS_TRUE if the timer was stopped (even if it was already stopped).
OS_FALSE if an error occurred.
```

## **Notes/Warnings**

- 1. You should examine the return value to make sure what you get from this function is valid.
- 2. You MUST NOT call this function from an ISR.
- 3. The callback function is **NOT** called if the timer is already stopped.

# OSVersion()

INT16U OSVersion(void);

Chapter	File	Called from	Code enabled by
3	OS_CORE.C	Task or ISR	N/A

OSVersion () obtains the current version of  $\mu C/OS\text{-}II$ .

## **Arguments**

none

### **Returned Value**

The version is returned as *x.yy* multiplied by 100. For example, v2.87 is returned as 287.

## **Notes/Warnings**

none

## **Example**

561