CSCE 586 - Design and Analysis of Algorithms

Date: Thursday 6th December, 2018

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Assignment: Final Exam - Take Home Portion, Version A

Documentation:

Problem 1 - Handout

Problem Statement:

Description of Algorithm:

Algorithm 1 Delete-Cycles(S, u)

```
Push u into S V \leftarrow V \cup u for Vertices v_i incident to u do if v \notin V then Delete - Cycles(S, v) else if v \in V and v \in S Delete edge u \rightarrow v end if end for Pop u from S return
```

Asymptotic Analysis:

Worst Case: Best Case: Average Case:

Chapter 5, Problem #4 Problem Statement: Description of Algorithm: Asymptotic Analysis: Worst Case: Best Case: Average Case:

Chapter 6, Problem #9

Problem Statement:

Description of Algorithm:

Algorithm 2 Pseudocode of Algorithm:

```
\begin{aligned} & \textbf{for } j = 1 \rightarrow n \ \textbf{do} \\ & OPT(n,j) \leftarrow min(x_n,s_j) \\ & \textbf{end for} \\ & \textbf{for } i = n-1 \rightarrow 1 \ \textbf{do} \\ & \textbf{for } j = 1 \rightarrow i \ \textbf{do} \\ & OPT(i,j) = max[OPT(i+1,1), min(x_i,s_j) + OPT(i+1,j+1)] \\ & \textbf{end for} \\ & \textbf{end for} \\ & \textbf{return } OPT(1,1) \end{aligned}
```

Asymptotic Analysis:

Worst Case: Best Case: Average Case:

Chapter 6, Problem #12

Problem Statement:

Description of Algorithm:

Algorithm 3 Dynamic Programming for Minimum Server Cost

```
\begin{array}{l} \operatorname{Let} OPT(0) = 0 \text{ and } \binom{1}{2} = 0 \\ \operatorname{for} 1 \leq j \leq (n-1) \operatorname{do} \\ OPT(j) \leftarrow c_j + \min_{0 \leq i < j} (OPT(i) + \binom{j-i}{2}) \\ \operatorname{end} \operatorname{for} \\ OPT(n) = OPT(n-1) + c_n. \\ \operatorname{return} OPT(n) \end{array}
```

Asymptotic Analysis:

Worst Case: Best Case: Average Case: