

## 1 Hidden Surface Removal Problem:

### 1.1 Which algorithmic paradigm will you use to solve this problem?

I will utilize Divide-and-Conquer to solve this problem.

### 1.2 Why did you chose the algorithmic paradigm selected above to solve this problem?

I believe this paradigm is suited to divide and conquer. Each line has similar characteristics, following the same equation  $y_i = a_i \cdot x + b_i$ . The problem can be divided into small, independent sub-problems using the slope of each line. The base case occurs when  $n \leq 3$ : because no 3 lines intersect at a single point, the resulting "uppermost" lines can be found in constant time.

### 1.3 Give an algorithm that takes $n$ lines as input, and in $O(n \log n)$ time returns all of the lines that are visible. Provide a clear description of the algorithm.

Let  $L$  be a set of lines,  $|L| = n$ , where  $L_i = m_i \cdot x + b_i$ .

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**Algorithm 1** Hidden-Surface-Removal: HSR( $L$ )

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Sort  $L$  by ascending slope, such that  $L_i$  has slope  $m_i$  and  $m_i < m_{i+1}$  for all  $i$ .

$|L| = n$

**if**  $n \leq 2$  **then**

**return** The set of lines  $L$ , and their intersection  $a$ .

**else if**  $n = 3$  **then**

    Let  $a$  = the intersection of  $L_1$  and  $L_3$

    Let  $b$  = the intersection of  $L_1$  and  $L_2$

**if**  $x_b < x_a$  **then**

        Let  $c$  = the intersection of  $L_2$  and  $L_3$ .

**return**  $L$  and  $\{b, c\}$

**else**

**return**  $L - \{L_2\}$  and  $\{a\}$

**end if**

**else**

$L, A = \text{HSR}(\{L_1, \dots, L_{\frac{n}{2}}\})$

$L', B = \text{HSR}(\{L_{\frac{n}{2}+1}, \dots, L_n\})$

**end if**

Merge  $A$  and  $B$  into  $C$  by increasing  $x$  coordinate

Find the first element in  $C$  for which  $L' > L$ , denote this element as  $c_k$ .

Let  $L_i \in L$  be the uppermost line in  $L$  immediately before  $c_k$ .

Let  $L_j \in L'$  be the uppermost line in  $L'$  immediately after  $c_k$ .

Let point  $p_{int}$  be the intersection of lines  $L_i$  and  $L_j$ .

$L = \{L_1, L_2, \dots, L_i\} \cup \{L'_j, L'_{j+1}, \dots, L_n\}$

$C = \{A_1, \dots, A_{i-1}\} \cup p_{int} \cup \{B_j, \dots, B_{n-1}\}$

**return**  $L$  and  $C$

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**1.4 Perform asymptotic analysis of your algorithm's running time. Also, consider the run time performance of a best case, worst case, and average case input model scenario.**

**Worst Case:**

**Best Case:**

**Average Case:**

**1.5 Provide a proof that your algorithm works correctly:**

## 2 Bipartite Matching Problem

- 2.1 Which algorithmic paradigm best describes this algorithm?
- 2.2 Why did you choose the algorithmic paradigm selected above?
- 2.3 Give an example of a bipartite graph  $G$  for which this algorithm does not return the maximum matching.
- 2.4 Let  $M$  and  $M'$  be matchings in a bipartite graph  $G$ . Suppose that  $|M'| > 2 \cdot |M|$ . Show that there is an edge  $e' \in M'$  such that  $M \cup e'$  is a matching in  $G$ .
- 2.5 Using the previous claim (and your supporting proof) to further prove that the algorithm is optimal or that the algorithm is  $\rho$ -optimal approximate (in this case be sure to derive the value of  $\rho$  as part of your proof).

### 3 Number Partitioning Problem

#### 3.1 Problem Statement:

Show that the *Number Partitioning* is NP-complete using the *Subset Sum* problem.

**Show *Number Partitioning*  $\in$  NP**

Let  $Y = \text{Number Partitioning}$

**Choose an NP-Complete problem  $X$ :**

Let *Subset Sum* be  $X$ .

**Prove that  $X \leq_P Y$**