

Finite field

In [mathematics](#), a **finite field** or **Galois field** (so-named in honor of [Évariste Galois](#)) is a [field](#) that contains a finite number of [elements](#). As with any field, a finite field is a [set](#) on which the operations of multiplication, addition, subtraction and division are defined and satisfy certain basic rules. The most common examples of finite fields are given by [the integers mod *p*](#) when *p* is a prime number

Finite fields are fundamental in a number of areas of mathematics and computer science, including [number theory](#), [algebraic geometry](#), [Galois theory](#), [finite geometry](#), [cryptography](#) and [coding theory](#).

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Definitions, first examples, and basic properties

The number of elements of a finite field is called its *order*. A finite field of order *q* exists if and only if the order *q* is a [prime power](#) *p*^{*k*} (where *p* is a [prime number](#) and *k* is a positive integer). All finite fields of a given order are [isomorphic](#). In a field of order *p*^{*k*}, adding *p* copies of any element always results in zero; that is, the [characteristic](#) of the field is *p*.

In a finite field of order *q*, the [polynomial](#) *X*^{*q*} − *X* has all *q* elements of the finite field as [roots](#). The non-zero elements of a finite field form a [multiplicative group](#). This group is [cyclic](#), so all non-zero elements can be expressed as powers of a single element called [primitive element](#) of the field (in general there will be several primitive elements for a given field.)

A finite field is a finite set on which the four operations multiplication, addition, subtraction and division (excluding division by zero) are defined, satisfying the rules of arithmetic known as the [field axioms](#). The simplest examples of finite fields are the fields of prime order: for each [prime number](#) *p*, the field GF(*p*) (also denoted **Z**/*p***Z**, **F**_{*p*}, or **F**_{*p*}) of order (that is, size) *p* may be constructed as the [integers modulo *p*](#). A more general algebraic structure that satisfies all the other axioms of a field, but whose multiplication is not required to be commutative, is called a [division ring](#) (or sometimes *skewfield*), however by [Wedderburn's little theorem](#) any finite division ring must be commutative, and hence a finite field.

The elements of such a field may be represented by integers in the range 0, ..., *p* − 1. The sum, the difference and the product are computed by taking the [remainder](#) by *p* of the integer result. The multiplicative inverse of an element may be computed by using the [extended Euclidean algorithm](#) ([§ Extended Euclidean algorithm § Modular integer](#)).

Let *F* be a finite field. For any element *x* in *F* and any [integer](#) *n*, let us denote by *n* · *x* the sum of *n* copies of *x*. The least positive *n* such that *n* · 1 = 0 must exist and is a prime number; it is called the *characteristic* of the field.

If the characteristic of *F* is *p*, one can define multiplication of an element *k* of GF(*p*) by an element *x* of *F* (***k***, ***x***) ↦ ***k*** · ***x*** by choosing an integer representative for *k* and using repeated addition. This multiplication makes *F* into a GF(*p*)-[vector space](#). It follows that the number of elements of *F* is *p*^{*n*} for some integer *n*.

For every prime number *p* and every positive integer *n*, there exist finite fields of order *p*^{*n*}, and all fields of this order are [isomorphic](#) (see [§ Existence and uniqueness](#) below). One may therefore identify all fields of order *p*^{*n*}, which are therefore unambiguously denoted **F**_{*p*^{*n*}}, **F**_{*p*^{*n*}} or GF(*p*^{*n*}), where the letters GF stand for "Galois field"^[1]

The [identity](#)

$$(x + y)^p = x^p + y^p$$

(sometimes called the [freshman's dream](#)) is true (for every *x* and *y*) in a field of characteristic *p*. (This follows from each [binomial coefficient](#) of the expansion of (*x* + *y*)^{*p*}, except the first and the last, being a multiple of *p*).

For each element *x* in the field GF(*p*) for a prime number *p*, one has *x*^{*p*} = *x* (This is an immediate consequence of [Fermat's little theorem](#), and this may be proved as follows: the equality is trivially true for *x* = 0 and *x* = 1; one obtains the result for the other elements of GF(*p*) by applying the above identity to *x* and 1, where *x* successively takes the values 1, 2, ..., *p* − 1 modulo *p*.) This implies the equality

$$X^p - X = \prod_{a \in \text{GF}(p)} (X - a)$$

for polynomials over $\text{GF}(p)$. More generally, every element in $\text{GF}(p^n)$ satisfies the polynomial equation $x^{p^n} - x = 0$.

Any finite field extension of a finite field is separable and simple. That is, if E is a finite field and F is a subfield of E , then E is obtained from F by adjoining a single element whose minimal polynomial is separable. To use a jargon, finite fields are perfect.

Existence and uniqueness

Let $q = p^n$ be a prime power, and F be the splitting field of the polynomial

$$P = X^q - X$$

over the prime field $\text{GF}(p)$. This means that F is a finite field of lowest order, in which P has q distinct roots (the roots are distinct, as the formal derivative of P is equal to -1). The above identity shows that the sum and the product of two roots of P are roots of P , as well as the multiplicative inverse of a root of P . In other word, the roots of P form a field of order q , which is equal to F by the minimality of the splitting field.

The uniqueness up to isomorphism of splitting fields implies thus that all fields of order q are isomorphic.

In summary, we have the following classification theorem first proved in 1893 by E. H. Moore^[2]

The order of a finite field is a prime power. For every prime power q there are fields of order q , and they are all isomorphic. In these fields, every element satisfies

$$x^q = x,$$

and the polynomial $X^q - X$ factors as

$$X^q - X = \prod_{a \in F} (X - a).$$

It follows that $\text{GF}(p^n)$ contains a subfield isomorphic to $\text{GF}(p^m)$ if and only if m is a divisor of n ; in that case, this subfield is unique. In fact, the polynomial $X^{p^m} - X$ divides $X^{p^n} - X$ if and only if m is a divisor of n .

Explicit construction of finite fields

Non-prime fields

Given a prime power $q = p^n$ with p prime and $n > 1$, the field $\text{GF}(q)$ may be explicitly constructed in the following way. One chooses first an irreducible polynomial P in $\text{GF}(p)[X]$ of degree n (such an irreducible polynomial always exists). Then the quotient ring

$$\text{GF}(q) = \text{GF}(p)[X]/(P)$$

of the polynomial ring $\text{GF}(p)[X]$ by the ideal generated by P is a field of order q .

More explicitly, the elements of $\text{GF}(q)$ are the polynomials over $\text{GF}(p)$ whose degree is strictly less than n . The addition and the subtraction are those of polynomials over $\text{GF}(p)$. The product of two elements is the remainder of the Euclidean division by P of the product in $\text{GF}(p)[X]$. The multiplicative inverse of a non-zero element may be computed with the extended Euclidean algorithm; see Extended Euclidean algorithm § Simple algebraic field extensions

Except in the construction of $\text{GF}(4)$, there are several possible choices for P , which produce isomorphic results. To simplify the Euclidean division, for P one commonly chooses polynomials of the form

$$X^n + aX + b,$$

which make the needed Euclidean divisions very efficient. However, for some fields, typically in characteristic 2, irreducible polynomials of the form $X^n + aX + b$ may not exist. In characteristic 2, if the polynomial $X^n + X + 1$ is reducible, it is recommended to choose $X^n + X^k + 1$ with the lowest possible k that makes the polynomial irreducible. If all these trinomials are reducible, one chooses "pentanomials" $X^n + X^a + X^b + X^c + 1$, as polynomials of degree greater than 1, with an even number of terms, are never irreducible in characteristic 2, having 1 as a root.^[3]

In the next sections, we will show how this general construction method works for small finite fields.

Field with four elements

Over $\text{GF}(2)$, there is only one irreducible polynomial of degree 2:

$$X^2 + X + 1$$

Therefore, for $\text{GF}(4)$ the construction of the preceding section must involve this polynomial, and

$$\text{GF}(4) = \text{GF}(2)[X]/(X^2 + X + 1).$$

If one denotes a a root of this polynomial in $\text{GF}(4)$, the tables of the operations in $\text{GF}(4)$ are the following. There is no table for subtraction, as, in every field of characteristic 2, subtraction is identical to addition. In the third table, for the division $x \div y$, x must be read on the left, and y on the top.

+	0	1	a	1+a
0	0	1	a	1+a
1	1	0	1+a	a

×	0	1	a	1+a
0	0	0	0	0
1	0	1	a	1+a

x/y	0	1	a	1+a
0	–	0	0	0
1	–	1	1+a	a

a	a	1+a	0	1
1+a	1+a	a	1	0

a	0	a	1+a	1
1+a	0	1+a	1	a

a	−	a	1	1+a
1+a	−	1+a	a	1

GF(p²) for an odd prime p

For applying the above general construction of finite fields in the case of GF(*p*²), one has to find an irreducible polynomial of degree 2. For *p* = 2, this has been done in the preceding section. If *p* is an odd prime, there are always irreducible polynomials of the form *x*² − *r*, with *r* in GF(*p*).

More precisely, the polynomial *x*² − *r* is irreducible over GF(*p*) if and only if *r* is a quadratic non-residue modulo *p* (this is almost the definition of a quadratic non-residue). There are $\frac{p-1}{2}$ quadratic non-residues modulo *p*. For example, 2 is a quadratic non-residue for *p* = 3, 5, 11, 13, ..., and 3 is a quadratic non-residue for *p* = 5, 7, 17, If *p* ≡ 3 mod 4, that is *p* = 3, 7, 11, 19, ..., one may choose −1 ≡ *p* − 1 as a quadratic non-residue, which allows us to have a very simple irreducible polynomial *x*² + 1.

Having chosen a quadratic non-residue *r*, let *α* be a symbolic square root of *r*, that is a symbol which has the property *α*² = *r*, in the same way as the complex number *i* is a symbolic square root of −1. Then, the elements of GF(*p*²) are all the linear expressions

***a* + *bα*,**

with *a* and *b* in GF(*p*). The operations on GF(*p*²) are defined as follows (the operations between elements of GF(*p*) represented by Latin letters are the operations in GF(*p*)):

−(a + bα) = −a + (−b)α
(a + bα) + (c + dα) = (a + c) + (b + d)α
(a + bα)(c + dα) = (ac + rbd) + (ad + bc)α
(a + bα)^{−1} = a(a² − rb²)^{−1} + (−b)(a² − rb²)^{−1}α

GF(8) and GF(27)

The polynomial

***X*³ − *X* − 1**

is irreducible over GF(2) and GF(3), that is, it is irreducible modulo 2 and 3 (to show this it suffices to show that it has no root in GF(2) nor in GF(3)). It follows that the elements of GF(8) and GF(27) may be represented by expressions

***a* + *bα* + *cα*²,**

where *a*, *b*, *c* are elements of GF(2) or GF(3) (respectively), and *α* is a symbol such that

***α*³ = *α* + 1.**

The addition, additive inverse and multiplication on GF(8) and GF(27) may thus be defined as follows; in following formulas, the operations between elements of GF(2) or GF(3), represented by Latin letters, are the operations in GF(2) or GF(3), respectively:

***−(a + bα + cα*²*) = −a + (−b)α + (−c)α*² (for GF(8), this operation is the identity)**
***(a + bα + cα*²*) + (d + eα + fα*²*) = (a + d) + (b + e)α + (c + f)α*²**
***(a + bα + cα*²*)(d + eα + fα*²*) = (ad + bf + ce) + (ae + bd + bf + ce + cf)α + (af + be + cd + cf)α*²**

GF(16)

The polynomial

***X*⁴ + *X* + 1**

is irreducible over GF(2), that is, it is irreducible modulo 2. It follows that the elements of GF(16) may be represented by expressions

***a* + *bα* + *cα*² + *dα*³,**

where *a*, *b*, *c*, *d* are either 0 or 1 (elements of GF(2)), and *α* is a symbol such that

***α*⁴ = *α* + 1.**

As the characteristic of GF(2) is 2, each element is its additive inverse in GF(16). The addition and multiplication on GF(16) may be defined as follows; in following formulas, the operations between elements of GF(2), represented by Latin letters are the operations in GF(2).

***(a + bα + cα*² *+ dα*³*) + (e + fα + gα*² *+ hα*³*) = (a + e) + (b + f)α + (c + g)α*² *+ (d + h)α*³**
(a + bα + cα*² *+ dα*³*)(e + fα + gα*² *+ hα*³*) = (ae + bh + cg + df) + (af + be + bh + cg + df + ch + dg)α +
***(ag + bf + ce + ch + dg + dh)α*² *+ (ah + bg + cf + de + dh)α*³**

Multiplicative structure

The set of non-zero elements in GF(*q*) is an abelian group under the multiplication, of order *q* − 1. By Lagrange's theorem there exists a divisor *k* of *q* − 1 such that *x*^{*k*} = 1 for every non-zero *x* in GF(*q*). As the equation *x*^{*k*} = 1 has at most *k* solutions in any field, *q* − 1 is the lowest possible value for *k*. The structure theorem of finite abelian groups implies that this multiplicative group is cyclic, that all non-zero elements are powers of single element. In summary:

The multiplicative group of the non-zero elements in $\text{GF}(q)$ is cyclic, and there exist an element a , such that the $q - 1$ non-zero elements of $\text{GF}(q)$ are $a, a^2, \dots, a^{q-2}, a^{q-1} = 1$.

Such an element a is called a primitive element. Unless $q = 2, 3$, the primitive element is not unique. The number of primitive elements is $\varphi(q - 1)$ where φ is Euler's totient function.

The result above implies that $x^q = x$ for every x in $\text{GF}(q)$. The particular case where q is prime is Fermat's little theorem.

Discrete logarithm

If a is a primitive element in $\text{GF}(q)$, then for any non-zero element x in F , there is a unique integer n with $0 \leq n \leq q - 2$ such that

$$x = a^n.$$

This integer n is called the discrete logarithm of x to the base a .

While a^n can be computed very quickly, for example using exponentiation by squaring, there is no known efficient algorithm for computing the inverse operation, the discrete logarithm. This has been used in various cryptographic protocols see Discrete logarithm for details.

When the nonzero elements of $\text{GF}(q)$ are represented by their discrete logarithms, multiplication and division are easy, as they reduce to addition and subtraction modulo $q - 1$. However, addition amounts to computing the discrete logarithm of $a^m + a^n$. The identity

$$a^m + a^n = a^n(a^{m-n} + 1)$$

allows one to solve this problem by constructing the table of the discrete logarithms of $a^n + 1$, called Zech's logarithms for $n = 0, \dots, q - 2$ (it is convenient to define the discrete logarithm of zero as being $-\infty$).

Zech's logarithms are useful for large computations, such as linear algebra over medium-sized fields, that is, fields that are sufficiently large for making natural algorithms inefficient, but not too large, as one has to pre-compute a table of the same size as the order of the field.

Roots of unity

Every nonzero element of a finite field is a root of unity, as $x^{q-1} = 1$ for every nonzero element of $\text{GF}(q)$.

If n is a positive integer, an n th **primitive root of unity** is a solution of the equation $x^n = 1$ that is not a solution of the equation $x^m = 1$ for any positive integer $m < n$. If a is a n th primitive root of unity in a field F , then F contains all the n roots of unity, which are $1, a, a^2, \dots, a^{n-1}$.

The field $\text{GF}(q)$ contains a n th primitive root of unity if and only if n is a divisor of $q - 1$; if n is a divisor of $q - 1$, then the number of primitive n th roots of unity in $\text{GF}(q)$ is $\varphi(n)$ (Euler's totient function). The number of n th roots of unity in $\text{GF}(q)$ is $\gcd(n, q - 1)$.

In a field of characteristic p , every (np) th root of unity is also an n th root of unity. It follows that primitive (np) th roots of unity never exist in a field of characteristic p .

On the other hand, if n is coprime to p , the roots of the n th cyclotomic polynomial are distinct in every field of characteristic p , as this polynomial is a divisor of $X^n - 1$, which has 1 as formal derivative. It follows that the n th cyclotomic polynomial factors over $\text{GF}(p)$ into distinct irreducible polynomials that have all the same degree, say d , and that $\text{GF}(p^d)$ is the smallest field of characteristic p that contains the n th primitive roots of unity.

Example: GF(64)

The field $\text{GF}(64)$ has several interesting properties that smaller fields do not share: it has two subfields such that neither is contained in the other; not all generators (elements with minimal polynomial of degree 6 over $\text{GF}(2)$) are primitive elements; and the primitive elements are not all conjugate under the Galois group.

The order of this field being 2^6 , and the divisors of 6 being 1, 2, 3, 6, the subfields of $\text{GF}(64)$ are $\text{GF}(2)$, $\text{GF}(2^2) = \text{GF}(4)$, $\text{GF}(2^3) = \text{GF}(8)$, and $\text{GF}(64)$ itself. As 2 and 3 are coprime, the intersection of $\text{GF}(4)$ and $\text{GF}(8)$ in $\text{GF}(64)$ is the prime field $\text{GF}(2)$.

The union of $\text{GF}(4)$ and $\text{GF}(8)$ has thus 10 elements. The remaining 54 elements of $\text{GF}(64)$ generate $\text{GF}(64)$ in the sense that no other subfield contains any of them. It follows that they are roots of irreducible polynomials of degree 6 over $\text{GF}(2)$. This implies that, over $\text{GF}(2)$, there are exactly $9 = \frac{54}{6}$ irreducible monic polynomials of degree 6. This may be verified by factoring $X^{64} - X$ over $\text{GF}(2)$.

The elements of $\text{GF}(64)$ are primitive n th roots of unity for some n dividing 63. As the 3rd and the 7th roots of unity belong to $\text{GF}(4)$ and $\text{GF}(8)$, respectively, the 54 generators are primitive n th roots of unity for some n in $\{9, 21, 63\}$. Euler's totient function shows that there are 6 primitive 9th roots of unity, 12 primitive 21st roots of unity, and 36 primitive 63rd roots of unity. Summing these numbers, one finds again 54 elements.

By factoring the cyclotomic polynomial over $\text{GF}(2)$, one finds that:

- The six primitive 9th roots of unity are roots of

$$X^6 + X^3 + 1,$$

and are all conjugate under the action of the Galois group.

- The twelve primitive 21st roots of unity are roots of

$$(X^6 + X^4 + X^2 + X + 1)(X^6 + X^5 + X^4 + X^2 + 1).$$

They form two orbits under the action of the Galois group. As the two factors are reciprocal to each other, a root and its (multiplicative) inverse do not belong to the same orbit.

- The 36 primitive elements of $\text{GF}(64)$ are the roots of

$$(X^6 + X^4 + X^3 + X + 1)(X^6 + X + 1)(X^6 + X^5 + 1)(X^6 + X^5 + X^3 + X^2 + 1)(X^6 + X^5 + X^2 + X + 1)(X^6 + X^5$$

They split into 6 orbits of 6 elements under the action of the Galois group.

This shows that the best choice to construct GF(64) is to define it as GF(2)[X]/(X⁶ + X + 1). In fact, this generator is a primitive element, and this polynomial is the irreducible polynomial that produces the easiest Euclidean division.

Frobenius automorphism and Galois theory

In this section, *p* is a prime number and *q* = *p*^{*n*} is a power of *p*.

In GF(*q*), the identity (*x* + *y*)^{*p*} = *x*^{*p*} + *y*^{*p*} implies that the map

φ : x ↦ x^p

is a GF(*p*)-linear endomorphism and a field automorphism of GF(*q*), which fixes every element of the subfield GF(*p*). It is called the Frobenius automorphism after Ferdinand Georg Frobenius.

Denoting by φ^{*k*} the composition of φ with itself *k* times, we have

φ^k : x ↦ x^{p^k}.

It has been shown in the preceding section that φ^{*n*} is the identity: For 0 < *k* < *n*, the automorphism φ^{*k*} is not the identity as, otherwise, the polynomial

X^{p^k} − X

would have more than *p*^{*k*} roots.

There are no other GF(*p*)-automorphisms of GF(*q*). In other words, GF(*p*^{*n*}) has exactly *n* GF(*p*)-automorphisms, which are

Id = φ⁰, φ, φ², . . . , φ^{*n*−1}.

In terms of Galois theory, this means that GF(*p*^{*n*}) is a Galois extension of GF(*p*), which has a cyclic Galois group.

The fact that the Frobenius map is surjective implies that every finite field is perfect.

Polynomial factorization

If *F* is a finite field, a non-constant monic polynomial with coefficients in *F* is irreducible over *F*, if it is not the product of two non-constant monic polynomials, with coefficients in *F*.

As every polynomial ring over a field is a unique factorization domain, every monic polynomial over a finite field may be factored in a unique way (up to the order of the factors) into a product of irreducible monic polynomials.

There are efficient algorithms for testing polynomial irreducibility and factoring polynomials over finite field. They are a key step for factoring polynomials over the integers or the rational numbers. At least for this reason, every computer algebra system has functions for factoring polynomials over finite fields, or at least, over finite prime fields.

Irreducible polynomials of a given degree

The polynomial

X^{*q*} − X

factors into linear factors over a field of order *q*. More precisely, this polynomial is the product of all monic polynomials of degree one over a field of order *q*.

This implies that, if *q* = *p*^{*n*} then X^{*q*} − X is the product of all monic irreducible polynomials over GF(*p*), whose degree divides *n*. In fact, if *P* is an irreducible factor over GF(*p*) of X^{*q*} − X, its degree divides *n*, as its splitting field is contained in GF(*p*^{*n*}). Conversely, if *P* is an irreducible monic polynomial over GF(*p*) of degree *d* dividing *n*, it defines a field extension of degree *d*, which is contained in GF(*p*^{*n*}), and all roots of *P* belong to GF(*p*^{*n*}), and are roots of X^{*q*} − X; thus *P* divides X^{*q*} − X. As X^{*q*} − X does not have any multiple factor, it is thus the product of all the irreducible monic polynomials that divide it.

This property is used to compute the product of the irreducible factors of each degree of polynomials over GF(*p*); see Distinct degree factorization

Number of monic irreducible polynomials of a given degree over a finite field

The number *N*(*q*, *n*) of monic irreducible polynomials of degree *n* over GF(*q*) is given by^[4]

N(q,n) = 1/n ∑_{d|n} μ(d) q^{n/d},

where μ is the Möbius function. This formula is almost a direct consequence of above property on X^{*q*} − X.

By the above formula, the number of irreducible (not necessarily monic) polynomials of degree *n* over GF(*q*) is (q − 1)*N*(*q*, *n*).

A (slightly simpler) lower bound for *N*(*q*, *n*) is

N(q,n) ≥ 1/n (qⁿ − ∑_{p|n, p prime} q^{n/p}).

One may easily deduce that, for every *q* and every *n*, there is at least one irreducible polynomial of degree *n* over GF(*q*). This lower bound is sharp for *q* = *n* = 2.

Applications

In cryptography, the difficulty of the discrete logarithm problem in finite fields or in elliptic curves is the basis of several widely used protocols, such as the Diffie–Hellman protocol. For example, in 2014 a secure internet connection to Wikipedia involved the elliptic curve Diffie–Hellman protocol (ECDHE) over a large finite field.^[5] In coding theory, many codes are constructed as subspaces of vector spaces over finite fields.

Finite fields are widely used in number theory, as many problems over the integers may be solved by reducing them modulo one or several prime numbers. For example, the fastest known algorithms for polynomial factorization and linear algebra over the field of rational numbers proceed by reduction modulo one or several primes, and then reconstruction of the solution by using Chinese remainder theorem, Hensel lifting or the LLL algorithm.

Similarly many theoretical problems in number theory can be solved by considering their reductions modulo some or all prime numbers. See, for example, Hasse principle. Many recent developments of algebraic geometry were motivated by the need to enlarge the power of these modular methods. Wiles' proof of Fermat's Last Theorem is an example of a deep result involving many mathematical tools, including finite fields.

Extensions

Algebraic closure

A finite field **F** is not algebraically closed. To demonstrate this, consider the polynomial

f(T) = 1 + \prod_{\alpha \in \mathbf{F}} (T - \alpha),

which has no roots in **F**, since f(α) = 1 for all α in **F**.

The direct limit of the system:

{\mathbf{F}_p, \mathbf{F}_{p^2}, ..., \mathbf{F}_{p^n}, ...},

with inclusion, is an infinite field. It is the algebraic closure of all the fields in the system, and is denoted by $\overline{\mathbf{F}_p}$.

The inclusions commute with the Frobenius map, as it is defined the same way on each field ($x \mapsto x^p$), so the Frobenius map defines an automorphism of $\overline{\mathbf{F}_p}$, which carries all subfields back to themselves. In fact \mathbf{F}_{p^n} can be recovered as the fixed points of the *n*th iterate of the Frobenius map.

However unlike the case of finite fields, the Frobenius automorphism on $\overline{\mathbf{F}_p}$ has infinite order, and it does not generate the full group of automorphisms of this field. That is, there are automorphisms of $\overline{\mathbf{F}_p}$ which are not a power of the Frobenius map. However, the group generated by the Frobenius map is a dense subgroup of the automorphism group in the Krull topology. Algebraically, this corresponds to the additive group **Z** being dense in the profinite integers (direct product of the *p*-adic integers over all primes *p*, with the product topology).

If we actually construct our finite fields in such a fashion that \mathbf{F}_{p^n} is contained in \mathbf{F}_{p^m} whenever *n* divides *m*, then this direct limit can be constructed as the union of all these fields. Even if we do not construct our fields this way we can still speak of the algebraic closure, but some more delicacy is required in its construction.

Wedderburn's little theorem

A division ring is a generalization of field. Division rings are not assumed to be commutative. There are no non-commutative finite division rings: Wedderburn's little theorem states that all finite division rings are commutative, hence finite fields. The result holds even if we relax associativity and consider alternative rings by the Artin–Zorn theorem

Relationship to other commutative ring classes

Finite fields appear in the following chain of inclusions:

commutative rings \supset integral domains \supset integrally closed domains \supset GCD domains \supset unique factorization domains \supset principal ideal domains \supset Euclidean domains \supset fields \supset finite fields

See also

- Quasi-finite field
- Field with one element
- Finite field arithmetic
- Trigonometry in Galois fields
- Finite ring
- Finite group
- Elementary abelian group
- Hamming space

Notes

1. This notation was introduced by E. H. Moore in an address given in 1893 at the International Mathematical Congress held in Chicago. Mullen & Panario 2013 p. 10.

2. Moore, E. H. (1896), "Tables of Galois fields of order < 1000" *Bulletin of the American Mathematical Society* **16**(4): 188–206, doi:10.1090/S0002-9904-1910-01888-7

3. *Recommended Elliptic Curves for Government Use* (<http://csrc.nist.gov/groups/ST/toolkit/documents/dss/NISTReCurve.pdf>) (PDF), National Institute of Standards and Technology, July 1999, p. 3

4. Jacobson 2009 §4.13

5. This can be verified by looking at the information on the page provided by the browser

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External links

- Finite Fields at Wolfram research.

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