

Warmup: any concerns?

Library code

```
/** Returns the positive
 * square root of `a`.
 * If the argument is
 * NaN or less than
 * zero, the result is
 * NaN. */
public static double
    sqrt(double a) { ... }
```

Client code

```
double a1 = 16.0;
double a2 = 25.0;
double s1 =
    Math.sqrt(a1 - a2);
assert !(s1 < 0);
```

Warmup: any concerns?

Note: NaN is not less than, greater than, or equal to any value (even itself)

- A. The client code is buggy – it violates a *precondition*
- B. The client code is buggy – it asserts something not guaranteed by the *postcondition*
- C. The library's *specification* is incomplete
- D. The code shown here is fine

```
/** Returns the positive square  
 * root of `a`. If the argument  
 * is NaN or less than zero,  
 * the result is NaN. */
```

```
public static double  
  sqrt(double a) { ... }
```

```
// Client code  
double a1 = 16.0;  
double a2 = 25.0;  
double s1 =  
  Math.sqrt(a1 - a2);  
assert !(s1 < 0);
```



CS 2110

Lecture 5

Interfaces, subtyping,
polymorphism



Coming up

A2 should be released today

A1 is being graded now

Expect feedback Thurs

iClickers



Interfaces

“Interfaces”

- Mechanism by which two parties work together, decided ahead of time
 - HDMI cable: interface between laptop and projector
 - ¼”-20 screw: interface between tripods and camera gear
 - API: interface between client programmers and classes
- Parties don’t need to know each other’s details
 - Forms an “**abstraction barrier**”
- (Client) interface = a class’s `public` methods

Abstraction reminder: separate “what” from “how”



“The power of abstraction... is the secret sauce of the internet. And, indeed, all of computer science.”

- *The Economist*, Jan 29, 2024

Discuss

In a class representing a cafeteria, someone has written a public method to predict the wait time.

1.

2.

What do you (the client) need to know in order to make use of this method?

3.

4.

Java interfaces

- Guarantee to clients what a type *can do*, without committing to details (i.e. fields)
 - Method signatures
 - Method return types
 - Method specs
- Method declarations are implicitly **public**

```
/** Closed interval on real
 * number line. */
public interface Interval {
    /** Return left endpoint. */
    double left();
    /** Return right endpoint. */
    double right();
    /** Whether x is contained
     * in this interval. */
    boolean contains(double x);
}
```

Client code 1

```
static boolean intervalsOverlap(Interval i1,  
                                Interval i2) {  
    return i1.left() <= i2.right() &&  
           i2.left() <= i1.right();  
}
```

Implementer's code: Option A

```
public class TwoPtInterval
    implements Interval {
    /** Left endpoint. */
    private double left;
    /** Right endpoint. */
    private double right;

    @Override
    double left() {
        return left;
    }
}
```

```
    @Override
    double right() {
        return right;
    }

    @Override
    boolean contains(
        double x) {
        return left <= x &&
            x <= right;
    }
}
```

Implementer's code: Option B

```
public class CenterInterval
    implements Interval {
    /** Midpoint. */
    private double center;
    /** Width. */
    private double width;

    @Override
    double left() {
        return center - width/2;
    }
}
```

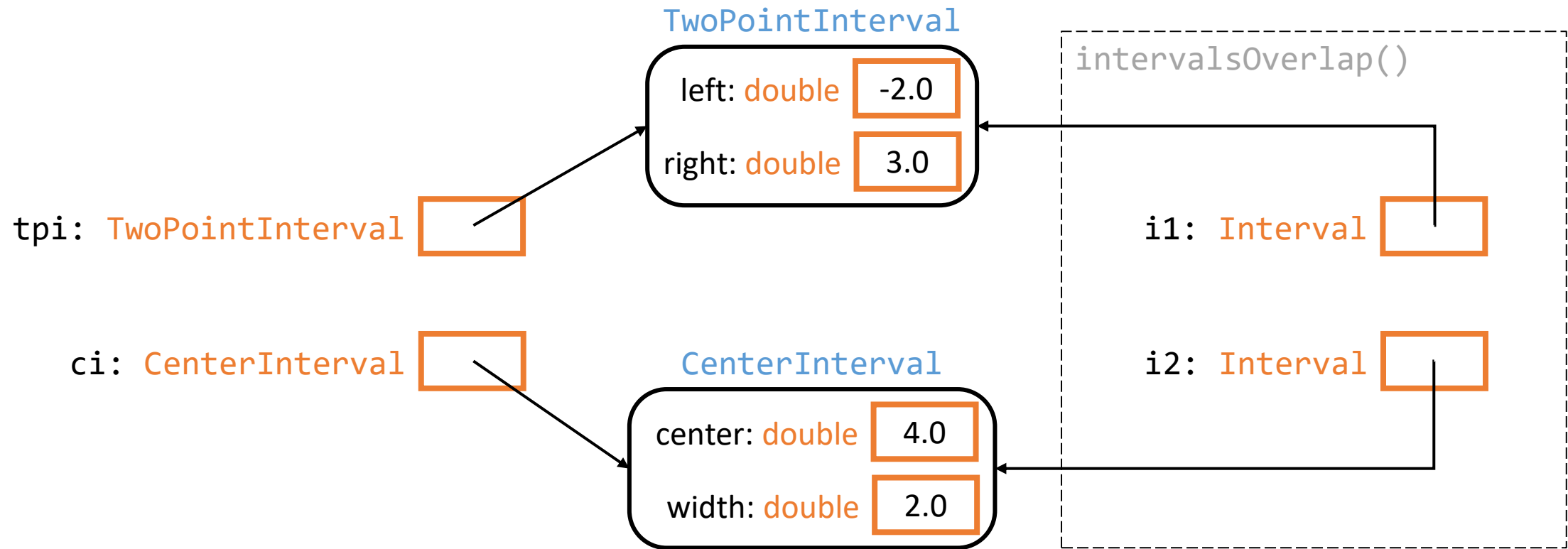
```
@Override
double right() {
    return center + width/2;
}
```

```
@Override
boolean contains(
    double x) {
    return abs(x - center)
        <= width/2;
}
}
```

Client code 2

```
TwoPtInterval tpi = new TwoPtInterval(-2, 3);  
CenterInterval ci = new CenterInterval(4, 2);  
  
// Is this allowed?  
boolean overlap = intervalsOverlap(tpi, ci);
```

Object diagram



Subtypes

- A `TwoPtInterval` can do anything an `Interval` can do
 - Behavior bound by same specifications
- Therefore, a `TwoPtInterval` can be used anywhere an `Interval` is expected
- Implementing an `interface` establishes a **subtype** relationship
 - `TwoPtInterval <: Interval`
 - `CenterInterval <: Interval`

Subtype compatibility

- Assignment
 - T $x = \text{expr}$; is allowed if the type of expr is a subtype of T
 - Argument passing
 - `void foo(T x);`
`foo(expr)`
is allowed if the type of expr is a subtype of T
 - Returning
 - T `bar() { return expr ; }` is allowed if the type of expr is a subtype of T
- “If x can store an animal, it can store a cat.”
 - “If `foo()` expects a bird, it can work with a robin.”
 - “If `bar()` says it will return an insect, it’s allowed to give me a cricket.”

Quick check

Suppose `Foo <: Bar` .

Consider these method declarations:

- `Foo f();`
- `Bar g();`

Which of these is allowed?

A. `Foo f = g();`

B. `Bar b = f();`

C. Neither

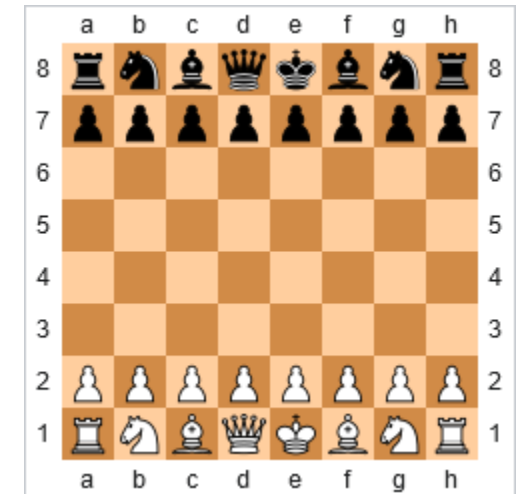




(subtype) Polymorphism

Variations in behavior

- The Interval interface abstracted over state, but both implementations behaved identically
- Sometimes, behavior specifications leave room for variation
- Example: chess pieces



Chess piece interface

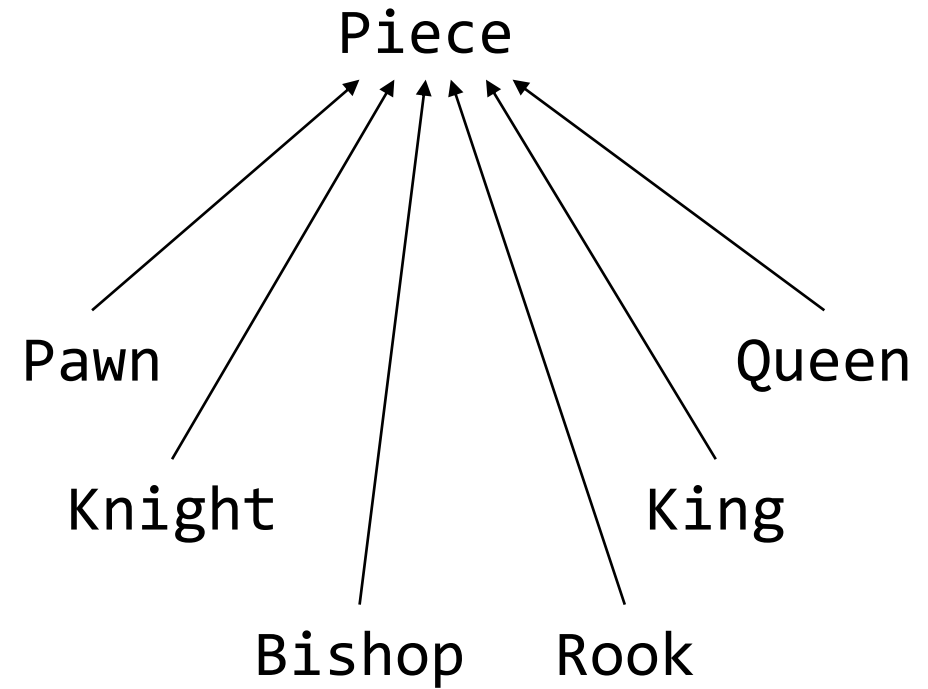
```
public interface Piece {  
    /** Return whether this piece is able to move to  
     * location (`dstRow`, `dstCol`) from its current  
     * position, given board config. `board`.  
     * Requires dstRow, dstCol in [0..7]. */  
    boolean legalMove(int dstRow, int dstCol,  
                      Board board);  
}
```

Chess board interface

```
public interface Board {  
    /** Return 0 if position (`row`, `col`) is empty,  
     * 1 if occupied by a white piece, 2 if occupied  
     * by a black piece. Requires row, col in  
     * [0..7]. */  
    int playerAt(int row, int col);  
}
```

Type hierarchy

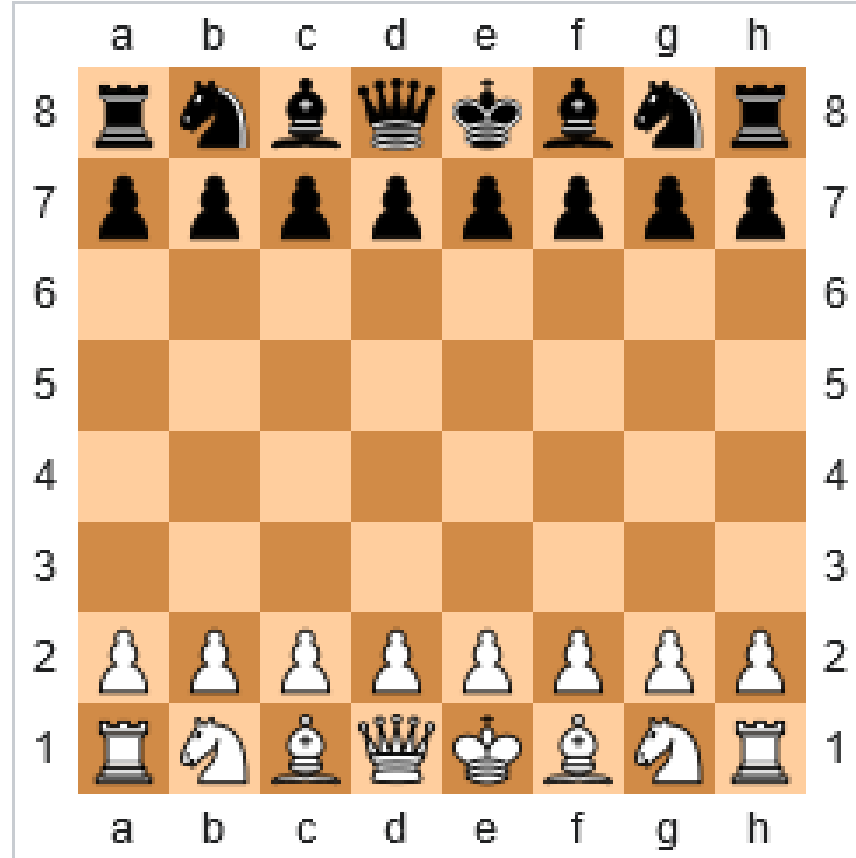
- Pawn <: Piece
- Knight <: Piece
- Bishop <: Piece
- Rook <: Piece
- Queen <: Piece
- King <: Piece



Knight

```
public class Knight
    implements Piece {
    private int row;
    private int col;
    private int player;
    @Override
    public boolean legalMove(
        int dstRow,
        int dstCol,
        Board board) {
```

```
    int dx = abs(row-dstRow);
    int dy = abs(col-dstCol);
    return board.playerAt(
        dstRow, dstCol) != player
        && ((dx==1 && dy==2) ||
            (dx==2 && dy==1));
    }}
```



King

```
public class King
    implements Piece {
    private int row;
    private int col;
    private int player;
    private boolean hasMoved;
    @Override
    public boolean legalMove(
        int dstRow
        int dstCol,
        Board board) {
```

```
    int dx = abs(row-dstRow);
    int dy = abs(col-dstCol);
    return board.playerAt(
        dstRow, dstCol) != player
        && (dx <= 1 && dy <= 1
            || !hasMoved &&
                canCastle(board));
}
    public boolean canCastle(
        Board board) { ...
    }}
```

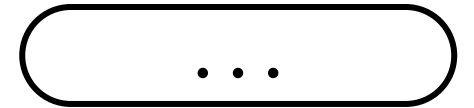
Object diagram

```
Piece pickNextPiece() {...}  
// ...  
Piece p;  
while (!gameOver) {  
    p = pickNextPiece();  
    // assign r, c  
    if (p.legalMove(r, c)) {  
        // ...  
    }  
}
```

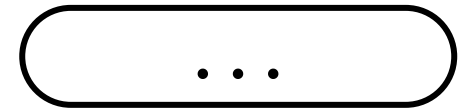
p: Piece



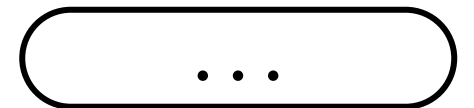
Pawn



King



Knight



Static vs. dynamic type

- While the program is running, the type of the object referenced by `p` could change, but it will always be a subtype of `Piece`
- **Static type**: types declared for variables & return values, derived for expressions (compile-time)
- **Dynamic type**: the type of an object being referenced (runtime)
- Behavior is determined by *dynamic* type
 - “**Dynamic dispatch**”

Poll

Should a client be able to call `p.canCastle()` when the *dynamic type* of the object referenced by `Piece p`` is a `King`?

- A. Yes
- B. No
- C. Only if they know more than the compiler



Compile-time reference rule

- Client can only request behavior supported by the target's **static** type
 - Guarantees that requested method will exist (unless target is null)
 - Compiler does not reason about **dynamic types**, even if “obvious”
 - **Piece** p = **new** King(...); // Static type of p is Piece

**Most important rule
in the course!**

Factories

```
/** Create a Piece of the
 * type specified by `code`
 * (algebraic notation). */
static Piece makePiece(
    char code) {
    if (code == 'K') {
        return new King(...);
    } else if (code == 'N') {
        return new Knight(...);
    } else // ...
}
```

Why declare with less-specific static types?

- Might not know which constructor will be called!

```
Piece p1 = makePiece('K');
// Static type of p1?
// Dynamic type of p1?
```

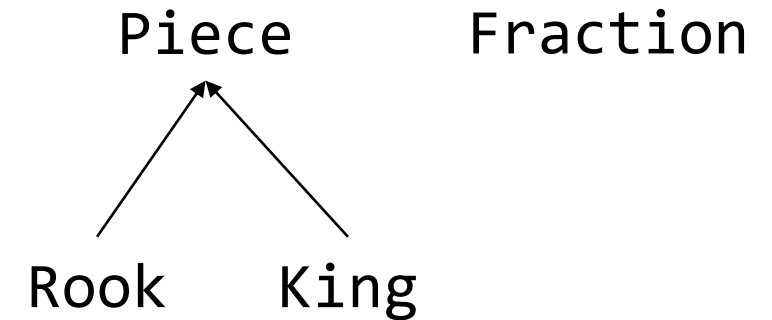
```
char code = in.readChar();
Piece p2 = makePiece(code);
```

Casting between reference types

- **Dynamic dispatch** means an object will behave according to its **dynamic type** at runtime
- What if you think you know an object's dynamic type more specifically than its static type?
 - Can **cast** an expression to recover access to additional behavior
 - **King** k = (King) p;
- Casting is *checked* at runtime, but otherwise has no effect on the object
 - Just used to massage static types
 - If you're wrong, will get a runtime error (ClassCastException)

Casting examples

- `k = (King) p;`
 - Allowed; could fail at runtime
- `p = (Piece) k;`
 - Allowed but unnecessary; will always succeed
 - `p = k;` allowed by subtype substitution rule
- `k = (King) r; k = (King)f;`
 - Not allowed (impossible)
- `k = (King)(Piece) r;`
 - Allowed; will fail at runtime



`p: Piece` `k: King`

`f: Fraction` `r: Rook`

Checking dynamic type before casting

- `instanceof` operator
 - ``expr instanceof T`` is true if the `dynamic type` of `expr` is a subtype of `T`

```
if (p instanceof King) {  
    King k = (King) p;  
    if (k.canCastle()) {...}  
}
```

- Runtime type queries are useful in *some* circumstances, but are usually a sign of poor OO design
- When possible, specify common behavior in a supertype, then leverage dynamic dispatch for polymorphism

Checkpoint

- Compile-time errors
 - Red underlines in IntelliJ
 - Red icons in “Problems” panel
 - Compilation errors from smoketester
- Runtime errors
 - Exception backtrace when program is run
- Spec violations are *bugs* but might not prevent compilation or produce an exception
 - Defensive programming: turn spec violations into runtime errors

Any problems?

A: Compile-time error
B: Runtime error
C: Spec violated
D: Okay

Implementer

```
public class DynArray {  
    /** Double capacity */  
    private void incCap() {  
        ...  
    }  
}
```

Client

```
DynArray a =  
    new DynArray();  
a.incCap();
```



Any problems?

A: Compile-time error
B: Runtime error
C: Spec violated
D: Okay

Implementer

```
interface Phone {  
    void call(int[] num);  
}  
class iPhone  
    implements Phone {  
    public void call(  
        int[] num) {...}  
    public void takePic() {...}  
}
```

Client

```
Phone p;  
p = new iPhone();  
p.takePic();
```



ORGANIZATION CHART of THE TABULATING MACHINE CO.

BOARD OF DIRECTORS - C-T-R- CO.

Alfred DeBuys	Clarence P. King
George W. Fairchild	Stacy C. Richmond
Charles R. Flint	Joseph E. Rogers
A. Ward Ford	Christopher D. Smithers
Oscar L. Gubelman	Thomas J. Watson
Samuel M. Hastings	George I. Wilber
John W. Herbert	Rollin S. Woodruff
Joel S. Coffin	

OFFICERS-C-T-R-CO.

Thomas J. Watson - Pres. & Genl. Mgr.
George W. Fairchild - Vice-President
James S. Ogsbury - Secy & Treasurer

COMPUTING-TABULATING-RECORDING CO.
Offices - 50 Broad St. - New York City

THE TABULATING MACHINE CO.

General Offices - 50 Broad St.
New York City

DIRECTORS

George M. Bond James S. Ogsbury
George W. Fairchild Gershom Smith
Thomas J. Watson

FACTORIES - WASHINGTON, D. C.
- ENDICOTT, N. Y.
- DAYTON, O.

THOMAS J. WATSON *President*
R. L. Houston *General Manager*

OFFICERS

Thomas J. Watson - President
Gershom Smith - Vice-President
R. L. Houston - Treasurer
W. D. Jones - Asst. Treasurer
James S. Ogsbury - Secretary
O. E. Braitmayer - Asst. Secretary

