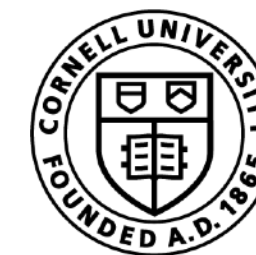




Market Clearing Price

NETWORKS INFO 2040 / CS 2850 / ECON 2040 / SOC 2090



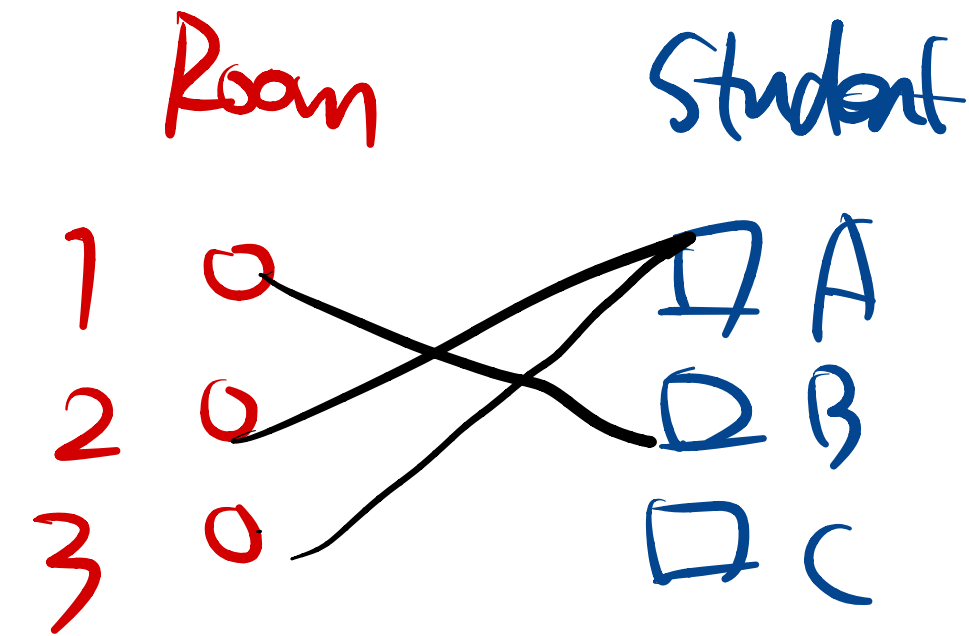
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College of Computing
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PS 2. Due time 3.30 PM

Recap

Bipartite Graph

2 types of nodes



Assignment Problem

Perfect Matching (P.M.)

Neighbor set $N(S)$: {set of nodes that connect to at least one node in S }

$$f(x) = 2x \quad \text{number} \rightarrow \text{number}$$

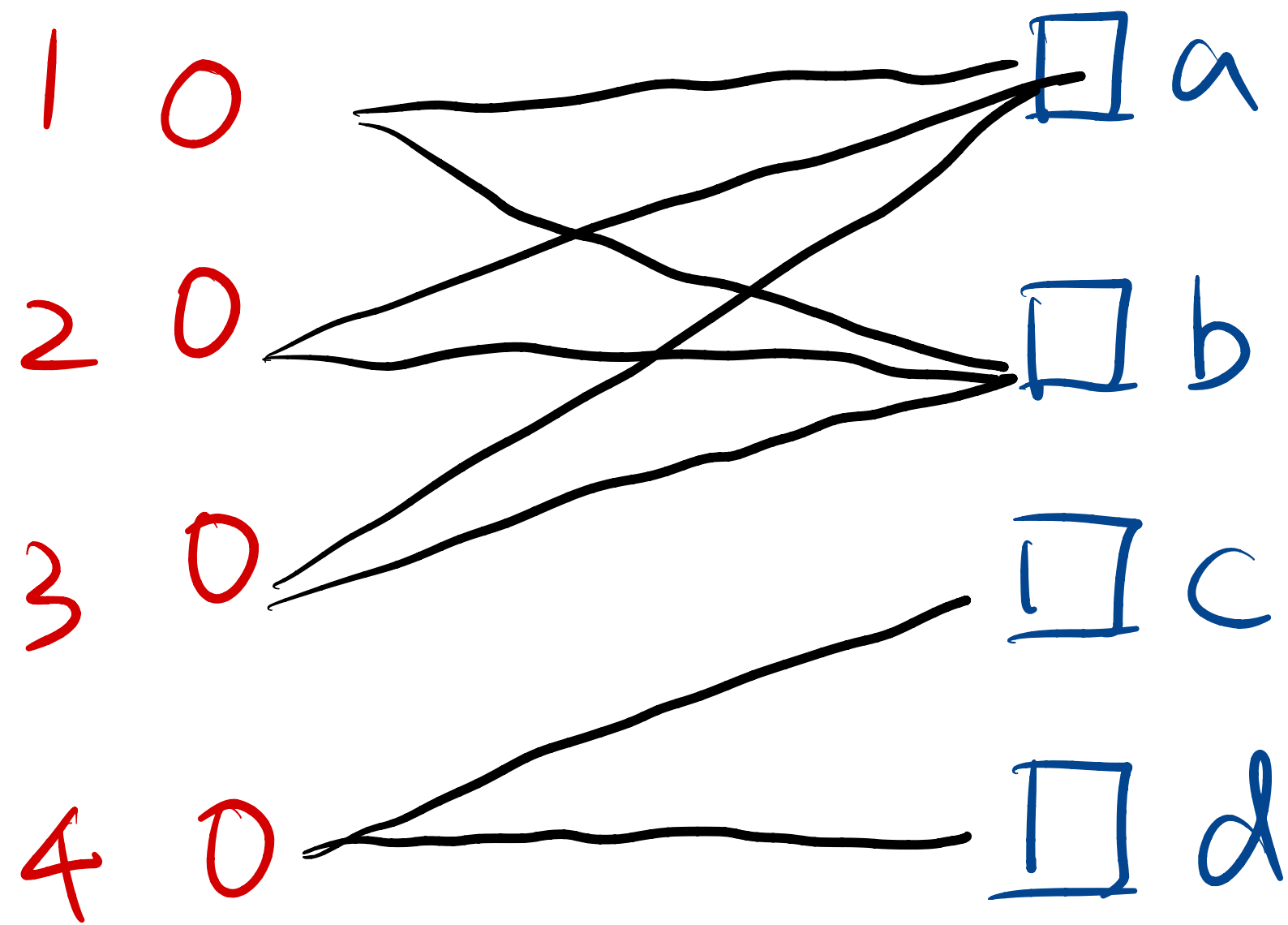
Function / Mapping on sets: $S = \{1, 2\} \quad N(S) = \{A, B\}$

We call a set S is **constricted** if S contains strictly more nodes than $N(S)$

$$S = \{2, 3\} \quad N(S) = \{A\} \quad |S| = 2 > |N(S)| = 1$$

Having a constricted set \Rightarrow Not having a P.M.

iClicker Question



$$S = \{c, d\}$$

$$N(S) = \{a, b\}$$

$$|S| = 2 > |N(S)| = 1$$

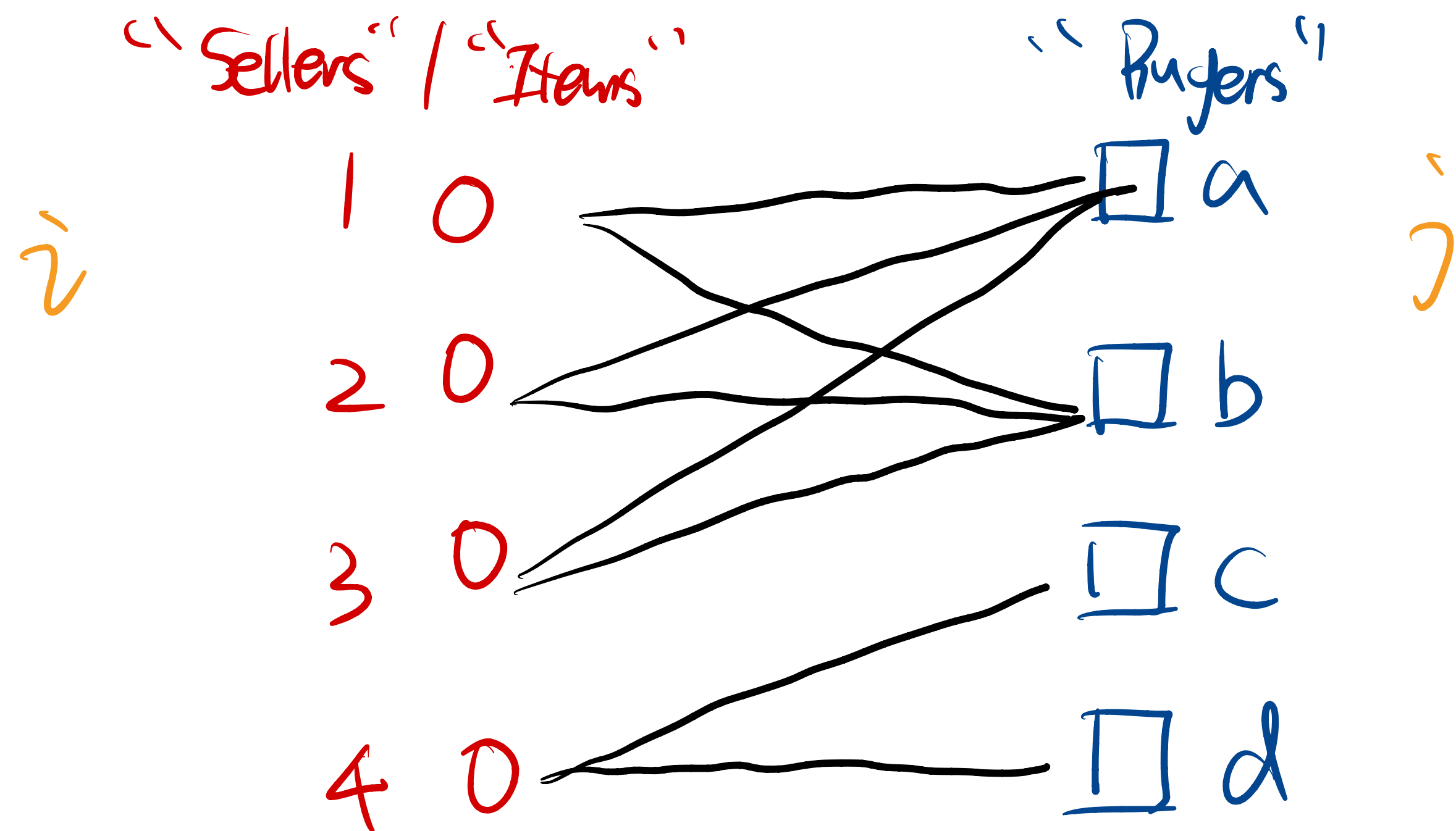
so S is a constriction set

\Rightarrow There's no P.M.

Does the graph has a perfect matching?

A Yes

B No



"Y/N"

Numerical ratings

Valuation

Each buyer can have a **valuation** for each item (seller)

✓ $V_{ij} \geq 0$

How much it worth

$i \in \{1, 2, 3, 4\}$

$j \in \{a, b, c, d\}$

Valuation

Sellers / Items

1 0

2 0

3 0

Buyers

☐ A

☐ B

☐ C

$$V_{1A}=10 \quad V_{2A}=3 \quad V_{3A}=8$$

$$V_{1B}=0 \quad V_{2B}=5 \quad V_{3B}=6$$

$$V_{1C}=12 \quad V_{2C}=2 \quad V_{3C}=1$$

10, 3, 8

0, 5, 6

12, 2, 1

What if each buyer simply maximizes their own valuation?

Buyer A: Choose item # 1

B: item # 3

C: item # 1

Valuation under some price?

"Price"



	Sellers / Items	
P_1	1	0
P_2	2	0
P_3	3	0

Buyers
☐ A
☐ B
☐ C

$$\begin{array}{lll} V_{1A}=10 & V_{2A}=3 & V_{3A}=8 \\ \hline V_{1B}=0 & V_{2B}=5 & V_{3B}=6 \\ \hline V_{1C}=12 & V_{2C}=2 & V_{3C}=1 \end{array}$$

each buyer simply maximizes their own valuation?

What if each buyer simply maximizes their own ~~valuation~~? Payoff = Valuation - price

$$(i, j) \quad V_{ij} - p_i$$

Ⓐ

Buyer A { Used to if no price
 $\max \{ V_{1A}, V_{2A}, V_{3A} \}$

Now: price

$$y = \max \{ V_{1A} - p_1, V_{2A} - p_2, V_{3A} - p_3 \}$$

if $y < 0$, buy nothing

if $y \geq 0$ buy item(s) that let me achieve y

Valuation under some price?

Price		Sellers / Items	Buyers
$P_1 = 9$	P_1	1 0	<input type="checkbox"/> A
$P_2 = 1$	P_2	2 0	<input type="checkbox"/> B
$P_3 = 2$	P_3	3 0	<input type="checkbox"/> C

Valuation

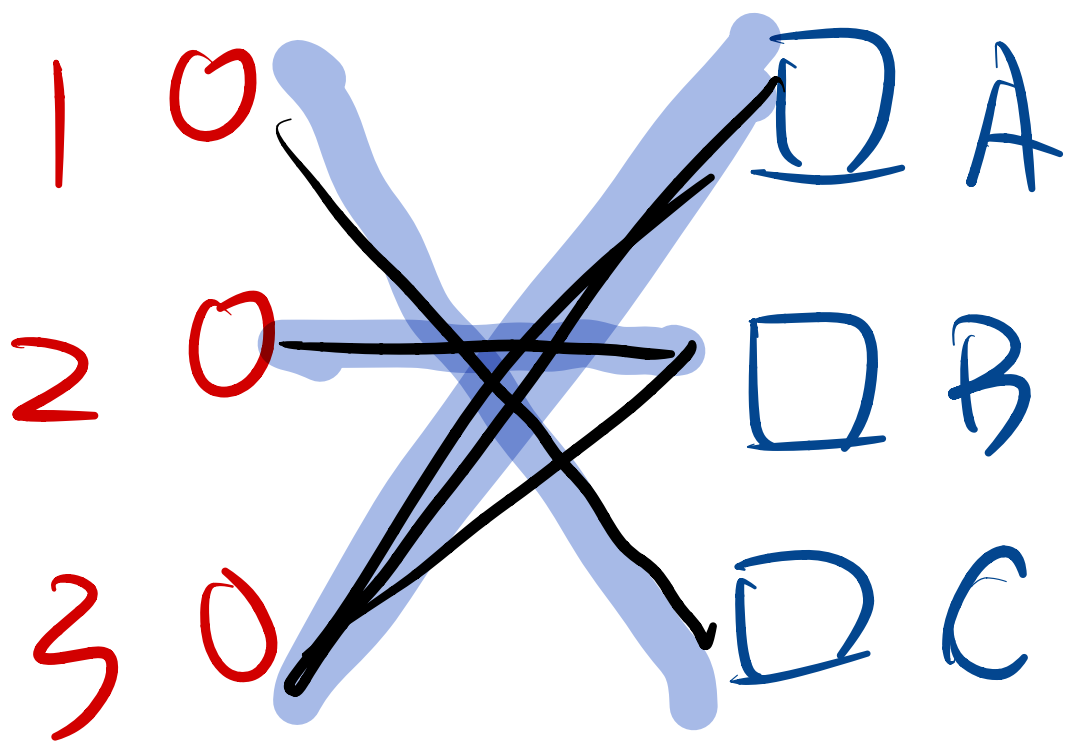
$V_{1A} = 10$	$V_{2A} = 3$	$V_{3A} = 8$
$V_{1B} = 0$	$V_{2B} = 5$	$V_{3B} = 6$
$V_{1C} = 12$	$V_{2C} = 2$	$V_{3C} = 1$

Payoff = Val - Pri

10 - 9	3 - 1	8 - 2
0 - 9	5 - 1	6 - 2
12 - 9	2 - 1	1 - 2

each buyer simply maximizes their own valuation?

- Buyer A: (3)
- Buyer B: (2) (3)
- Buyer C: (1)



⇓

1	2	6
-9	4	4
3	1	-1

“Preferred sellers”

“Preferred seller graph”

Market-clearing Price (MCP)

For a set of prices, we define the **preferred-seller graph** on buyers and sellers by simply constructing an edge between each buyer and her preferred seller or sellers.

A set of prices is **market clearing** if the resulting preferred-seller graph has a perfect matching.

Price

$P_1 = 9$

$P_2 = 1$

$P_3 = 2$

P_1

P_2

P_3

1

2

3

0

0

0

Sellers / Items

\square A

\square B

\square C

\circ

\circ

\circ

Buyers

Valuation

$V_{1A} = 10$

$V_{2A} = 3$

$V_{3A} = 8$

$V_{1B} = 0$

$V_{2B} = 5$

$V_{3B} = 6$

$V_{1C} = 12$

$V_{2C} = 2$

$V_{3C} = 1$

Buyer A: (3)

Buyer B: (2) (3)

Buyer C: (1)

“Preferred sellers”

Price (9, 1, 2) is Market clearing

Price (0, 0, 0)

\circ

\circ

\circ

\circ

\circ

\circ

\square A

\square B

\square C

Market-clearing Price (MCP)

Two important properties:

1. MCP always exists Friday
2. MCP (and the resulting P.M on preferred seller graph) is “socially optimal”

“Prisoner's dilemma” Each person max their own payoff. \rightarrow total payoff is NOT max

MCP + P.M

total payoff is maximized

1	2	6
-9	4	4
3	1	-1

Total payoff. 13

6

Market-clearing Price (MCP)

Two important properties:

- 1. MCP always exists
- 2. **MCP (and the resulting P.M on preferred seller graph) is “socially optimal”**

<div><div>1</div><div>-9</div><div>3</div></div>	<div>2</div> <div><div>4</div><div>1</div></div>	<div><div>6</div><div>4</div><div>-1</div></div>
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Total payoff. 13
6

Buyer A

≥

✓ True for Buyer B

⇒ Add up:

≥

Market-clearing Price (MCP)

Two important properties:

- 1. MCP always exists
- 2. **MCP (and the resulting P.M on preferred seller graph) is “socially optimal”**

Payoff

1	2	6
-9	4	4
3	1	-1

Valuation

10	3.8
0	5
12	2

Max total payoff ✓
Max total valuation ?

Total Payoff	=	Total Valuation
13		25
6		18

Total Price
$P_1 + P_2 + \dots$

12 → “Constant”
Give a set of prices

$$P_1 + P_2 + P_3 = 12$$