

Team notebook

CodeRats al fallo

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1 Codes

1.1 Algebra

1.1.1 FFT

```
// Multiplicacion de polinomios en  $O(n \log n)$ 
```

```
const double PI = acos(-1.0);
```

```
namespace fft {
    struct pt {
        double r, i;
        pt(double r = 0.0, double i = 0.0) : r(r), i(i) {}
        pt operator + (const pt &b) { return pt(r+b.r, i+b.i); }
        pt operator - (const pt &b) { return pt(r-b.r, i-b.i); }
        pt operator * (const pt &b) { return pt(r*b.r - i*b.i,
            r*b.i + i*b.r); }
    };
    vector<int> rev;
```

```
void fft(vector<pt> &y, int on) {
    int n = y.size();
    for (int i = 1; i < n; i++)
        if (i < rev[i]) swap(y[i], y[rev[i]]);
    for (int m = 2; m <= n; m <= 1) {
        double ang = -on * 2 * PI / m;
        pt wm(cos(ang), sin(ang));
        for (int k = 0; k < n; k += m) {
            pt w(1, 0);
```

```
                for (int j = 0; j < m / 2; j++) {
                    pt u = y[k + j];
                    pt t = w * y[k + j + m / 2];
                    y[k + j] = u + t;
                    y[k + j + m / 2] = u - t;
                    w = w * wm;
                }
            }
        }
    if (on == -1) for (int i = 0; i < n; i++) y[i].r /= n;
}
```

```
vector<ll> mul(vector<ll> &a, vector<ll> &b) {
    int n = 1, t = 0, la = a.size(), lb = b.size();
    for (; n <= (la+lb+1); n <= 1, t++); t = 1<<(t-1);
    vector<pt> x1(n), x2(n);
    rev.assign(n, 0);
    for (int i = 0; i < n; i++) rev[i] = rev[i >> 1] >> 1 |
        (i & 1 ? t : 0);
    for (int i = 0; i < la; i++) x1[i] = pt(a[i], 0);
    for (int i = 0; i < lb; i++) x2[i] = pt(b[i], 0);
    fft(x1, 1); fft(x2, 1);
    for (int i = 0; i < n; i++) x1[i] = x1[i] * x2[i];
    fft(x1, -1);
    vector<ll> ans(n);
    for (int i = 0; i < n; i++) ans[i] = x1[i].r + 0.5;
    return ans;
}
}
```

1.1.2 Linear Recurrence

```
/*
```

Calcula el n-esimo termino de una recurrencia lineal (que depende de los k terminos anteriores).

* Llamar init(k) en el main una unica vez si no es necesario inicializar las matrices multiples veces.

Este ejemplo calcula el fibonacci de n como la suma de los k terminos anteriores de la secuencia (En la secuencia comun k es 2).

Agregar Matrix Multiplication con un constructor vacio.

```

*/
matrix F, T;

void init(int k) {
    F = {k, 1}; // primeros k terminos
    F[k-1][0] = 1;
    T = {k, k}; // fila k-1 = coeficientes: [c_k, c_{k-1}, ..., c_1]
    for (int i = 0; i < k-1; i++) T[i][i+1] = 1;
    for (int i = 0; i < k; i++) T[k-1][i] = 1;
}

/// O(k^3 log(n))
int fib(ll n, int k = 2) {
    init(k);
    matrix ans = pow(T, n+k-1) * F;
    return ans[0][0];
}

```

1.1.3 Matrix Multiplication

// Estructura para realizar operaciones de multiplicacion y exponenciacion modular sobre matrices.

```
const int mod = 1e9+7;
```

```
struct matrix {
```

```

vector<vector<int>>> v;
int n, m;

matrix(int n, int m, bool o = false) : n(n), m(m), v(n,
    vector<int>(m)) {
    if (o) while (n--) v[n][n] = 1;
}

matrix operator * (const matrix &o) {
    matrix ans(n, o.m);
    for (int i = 0; i < n; i++)
        for (int k = 0; k < m; k++) if (v[i][k])
            for (int j = 0; j < o.m; j++)
                ans[i][j] = (1ll*v[i][k]*o.v[k][j] +
                    ans[i][j]) % mod;
    return ans;
}

vector<int>& operator[] (int i) { return v[i]; }
};

matrix pow(matrix b, ll e) {
    matrix ans(b.n, b.m, true);
    while (e) {
        if (e&1) ans = ans*b;
        b = b*b;
        e /= 2;
    }
    return ans;
}

```

1.2 Data Structures

1.2.1 Fenwick Tree

```
int lso(int n) {return (n & (-n));}

// las consultas estn indexadas en 1
struct FenwickTree{
    vector<int> ft;
    FenwickTree(int m) {ft.assign(m + 1, 0);};
    int rsq(int j) {
        int sum = 0;
        for(; j; j -= lso(j)) sum += ft[j];
        return sum;
    }
    void upd(int i, int v) {
        for(; i < ft.size(); i += lso(i)) ft[i] += v;
    }
};
```

1.2.2 Ordered Set

```
/*
Estructura de datos basada en politicas. Funciona como un set<>
pero es internamente indexado, cuenta con dos funciones
adicionales.
.find_by_order(k) -> Retorna un iterador al k-esimo elemento, si
k >= size() retona .end()
.order_of_key(x) -> Retorna cuantos elementos hay menores (<)
que x
*/

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
```

```
typedef tree<int, null_type, less<int>, rb_tree_tag,
tree_order_statistics_node_update> ordered_set;
```

1.2.3 Segment Tree(Lazy)

```
const int MOD = 1e9 + 7;
// modificar los struct
// operadores * y +
// en node + = merge
// en operation * = merge
// node * operation = como afecta la operacion al nodo
struct operation{
    bool swapped;
    operation() : swapped(false) {}
    operation(bool s) : swapped(s) {}

    operation operator * (const operation &rhs) const {
        operation res;
        res.swapped = (this -> swapped) ^ rhs.swapped;
        return res;
    }

    bool id() {
        return !swapped;
    }

    void clear() {
        swapped = false;
    }
};

struct node {
```

```

int imx, imn, mx, mn;
node() { imx = imn = mx = mn = 0; }
node(int imx, int imn, int mx, int mn) : imx(imx), imn(imn),
mx(mx), mn(mn) { }

node operator + (const node &rhs) const {
    node res;
    res.mx = max(this -> mx, rhs.mx);
    res.mn = min(this -> mn, rhs.mn);
    if(this -> mx >= rhs.mx) res.imx = (this -> imx);
    else res.imx = rhs.imx;
    if(this -> mn <= rhs.mn) res.imn = (this -> imn);
    else res.imn = rhs.imn;
    return res;
}

node operator * (const operation &rhs) const {
    node res;
    if(rhs.swapped) {
        res.imx = imn;
        res.imn = imx;
        res.mx = MOD - mn;
        res.mn = MOD - mx;
    } else {
        res.imx = imx;
        res.imn = imn;
        res.mx = mx;
        res.mn = mn;
    }
    return res;
}

};

const int N = 4e6 + 20;
node tree[N * 4];

```

```

operation lazy[N * 4];
pii arr[N];

void build(int v, int tl, int tr) {
    lazy[v].clear();
    if(tl == tr) {
        tree[v] = node(arr[tl].second, arr[tl].second,
            arr[tl].first, arr[tl].first);
        return;
    }
    int tm = (tl + tr) / 2;
    build(v * 2, tl, tm);
    build(v * 2 + 1, tm + 1, tr);
    tree[v] = tree[v * 2] + tree[v * 2 + 1];
}

void push(int v) {
    if(lazy[v].id()) return;
    tree[v * 2] = tree[v * 2] * lazy[v];
    tree[v * 2 + 1] = tree[v * 2 + 1] * lazy[v];
    lazy[v * 2] = lazy[v] * lazy[v * 2];
    lazy[v * 2 + 1] = lazy[v] * lazy[v * 2 + 1];
    lazy[v].clear();
}

void update(int v, int tl, int tr, int l, int r, operation &op) {
    if(l > r) return;
    if(l == tl && r == tr) {
        tree[v] = tree[v] * op;
        lazy[v] = op * lazy[v];
        return;
    }
    push(v);
    int tm = (tl + tr) / 2;
    update(v * 2, tl, tm, l, min(r, tm), op);
}

```

```

        update(v * 2 + 1, tm + 1, tr, max(l, tm + 1), r, op);
        tree[v] = tree[v * 2] + tree[v * 2 + 1];
    }

node query(int v, int tl, int tr, int l, int r) {
    if(l > r) return node();
    if(l == tl && r == tr) return tree[v];
    push(v);
    int tm = (tl + tr) / 2;
    return query(v * 2, tl, tm, l, min(r, tm))
        + query(v * 2 + 1, tm + 1, tr, max(tm + 1, l), r);
}

```

1.2.4 Segment Tree

```

const int N = 100000 + 5;
const long long inf = 1e18 + 10;

```

```

struct node {
    long long sum;
    long long maxi;
    node(){
        sum = 0;
        maxi = -inf;
    }

    node(long long x) {
        sum = maxi = x;
    }

    node operator + (const node &rhs) const {
        node q;
        q.sum = sum + rhs.sum;
        q.maxi = max(maxi, rhs.maxi);
    }
}

```

```

        return q;
    }
};

int n;
int q;
node NIL;
long long a[N];
node st[4 * N];

void build(int pos = 1, int l = 1, int r = n) {
    if(l == r) {
        st[pos] = node(a[l]);
        return;
    }
    int mi = (l + r) / 2;
    build(2 * pos, l, mi);
    build(2 * pos + 1, mi + 1, r);
    st[pos] = st[2 * pos] + st[2 * pos + 1];
}

void update(int x, int y, int pos = 1, int l = 1, int r = n) {
    if(st[pos].maxi <= 1) return; // Funcion Potencial
    sqrt(1) = 1
    if(y < l or r < x) return;
    if(l == r) {
        // to change
        st[pos].sum = sqrt(st[pos].sum);
        st[pos].maxi = st[pos].sum;
        return;
    }
    int mi = (l + r) / 2;
    update(x, y, 2 * pos, l, mi);
    update(x, y, 2 * pos + 1, mi + 1, r);
    st[pos] = st[2 * pos] + st[2 * pos + 1];
}

```

```

}

node query(int x, int y, int pos = 1, int l = 1, int r = n) {
    if(y < l or r < x) return NIL;
    if(x <= l and r <= y) return st[pos];
    int mi = (l + r) / 2;
    return query(x, y, 2 * pos, l, mi) + query(x, y, 2 * pos
        + 1, mi + 1, r);
}

int main() {
    build();
    update(1, r);
    query(1, r).sum;
    query(1, r).maxi;
}

```

1.2.5 Sparse Table

Estructura de datos que permite procesar consultas por rangos.

```

const int MX = 1e5+5;
const int LG = log2(MX)+1;
int spt[LG][MX];
int arr[MX];
int n;

void build() {
    for (int i = 0; i < n; i++) spt[0][i] = arr[i];
    for (int j = 0; j < LG-1; j++)
        for (int i = 0; i+(1<<(j+1)) <= n; i++)
            spt[j+1][i] = min(spt[j][i], spt[j][i+(1<<j)]);
}

```

```

int rmq(int i, int j) {
    int k = 31-__builtin_clz(j-i+1);
    return min(spt[k][i], spt[k][j-(1<<k)+1]);
}

```

1.3 Flows

1.3.1 Blossom

// Halla el mximo match en un grafo general $O(E * v^2)$

```

struct network {
    struct struct_edge {
        int v; struct_edge * n;
    };

    typedef struct_edge* edge;

    int n;
    struct_edge pool[MAXE]; ///2*n*n;
    edge top;
    vector<edge> adj;
    queue<int> q;
    vector<int> f, base, inq, inb, inp, match;
    vector<vector<int>> ed;

    network(int n) : n(n), match(n, -1), adj(n), top(pool),
        f(n), base(n),
        inq(n), inb(n), inp(n), ed(n, vector<int>(n))
        {}

    void add_edge(int u, int v) {
        if(ed[u][v]) return;
        ed[u][v] = 1;
    }
}

```

```

    top->v = v, top->n = adj[u], adj[u] = top++;
    top->v = u, top->n = adj[v], adj[v] = top++;
}

```

```

int get_lca(int root, int u, int v) {
    fill(inp.begin(), inp.end(), 0);
    while(1) {
        inp[u = base[u]] = 1;
        if(u == root) break;
        u = f[ match[u] ];
    }
    while(1) {
        if(inp[v = base[v]]) return v;
        else v = f[ match[v] ];
    }
}

```

```

void mark(int lca, int u) {
    while(base[u] != lca) {
        int v = match[u];
        inb[ base[u] ] = 1;
        inb[ base[v] ] = 1;
        u = f[v];
        if(base[u] != lca) f[u] = v;
    }
}

```

```

void blossom_contraction(int s, int u, int v) {
    int lca = get_lca(s, u, v);
    fill(inb.begin(), inb.end(), 0);
    mark(lca, u); mark(lca, v);
    if(base[u] != lca) f[u] = v;
    if(base[v] != lca) f[v] = u;
    for(int u = 0; u < n; u++){
        if(inb[base[u]]) {

```

```

            base[u] = lca;
            if(!inq[u]) {
                inq[u] = 1;
                q.push(u);
            }
        }
    }

    int bfs(int s) {
        fill(inq.begin(), inq.end(), 0);
        fill(f.begin(), f.end(), -1);
        for(int i = 0; i < n; i++) base[i] = i;
        q = queue<int>();
        q.push(s);
        inq[s] = 1;
        while(q.size()) {
            int u = q.front(); q.pop();
            for(edge e = adj[u]; e; e = e->n) {
                int v = e->v;
                if(base[u] != base[v] && match[u] != v) {
                    if((v == s) || (match[v] != -1 && f[match[v]]
                        != -1)){
                        blossom_contraction(s, u, v);
                    }else if(f[v] == -1) {
                        f[v] = u;
                        if(match[v] == -1) return v;
                        else if(!inq[match[v]]) {
                            inq[match[v]] = 1;
                            q.push(match[v]);
                        }
                    }
                }
            }
        }
    }
}

```



```

        return -1;
    }

    int doit(int u) {
        if(u == -1) return 0;
        int v = f[u];
        doit(match[v]);
        match[v] = u; match[u] = v;
        return u != -1;
    }

    /// (i < net.match[i]) => means match
    int maximum_matching() {
        int ans = 0;
        for(int u = 0; u < n; u++)
            ans += (match[u] == -1) && doit(bfs(u));
        return ans;
    }
};

```

1.3.2 Dinic

Halla el flujo mximo $O(E * V^2)$

```

struct edge { int v, cap, inv, flow; };

struct network {
    int n, s, t;
    vector<int> lvl;
    vector<vector<edge>> g;

    network(int n) : n(n), lvl(n), g(n) {}

    void add_edge(int u, int v, int c) {

```

```

        g[u].push_back({v, c, (int)g[v].size(), 0});
        g[v].push_back({u, 0, (int)g[u].size()-1, c});
    }

    bool bfs() {
        fill(lvl.begin(), lvl.end(), -1);
        queue<int> q;
        lvl[s] = 0;
        for(q.push(s); q.size(); q.pop()) {
            int u = q.front();
            for(auto &e : g[u]) {
                if(e.cap > 0 && lvl[e.v] == -1) {
                    lvl[e.v] = lvl[u]+1;
                    q.push(e.v);
                }
            }
        }
        return lvl[t] != -1;
    }

    int dfs(int u, int nf) {
        if(u == t) return nf;
        int res = 0;
        for(auto &e : g[u]) {
            if(e.cap > 0 && lvl[e.v] == lvl[u]+1) {
                int tf = dfs(e.v, min(nf, e.cap));
                res += tf; nf -= tf; e.cap -= tf;
                g[e.v][e.inv].cap += tf;
                g[e.v][e.inv].flow -= tf;
                e.flow += tf;
                if(nf == 0) return res;
            }
        }
        if(!res) lvl[u] = -1;
        return res;
    }

```

```

    }

    int max_flow(int so, int si, int res = 0) {
        s = so; t = si;
        while(bfs()) res += dfs(s, INT_MAX);
        return res;
    }
};

```

1.3.3 Hungarian

```

typedef ll T;
const T inf = 1e18;

struct hung {
    int n, m;
    vector<T> u, v; vector<int> p, way;
    vector<vector<T>> g;

    hung(int n, int m):
        n(n), m(m), g(n+1, vector<T>(m+1, inf-1)),
        u(n+1), v(m+1), p(m+1), way(m+1) {}

    void set(int u, int v, T w) { g[u+1][v+1] = w; }

    T assign() {
        for (int i = 1; i <= n; ++i) {
            int j0 = 0; p[0] = i;
            vector<T> minv(m+1, inf);
            vector<char> used(m+1, false);
            do {
                used[j0] = true;
                int i0 = p[j0], j1; T delta = inf;
                for (int j = 1; j <= m; ++j) if (!used[j]) {

```

```

                    T cur = g[i0][j] - u[i0] - v[j];
                    if (cur < minv[j]) minv[j] = cur, way[j] = j0;
                    if (minv[j] < delta) delta = minv[j], j1 = j;
                }
                for (int j = 0; j <= m; ++j)
                    if (used[j]) u[p[j]] += delta, v[j] -= delta;
                    else minv[j] -= delta;
                j0 = j1;
            } while (p[j0]);
            do {
                int j1 = way[j0]; p[j0] = p[j1]; j0 = j1;
            } while (j0);
        }
        return -v[0];
    }
};

```

1.3.4 Maximum Bipartite Matching

// Halla el maximo match en un grafo bipartito $O(|E|*|V|)$

```

struct mbm {
    int l, r;
    vector<vector<int>> g;
    vector<int> match, vis;

    mbm(int l, int r) : l(l), r(r), g(l) {}

    void add_edge(int l, int r) {
        g[l].push_back(r);
    }

    bool dfs(int u) {
        for (auto &v : g[u]) {

```

```

        if (vis[v]++) continue;
        if (match[v] == -1 || dfs(match[v])) {
            match[v] = u;
            return true;
        }
    }
    return false;
}

int max_matching() {
    int ans = 0;
    match.assign(r, -1);
    for (int u = 0; u < l; ++u) {
        vis.assign(r, 0);
        ans += dfs(u);
    }
    return ans;
}

};

// Hopcroft Karp: O(E * sqrt(V))

const int INF = INT_MAX;

struct mbm {
    vector<vector<int>> g;
    vector<int> d, match;
    int nil, l, r;
    /// u -> 0 to l, v -> 0 to r
    mbm(int l, int r) : l(l), r(r), nil(l+r), g(l+r),
                      d(l+r, INF), match(l+r, l+r) {}

    void add_edge(int a, int b) {
        g[a].push_back(l+b);
        g[l+b].push_back(a);
    }
};

```

```

    }

bool bfs() {
    queue<int> q;
    for(int u = 0; u < l; u++) {
        if(match[u] == nil) {
            d[u] = 0;
            q.push(u);
        } else {
            d[u] = INF;
        }
    }
    d[nil] = INF;
    while(q.size()) {
        int u = q.front(); q.pop();
        if(u == nil) continue;
        for(auto v : g[u]) {
            if(d[ match[v] ] == INF) {
                d[ match[v] ] = d[u]+1;
                q.push(match[v]);
            }
        }
    }
    return d[nil] != INF;
}

bool dfs(int u) {
    if(u == nil) return true;
    for(int v : g[u]) {
        if(d[ match[v] ] == d[u]+1 && dfs(match[v])) {
            match[v] = u; match[u] = v;
            return true;
        }
    }
    d[u] = INF;
}

```

```

        return false;
    }

    int max_matching() {
        int ans = 0;
        while(bfs()) {
            for(int u = 0; u < 1; u++) {
                ans += (match[u] == nil && dfs(u));
            }
        }
        return ans;
    }
};

```

1.3.5 MinCost MaxFlow

Dado un grafo, halla el flujo maximo y el costo minimo entre el source s y el sink t.

```

struct edge {
    int u, v, cap, flow, cost;
    int rem() { return cap - flow; }
};

const int inf = 1e9;
const int MX = 405; //Cantidad maxima TOTAL de nodos
vector<int> g[MX]; //Lista de adyacencia
vector<edge> e; //Lista de aristas
vector<bool> in_queue; //Marca los nodos que estan en cola
vector<int> pre, dist, cap; //Almacena el nodo anterior, la
    distancia y el flujo de cada nodo
int mxflow, mncost; //Flujo maximo y costo minimo
int N; //Cantidad TOTAL de nodos

```

```

void add_edge(int u, int v, int cap, int cost) {
    g[u].push_back(e.size());
    e.push_back({u, v, cap, 0, cost});
    g[v].push_back(e.size());
    e.push_back({v, u, 0, 0, -cost});
}

void flow(int s, int t) {
    mxflow = mncost = 0;
    in_queue.assign(N, false);
    while (true) {
        dist.assign(N, inf); dist[s] = 0;
        cap.assign(N, 0); cap[s] = inf;
        pre.assign(N, -1); pre[s] = 0;
        queue<int> q; q.push(s);
        in_queue[s] = true;

        while (q.size()) {
            int u = q.front(); q.pop();
            in_queue[u] = false;
            for (int &id : g[u]) {
                edge &ed = e[id];
                int v = ed.v;
                if (ed.rem() && dist[v] > dist[u]+ed.cost) {
                    dist[v] = dist[u]+ed.cost;
                    cap[v] = min(cap[u], ed.rem());
                    pre[v] = id;
                    if (!in_queue[v]) {
                        q.push(v);
                        in_queue[v] = true;
                    }
                }
            }
        }
    }
    if (pre[t] == -1) break;
}

```

```

        mxflow += cap[t];
        mncost += cap[t] * dist[t];
        for (int v = t; v != s; v = e[pre[v]].u) {
            e[pre[v]].flow += cap[t];
            e[pre[v]^1].flow -= cap[t];
        }
    }
}

void init() {
    e.clear();
    for (int i = 0; i <= N; i++) {
        g[i].clear();
    }
}

// O(V * E * 2 * log(E))
template <class type>
struct mcmf {
    struct edge { int u, v, cap, flow; type cost; };

    int n;
    vector<edge> ed;
    vector<vector<int>> g;
    vector<int> p;
    vector<type> d, phi;

    mcmf(int n) : n(n), g(n), p(n), d(n), phi(n) {}

    void add_edge(int u, int v, int cap, type cost) {
        g[u].push_back(ed.size());
        ed.push_back({u, v, cap, 0, cost});
        g[v].push_back(ed.size());
        ed.push_back({v, u, 0, 0, -cost});
    }
}

```

```

bool dijkstra(int s, int t) {
    fill(d.begin(), d.end(), INF);
    fill(p.begin(), p.end(), -1);
    set<pair<type, int>> q;
    d[s] = 0;
    for(q.insert({d[s], s}); q.size();) {
        int u = (*q.begin()).second; q.erase(q.begin());
        for(auto v : g[u]) {
            auto &e = ed[v];
            type nd = d[e.u] + e.cost + phi[e.u] - phi[e.v];
            if(0 < (e.cap - e.flow) && nd < d[e.v]) {
                q.erase({d[e.v], e.v});
                d[e.v] = nd; p[e.v] = v;
                q.insert({d[e.v], e.v});
            }
        }
    }
    for(int i = 0; i < n; i++) phi[i] = min(INF, phi[i] + d[i]);
    return d[t] != INF;
}

pair<int, type> max_flow(int s, int t) {
    type mc = 0;
    int mf = 0;
    fill(phi.begin(), phi.end(), 0);
    while(dijkstra(s, t)) {
        int flow = INF;
        for(int v = p[t]; v != -1; v = p[ ed[v].u ])
            flow = min(flow, ed[v].cap - ed[v].flow);
        for(int v = p[t]; v != -1; v = p[ ed[v].u ]) {
            edge &e1 = ed[v];
            edge &e2 = ed[v^1];
            mc += e1.cost * flow;
            e1.flow += flow;
            e2.flow -= flow;
        }
    }
}

```

```

    }
    mf += flow;
}
return {mf, mc};
}
};

```

1.3.6 Weighted Matching

Halla el mximo match con pesos $O(V^3)$

```

typedef int type;
struct matching_weighted {
    int l, r;
    vector<vector<type>> c;
    matching_weighted(int l, int r) : l(l), r(r), c(l,
        vector<type>(r)) {
        assert(l <= r);
    }

    void add_edge(int a, int b, type cost) { c[a][b] = cost; }

    type matching() {
        vector<type> v(r), d(r); // v: potential
        vector<int> ml(l, -1), mr(r, -1); // matching pairs
        vector<int> idx(r), prev(r);
        iota(idx.begin(), idx.end(), 0);
        auto residue = [&](int i, int j) { return c[i][j]-v[j]; };
        for(int f = 0; f < l; ++f) {
            for(int j = 0; j < r; ++j) {
                d[j] = residue(f, j);
                prev[j] = f;
            }
            type w;

```

```

            int j, l;
            for (int s = 0, t = 0;;) {
                if(s == t) {
                    l = s;
                    w = d[ idx[t++] ];
                    for(int k = t; k < r; ++k) {
                        j = idx[k];
                        type h = d[j];
                        if (h <= w) {
                            if (h < w) t = s, w = h;
                            idx[k] = idx[t];
                            idx[t++] = j;
                        }
                    }
                    for (int k = s; k < t; ++k) {
                        j = idx[k];
                        if (mr[j] < 0) goto aug;
                    }
                }
                int q = idx[s++], i = mr[q];
                for (int k = t; k < r; ++k) {
                    j = idx[k];
                    type h = residue(i, j) - residue(i, q) + w;
                    if (h < d[j]) {
                        d[j] = h;
                        prev[j] = i;
                        if(h == w) {
                            if(mr[j] < 0) goto aug;
                            idx[k] = idx[t];
                            idx[t++] = j;
                        }
                    }
                }
            }
        }
        aug:

```

```

    for (int k = 0; k < l; ++k)
        v[ idx[k] ] += d[ idx[k] ] - w;
    int i;
    do {
        mr[j] = i = prev[j];
        swap(j, ml[i]);
    } while (i != f);
}
type opt = 0;
for (int i = 0; i < l; ++i)
    opt += c[i][ml[i]]; // (i, ml[i]) is a solution
return opt;
}
};

```

1.4 Graphs

1.4.1 2SAT

```

struct SATSolver {
    // Assumes that nodes are 0-indexed
    int n;
    int m;
    vector<bool> vis;
    vector<int> comp;
    vector<int> order;
    vector<int> component;
    vector<vector<int>> G;
    vector<vector<int>> Gt;

    SATSolver(int n, int m) : n(n), m(m) {
        // X_i = 2i
        // ~X_i = 2i + 1
        comp.resize(2 * n);
    }
}

```

```

        vis.resize(2 * n, false);
        G.resize(2 * n, vector<int>());
        Gt.resize(2 * n, vector<int>());
    }

    void add_edge(int u, int v) {
        // u OR v
        G[u ^ 1].emplace_back(v);
        G[v ^ 1].emplace_back(u);
        Gt[v].emplace_back(u ^ 1);
        Gt[u].emplace_back(v ^ 1);
    }

    void DFS1(int u) {
        vis[u] = true;
        for(int v : G[u]) {
            if(vis[v]) continue;
            DFS1(v);
        }
        order.emplace_back(u);
    }

    void DFS2(int u) {
        vis[u] = true;
        for(int v : Gt[u]) {
            if(vis[v]) continue;
            DFS2(v);
        }
        component.emplace_back(u);
    }

    void get_scc() {
        for(int i = 0; i < 2 * n; i++) {
            if(vis[i]) continue;
            DFS1(i);
        }
    }
}

```

```

    }
    reverse(order.begin(), order.end());
    fill(vis.begin(), vis.end(), false);
    int component_id = 0;
    for(int u : order) {
        if(vis[u]) continue;
        component.clear();
        DFS2(u);
        for(int x : component) comp[x] =
            component_id;
        component_id += 1;
    }

    vector<int> solve() {
        vector<int> res(n);
        get_scc();
        for(int i = 0; i < n; i++) {
            int val = 2 * i;
            if(comp[val] == comp[val ^ 1]) return
                vector<int>();
            if(comp[val] < comp[val ^ 1]) res[i] = 0;
            else res[i] = 1;
        }
        return res;
    }
};

```

1.4.2 Articulation Points

```

int n;
int m;
int timer;
int low[N];

```

```

int tin[N];
bool vis[N];
bool cut_point[N];
vector<int> G[N];

void DFS(int u, int p = -1) {
    tin[u] = low[u] = timer++;
    vis[u] = true;
    int children = 0;
    for(int v : G[u]) {
        if(v == p) continue;
        if(vis[v]) {
            // Es una backedge, aporta a low[u]
            low[u] = min(low[u], tin[v]);
        }
        else {
            // Es una tree-edge, verificar si u es cut
            // point
            DFS(v, u);
            // Minimizamos el low del padre con el del
            // hijo
            low[u] = min(low[u], low[v]);
            // Ya tenemos procesado low[v]
            if(p != -1 and low[v] >= tin[u]) {
                // u es articulacion si low[v] >=
                // tin[u] cuando u no es la raz
                cut_point[u] = true;
            }
            children++;
        }
    }
    // Si es la raiz, es articulacion si tiene 2 o ms hijos
    if(p == -1 and children > 1) cut_point[u] = true;
}

```


1.4.3 Bellman Ford

```

int n;
int m;
int D[N]; // D[u] : Minima distancia de src a u usando <= k
           aristas en la k-sima iteracin
bool vis[N]; // vis[u] : El nodo se ha vuelto alcanzable por src
vector<tuple<int, int, int>> edges;

// retorna true si no hay ciclos negativos
bool bellman_ford(int src) {
    for(int i = 0; i < n; i++) {
        D[i] = -1;
        vis[i] = false;
    }
    D[src] = 0;
    vis[src] = true;
    for(int i = 1; i < n; i++) {
        for(auto e : edges) {
            int u, v, w;
            tie(u, v, w) = e;
            if(not vis[u]) continue;
            if(not vis[v] or D[v] > D[u] + w) {
                D[v] = D[u] + w;
                vis[v] = true;
            }
        }
    }
    for(auto e : edges) {
        int u, v, w;
        tie(u, v, w) = e;
        if(not vis[u]) continue;
        if(not vis[v] or D[v] > D[u] + w) {
            return false; // Ciclo negativo alcanzable
                           por src
        }
    }
}

```

```

    }
    }
    return true;
}

```

1.4.4 Bridges

```

int n; // para guardar el numero de vertices
int m; // aristas
int k; // numero de componentes 2-conexas en el arbol final
int C[N]; // componente de cada vertice
int to[N], from[N]; // vertice de salida y llegada para cada
                    arista
int timer; // para el DFS
int low[N];
int tin[N];
bool vis[N];
bool bridge[N]; // es puente ?
vector<int> T[N]; // arbol final
vector<pair<int, int>> G[N]; // first = v, second = id de arista

void DFS(int u, int p = -1) {
    vis[u] = true;
    low[u] = tin[u] = timer++;
    for(auto edge : G[u]) {
        int v, e;
        tie(v, e) = edge;
        if(v == p) continue;
        if(vis[v]) {
            low[u] = min(low[u], tin[v]);
        }
        else {
            DFS(v, u);
            low[u] = min(low[u], low[v]);
        }
    }
}

```

```

        if(low[v] > tin[u]) {
            bridge[e] = true;
        }
    }
}

void build_tree() {
    DFS(1);
    for(int i = 1; i <= n; i++) vis[i] = false;
    k = 0;
    for(int i = 1; i <= n; i++) {
        if(vis[i]) continue;
        queue<int> Q;
        Q.emplace(i);
        vis[i] = true;
        while(!Q.empty()) {
            int u = Q.front(); Q.pop();
            C[u] = k;
            for(auto edge : G[u]) {
                int v, e;
                tie(v, e) = edge;
                if(bridge[e]) continue;
                if(vis[v]) continue;
                Q.emplace(v);
                vis[v] = true;
            }
        }
        k += 1;
    }
    for(int i = 1; i <= m; i++) {
        if(not bridge[i]) continue;
        int u = C[to[i]], v = C[from[i]];
        T[u].emplace_back(v);
        T[v].emplace_back(u);
    }
}

```

```

    }
}

```

1.4.5 Dijkstra

```

int n;
int m;
int D[N];
vector<pair<int, int>> G[N];

// las distancias mnimas se guardan en D
// indexa en 0
void Dijkstra(int src) {
    for(int i = 0; i < n; i++) D[i] = -1;
    D[src] = 0;
    priority_queue<pair<int, int>, vector<pair<int, int>>,
        greater<pair<int, int>>> Q; // Min PQ
    Q.emplace(0, src);
    while(!Q.empty()) {
        int dis, u;
        tie(dis, u) = Q.top(); Q.pop();
        if(dis != D[u]) continue; // Verificacion de que u
                                // no ha sido visitado todavia
        for(auto e : G[u]) {
            int v, w;
            tie(v, w) = e;
            if(D[v] == -1 or D[v] > D[u] + w) {
                D[v] = D[u] + w;
                Q.emplace(D[v], v);
            }
        }
    }
}

```

1.4.6 Floyd Warshall

```

const int N = 100 + 5;
const int inf = 2e9 + 10;

int n;
int m;
int d[N][N];

/*
 * inicializar el arreglo d con INF, a menos que i == j
   for(int i = 0; i < n; i++) {
       for(int j = 0; j < n; j++) {
           d[i][j] = i == j ? 0 : inf;
       }
   }
 * para escoger siempre la menor
 * arista en caso de aristas mltiples
   for(int i = 0; i < m; i++) {
       int u, v, w;
       scanf("%d %d %d", &u, &v, &w);
       d[u][v] = min(d[u][v], w);
   }
 */

bool floyd_warshall(){
    for(int k = 0; k < n; k++) {
        for(int i = 0; i < n; i++) {
            for(int j = 0; j < n; j++) {
                if(d[i][k] == inf || d[k][j] ==
                    inf) continue;

```

```

                if(d[i][j] > d[i][k] + d[k][j])
                    d[i][j] = d[i][k] + d[k][j];
            }
        }

        // Termina Floyd-Warshall
        // comprobacion de ciclos negativos
        for(int i = 0; i < n; i++) {
            if(d[i][i] < 0) {
                return false;
            }
        }
        return true;
    }
}

```

1.4.7 Heavy Light Decomposition

```

/*

Para inicializar llamar build().
Agregar Segment Tree con un constructor vacio,
actualizaciones puntuales y declarar el valor neutro de forma
global.

Para consultas sobre aristas guardar el valor de cada arista
en su nodo hijo y cambiar pos[u] por pos[u]+1 en la linea 54.
*/

typedef int T; //tipo de dato del segtree
const int MX = 1e5+5;
vector<int> g[MX];
int par[MX], dep[MX], sz[MX];
int pos[MX], top[MX], value[MX];
vector<T> arr;

```

```

int idx;

int pre(int u, int p, int d) {
    par[u] = p; dep[u] = d;
    int aux = 1;
    for (auto &v : g[u]) {
        if (v != p) {
            aux += pre(v, u, d+1);
            if (sz[v] >= sz[g[u][0]]) swap(v, g[u][0]);
        }
    }
    return sz[u] = aux;
}

void hld(int u, int p, int t) {
    arr[idx] = value[u]; //vector para inicializar el segtree
    pos[u] = idx++;
    top[u] = t < 0 ? t = u : t;
    for (auto &v : g[u]) {
        if (v != p) {
            hld(v, u, t);
            t = -1;
        }
    }
}

segtree sgt;

void build(int n, int root) {
    idx = 0;
    arr.resize(n);
    pre(root, root, 0);
    hld(root, root, -1);
    sgt = segtree(arr);
}

```

```

T query(int u, int v) {
    T ans = neutro;
    while (top[u] != top[v]) {
        if (dep[top[u]] > dep[top[v]]) swap(u, v);
        ans = min(ans, sgt.query(pos[top[v]], pos[v]));
        v = par[top[v]];
    }
    if (dep[u] > dep[v]) swap(u, v);
    ans = min(ans, sgt.query(pos[u], pos[v]));
    return ans;
}

void upd(int u, T val) {
    sgt.upd(pos[u], val);
}

```

1.4.8 Prim

// Prim clsico, retorna el MST de un grafo

```

int n;
int m;
int q[N];
bool vis[N];
int wedge[N];
vector<pair<int, int>> G[N];

int Prim(int src) {
    memset(wedge, -1, sizeof wedge);
    wedge[src] = 0;
    priority_queue<pair<int, int>, vector<pair<int, int>>,
        greater<pair<int, int>>> Q;
    Q.emplace(0, src);
}

```

```

while(!Q.empty()) {
    int we, u;
    tie(we, u) = Q.top(); Q.pop();
    vis[u] = true;
    for(auto e : G[u]) {
        int v, w;
        tie(v, w) = e;
        if(wedge[v] == -1 or wedge[v] > w) {
            wedge[v] = w;
            Q.emplace(wedge[v], v);
        }
    }
}

for(int i = 1; i <= n; i++) {
    if(not vis[i]) return -1;
}

return accumulate(wedge + 1, wedge + n + 1, 0);
}

/*
 *
 * Caso especial para un grafo completo
 * no es posible construir el grafo (memoria)
 * pero es posible un algoritmo  $O(n^2)$ 
 *
 */
int n;
int x[N];
int y[N];
bool vis[N];
int edge[N]; // Minima arista que cruza desde S hasta mi nodo

int dis(int i, int j) {

```

```

    return abs(x[i] - x[j]) + abs(y[i] - y[j]);
}

int Prim(int src) {
    vis[src] = true;
    for(int i = 0; i < n; i++) edge[i] = dis(src, i);
    for(int i = 1; i < n; i++) {
        // El nodo al que llega la arista ligera es al
        // argmin(edge[i]) pero con vis[i] = false
        int v = -1;
        for(int j = 0; j < n; j++) {
            if(vis[j]) continue;
            if(v == -1 or edge[v] > edge[j]) v = j;
        }
        vis[v] = true;
        for(int j = 0; j < n; j++) {
            if(vis[j]) continue;
            edge[j] = min(edge[j], dis(v, j));
        }
    }
    return accumulate(edge, edge + n, 0);
}

```

1.4.9 SCC

```

int n;
int m;
bool in[N];
bool vis[N];
int comp[N];
vector<int> order;
vector<int> component;
vector<int> G[2][N]; // en G[1] es el grafo transpuesto

```

```

void DFS(int id, int u) {
    vis[u] = id ^ 1;
    for(int v : G[id][u]) {
        if(vis[v] == (id ^ 1)) continue;
        DFS(id, v);
    }
    if(id == 0) order.emplace_back(u);
    else component.emplace_back(u);
}

int solve() {
    order.clear();
    component.clear();
    for(int i = 0; i < n; i++) {
        if(vis[i]) continue;
        DFS(0, i);
    }
    reverse(order.begin(), order.end());
    vector<vector<int>> res;
    for(int u : order) {
        if(not vis[u]) continue;
        component.clear();
        DFS(1, u);
        for(int x : component) comp[x] = res.size();
        res.emplace_back(component);
    }
    for(int i = 0; i < n; i++) {
        for(int v : G[0][i]) {
            if(comp[i] == comp[v]) continue;
            in[comp[v]] = true;
        }
    }
    int cnt = 0;
    for(int i = 0; i < res.size(); i++) cnt += !in[i];
    return cnt;
}

```

```

}

```

1.4.10 Topological Sort

```

vector<int> toposort(int n, int m, vector<vector<int>> &G) {
    vector<int> in_degree(n, 0);
    for(int i = 0; i < n; i++) {
        for(int v : G[i]) in_degree[v] += 1;
    }
    queue<int> Q;
    vector<int> res;
    for(int i = 0; i < n; i++) {
        if(in_degree[i] == 0) {
            Q.emplace(i);
        }
    }
    while(!Q.empty()) {
        int u = Q.front(); Q.pop();
        res.emplace_back(u);
        for(int v : G[u]) {
            in_degree[v] -= 1;
            if(in_degree[v] == 0) {
                Q.emplace(v);
            }
        }
    }
    return res.size() < n ? vector<int>() : res;
}

```

1.5 Number theory

1.5.1 Chinese Remainder Theorem

```

/*
Encuentra un x tal que para cada i : x es congruente con A_i mod
M_i
Devuelve {x, lcm}, donde x es la solucion con modulo lcm (lcm =
LCM(M_0, M_1, ...)). Dado un k : x + k*lcm es solucion
tambien.
Si la solucion no existe o la entrada no es valida devuelve {-1,
-1}
Agregar Extended Euclides.
*/

```

```

pair<int, int> crt(vector<int> A, vector<int> M) {
    int n = A.size(), ans = A[0], lcm = M[0];
    for (int i = 1; i < n; i++) {
        int d = euclid(lcm, M[i]);
        if ((A[i] - ans) % d) return {-1, -1};
        int mod = lcm / d * M[i];
        ans = (ans + x * (A[i] - ans) / d % (M[i] / d) * lcm) %
            mod;
        if (ans < 0) ans += mod;
        lcm = mod;
    }
    return {ans, lcm};
}

```

1.5.2 Lineal Sieve

```

const int N = 100000000 + 5;
vector<int> primes;
bitset<N> composite;

void lineal(int n){
    for(int i = 2; i <= n; i++) {

```

```

        if(not composite[i]) primes.emplace_back(i);
        for(int p : primes) {
            if(i * p > n) break;
            composite[i * p] = true;
            if(i % p == 0) break;
        }
    }
}

```

1.5.3 Miller Rabin

```

// El algoritmo de Miller-Rabin determina si un numero es primo
o no. Agregar Modular Exponentiation (para m ll) y Modular
Multiplication.

```

```

/// O(log^3(n))
bool test(ll n, int a) {
    if (n == a) return true;
    ll s = 0, d = n-1;
    while (d%2 == 0) s++, d /= 2;
    ll x = expmod(a, d, n);
    if (x == 1 || x+1 == n) return true;
    for (int i = 0; i < s-1; i++) {
        x = mulmod(x, x, n);
        if (x == 1) return false;
        if (x+1 == n) return true;
    }
    return false;
}

```

```

bool is_prime(ll n) {
    if (n == 1) return false;
    int ar[] = {2,3,5,7,11,13,17,19,23};
    for (auto &p : ar) if (!test(n, p)) return false;

```

```

    return true;
}

```

1.5.4 Mobius

```

/*
La funcion mu de Mobius devuelve 0 si n es divisible por algun
cuadrado (x^2).
Si n es libre de cuadrados entonces devuelve 1 o -1 si n tiene
un numero par o impar de factores primos distintos.
* Calcular Mobius para todos los numeros menores o iguales a MX
con Sieve of Eratosthenes.
*/
const int MX = 1e6;
short mu[MX+1] = {0, 1};
/// O(MX log(log(MX)))
void mobius() {
    for (int i = 1; i <= MX; i++) {
        if (!mu[i]) continue;
        for (int j = i*2; j <= MX; j += i) {
            mu[j] -= mu[i];
        }
    }
}

```

1.5.5 Modular Multiplication

```

/* Calcula (a*b) % m sin overflow cuando m es ll. */

/// O(1)
ll mulmod(ll a, ll b, ll m) {
    ll r = a*b-(ll)((long double)a*b/m+.5)*m;

```

```

    return r < 0 ? r+m : r;
}

```

1.5.6 Natural Sieve

```

const int N = 100000000 + 5;
bitset<N> composite;

void natural(int n){
    // (WARNING) Todos los pares son primos
    for(int i = 3; i * i <= n; i += 2) {
        if(not composite[i]) {
            for(int j = i * i; j <= n; j += i)
                composite[j] = true;
        }
    }
}

```

1.5.7 Pollard Rho

```

/*
La funcion Rho de Pollard calcula un divisor no trivial de n.
Agregar Modular Multiplication.
*/
ll gcd(ll a, ll b) { return a ? gcd(b%a, a) : b; }

ll rho(ll n) {
    if (!(n&1)) return 2;
    ll x = 2, y = 2, d = 1;
    ll c = rand() % n + 1;
    while (d == 1) {
        x = (mulmod(x, x, n) + c) % n;

```

```

        y = (mulmod(y, y, n) + c) % n;
        y = (mulmod(y, y, n) + c) % n;
        d = gcd(abs(x-y), n);
    }
    return d == n ? rho(n) : d;
}

* Version optimizada

ll add(ll a, ll b, ll m) { return (a += b) < m ? a : a-m; }

ll rho(ll n) {
    static ll s[MX];
    while (1) {
        ll x = rand()%n, y = x, c = rand()%n;
        ll *px = s, *py = s, v = 0, p = 1;
        while (1) {
            *py++ = y = add(mulmod(y, y, n), c, n);
            *py++ = y = add(mulmod(y, y, n), c, n);
            if ((x = *px++) == y) break;
            ll t = p;
            p = mulmod(p, abs(y-x), n);
            if (!p) return gcd(t, n);
            if (++v == 26) {
                if ((p = gcd(p, n)) > 1 && p < n) return p;
                v = 0;
            }
        }
        if (v && (p = gcd(p, n)) > 1 && p < n) return p;
    }
}

```

1.6 Strings

1.6.1 Aho Corasick

```

/*
    El trie (o prefix tree) guarda un diccionario de strings
    como un arbol enraizado.
    Aho corasick permite encontrar las ocurrencias de todos los
    strings del trie en un string s.
*/
const int alpha = 26; //cantidad de letras del lenguaje
const char L = 'a'; //primera letra del lenguaje

struct node {
    int next[alpha], end;
    //int link, exit, cnt; //para aho corasick
    int& operator[](int i) { return next[i]; }
};

vector<node> trie = {node()};

void add_str(string &s, int id = 1) {
    int u = 0;
    for (auto ch : s) {
        int c = ch-L;
        if (!trie[u][c]) {
            trie[u][c] = trie.size();
            trie.push_back(node());
        }
        u = trie[u][c];
    }
    trie[u].end = id; //con id > 0
    //trie[u].cnt++; //para aho corasick
}

```

```
// aho corasick
void build_ac() {
    queue<int> q; q.push(0);
    while (q.size()) {
        int u = q.front(); q.pop();
        for (int c = 0; c < alpha; ++c) {
            int v = trie[u][c];
            if (!v) trie[u][c] = trie[trie[u].link][c];
            else q.push(v);
            if (!u || !v) continue;
            trie[v].link = trie[trie[u].link][c];
            trie[v].exit = trie[trie[v].link].end ?
                trie[v].link : trie[trie[v].link].exit;
            trie[v].cnt += trie[trie[v].link].cnt;
        }
    }
}

vector<int> cnt; //cantidad de ocurrencias en s para cada patron

void run_ac(string &s) {
    int u = 0, sz = s.size();
    for (int i = 0; i < sz; ++i) {
        int c = s[i]-L;
        while (u && !trie[u][c]) u = trie[u].link;
        u = trie[u][c];
        int x = u;
        while (x) {
            int id = trie[x].end;
            if (id) cnt[id-1]++;
            x = trie[x].exit;
        }
    }
}
```

1.6.2 Hashing

```
/**
Convierte el string en un polinomio, en  $O(n)$ , tal que podemos
comparar substrings como valores numericos en  $O(1)$ .
Primero llamar calc_xpow() (una unica vez) con el largo maximo
de los strings dados.
*/
inline int add(int a, int b, const int &mod) { return a+b >= mod
    ? a+b-mod : a+b; }
inline int sbt(int a, int b, const int &mod) { return a-b < 0 ?
    a-b+mod : a-b; }
inline int mul(int a, int b, const int &mod) { return 1ll*a*b %
    mod; }

const int X[] = {257, 359};
const int MOD[] = {(int)1e9+7, (int)1e9+9};
vector<int> xpow[2];

struct hashing {
    vector<int> h[2];

    hashing(string &s) {
        int n = s.size();
        for (int j = 0; j < 2; ++j) {
            h[j].resize(n+1);
            for (int i = 1; i <= n; ++i) {
                h[j][i] = add(mul(h[j][i-1], X[j], MOD[j]),
                    s[i-1], MOD[j]);
            }
        }
    }

    //Hash del substring en el rango [i, j)
    ll value(int l, int r) {
```

```

    int a = sbt(h[0][r], mul(h[0][1], xpow[0][r-1], MOD[0]),
        MOD[0]);
    int b = sbt(h[1][r], mul(h[1][1], xpow[1][r-1], MOD[1]),
        MOD[1]);
    return (1ll(a)<<32) + b;
}
};

void calc_xpow(int mxlen) {
    for (int j = 0; j < 2; ++j) {
        xpow[j].resize(mxlen+1, 1);
        for (int i = 1; i <= mxlen; ++i) {
            xpow[j][i] = mul(xpow[j][i-1], X[j], MOD[j]);
        }
    }
}

```

1.6.3 Prefix Function

```

vector<int> prefix_function(string s) {
    int n = (int)s.length();
    vector<int> pi(n);
    for (int i = 1; i < n; i++) {
        int j = pi[i-1];
        while (j > 0 && s[i] != s[j])
            j = pi[j-1];
        if (s[i] == s[j])
            j++;
        pi[i] = j;
    }
    return pi;
}

```

1.6.4 Suffix Array

```

const int MAXL = 300;

struct suffixArray {
    string s;
    int n, MX;
    vector<int> ra, tra, sa, tsa, lcp;

    suffixArray(string &s) {
        s = _s+"$";
        n = s.size();
        MX = max(MAXL, n)+2;
        ra = tra = sa = tsa = lcp = vector<int>(n);
        build();
    }

    void radix_sort(int k) {
        vector<int> cnt(MX, 0);
        for (int i = 0; i < n; i++)
            cnt[(i+k < n) ? ra[i+k]+1 : 1]++;
        for (int i = 1; i < MX; i++)
            cnt[i] += cnt[i-1];
        for (int i = 0; i < n; i++)
            tsa[cnt[(sa[i]+k < n) ? ra[sa[i]+k] : 0]++] = sa[i];
        sa = tsa;
    }

    void build() {
        for (int i = 0; i < n; i++)
            ra[i] = s[i], sa[i] = i;
        for (int k = 1, r; k < n; k <= 1) {
            radix_sort(k);
            radix_sort(0);
            tra[sa[0]] = r = 0;

```

```

    for (int i = 1; i < n; i++) {
        if (ra[sa[i]] != ra[sa[i-1]] || ra[sa[i]+k] !=
            ra[sa[i-1]+k]) ++r;
        tra[sa[i]] = r;
    }
    ra = tra;
    if (ra[sa[n-1]] == n-1) break;
}

int& operator[] (int i) { return sa[i]; }

void build_lcp() {
    lcp[0] = 0;
    for (int i = 0, k = 0; i < n; i++) {
        if (!ra[i]) continue;
        while (s[i+k] == s[sa[ra[i]-1]+k]) k++;
        lcp[ra[i]] = k;
        if (k) k--;
    }
}

//Longest Common Substring: construir el suffixArray s = s1
+ "#" + s2 + "$" y m = s2.size()
pair<int, int> lcs() {
    int mx = -1, ind = -1;
    for (int i = 1; i < n; i++) {
        if (((sa[i] < n-m-1) != (sa[i-1] < n-m-1)) && mx <
            lcp[i]) {
            mx = lcp[i]; ind = i;
        }
    }
    return {mx, ind};
}
};

```

1.6.5 Suffix Automaton

```

struct suffixAutomaton {
    struct node {
        int len, link; bool end;
        map<char, int> next;
        int cnt; ll in, out;
    };

    vector<node> sa;
    int last; ll subtrs = 0;

    suffixAutomaton() {}
    suffixAutomaton(string &s) {
        sa.reserve(s.size()*2);
        last = add_node();
        sa[0].link = -1;
        sa[0].in = 1;
        for (char &c : s) add_char(c);
        for (int p = last; p; p = sa[p].link) sa[p].end = 1;
    }

    int add_node() { sa.pb({}); return sa.size()-1; }

    void add_char(char c) {
        int u = add_node(), p = last;
        sa[u].len = sa[last].len + 1;
        while (p != -1 && !sa[p].next.count(c)) {
            sa[p].next[c] = u;
            sa[u].in += sa[p].in;
            subtrs += sa[p].in;
            p = sa[p].link;
        }
        if (p != -1) {
            int q = sa[p].next[c];

```

```

    if (sa[p].len + 1 != sa[q].len) {
        int clone = add_node();
        sa[clone] = sa[q];
        sa[clone].len = sa[p].len + 1;
        sa[clone].in = 0;
        sa[q].link = sa[u].link = clone;
        while (p != -1 && sa[p].next[c] == q) {
            sa[p].next[c] = clone;
            sa[q].in -= sa[p].in;
            sa[clone].in += sa[p].in;
            p = sa[p].link;
        }
        } else sa[u].link = q;
    }
    last = u;
}

void run(string &s) {
    int u = 0;
    for (int i = 0; i < s.size(); ++i) {
        while (u && !sa[u].next.count(s[i])) u = sa[u].link;
        if (sa[u].next.count(s[i])) u = sa[u].next[s[i]];
    }
}

int match_str(string &s) {
    int u = 0, n = s.size();
    for (int i = 0; i < n; ++i) {
        if (!sa[u].next.count(s[i])) return 0;
        u = sa[u].next[s[i]];
    }
    return count_occ(u);
}

int count_occ(int u) {

```

```

    if (sa[u].cnt != 0) return sa[u].cnt;
    sa[u].cnt = sa[u].end;
    for (auto &v : sa[u].next)
        sa[u].cnt += count_occ(v.S);
    return sa[u].cnt;
}

ll count_paths(int u) {
    if (sa[u].out != 0) return sa[u].out;
    for (auto &v : sa[u].next)
        sa[u].out += count_paths(v.S) + 1;
    return sa[u].out;
}

node& operator[](int i) { return sa[i]; }
};

```

1.6.6 Trie

```

const int N = 1e6 + 100;
int trie[N][26]; // N = suma de longitudes
bool stop[N];
int ct = 0;

void insert(string word) {
    int node = 0;
    for(int i = 0; i < (int)word.size(); i++) {
        if(!trie[node][word[i] - 'a']) trie[node][word[i] - 'a'] = ++ct;
        node = trie[node][word[i] - 'a'];
    }
    stop[node] = true;
}

```

1.6.7 Z Function

```
vector<int> z_function(string s) {  
    int n = s.size();  
    vector<int> z(n);  
    int l = 0, r = 0;  
    for(int i = 1; i < n; i++) {  
        if(i < r) {  
            z[i] = min(r - i, z[i - l]);  
        }  
        while(i + z[i] < n && s[z[i]] == s[i + z[i]]) {
```

```
            z[i]++;  
        }  
        if(i + z[i] > r) {  
            l = i;  
            r = i + z[i];  
        }  
    }  
    return z;  
}
```
