

MAC0331 - Lista 1

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Q 1:

Algorithm details: if a point is intersecting a line segment, I will say that the point is in the region on the right of the segment.

Both the segment insertion and the point look-up are based on a routine (1) that checks if a point is strictly to the left of the segment (the same used in class). Note that (1) is a *constant time* operation.

The tree is going to be a self-balancing binary search tree, and the insertion algorithm uses **any** point of the segment being inserted and (1) to decide if insertion is on the left/right of the segment it is compared to. Since it is a self-balancing binary tree, n consecutive insertions take time in $O(n \lg n)$.

In the same way, a single point can be checked against the segments on the binary tree. Since it is a balanced binary tree, the height is in $O(\lg n)$, and, therefore, determining the region a point belongs to is also in $O(\lg n)$ since (1) is a *constant time* operation.

Q 4:

This is exactly what I did in PE01. I will solve it with the sweep line method as described in class. Following is python code that solves the problem. Segment takes the polygon and returns a list of its segments. `verify_intersection` returns whether the two segments intersect.

```
1  def treat_left(s, bst):
2      bst.insert(s)
3      ns = bst.get_neighbours(s)
4      ret = verify_intersection(s, ns[0])
5      if (ret) : return True
6      ret = verify_intersection(s, ns[1])
7      if (ret) : return True
8
9  def treat_right (s, bst):
10     ret = False
11     bst.remove(s)
12     ns = bst.get_neighbours(s)
13     if(ns[0] and ns[1]):
14         ret = verify_intersection(ns[0], ns[1])
15     return ret
16 def Scanline (segments):
17     segments = segment(P1) + segment(P2)
18     segments = sorted(segments, key=functools.cmp_to_key(compare_segments))
19     heap, hmap = make_event_points(segments)
20     bst = balanced_binary_tree(Node_Seg)
21     while (not heap.empty()):
22         pt = heap.get()
23         for seg in pt.left:
24             if(treat_left(segments[seg], bst)):
25                 return True
```

```

26             for seg in pt.right:
27                 if(treat_right(segments[seg], bst)):
28                     return True
29         return False

```

Q5:

This is another application of the same sweep line algorithm above. I will make a change on the `verify_intersection()`, `segment()` methods. The only line changed is line 17.

Now, `verify_intersection()` will check $\Delta x^2 + \Delta y^2 < \Delta r^2$ to return whether there is an intersection. `segment()` will, instead of getting all the edges of the polygon, generate a single segment given by $[(x - r, y), (x + r, y)]$. Therefore, line 17 will be:

17. `segments = [segment(D) for D in disks]`

This algorithm is in $O(n \lg n)$ because sorting the disks is in that class, and there is a linear amount ($2n$) of event points generated (each event is in $O(\lg n)$).