```
mirror_mod.mirror_object
peration == "MIRROR_X":
mirror_mod.use_x = True
mirror_mod.use_y = False
mlrror_mod.use_z = False
 _operation == "MIRROR_Y"
irror_mod.use_x = False
mirror_mod.use_y = True
 mirror_mod.use_z = False
  operation == "MIRROR Z"
  lrror mod.use_x = False
  Irror mod.use y = False
  Irror mod.use z = True
  election at the end -add
   ob select= 1
    pplied Programming I
   bpy.context.selected_ob
   ata.objects[one.name].sel
  Programming computers in a nutshell.
     OPERATOR CLASSES
```

```
mirror_object
 peration == "MIRROR_X":
irror_mod.use_x = True
"Irror_mod.use_y = False
"Irror_mod.use_z = False
 _operation == "MIRROR_Y"
lrror_mod.use_x = False
irror_mod.use_y = True
 irror_mod.use_z = False
  operation == "MIRROR_Z";
  rror_mod.use_x = False
rror_mod.use_y = False
rror_mod.use_rnturoduction to
          ed systems programming
   rror ob.select = 0
 bpy.context.selected_obj
  mata.objects[one.name].sel
   Focus on system programming and RTOS usage
   X mirror to the selected
  ject.mirror_mirror_x"
```

Today's objectives

- Using an RTOS environment
- Working with files
- Command line interface
 - Starting a webserver

on the microcontroller ESP32

RTOS

Real-time operating system

Operating system's tasks



- An operating system (OS) controls the shared use of hardware between multiple applications.
- It provides resource management of CPU allocation, memory, storage.
- And offers application support by
 - Hardware abstraction,
 - Service provision.
- The kernel is the core component of an OS. It handles the creation, scheduling and termination of processes (applications).

General purpose operating systems

General purpose operating systems (e.g., Windows, Linux, macOS) provide:

- File system management
- Security and access control
- User interface, e.g., command line interface (CLI)
- Networking

Real-time operating system (RTOS)



- An RTOS provides scheduling of processes ensuring deterministic timing. It guarantees low latency (quick response).
- Operating systems for microcontroller can use only little resources as the microcontroller has limited resources.
- RTOS are reduced to the essentials and use only a small footprint.
- Essentials of an RTOS are scheduler, queues and semaphores.
- Scheduling of processes is **priority-based**.
- Configuration of an RTOS is done at compile time.

Advantages of using RTOS on microcontroller

An RTOS provides a framework for

- Scheduling
- Handling of timer and interrupts
- Usually different modules for specific functions.

Maintenance and portability of application is simple if the RTOS is maintained by its provider.

FreeRTOS

Example of an RTOS

FreeRTOS



- <u>FreeRTOS</u> is an example of an RTOS. It is very popular and supports many microcontroller architectures.
- It has a small kernel and is offered under the MIT open source license.
- There are two further variants of FreeRTOS, such as OpenRTOS and SafeRTOS. These are with commercial licenses and further features.

Core features of FreeRTOS

- FreeRTOS provides preemptive and co-operative scheduling of tasks ('kind of programs that should share the microcontroller').
- Priorities of tasks can be flexible managed.
- Important data structures such as queues, semaphores and mutexes are provided.
- Heap and stack management is provided. Stack overflow checking is possible.

FreeRTOS: task

- A task is a simple C function. There can be many tasks running on the system. The tasks will share the CPU(s). The scheduler determines the order of running the tasks.
- Every task get its own stack.
- A task function has no return value and takes a void pointer as input:

void aTaskFunction(void *pvParameters);

FreeRTOS: task structure

```
void aTask(void *pvParameters) {
 /* task setup before the infinite loop */
 /* each instance of a task has it own local variables */
 int aVariable = 0;
 /* this is main / infinite loop of the task: */
 while(1) {
 /* task codehere */
```

FreeRTOS: task generation

 A task has to be generated and register to be scheduled by FreeRTOS. This is done by the function xTaskCreate:

```
BaseType_t xTaskCreate( /* return pdPASS if ok*/

TaskFunction_t pxTaskCode, /* C function of task */

const char * const pcName, /* name of task */

const configSTACK_DEPTH_TYPE usStackDepth, /* size of stack */

void * const pvParameters, /* pointer to parameters for task */

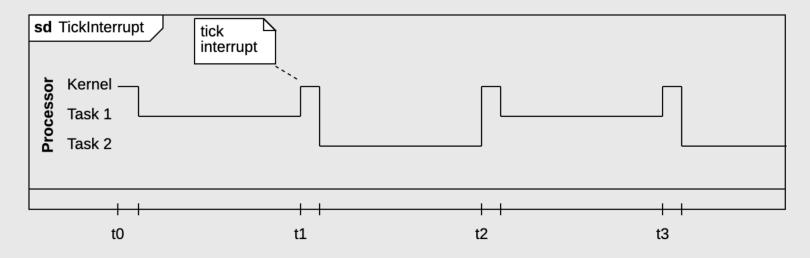
UBaseType_t uxPriority, /* priority of task */

TaskHandle_t * const pxCreatedTask ); /* handle of task */
```

FreeRTOS: task creation example

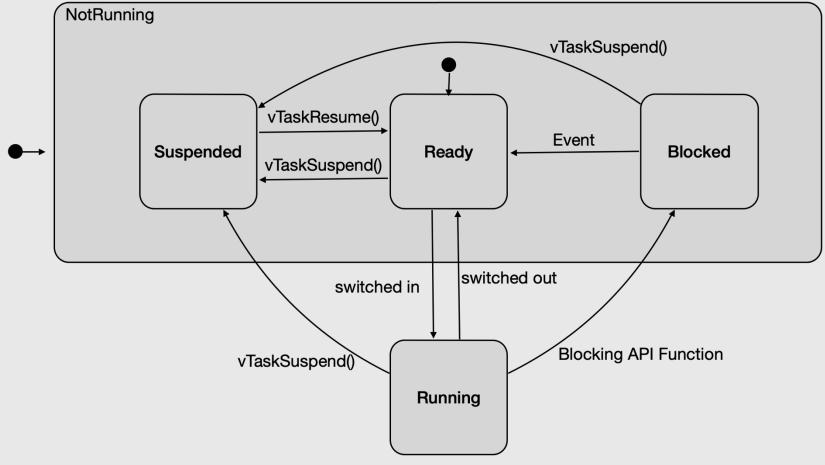
```
/* task function to be used */
void vTask1( void *pvParameters ) {
 int param = (int) pvParameters; /* cast and result: 1 */
 for(;;) {
/* generate task for vTask1, name: "Task 1", stack size: 4096
parameter: 1, priority: 2 and no taskhandle (NULL) */
xTaskCreate(vTask1, "Task 1", 4096, (void*) 1, 2, NULL );
```

FreeRTOS: scheduling example



- Two tasks with similar priority are scheduled round-robin.
- After every time slice, the kernel takes some control by the tick interrupt.

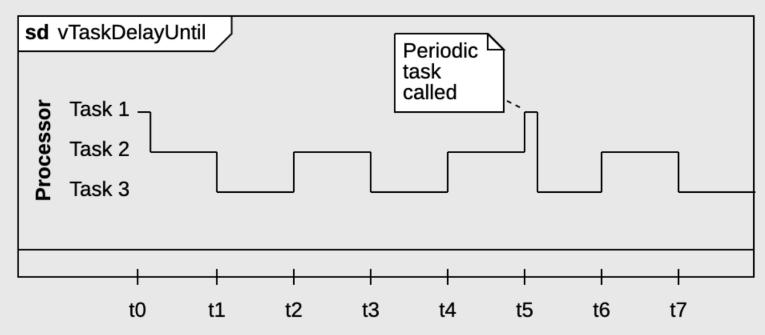
FreeRTOS: States of a task



FreeRTOS: States of a task

- Not-Running:
 - Blocked: wait for time event or synchronization event
 - Suspended: stopped of being scheduled by API call
 - Ready: can be scheduled
- Running: is executed

FreeRTOS: scheduling using vTaskDelayUntil



- Task 1 has higher priority than Task 2 and Task 3
- Task 1 blocks itself for 5 time slices by calling vTaskDelayUntil.

FreeRTOS: example scheduling

```
/* Task 1 has higher priority (2) than Task 2 and Task 3 (0)*/
xTaskCreate( vTask1, "Task 1", 4096, NULL, 2, NULL);
xTaskCreate( vTask2, "Task 2", 4096, NULL, 0, NULL);
xTaskCreate( vTask3, "Task 3", 4096, NULL, 0, NULL);
void vTask1( void *pvParameters ) {
 TickType t xLastWakeTime; /* a point in time given in ticks */
 const TickType t xDelay = pdMS TO TICKS(5000);
 xLastWakeTime = xTaskGetTickCount(); /* current time */
 for(;;) {
  /* periodic (it steps now back) for xDelay */
 vTaskDelayUntil( &xLastWakeTime, xDelay ); /* wait until (xLastW..T.. +xDelay) */
 }}
```

FreeRTOS: scheduling

- Fixed priority pre-emptive scheduling
 - Only the task itself can change priority
 - Pre-emption
 - Highest priority first
 - In case of same priority, round-robing
- Co-operative scheduling
 - Tasks need to release themselves

FreeRTOS: queues

- Queues are a means to communicate between tasks.
- Queues are used in FIFO mode (first in, first out)
- Accessing a queue has an optional waiting time
 - A task will be blocked until queue access is possible or time is over.
 - The task with highest priority accesses first in case of more tasks trying to access the queue.

FreeRTOS: queues usage / API

Generation of a queue:

```
QueueHandle_t xQueueCreate( /* return queue handle */
const UBaseType_t uxQueueLength, /* number of items */
const UBaseType_t uxItemSize ); /* size of an item in bytes */
```

Send to / receive from a queue:

FreeRTOS: queue example excerpt

```
/* Create a queue for uint32 t and store handle in blinkQueue */
blinkQueue = xQueueCreate( QUEUE LENGTH, sizeof( uint32 t ) );
... /* in a sender task, */
void senderTask( void *pvParameters ) {
 u int32 t data = 10;
... /* send the number stored at address data, the value 10 */
             xQueueSend(blinkQueue, &data, 0);
...}
... /* in a receiver task, */
void receiverTask( void *pvParameters ) {
u int32 t recValue = 10;
... /* if available, take a value from queue, wait max. 2s, store in recValue
    xQueueReceive(blinkQueue, &recValue, pdMS TO TICKS(2000));
...}
```



Your task: Work with queues in the example project

Group work, time: 20 Minutes.

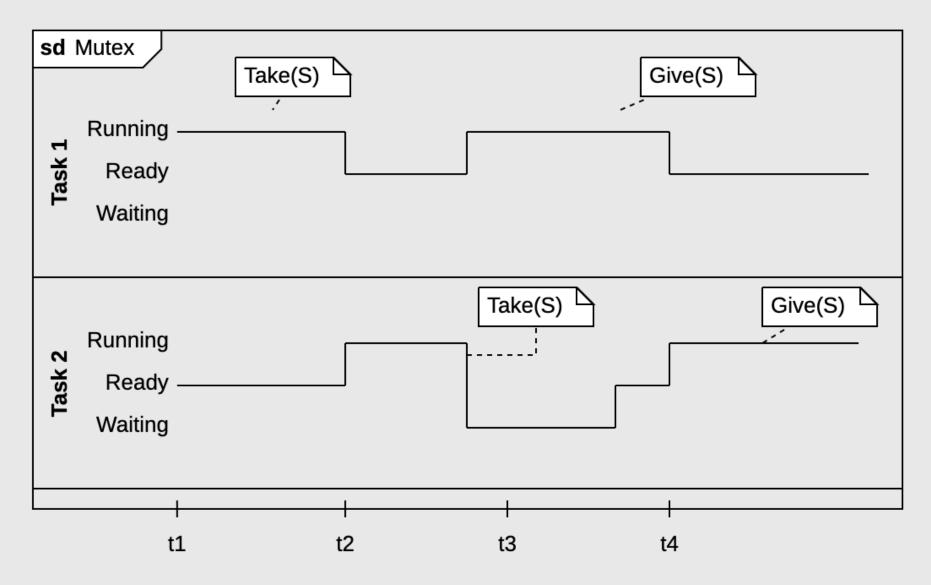
- 1. Open folder 'ap01/esp/freertos_queues_espidf' and run it once.
- 2. Change the program so that the LED blinks once quickly for each received message. Use vTaskDelay in the receiver task.
- 3. Send two or three messages (the LED should two-/three-times blink).

Mutexes



- When at least two tasks access a shared resource, it can be a problem (called race condition). Hence, the critical sections (with the resource access) have to be exclusively used.
- A mutex (**mutual exclusion**) is a mean to allow only one task accessing the shared resource:
 - When a process uses the resource, it **takes** the mutex. When the process is done with the resource, it **gives** back the mutex.
 - Another task trying to access the resource cannot take the mutex and has
 to wait until it is given back. Only one task at a time can have the mutex.
- Side note: priority inversion is usually prevented by priority inheritance.

Mutex example (tasks of same priority)

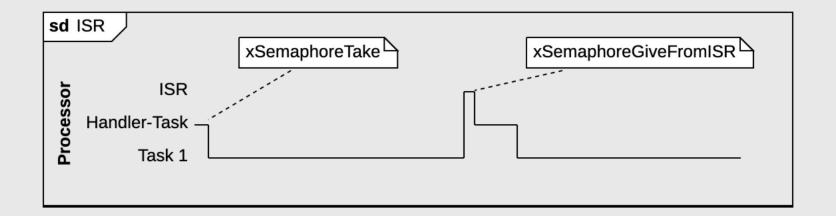


Semaphores



- Semaphores are 'counting mutexes' (or mutexes are 'binary semaphores').
- Semaphores can be used to
 - allow access control if there a shared resource can be used by a limited number of tasks,
 - synchronize tasks.

FreeRTOS: semaphore for ISR handling



- Binary semaphores should be used to interrupt handling in FreeRTOS so that interrupt service routines can be short.
- A high priority task is activated by the semaphore to execute the code for the handler.



Your task: Understand the example project

Group work, time: 10 Minutes.

- 1. Open folder 'ap01/esp/freertos_semaphore_espidf' and run it once.
- 2. Find the API calls of the previous slide and discuss / explain them in your group.

Filesystems

Filesystem



- A filesystem is the component of an operating system that organizes the storage of data on storage devices.
- This includes how files and folders are organized in a (usually) hierarchical structure and how they are mapped on the device.
- Nowadays, a filesystem provides access control and data integrity as well.
- File systems are mounted by the system for usage and unmounted after usage. They can be mounted read-only.
- Examples are NTFS, FAT, ext4, APFS.

Filesystem on ESP32



- Filesystems on microcontrollers, such as the ESP32 groups, are usually mapped to be used in flash memory or on SD cards.
- The flash memory of a microcontroller is partitioned into different **partitions** (areas) so that different partition contain different data (e.g., program code binaries and filesystems).
- A partition table describes how the memory is configured.
- The application can **mount** a filesystem to use it.

Filesystem on ESP32

- Among other filesystems for embedded systems, ESP32 supports the FAT filesystem.
- Using a partition tool, filesystems can be created as part of the binary firmware to the microcontroller.
- The flash utility allows not only write a binary but also read it so that data of the filesystem can be retrieved.

Filesystem on ESP32: partition table (1)

1. Step: create a partition table file with filesystem FAT for data:

```
# file partition.csv

# Name, Type, SubType, Offset, Size, Flags
nvs, data, nvs, 0x9000, 0x6000,
phy_init, data, phy, 0xf000, 0x1000,
factory, app, factory, 0x10000, 1M,
storage, data, fat, , 1M,
```

Filesystem on ESP32: partition table (1)

2. Step: configure project for partition table (in sdkconfig.default):

```
CONFIG_PARTITION_TABLE_CUSTOM=y

CONFIG_PARTITION_TABLE_CUSTOM_FILENAME="partitions.csv"

CONFIG_PARTITION_TABLE_FILENAME="partitions.csv"

CONFIG_ESPTOOLPY_FLASHSIZE_8MB=y
```

Filesystem on ESP32: partition table (2)

2. Step: configure project for partition table (in sdkconfig.default):

```
CONFIG_PARTITION_TABLE_CUSTOM=y

CONFIG_PARTITION_TABLE_CUSTOM_FILENAME="partitions.csv"

CONFIG_PARTITION_TABLE_FILENAME="partitions.csv"

CONFIG_ESPTOOLPY_FLASHSIZE_4MB=y
```

Filesystem on ESP32: files for partition (3)

3. Create a folder with file that should be in the FAT partition, e.g., folder 'fatfs_image/test.txt'. And configure CMakeLists.txt:

```
# Create a FATFS image from the contents of the 'fatfs_image' directory

# that fits the partition named 'storage'. FLASH_IN_PROJECT indicates that

# the generated image should be flashed when the entire project is flashed to

# the target with 'idf.py -p PORT flash'.

set(image ../fatfs_image)

fatfs create spiflash image(storage ${image} FLASH IN PROJECT)
```

Filesystem on ESP32: mount in application (4)

```
const char *base path = "/fat"; /* base path to mount partition */
const esp_vfs_fat_mount_config_t mount_config = {
 .max_files = 4, /* no. that can be opened */
 .format if mount failed = false, /* s. API documentation */
 .allocation_unit_size = CONFIG_WL_SECTOR_SIZE
esp_err_t err = esp_vfs_fat_spiflash_mount_rw_wl(
  base_path, /* base, e.g., root file system, mount point */
  "storage", /* partition to mount, see partition table */
  &mount_config, /* configuration, see above */
  &s wl handle); /* drive handle (for low-level use) */
```

Filesystem on ESP32: accessing files (5)

- Files on the filesystem can be accessed by the POSIX (Portable Operating System Interface) standard.
- Besides a lot of other things, POSIX allows a standard C API for file access (stdio.h):
 - Open and close a file by using its file name.
 - Reading from and writing to a file (when it is open).

Filesystem on ESP32: reading files (5)

```
const char *filename_read = "/fat/test.txt"; /* file name */
char line[128];
                               /* buffer for text lines */
FILE *f;
                           /* file handle */
f = fopen(filename_read, "r"); /* open file for read */
if (f == NULL) { return; } /* error check*/
                         /* read line by line until end and print */
while (fgets(line, sizeof(line), f) != NULL) {/* read character by character with fgets */
 printf("%s", line);
                             /* close file */
fclose(f);
```

Filesystem on ESP32: writing files (5)

Filesystem on ESP32: reading the partition

 The partition can be retrieved from the microcontroller and translated into files by:

parttool.py --port /dev/tty.usbmodem1101 read_partition --partition-type=data --partition-subtype=fat --output flash.img

\$IDF_PATH/components/fatfs/fatfsparse.py flash.img



Your task: Make an application that saves sensor data.

Group work, time: 60 Minutes.

- 1. Open the folder 'ap01/esp/fatfs_espidf' and understand the code.
- 2. Open the folder 'ap01/esp/fatfs_mpu6050_espidf' and understand the code.
- 3. Write your code in the prepared setting between the /* TODO: begin */ and /* TODO: end */. The comments should help you how to do it.

Console

Console (or terminal)



- A console is a text-based interface to interact with the operating system.
- It implements a command line interface (CLI). Commands are typed in, executed and results displayed.
- Using a serial communication, a console can be easily implemented.

ESP32 and console

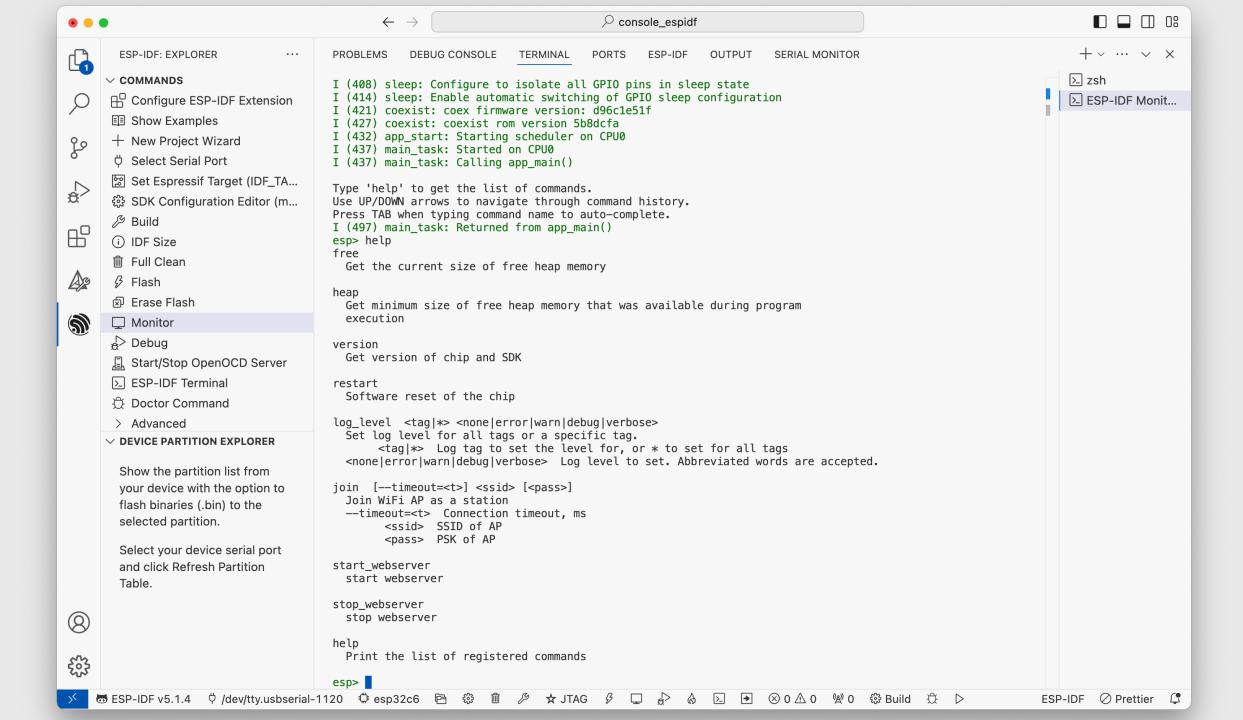
- ESP-IDF offers an API to implement a console.
- It provides basic commands that can be registered with a command 'word' to provide common used functionality.

ESP32: setting up a console

```
/* prepare console */
esp_console_repl_t *repl = NULL;
esp_console_repl_config_t repl_config = ESP_CONSOLE_REPL_CONFIG_DEFAULT();
/* register commands */
esp_console_register_help_command();
register_system_common();
register_wifi();
/* init console on uart */
esp_console_dev_uart_config_t hw_config = ESP_CONSOLE_DEV_UART_CONFIG_DEFAULT();
esp_console_new_repl_uart(&hw_config, &repl_config, &repl);
/* start console */
esp_console_start_repl(repl);
```

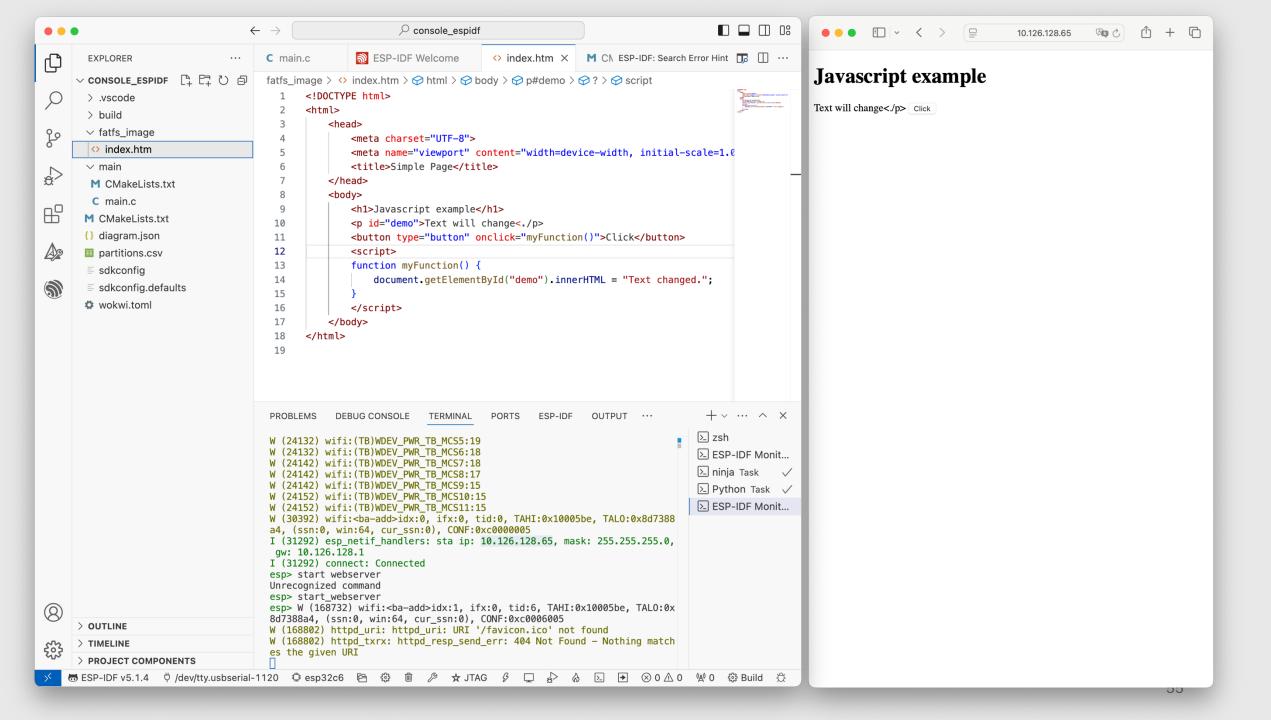
ESP32: register a console command

```
int start webserver func(int argc, char** argv);
void register start webserver() {
 /* structure defining the command */
 const esp_console_cmd_t start_webserver_cmd = {
  .command = "start webserver", /* command name */
  .help = "start webserver", /* info text */
  .hint = NULL, /* no hint text */
  .func = &start webserver func, /* callback function */
 };
 /* register command */
 esp_console_cmd_register(&start_webserver_cmd);
```



ESP32: essentials of running a webserver

```
esp err tindex get handler(httpd reg t *reg) {
 ... httpd resp send(reg, html page, HTTPD RESP USE STRLEN); ...
httpd uri t uri index = {
 .uri = "/", /* uniform resource identifier */
 .method = HTTP GET, /* http method: for request GET */
 .handler = index get handler, /* callback function */
 .user ctx = NULL /* no user context here */ };
httpd config t config = HTTPD DEFAULT CONFIG();
if (httpd start(&web server handle, &config) == ESP OK) {
 httpd register uri handler(web server handle, &uri index);
```



Today's summary

- Using FreeRTOS as example of an RTOS
- Working with files on a microcontroller (ESP32)
 - Partitions
 - Mounting
 - Standardized file I/O
- Console (command line interface) (ESP32)
 - Using standard functionality of system API
 - Running a web server

As always...
Questions?

