

Question 1

Correct

Mark 1.00 out of

1.00

Flag question

Do there exists some prefix binary code containing words of the following lengths?

[3,3,3,5,5,6,6,7,8,8,9,9,9]

Select one:

☒ Yes ✓

The given lengths do satisfy Kraft inequality

☐ No

Your answer is correct.

The correct answer is: Yes

```
def prefix_code(lst):
    output = []
    sumatorio = 0
    for i in lst:
        sumatorio += 2**(-i)
    if (sumatorio > 1):
        return output
    else:
        for index, i in enumerate(lst):
            if (index is 0):
                output.append(bin(0)[2:].zfill(i))
            else:
                for j in range(0, 2**i):
                    prefix = False
                    aux = bin(j)[2:].zfill(i)
                    for binword in output:
                        prefix = (aux.startswith(binword)) or (binword.startswith(aux))
                    if prefix:
                        break
                    if not prefix:
                        output.append(aux)
                        break
        return output

prova = [3,3,3,5,5,6,6,7,8,8,9,9,9]
binarywords = prefix_code(prova)
if not binarywords:
    print("No")
else:
    print("Yes")
    [print(binword) for binword in binarywords]
...
Yes
000
001
010
01100
01101
011100
011101
0111100
01111010
01111011
011111000
011111001
011111010
...
```

Question 2

Partially correct

Mark 0.67 out of

1.00

Remove flag

Of the following sets of binary words; which ones are codes?

Select one or more:

- ☐ ['00', '11', '001', '111', '00111', '01111', '10000']
- ☒ ['000', '001', '010', '0110', '01110', '01111', '10000'] ✓ This is a prefix set
- ☐ ['00', '01', '10', '11', '00011']
- ☒ ['00', '010', '011', '1000', '10010', '10011', '10100'] ✓ This is a prefix set
- ☐ ['00', '01', '10', '110', '111', '10101']
- ☐ ['00', '001']

You have correctly selected 2.

The correct answer is: ['00', '010', '011', '1000', '10010', '10011', '10100'], ['000', '001', '010', '0110', '01110', '01111', '10000'], ['00', '001']

```
def prefix_code(code):
    for binword in code:
        for aux in code:
            if (binword != aux):
                prefix = (aux.startswith(binword)) or (binword.startswith(aux))
                if (prefix):
                    return('No')
        return('Yes')

code = ['00', '11', '001', '111', '00111', '01111', '10000']
print(prefix_code(code))
code = ['000', '001', '010', '0110', '01110', '01111', '10000']
print(prefix_code(code))
code = ['00', '01', '10', '11', '00011']
print(prefix_code(code))
code = ['00', '010', '011', '1000', '10010', '10011', '10100']
print(prefix_code(code))
code = ['00', '01', '10', '110', '111', '10101']
print(prefix_code(code))
code = ['00', '001']
print(prefix_code(code))
'''
No
Yes
No
Yes
No
No
'''
```

Question 3

Correct

Mark 1.00 out of

1.00

Flag question

The **Shannon source code theorem** gives

Select one:

- ☒ Lower and upper bounds for the mean length of an optimal code ✓ Yes, this is exactly what it does
- ☐ A criterion to decide whether a set of binary words is a code or not
- ☐ A method to build optimal codes
- ☐ A formula for the entropy of a source

Your answer is correct.

The correct answer is: Lower and upper bounds for the mean length of an optimal code

Question 4

Correct

Mark 1.00 out of

1.00

Flag question

Which of the following sources has the largest entropy?

Select one:

- ☐ A source of 8 letters with different probabilities
- ☐ The 3-extension of a binary source with different probabilities
- ☐ A source of 8 letters with equal probabilities
- ☒ The 4-extension of a binary source with equal probabilities ✓ Correct. The entropy is 4, the largest possible among all examples
- ☐ The 4-extension of a binary source with different probabilities

Your answer is correct.

The correct answer is: The 4-extension of a binary source with equal probabilities

```
from math import log2
from itertools import product

def source(str):
    output = []
    for char in set(str):
        output.append((char, str.count(char) / len(str)))
    return output

def source_extension(src,k):
    letters = []
    output = []
    for tupla in src:
        letters.append(tupla[0])
    permutations = [''.join(i) for i in product(letters, repeat = k)]
    for perm in permutations:
        suma = 1
        for let in letters:
            proplet = [tupla[1] for tupla in src if tupla[0] == let]
            suma *= proplet[0]**(perm.count(let))
        output.append((perm, suma))
    return output

def entropy_source(src):
    hx = 0.0
    for i in range(len(src)):
        pxi = src[i][1]
        hx += pxi*(log2(1.0/pxi))
    return hx

s = [('a',0.03), ('b',0.05), ('c',0.08), ('d',0.09),
      ('e',0.1), ('f',0.15), ('g',0.2), ('h',0.3)]
print('A source of 8 letters with different probabilities',entropy_source(s))

string = '00000100000000000100'
s = source(string)
source_ext = source_extension(s,3)
print('The 3-extension of a binary source with different probabilities',entropy_source(source_ext))
```

```

s = [('a',0.125), ('b',0.125), ('c',0.125), ('d',0.125),
      ('e',0.125), ('f',0.125), ('g',0.125), ('h',0.125)]
print('A source of 8 letters with equal probabilities',entropy_source(s))

string = '0000011111'
s = source(string)
source_ext = source_extension(s,4)
print('The 4-extension of a binary source with equal probabilities',entropy_source(source_ext))

string = '0000010000000000100'
s = source(string)
source_ext = source_extension(s,4)
print('The 4-extension of a binary source with different probabilities',entropy_source(source_ext))

'''
A source of 8 letters with different probabilities 2.700238463388872
The 3-extension of a binary source with different probabilities 1.406986780767844
A source of 8 letters with equal probabilities 3.0
The 4-extension of a binary source with equal probabilities 4.0
The 4-extension of a binary source with different probabilities 1.875982374357125
'''

```

Question 5

Partially correct

Mark 0.33 out of 1.00

Flag question

Looking at a long string of binary digits we see that

- The probability that a digit is a one is $\frac{1}{4}$
- The probability that two consecutive digits are both ones is $\frac{1}{10}$
- The probability that after a zero comes another zero is $\frac{4}{5}$

Which of the following are true?

Select one or more:

- ☐ The probability that a digit is a zero is $\frac{15}{16}$
- ☒ The probability that after a one comes a one is $\frac{2}{5}$
- ☐ The probability of two consecutive zeros is $\frac{13}{15}$
- ☐ The probability of the pair of consecutive symbols zero-one is $\frac{3}{20}$
- ☐ The probability that after a zero comes a one is $\frac{1}{5}$
- ☐ The probability that after a one comes a zero is $\frac{3}{10}$

Your answer is partially correct.

If X, Y are the random variables corresponding to pick randomly two consecutive digits,

- $P(X=1) = \frac{1}{4}$
- $P(X=1, Y=1) = p_2$
- $P(Y=0|X=0) = p_3$

With the formulas and properties of probability theory for pairs of random variables, one can compute:

- $P(X=0) = 1 - P(X=1)$
- $P(X=0, Y=0) = P(Y=0|X=0) \cdot P(X=0)$
- $P(Y=1|X=0) = 1 - P(Y=0|X=0)$
- $P(X=0, Y=1) = P(Y=1|X=0) \cdot P(X=0)$
- $P(X=1, Y=0) = 1 - P(X=0, Y=0) - P(X=1, Y=1) - P(X=0, Y=1)$

and answer all the questions

The correct answer is: The probability that after a zero comes a one is $\frac{1}{5}$, The probability that after a one comes a one is $\frac{2}{5}$, The probability of the pair of consecutive symbols zero-one is $\frac{3}{20}$

```
from fractions import Fraction
```

```
probX1 = Fraction(1,4)
probX1Y1 = Fraction(1,10)
probY0_X0 = Fraction(4,5)
```

```
probX0 = 1 - probX1
probX0Y0 = probY0_X0 * probX0
probY1_X0 = 1 - probY0_X0
probX0Y1 = probY1_X0 * probX0
probX1Y0 = 1 - probX0Y0 - probX1Y1 - probX0Y1
probY0_X1 = probX1Y0 / probX1
probY1_X1 = 1 - probY0_X1
```

```
print('The probability that a digit is a zero is',probX0)
print('The probability that after a one comes a one is',probY1_X1)
print('The probability of two consecutive zeros is',probX0Y0)
print('The probability of the pair of consecutive symbols zero-one is',probX0Y1)
print('The probability that after a zero comes a one is',probY1_X0)
print('The probability that after a one comes a zero is',probY0_X1)
```

```
'''
The probability that a digit is a zero is 3/4
The probability that after a one comes a one is 2/5
The probability of two consecutive zeros is 3/5
The probability of the pair of consecutive symbols zero-one is 3/20
The probability that after a zero comes a one is 1/5
The probability that after a one comes a zero is 3/5
'''
```

Question 6

Not answered

Marked out of 1.00

Flag question

What knowledge about the outcome of the experiment of throwing two dice contains more information?

Select one:

- ☐ The product is (strictly) larger than 10
- ☐ The value of the sum
- ☐ The sum is (strictly) larger than 10
- ☐ The sum is equal to 7
- ☐ The outputs have different parity

Your answer is incorrect.

The probability that the sum is larger than 10 is 3/36 (favorable cases: (5,6), (6,5) and (6,6)), and the corresponding information is $\log_2(12) \approx 3.58496$...

The correct answer is: The sum is (strictly) larger than 10

Question 7

Correct

Mark 1.00 out of 1.00

Flag question

Who is the father of **Information Theory** ?

Select one:

- ☒ Claude Shannon ✓
- ☐ David Huffman
- ☐ Alan Turing
- ☐ Peter Arithmetic

Your answer is correct.

Information Theory was created in 1947 by engineer and mathematician Claude Shannon

The correct answer is: Claude Shannon

Question 8

Not answered

Marked out of 1.00

Flag question

A sensor in a factory takes measures at regular intervals of time, represented by integers n in the range from 0 to $N-1$. Due to the properties of the system for every $n=0,1,2,\dots,N-3,N-2$ the probability of measuring the value n is twice the probability of measuring the next value $n+1$.

The chief engineer of the factory proposes the three following ways to binary encode the information provided by the sensor:

- (A) Use a block code with words of length k with k the smallest integer such that $N \leq 2^k$;
- (B) Encode every integer n with the binary word consisting of n ones followed by a zero;
- (C) Encode every integer n with the binary word consisting of n ones followed by a zero, except the largest integer $n=N-1$, which is encoded just as the word consisting of n ones without a zero at the end.

Answer for true or false:

The code in (A) is a Huffman code for the source described

Choose...

(A) and (B) are prefix sets, but (C) is not prefix, and consequently it is not a code

Choose...

The code (C) is a Huffman code for the source described

Choose...

For $N=4$ the entropy of the source is $\log(15)-34/15$

Choose...

For $N=8$ the probability of $n=7$ is $1/255$

Choose...

When N tends to infinity the probability of $n=1$ tends to $1/4$

Choose...

For $N=8$ the probability of $n=7$ is $1/7$

Choose...

The code (B) is a Huffman code for the source described

Choose...

For $N=4$ the entropy of the source is $\log(15)+16/15$

Choose...

Your answer is incorrect.

The correct answer is: The code in (A) is a Huffman code for the source described – False, (A) and (B) are prefix sets, but (C) is not prefix, and consequently it is not a code – False, The code (C) is a Huffman code for the source described – True, For $N=4$ the entropy of the source is $\log(15)-34/15$ – True, For $N=8$ the probability of $n=7$ is $1/255$ – True, When N tends to infinity the probability of $n=1$ tends to $1/4$ – True, For $N=8$ the probability of $n=7$ is $1/7$ – False, The code (B) is a Huffman code for the source described – False, For $N=4$ the entropy of the source is $\log(15)+16/15$ – False

Question 9

Correct

Mark 1.00 out of

1.00

Flag question

An arithmetic code working with integers of 11 bits is used to encode information produced by a source with an alphabet of three letters and probabilities

11	16765	65
50436	16812	25218

Is the precision 11 enough for working with that source?

Select one:

- ☒ No ✓
- ☐ Yes

Your answer is correct.

The smallest good precision is 15, the base 2 logarithm of the inverse of the smallest probability (integer part by excess) plus 2

The correct answer is: No

```

from math import log2, ceil
k = 11
probs = [11/50436, 16765/16812, 65/25218]
print('The smallest good precision is',ceil(log2(1/min(probs))+2))
if k >= ceil(log2(1/min(probs))+2):
    print("Yes")
else:
    print("No")

'''
The smallest good precision is 15
No
'''

```

Question 10

Correct

Mark 1.00 out of 1.00

Flag question

For which probability distribution the following binary code is a Huffman code ? (codewords and probabilities are given in the same order)

[1,'0111','0110','01011','01010','01000','01001','0011','0010','0000','0001']

Select one:

- ☒ [0.9,0.01,0.01,0.01,0.01,0.01,0.01,0.01,0.01,0.01,0.01]
- ☐ [0.07,0.08,0.08,0.08,0.08,0.09,0.1,0.1,0.1,0.11,0.11]
- ☐ [1/11,1/11,1/11,1/11,1/11,1/11,1/11,1/11,1/11,1/11,1/11]

Your answer is correct.

The correct answer is: [0.9,0.01,0.01,0.01,0.01,0.01,0.01,0.01,0.01,0.01,0.01]

```

import queue

class HuffmanNode(object):
    def __init__(self, left, right, root):
        self.left = left
        self.right = right
        self.root = root

def fromRootToLeafs(node, binarywords, aux):
    if (node.left == None and node.right == None):
        binarywords.append(aux)
    else:
        if (node.right != None):
            aux += '0'
            fromRootToLeafs(node.right, binarywords, aux)
            aux = aux[:len(aux)-1]
        if (node.left != None):
            aux += '1'
            fromRootToLeafs(node.left, binarywords, aux)

def huffman_code(src):
    id_desempate = 0
    q = queue.PriorityQueue()
    q2 = queue.PriorityQueue()
    for tupla in src:
        leafNode = HuffmanNode(None, None, tupla[1])
        q.put((tupla[1], id_desempate, leafNode))
        q2.put((tupla[1], id_desempate, leafNode))
        id_desempate += 1
    while q.qsize() > 1:
        leftNode, rightNode = q.get()[2], q.get()[2]
        parentNode = HuffmanNode(leftNode, rightNode, leftNode.root+rightNode.root)
        q.put((parentNode.root, id_desempate, parentNode))
        id_desempate += 1
    huffmanTree = q.get()[2]
    binarywords = []
    fromRootToLeafs(huffmanTree, binarywords, '')
    binarywords.sort(key = len)
    mean_length = 0.0
    for binword in reversed(binarywords):
        mean_length += q2.get()[0] * len(binword)
    return binarywords, mean_length

```

```

bincode = ['1','0111','0110','01011','01010','01000','01001','0011','0010','0000','0001']
for i in bincode:
    print(len(i), end = " ")
print('')

s1 = [('a',0.9), ('b',0.01), ('c',0.01), ('d',0.01),
      ('e',0.01), ('f',0.01), ('g',0.01), ('h',0.01),
      ('i',0.01), ('j',0.01), ('k',0.01)]
huff, mean = huffman_code(s1)
print('Huffman code for the source s1')
print(huff,',',mean)
for i in huff:
    print(len(i), end = " ")
print('')

s2 = [('a',0.07), ('b',0.08), ('c',0.08), ('d',0.08),
      ('e',0.08), ('f',0.09), ('g',0.1), ('h',0.1),
      ('i',0.1), ('j',0.11), ('k',0.11)]
huff, mean = huffman_code(s2)
print('Huffman code for the source s2')
print(huff,',',mean)
for i in huff:
    print(len(i), end = " ")
print('')

s3 = [('a',1/11), ('b',1/11), ('c',1/11), ('d',1/11),
      ('e',1/11), ('f',1/11), ('g',1/11), ('h',1/11),
      ('i',1/11), ('j',1/11), ('k',1/11)]
huff, mean = huffman_code(s3)
print('Huffman code for the source s3')
print(huff,',',mean)
for i in huff:
    print(len(i), end = " ")
print('')

'''
1 4 4 5 5 5 5 4 4 4 4
Huffman code for the source s1
['0', '1010', '1011', '1100', '1101', '1110', '1111', '10000', '10001', '10010', '10011'] ,
1.3399999999999999
1 4 4 4 4 4 4 5 5 5 5
Huffman code for the source s2
['011', '100', '101', '110', '111', '0000', '0001', '0010', '0011', '0100', '0101'] ,
3.4800000000000004
3 3 3 3 3 4 4 4 4 4 4
Huffman code for the source s3
['011', '100', '101', '110', '111', '0000', '0001', '0010', '0011', '0100', '0101'] ,
3.545454545454545
3 3 3 3 3 4 4 4 4 4 4
'''

```

Question

11

Correct

Mark 1.00 out of

1.00

Flag question

What is an optimal source code ?

Select one:

- ☐ A source code with mean length smaller than the entropy of the source plus one
- ☐ An arithmetic code with infinite precision (exact numbers)
- ☒ A source code with the smallest mean length among all possible source codes ✓ Yes, this is the definition
- ☐ A code with mean length smaller than the entropy of the source

Your answer is correct.

The correct answer is: A source code with the smallest mean length among all possible source codes

Question 12

Correct

Mark 1.00 out of

1.00

Flag question

Consider the source with alphabet $\{a, b, c, d\}$ and uniform probability distribution (all letters have the same probability).

We encode the text 'cbaa' using arithmetic coding with infinite precision.

Of the following numbers, which one is an encoding of the given text?

Select one:

☐ 0.063

☐ 0.6911

☒ 0.5636 ✓

☐ 0.6096

Your answer is correct.

The correct answer is: 0.5636

```
src = [("a",0.25), ("b",0.25), ("c",0.25), ("d",0.25)]
m = 'cbaa'
cumulative_probs = [0]
aux = 0.0
for i in range(len(src)):
    aux += src[i][1]
    cumulative_probs.append(aux)
print('cumulative_probs ->', cumulative_probs)
alpha = 0
beta = 1
for char in m:
    subintervals = []
    for j in range(1, len(cumulative_probs)):
        aux = (alpha + (beta-alpha)*cumulative_probs[j-1],
              alpha + (beta-alpha)*cumulative_probs[j])
        subintervals.append(aux)
    ind = [src.index(tupla) for tupla in src if tupla[0] == char][0]
    print(subintervals)
    print(char, 'ocupa la posicion', ind, 'en src')
    alpha = subintervals[ind][0]
    beta = subintervals[ind][1]
    print('Nos quedamos con [' , subintervals[ind][0] , ',' , subintervals[ind][1] , ')')
print('¿Que valor de los siguientes puede representar el mensaje?')
opcionA = 0.063
if (subintervals[ind][0] <= opcionA < subintervals[ind][1]):
    print('opcionA')
opcionB = 0.6911
if (subintervals[ind][0] <= opcionB < subintervals[ind][1]):
    print('opcionB')
opcionC = 0.5636
if (subintervals[ind][0] <= opcionC < subintervals[ind][1]):
    print('opcionC')
opcionD = 0.6096
if (subintervals[ind][0] <= opcionD < subintervals[ind][1]):
    print('opcionD')

'''
cumulative_probs -> [0, 0.25, 0.5, 0.75, 1.0]
[(0, 0.25), (0.25, 0.5), (0.5, 0.75), (0.75, 1.0)]
c ocupa la posicion 2 en src
Nos quedamos con [ 0.5 , 0.75 )
[(0.5, 0.5625), (0.5625, 0.625), (0.625, 0.6875), (0.6875, 0.75)]
b ocupa la posicion 1 en src
Nos quedamos con [ 0.5625 , 0.625 )
[(0.5625, 0.578125), (0.578125, 0.59375), (0.59375, 0.609375), (0.609375, 0.625)]
a ocupa la posicion 0 en src
Nos quedamos con [ 0.5625 , 0.578125 )
[(0.5625, 0.56640625), (0.56640625, 0.5703125), (0.5703125, 0.57421875), (0.57421875, 0.578125)]
a ocupa la posicion 0 en src
Nos quedamos con [ 0.5625 , 0.56640625 )
¿Que valor de los siguientes puede representar el mensaje?
opcionC
'''
```

Pregunta 3

No s'ha respost

encara

Puntuat sobre 1,00

Marca la

pregunta

Edita la

pregunta aritmetica

texto -> X

Con el alfabeto ['a','b','c','d'] y probabilidades [0.25,0.25,0.25,0.25] se ha codificado aritmeticamente el mensaje bcba. ¿Qué valor de los siguientes puede representar el mensaje?

Trieu-ne una:

- ☐ a. 0.1422
- ☐ b. 0.9202
- ☐ c. 0.3926

```
src = [("a",0.25), ("b",0.25), ("c",0.25), ("d",0.25)]
m = 'bcba'
cumulative_probs = [0]
aux = 0.0
for i in range(len(src)):
    aux += src[i][1]
    cumulative_probs.append(aux)
print('cumulative_probs ->', cumulative_probs)
alpha = 0
beta = 1
for char in m:
    subintervals = []
    for j in range(1, len(cumulative_probs)):
        aux = (alpha + (beta-alpha)*cumulative_probs[j-1],
              alpha + (beta-alpha)*cumulative_probs[j])
        subintervals.append(aux)
    ind = [src.index(tupla) for tupla in src if tupla[0] == char][0]
    print(subintervals)
    print(char, 'ocupa la posicion', ind, 'en src')
    alpha = subintervals[ind][0]
    beta = subintervals[ind][1]
    print('Nos quedamos con [' + subintervals[ind][0] + ',' + subintervals[ind][1] + ')')
print('¿Que valor de los siguientes puede representar el mensaje?')
opcionA = 0.1422
if (subintervals[ind][0] <= opcionA < subintervals[ind][1]):
    print('opcionA')
opcionB = 0.9202
if (subintervals[ind][0] <= opcionB < subintervals[ind][1]):
    print('opcionB')
opcionC = 0.3926
if (subintervals[ind][0] <= opcionC < subintervals[ind][1]):
    print('opcionC')

'''
cumulative_probs -> [0, 0.25, 0.5, 0.75, 1.0]
[(0, 0.25), (0.25, 0.5), (0.5, 0.75), (0.75, 1.0)]
b ocupa la posicion 1 en src
Nos quedamos con [ 0.25 , 0.5 )
[(0.25, 0.3125), (0.3125, 0.375), (0.375, 0.4375), (0.4375, 0.5)]
c ocupa la posicion 2 en src
Nos quedamos con [ 0.375 , 0.4375 )
[(0.375, 0.390625), (0.390625, 0.40625), (0.40625, 0.421875), (0.421875, 0.4375)]
b ocupa la posicion 1 en src
Nos quedamos con [ 0.390625 , 0.40625 )
[(0.390625, 0.39453125), (0.39453125, 0.3984375), (0.3984375, 0.40234375), (0.40234375, 0.40625)]
a ocupa la posicion 0 en src
Nos quedamos con [ 0.390625 , 0.39453125 )
¿Que valor de los siguientes puede representar el mensaje?
opcionC
'''
```

Pregunta 4

No s'ha respost

encara

Puntuat sobre 1,00

Marca la pregunta

Edita la pregunta huffman y entropía

Una fuente S viene dada per la ddp $\left[\frac{19}{73}, \frac{6}{73}, \frac{17}{73}, \frac{6}{73}, \frac{25}{73} \right]$. Un código de Huffman asociado a S tiene una longitud media \tilde{l} . ¿Qué valores puede tomar \tilde{l} ?

Trieu-ne una:

- ☐ a. Entre 3.117 y 3.35
- ☐ b. Entre 3.35 y 4.117
- ☐ c. Entre 2.117 y 3.117

```
import queue

class HuffmanNode(object):
    def __init__(self, left, right, root):
        self.left = left
        self.right = right
        self.root = root

def fromRootToLeafs(node, binarywords, aux):
    if (node.left == None and node.right == None):
        binarywords.append(aux)
    else:
        if (node.right != None):
            aux += '0'
            fromRootToLeafs(node.right, binarywords, aux)
            aux = aux[:len(aux)-1]
        if (node.left != None):
            aux += '1'
            fromRootToLeafs(node.left, binarywords, aux)

def huffman_code(src):
    id_desempate = 0
    q = queue.PriorityQueue()
    q2 = queue.PriorityQueue()
    for tupla in src:
        leafNode = HuffmanNode(None, None, tupla[1])
        q.put((tupla[1], id_desempate, leafNode))
        q2.put((tupla[1], id_desempate, leafNode))
        id_desempate += 1
    while q.qsize() > 1:
        leftNode, rightNode = q.get()[2], q.get()[2]
        parentNode = HuffmanNode(leftNode, rightNode, leftNode.root+rightNode.root)
        q.put((parentNode.root, id_desempate, parentNode))
        id_desempate += 1
    huffmanTree = q.get()[2]
    binarywords = []
    fromRootToLeafs(huffmanTree, binarywords, '')
    binarywords.sort(key = len)
    mean_length = 0.0
    for binword in reversed(binarywords):
        mean_length += q2.get()[0] * len(binword)
    return binarywords, mean_length

src = [('a', 19/73), ('b', 6/73), ('c', 17/73), ('d', 6/73), ('e', 25/73)]
print(huffman_code(src))

'''
(['00', '01', '10', '110', '111'], 2.1643835616438354)
'''
```

Pregunta 5

No s'ha respost

encara

Puntuat sobre 1,00

Marca la

pregunta

Edita la

pregunta aritmetica

x->texto

Con el alfabeto ['a','b','c','d'] y probabilidades [0.25,0.25,0.25,0.25] se ha codificado un mensaje de longitud 5 con el valor 0.1094. ¿Cual es el mensaje?

Trieu-ne una:

- ☐ a. abcad
- ☐ b. abcd a
- ☐ c. aabcd

```
src = [("a",0.25), ("b",0.25), ("c",0.25), ("d",0.25)]
longitud = 5
valor = 0.1094
cumulative_probs = [0]
aux = 0.0
for i in range(len(src)):
    aux += src[i][1]
    cumulative_probs.append(aux)
print('cumulative_probs ->', cumulative_probs)
alpha = 0
beta = 1
mensaje = ''
for i in range(longitud):
    subintervals = []
    for j in range(1, len(cumulative_probs)):
        aux = (alpha + (beta-alpha)*cumulative_probs[j-1],
              alpha + (beta-alpha)*cumulative_probs[j])
        subintervals.append(aux)
    print(subintervals)
    for ind, subint in enumerate(subintervals):
        if subint[0] <= valor < subint[1]:
            mensaje += src[ind][0]
            alpha = subint[0]
            beta = subint[1]
            print(subint[0], '<=' , valor, '<' , subint[1])
            print('mensaje =', mensaje)
            print('Nos quedamos con [', alpha, ',', beta, ')')
            break
    print('¿Cual es el mensaje?', mensaje)
    ...
cumulative_probs -> [0, 0.25, 0.5, 0.75, 1.0]
[(0, 0.25), (0.25, 0.5), (0.5, 0.75), (0.75, 1.0)]
0 <= 0.1094 < 0.25
mensaje = a
Nos quedamos con [ 0 , 0.25 )
[(0.0, 0.0625), (0.0625, 0.125), (0.125, 0.1875), (0.1875, 0.25)]
0.0625 <= 0.1094 < 0.125
mensaje = ab
Nos quedamos con [ 0.0625 , 0.125 )
[(0.0625, 0.078125), (0.078125, 0.09375), (0.09375, 0.109375), (0.109375, 0.125)]
0.109375 <= 0.1094 < 0.125
mensaje = abd
Nos quedamos con [ 0.109375 , 0.125 )
[(0.109375, 0.11328125), (0.11328125, 0.1171875), (0.1171875, 0.12109375), (0.12109375, 0.125)]
0.109375 <= 0.1094 < 0.11328125
mensaje = abda
Nos quedamos con [ 0.109375 , 0.11328125 )
[(0.109375, 0.1103515625), (0.1103515625, 0.111328125), (0.111328125, 0.1123046875), (0.1123046875, 0.11328125)]
0.109375 <= 0.1094 < 0.1103515625
mensaje = abdaa
Nos quedamos con [ 0.109375 , 0.1103515625 )
¿Cual es el mensaje? abdaa
...
```

Pregunta 7

No s'ha respost

encara

Puntuat sobre 1,00

▼ Marca la pregunta

⚙ Edita la pregunta Kraft
palabras a añadir

¿Cuántas palabras de longitud máxima se pueden añadir a un código binario con palabras de longitudes [2,3,3,4,5,5,8,10,10,10]?

Trieu-ne una:

- ☐ a. 379
- ☐ b. 377
- ☐ c. 378

```
longitudes = (2,3,3,4,5,5,8,10,10,10)
sumatorio = 0
for li in longitudes:
    sumatorio += 2**(-li)
veces = 0
while sumatorio < 1:
    veces += 1
    sumatorio += 2**(-10)
print(veces)

'''
377
'''
```