

Discussion Disclosure

Sections of the assignment were discussed with:

- Munro Forgan

What Does The Code Do:

- Finds shortest route by distance
- Highlights route and prints out all roads and the roads distance's
 - Uses one way road correctly
 - Prints without duplicates
 - Follows intersection restrictions
- Highlights the map's articulation points

Sudocode:

A* Search Sudocode:

Input: A weighted graph representing the road network of Auckland, A start node representing the starting intersection, A goal node representing the destination intersection and a heuristic function for each node

Output: the shortest path from the starting intersection/node to the destination intersection/node

Initially while all of the nodes in the graph are unvisited and the fringe has a single element (the starting intersection)

While(the fringe is not empty and the route has not been found){

 Set current node to the item in the fringe with minimal f^*

 Add current node to the visited;

 If(current node is the destination){

 Set endNodeFound to true;

 }

 For(every road starting or ending in node*){

 If(the road is oneway and node* ends there){

 Check next road

 }

 Finds the node at the other end of the road and sets node* to that;

 Calculates the distance travelled to get to node* and sets g^*

 Calculates the remaining distance to the end node from node* and sets H^*

 If(node* has already been visited or the new distance to the end node is greater than $prev^*$ try the next neighbour);

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If(the fringe does not contain node* and the new distance to the endnode is smaller  
than the previous){
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```
    Add <node*,current node,g*,f*> to the fringe
```

```
}
```

```
}
```

```
}
```

Path cost and heuristic estimate

The path cost is calculated by adding the length of the next edge leading to the next node to the cumulative path cost from all previous nodes. The heuristic estimate is calculated by calculating the distance between the two supplied nodes. These two are combined into the heuristic function that creates an admissible monotonic heuristic. This heuristic is admissible because the structure of my heuristic $f = g + h$ is non decreasing and because it is non decreasing no revisit of any node will result in a smaller cost.

Generate set to store articulation points (AP)

findAllArticulationPoints{

Set all nodes depth to infinity;

Set number of subtrees of all nodes to 0;

Initializing a array containing all nodes in the graph as a list of unvisited nodes;

While(there is nodes left that are not visited){

Chooses a random node and sets as root;

Removes root from list of unvisited nodes;

For(all neighbouring roads){

Find node at end of road and set is as neighbour;

Checks if neighbour starts or terminates at root;

Removes neighbour from unvisited;

If(the neighbours depth has not been set){

Call recursive findAllArticulationPoints(neighbour, 1, root)

Increase number of the neighbours subtrees by 1

}

If(the number of the neighbours subtrees is greater than 1){

Add neighbour to AP

}

}

}

}

RecursiveFindArticulationPoints(currentNode, depth, parentNode){

Set depth of currentNode to depth;

Set reachBack to depth;

For(all neighbours of currentNode){

 Determines if the neighbour starts or ended at currentNode;

 If(the neighbour is the parent){

 skip this neighbour;

 }

 Remove neighbour from unvisited nodes;

 If(this neighbour's depth has been set){

 change reachback to the smallest of the two;

 }

 case 2: indirect alternative path: neighbour is an unvisited child in the same sub-tree

 calculate alternative paths of the child, which can also be reached by itself

 childReach = recursiveFindArticulationPoints(neighbour, depth + 1, node);

 set reachback to the smallest of childreach and reachback

 if(childReach is smaller than or the same size as depth){

 add currentNode to articulation points

 }

}

}

Testing

- Choosing straightforward routes the check that the algorithm chooses the correct route
 - Checking that the correct offramps and onramps are used on highways