

Nomenclature :

- ligne double : remplacement d'une variable (e_1, γ_2, \dots) par sa définition (ou inversement)
- ligne simple : application d'une des règles de la sémantique opérationnelle

Simplifications d'écriture

- $e_1 = \text{fun } n \rightarrow \text{if } n = 0 \text{ then } 1 \text{ else } n * fact (n - 1)$
- $e_3 = \text{if } n = 0 \text{ then } 1 \text{ else } n * fact (n - 1)$
- $\gamma_2 = \gamma :: \{fact \mapsto \langle \text{letrec } fact = e_1 \text{ in } e_1, \gamma \rangle\}$
- $\gamma_3 = \gamma_2 :: \{n \mapsto 2\}$
- $\gamma_4 = \gamma_2 :: \{n \mapsto 1\}$
- $\gamma_5 = \gamma_2 :: \{n \mapsto 0\}$

$$\frac{\frac{\frac{[A] \quad \gamma_2 \vdash 2 \Rightarrow 2 \quad \gamma_2 \vdash fact \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle \quad [B] \quad \gamma_2 :: \{n \mapsto 2\} \vdash e_3 \Rightarrow 2}{\gamma :: \{fact \mapsto \langle \text{letrec } fact = e_1 \text{ in } e_1, \gamma \rangle\} \vdash fact \ 2 \Rightarrow 2}}{\gamma \vdash \text{letrec } fact = \text{fun } n \rightarrow \text{if } n = 0 \text{ then } 1 \text{ else } n * fact (n - 1) \text{ in } fact \ 2 \Rightarrow 2}}$$

[A] : évaluation de $fact$

$$\frac{fact \in \gamma_2 \quad \gamma_2(fact) = \langle \text{letrec } fact = e_1 \text{ in } e_1, \gamma \rangle \quad \frac{\frac{\gamma_2 \vdash \text{fun } n \rightarrow e_3 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}{\gamma :: \{fact \mapsto \langle \text{letrec } fact = e_1 \text{ in } e_1, \gamma \rangle\} \vdash e_1 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}{\gamma \vdash \text{letrec } fact = e_1 \text{ in } e_1 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}{\gamma_2 \vdash fact \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}$$

[B] : On déroule un appel récursif

$$\frac{\frac{\frac{\gamma_3 \vdash n \Rightarrow 2 \quad \gamma_3 \vdash 0 \Rightarrow 0 \quad 2 \times 0 \in \text{dom}(=) \quad \text{false} = (2 = 0)}{\gamma_3 \vdash n = 0 \Rightarrow \text{false}} \quad \frac{[C] \quad \gamma_3 \vdash n \Rightarrow 2 \quad \gamma_3 \vdash fact (n - 1) \Rightarrow 1 \quad 2 \times 1 \in \text{dom}(*) \quad 2 = (2 * 1)}{\gamma_3 \vdash n * fact (n - 1) \Rightarrow 2}}{\frac{\gamma_3 \vdash \text{if } n = 0 \text{ then } 1 \text{ else } n * fact (n - 1) \Rightarrow 2}{\gamma_2 :: \{n \mapsto 2\} \vdash e_3 \Rightarrow 2}}$$

$[C]$: Appel de fonction

$$\frac{\frac{\gamma_3 \vdash n \Rightarrow 2 \quad \gamma_3 \vdash 1 \Rightarrow 1 \quad 2 \times 1 \in \text{dom}(-) \quad 1 = (2 - 1)}{\gamma_3 \vdash (n - 1) \Rightarrow 1} \quad \frac{[A'] \quad [B']}{\gamma_3 \vdash \text{fact} \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle \quad \gamma_2 :: \{n \mapsto 1\} \vdash e_3 \Rightarrow 1}$$

$[A']$: évaluation de fact , presque identique à $[A]$

$$\frac{\frac{\frac{\gamma_2 \vdash \text{fun } n \rightarrow e_3 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}{\gamma :: \{\text{fact} \mapsto \langle \text{letrec } \text{fact} = e_1 \text{ in } e_1, \gamma \rangle\} \vdash e_1 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}{\gamma \vdash \text{letrec } \text{fact} = e_1 \text{ in } e_1 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}{\frac{\text{fact} \in \gamma_3 \quad \gamma_3(\text{fact}) = \langle \text{letrec } \text{fact} = e_1 \text{ in } e_1, \gamma \rangle}{\gamma_3 \vdash \text{fact} \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}$$

$[B']$: On déroule un deuxième appel récursif

$$\frac{\frac{\frac{\gamma_4 \vdash n \Rightarrow 1 \quad \gamma_4 \vdash 0 \Rightarrow 0 \quad 1 \times 0 \in \text{dom}(=) \quad \text{false} = (1 = 0)}{\gamma_4 \vdash n = 0 \Rightarrow \text{false}} \quad \frac{[C']}{\gamma_4 \vdash n \Rightarrow 1 \quad \gamma_4 \vdash \text{fact} (n - 1) \Rightarrow 1 \quad 1 \times 1 \in \text{dom}(*) \quad 1 = (1 * 1)}}{\frac{\gamma_4 \vdash \text{if } n = 0 \text{ then } 1 \text{ else } n * \text{fact} (n - 1) \Rightarrow 1}{\gamma_2 :: \{n \mapsto 1\} \vdash e_3 \Rightarrow 1}}$$

$[C']$: Appel de fonction

$$\frac{\frac{\gamma_4 \vdash n \Rightarrow 1 \quad \gamma_4 \vdash 1 \Rightarrow 1 \quad 1 \times 1 \in \text{dom}(-) \quad 0 = (1 - 1)}{\gamma_4 \vdash (n - 1) \Rightarrow 0} \quad \frac{[A''] \quad [B'']}{\gamma_4 \vdash \text{fact} \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle \quad \gamma_2 :: \{n \mapsto 0\} \vdash e_3 \Rightarrow 1}}{\gamma_4 \vdash \text{fact } (n - 1) \Rightarrow 1}$$

$[A'']$: évaluation de *fact*, presque identique à $[A]$ et $[A']$

$$\frac{\frac{\frac{\gamma_2 \vdash \text{fun } n \rightarrow e_3 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}{\gamma :: \{\text{fact} \mapsto \langle \text{letrec } \text{fact} = e_1 \text{ in } e_1, \gamma \rangle\} \vdash e_1 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}{\gamma \vdash \text{letrec } \text{fact} = e_1 \text{ in } e_1 \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}{\frac{\text{fact} \in \gamma_4 \quad \gamma_4(\text{fact}) = \langle \text{letrec } \text{fact} = e_1 \text{ in } e_1, \gamma \rangle}{\gamma_4 \vdash \text{fact} \Rightarrow \langle \text{fun } n \rightarrow e_3, \gamma_2 \rangle}}$$

$[B'']$: On déroule un troisième appel récursif

$$\frac{\frac{\frac{\gamma_5 \vdash n \Rightarrow 0 \quad \gamma_5 \vdash 0 \Rightarrow 0 \quad 0 \times 0 \in \text{dom}(=) \quad \text{true} = (0 = 0)}{\gamma_5 \vdash n = 0 \Rightarrow \text{true}} \quad \gamma_5 \vdash 1 \Rightarrow 1}{\frac{\gamma_5 \vdash \text{if } n = 0 \text{ then } 1 \text{ else } n * \text{fact } (n - 1) \Rightarrow 1}{\gamma_2 :: \{n \mapsto 0\} \vdash e_3 \Rightarrow 1}}$$