

2) Proof L is a non regular language using the Pumping lemma:

a) $L = \{w \in \{0,1\}^*: w = 0^{n!}, n \geq 0\}$

b) $L = \{w \in \{0,1\}^*: w = 0^a 1^b 1^c 0^d, \text{ where } a + b = c + d \text{ and } a, b, c, d \geq 0\}$ [Be careful for (b): check if xyyz works]

c) $L = \{w \in \Sigma^*: w = a^i b^j, \text{ where } i > j, \text{ and } j \geq 0\}$

ii) Given the Context-Free Grammar, answer the following questions:

$$\begin{aligned}A &\rightarrow 1A \mid 1C \mid 0B \mid 00A \\B &\rightarrow 0A \mid 1B \mid 00B \\C &\rightarrow 0C0 \mid 0C1 \mid 1C0 \mid 1C1 \mid \varepsilon\end{aligned}$$

- (a) Give a leftmost derivation for the string 01011001. (3 points)
- (b) Sketch the parse tree corresponding to the derivation you gave in (a). (2 points)
- (c) Demonstrate that the given grammar is ambiguous by showing two more parse trees (apart from the one you already found in (b)) for the same string. (3 points)
- (d) Find a string w of length six such that w has exactly one parse tree in the grammar above. (1 point)
- (e) Design an unambiguous Context Free Grammar for the language represented by the given ambiguous grammar. (1 point)