

$L = \{w \in \{a, b, c, p, q, r, \#\}^* : a^i \# c^n p^{2x} q^y r^z b^j,$   
 where  $i = j + k$ ,  $y = 3x + z$ ,  $n$  is odd and  
 $i, j, k, n, x, y, z \geq 0\}$

$a^i \# c^n p^{2x} q^y r^z b^j$

$\Rightarrow a^{j+k} \# c^n p^{2x} q^{3x+z} r^z b^j$

$\Rightarrow a^j a^k \# c^n p^{2x} q^{3x} q^z r^z b^j$

$S \rightarrow a S b \mid T$

$T \rightarrow A B C$

$A \rightarrow a A c \mid x$

$x \rightarrow \#\# x \mid \#$

$B \rightarrow p p B q q q \mid \epsilon$

$C \rightarrow q C r \mid \epsilon$

CSE331  
Assignment - 02  
CFG Solution  
Faculty: KKP  
(IF you spot any error, please notify me)

3) Design a Context Free Grammar for the Language:

- a)  $L = \{w \in \{a,b,c,p,q,r,\# \}^* : a^i \#^n c^k p^{2x} q^y r^z b^j \text{ where } i=j+k, y=3x+z, n \text{ is odd and } i,j,k,n,x,y,z \geq 0\}$
- b)  $L = \{w \in \{0,1,2\}^* : w = 0^i 2^j 1^k, [\text{where } \dots \text{conditions} \dots] \}$

where...

- ~~i)  $i = k, i, k \geq 1$  and  $j \geq 2$~~
- ii)  $i = 3k, j \text{ is odd and } i, j, k \geq 0$
- iii)  $i$  is a multiple of two,  $k$  is two more than a multiple of 3,  $j = k+i$ , and  $i, j, k \geq 0$
- iv)  $i+j > k$  and  $i, j, k \geq 0$
- v)  $i+k$  is even,  $j = i+k$  and  $j \geq 1$
- c)  $L = \{w \in \{0,1\}^* : \text{the parity of 0s and 1s is different in } w\}$
- d)  $L = \{w \in \{0,1\}^* : \text{the number of 0s and 1s are different in } w\}$   
[Hint: First, try to solve for an equal number of 0s and 1s in  $w$ ]
- e)  $L = \{1^i 0 2^j 1^k \mid i, j, k \geq 0, 3i \geq 4k + 2, j \text{ is not divisible by three}\}$
- f) Recall that for a string  $w$ ,  $|w|$  denotes the length of  $w$ .  $\Sigma = \{0,1\}$   
 $L1 = \{w \in \Sigma^* : w \text{ contains exactly two 1s}\}$   
 $L2 = \{x\#y : x \in \Sigma^*, y \in L1, |x| = |y|\}$   
Construct a CFG for  $L2$ .
- g) Recall that for a string  $w$ ,  $|w|$  denotes the length of  $w$ .  $\Sigma = \{0,1\}$   
 $L1 = \{w \in \Sigma^* : w \text{ contains at least three 1s}\}$   
 $L2 = \{x\#y : x \in (\Sigma\Sigma)^*, y \in L1, |x| = |y|\}$   
Construct a CFG for  $L2$ .

## Assignment - 02

## CFG Solution

Faculty: KKP

(IF you spot any error, please notify me)

## 3) Design a Context Free Grammar for the Language:

a)  $L = \{w \in \{a,b,c,p,q,r,\# \}^* : a^i \#^n c^k p^{2x} q^y r^z b^j \text{ where } i=j+k, y=3x+z, n \text{ is odd and } i,j,k,n,x,y,z \geq 0\}$

b)  $L = \{w \in \{0,1,2\}^* : w = 0^i 2^j 1^k, [\text{where .....conditions.....} ] \}$

where..

i)  ~~$i = k, i, k \geq 1 \text{ and } j \geq 2$~~

ii)  $i = 3k, j \text{ is odd and } i,j,k \geq 0$

iii)  $i \text{ is a multiple of two, } k \text{ is two more than a multiple of 3, } j = k+1, \text{ and } i,j,k \geq 0$

iv)  $i+j > k \text{ and } i,j,k \geq 0$

v)  $i+k \text{ is even, } j = i+k \text{ and } j \geq 1$

c)  $L = \{w \in \{0,1\}^* : \text{the parity of 0s and 1s is different in } w\}$

d)  $L = \{w \in \{0,1\}^* : \text{the number of 0s and 1s are different in } w\}$

[Hint: First, try to solve for an equal number of 0s and 1s in  $w$ ]

e)  $L = \{1^i 02^j 1^k \mid i, j, k \geq 0, 3i \geq 4k + 2, j \text{ is not divisible by three}\}$

f) Recall that for a string  $w$ ,  $|w|$  denotes the length of  $w$ .  $\Sigma = \{0,1\}$

$L_1 = \{w \in \Sigma^* : w \text{ contains exactly two 1s}\}$

$L_2 = \{x\#y : x \in \Sigma^*, y \in L_1, |x| = |y|\}$

Construct a CFG for  $L_2$ .

g) Recall that for a string  $w$ ,  $|w|$  denotes the length of  $w$ .  $\Sigma = \{0,1\}$

$L_1 = \{w \in \Sigma^* : w \text{ contains at least three 1s}\}$

$L_2 = \{x\#y : x \in (\Sigma\Sigma)^*, y \in L_1, |x| = |y|\}$

Construct a CFG for  $L_2$ .

$$L = \{ w \in \{0,1,2\}^* : w = 0^i 2^j 1^k, \text{ where } i=k, i, k \geq 1, j \geq 2 \}$$

$$0^i 2^j 1^k \\ \Rightarrow 0^i 2^j 1^i$$

Solution:

$$\begin{aligned} S &\rightarrow OA1 \\ A &\rightarrow OA1 \mid 22B \\ B &\rightarrow 2B \mid \epsilon \end{aligned} \quad \begin{array}{l} i, k \geq 1 \\ \downarrow \\ j \geq 2 \end{array}$$

0221, 00022111, 022221  $\in L$   
22, 01, 021, 022  $\notin L$

0002222111



Another solution:

$$\begin{aligned} S &\rightarrow OS1 \mid O22A1 \\ A &\rightarrow 2A \mid \epsilon \end{aligned} \quad \begin{array}{l} i, k \geq 1 \\ \downarrow \\ j \geq 2 \end{array}$$

Another Solution:

$$\begin{aligned} S &\rightarrow OS1 \mid OA1 \\ A &\rightarrow 2A \mid 22 \end{aligned} \quad \begin{array}{l} i, k \geq 1 \\ \downarrow \\ j \geq 2 \end{array}$$



$L = \{ w \in \{0,1,2\}^* : w = 0^i 2^j 1^K, \text{ where } i=3K, j \text{ is odd and } i,j,K \geq 0 \}$

$$\begin{matrix} 0^i & 2^j & 1^K \\ \Rightarrow & 0^{3K} & 2^j & 1^K \end{matrix}$$

$K=0$ , 'odd 2s'

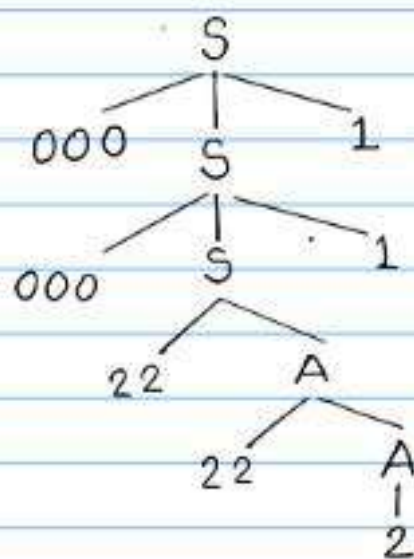
$K=1$ , 000 'odd 2s' 1

$K=2$ , 000000 'odd 2s' 11

Solution:

$S \rightarrow 000S1 \mid A$

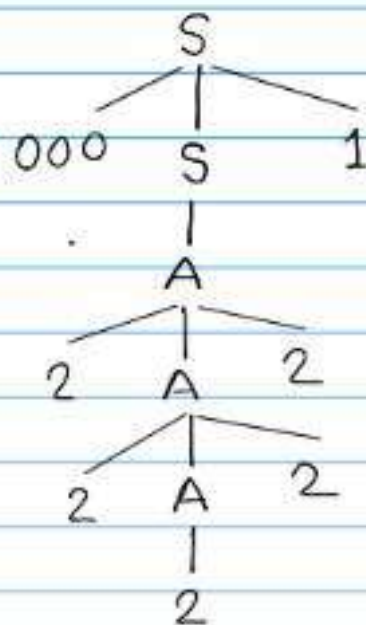
$A \rightarrow 22A12$



Another Solution:

$S \rightarrow 000S1 \mid A$

$A \rightarrow 2A212$



\* It is not mandatory to draw the parse tree in the q's of CFG. I have drawn for the understanding purpose.

$L = \{ w \in \{0,1,2\}^* : w = 0^i 2^j 1^K, i \text{ is multiple of two, } K \text{ is two more than multiple of three, } j = K + i, \text{ and } i, j, K \geq 0 \}$

Be careful with maintaining the same order

$$\begin{aligned}
 & \rightarrow 0^i 2^j 1^K \\
 & \Rightarrow 0^{2i} 2^{3K+2+2i} 1^{3K+2} \\
 & \Rightarrow \underbrace{0^{2i} 2^2 2^{2i}}_A \underbrace{2^{3K} 1^2 1^{3K}}_B
 \end{aligned}$$

$$S \rightarrow AB$$

$$A \rightarrow 00A22 \mid 22$$

$$B \rightarrow 222B111 \mid 11$$

Another solution:

$$\begin{aligned}
 & 0^i 2^j 1^K \\
 & \Rightarrow 0^{2i} 2^{3K+2+2i} 1^{3K+2} \\
 & \Rightarrow \underbrace{0^{2i} 2^{2i}}_A \underbrace{2^2}_B \underbrace{2^{3K} 1^{3K}}_C \underbrace{1^2}_D
 \end{aligned}$$

$$S \rightarrow ABCD$$

$$A \rightarrow 00A22 \mid \epsilon$$

$$B \rightarrow 22$$

$$C \rightarrow 222C111 \mid \epsilon$$

$$D \rightarrow 11$$

some 0s... some 2s... some 1s

some 0s... some 2s... some 1s

Be careful with  
maintaining the same order

Another solution:

$$\begin{aligned}
 & \rightarrow 0^i 2^j 1^K \\
 \Rightarrow & 0^{2i} 2^{3K+2+2i} 1^{3K+2} \\
 \Rightarrow & 0^{2i} 2^{3K} 2^2 2^{2i} 1^{3K} 1^2 \\
 \Rightarrow & 0^{2i} 2^{2i} 2^{3K} 2^2 2^2 1^2 1^{3K}
 \end{aligned}$$

$$S \rightarrow AB$$

$$A \rightarrow 00A22 \mid \epsilon$$

$$B \rightarrow 222B111 \mid 2211$$



$$L = \{ w \in \{0,1,2\}^* : w = 0^i 2^j 1^K, \text{ where } i+j > K \text{ and } i,j,K \geq 0 \}$$

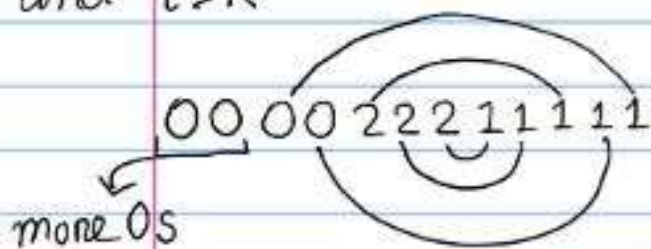
Let's first solve for  $i+j = K$

$$\begin{aligned} & 0^i 2^j 1^K \\ \Rightarrow & 0^i 2^j 1^{i+j} \\ \Rightarrow & 0^i 2^j 1^j 1^i \end{aligned}$$

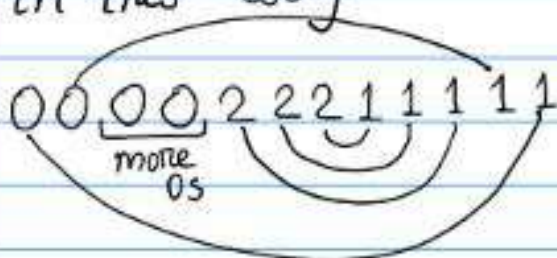
Now since  $i+j > K$ ,

$$0^i 2^j 1^j 1^i$$

there could be more  
0s and equal 2s & 1s  
means, in  $i+j > K$ ,  $j=K$   
and  $i > K$



you can also think  
in this way



$$\text{or } 0^i 2^j 1^j 1^i$$

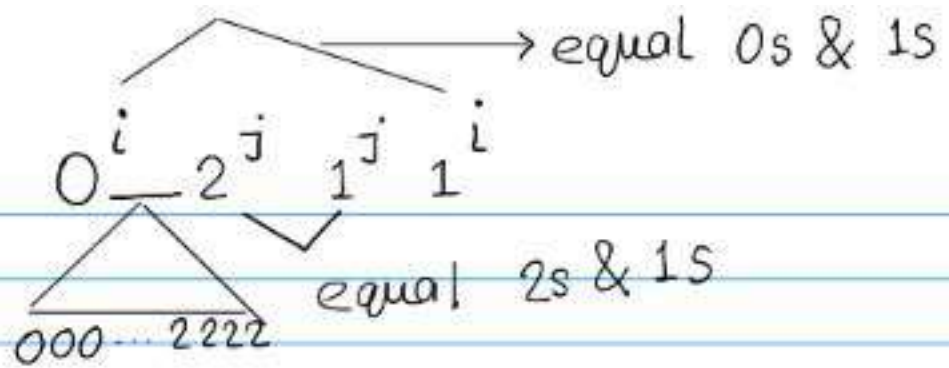
there could be more  
2s and equal 0s & 1s  
means, in  $i+j > K$ ,  
 $i=K$  and  $j > K$



or, there could be  
more 0s and 2s  
than 1s both,  
means in  $i+j > K$   
 $i > K$  and  $j > K$



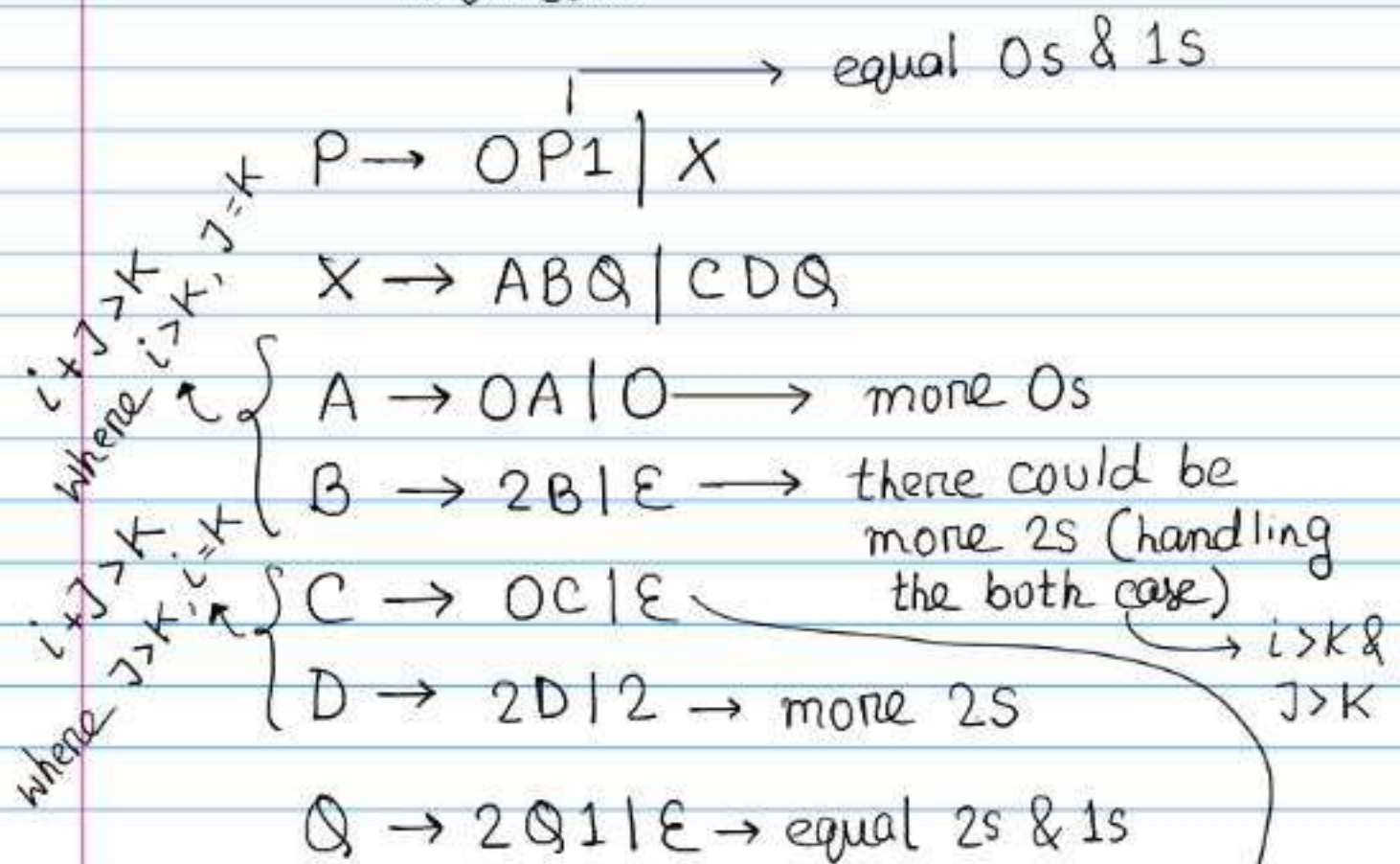
Too Summarize,



↳ either more 0s

↳ either more 2s

↳ or both



$P \rightarrow OP1 | X$

$X \rightarrow ABQ | CDQ$

$A \rightarrow OA | O \rightarrow$  more 0s

$B \rightarrow 2B | \epsilon \rightarrow$  there could be more 2s (handling the both case)

$C \rightarrow OC | \epsilon$

$D \rightarrow 2D | 2 \rightarrow$  more 2s

$Q \rightarrow 2Q1 | \epsilon \rightarrow$  equal 2s & 1s

this also handle the both case, can be skipped, since we already have handled the case previously.

$L = \{w \in \{0,1\}^* : \text{parity of number of 0s and 1s is different}\}$

Case 1: even 0s and odd 1s

Case 2: odd 0s and even 1s

So, the problem can be boiled down into  $L = \{\text{length of } w \text{ is odd}\}$

$S \rightarrow 00S \mid 01S \mid 10S \mid 11S \mid 0 \mid 1$

This can also be written as

$S \rightarrow XXS \mid X$

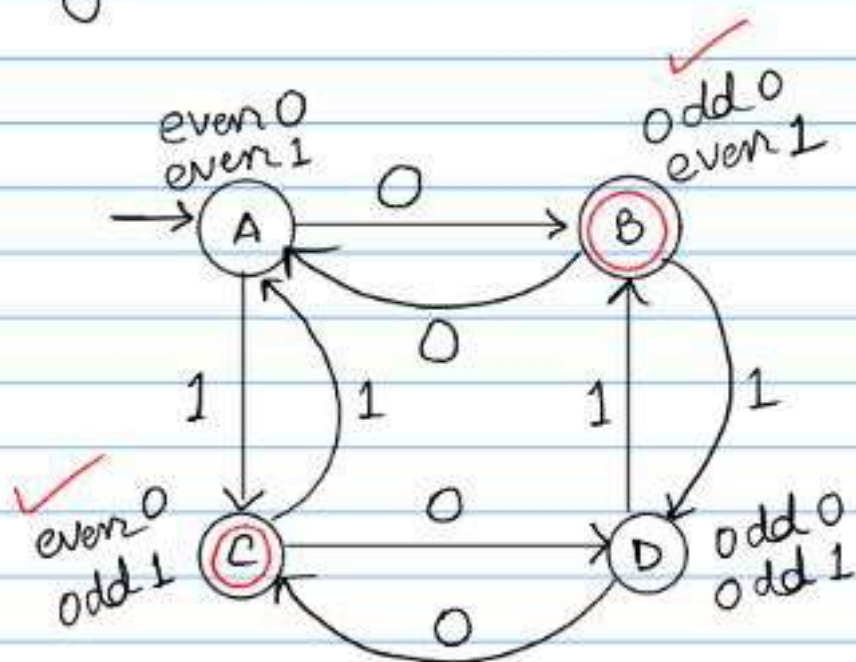
$X \rightarrow 0 \mid 1$

Another solution:

$S \rightarrow 0S0 \mid 0S1 \mid 1S0 \mid 1S1 \mid 0 \mid 1$

Another solution:

if you couldn't figure out the previous idea, then no worry. You may also recall the following DFA we had done in the class:



So another solution can be

$$A \rightarrow 0B \mid 1C$$

$$B \rightarrow 0A \mid 1D \mid \epsilon$$

$$C \rightarrow 0D \mid 1A \mid \epsilon$$

$$D \rightarrow 0C \mid 1B$$



$L1 = \{ w \in \{0,1\}^* : \text{the number of 0s and 1s are different in } w. \}$

Before solving  $L1$ , first we try to solve

$L2 = \{ w \in \{0,1\}^* : w \text{ contains equal numbers of 0s and 1s} \}$

$S \rightarrow 0S1 \mid 1S0 \mid \epsilon$

→ 0s & 1s are paired in pairs, so both the count of 0s & 1s will be same

However, this solution is partially correct. For example, 0110 can't be parsed.

If we take a string,  $w \in L2$ , and if it has equal numbers of

001011 10

↓

We can divide the string into two substrings, having equal 0s and 1s.

0s and 1s, then it means, in  $w$ , there are one or more substring in  $w$ , having equal 0s & 1s.


Now, recall the solution for valid parentheses. and let's fix the grammar.

$S \rightarrow 0S1 \mid 1S0 \mid SS \mid \epsilon$


□ Draw parse tree for 110101000011



Now, let's come back to our original question.  
let's consider a string having equal 0s & 1s

each block having equal { 


0s & 1s Now, if there is more 0s then having at least one additional 0 will be enough.



So, we can write

$T \rightarrow SOS$  where S produce equal 0s and 1s  
 $S \rightarrow 0S1 \mid 1S1 \mid SS \mid \epsilon$

However, there could be more than one additional 0 than 1s. So, those 0s should be parsed as well.



$T \rightarrow SOS$   
 $S \rightarrow 0S1 \mid 1S0 \mid SS \mid 0S \mid \epsilon$

However, we have handled only one case - more 0s than 1s. There could be more 1s than 0s also.

So, the final solution :

$$S \rightarrow A \mid B$$

$$A \rightarrow x0x$$

$$x \rightarrow 0x1 \mid 1x0 \mid xx \mid 0x \mid \epsilon$$

$$B \rightarrow y1y$$

$$y \rightarrow 0y1 \mid 1y0 \mid yy \mid 1y \mid \epsilon$$

Recall that for a string  $w$ ,  $|w|$  denotes the length of  $w$ .  $\Sigma = \{0,1\}$

$L_1 = \{w \in \Sigma^* : w \text{ contains exactly two 1s}\}$

$L_2 = \{x \# y : x \in \Sigma^*, y \in L_1, |x| = |y|\}$

Construct a CFG for  $L_2$ .

Before solving this problem, let's try to solve a few similar problems.

$$L = \{ w_1 \# w_2 \mid w_1, w_2 \in \{0,1\}^* \text{ and } |w_1| = |w_2| \}$$

1001#0010

$$S \rightarrow OSO \mid OS1 \mid 1SO \mid 1S1 \mid \#$$

We can also write it as

$$S \rightarrow XSX \mid \#$$

$$X \rightarrow 0 \mid 1$$

Now let's say,

$$L_1 = \{ w_1 \# w_2 \mid w_1 \in \{0,1\}^*, w_2 \in L_2 \text{ and } |w_1| = |w_2| \}$$

$$L_2 = \{ w \in \{0,1\}^* : w \text{ is even} \}$$

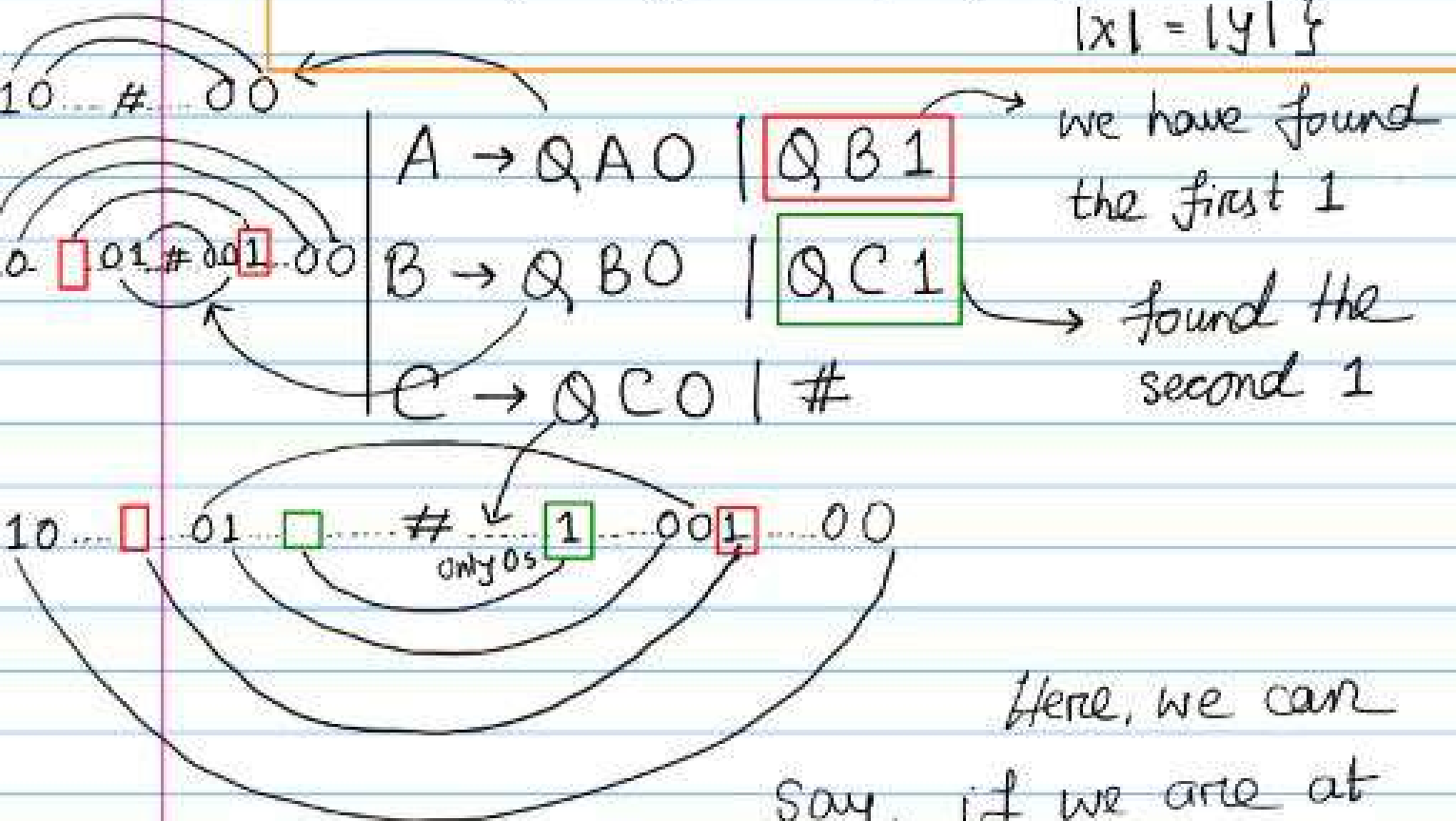
$$S \rightarrow XXSXX \mid \#$$

$$X \rightarrow 0 \mid 1$$

If you understood the previous two solutions, then we try to solve our original question.

$$L1 = \{w \in \{0,1\}^* : w \text{ contains exactly two 1s}\}$$

$$L2 = \{x \# y, x \in \{0,1\}^*, y \in L1 \text{ and } |x| = |y|\}$$



Here, we can say, if we are at production rule A, then have seen no 1 in y, if we are at the production rule B, then we have found exactly one 1 in y. Next, if we are at rule C, then we have seen exactly two 1s in the y.



$$L = \{ w \in \{0,1\}^* : 0^i 1^K \text{ where } i, K \geq 0 \text{ and } 3i \geq 4K + 2 \}$$

This means, if we have  $i$  amount of 0s and  $K$  amount of 1s, then  $(3 \times \text{total 0s})$  should be greater than or equal to  $(4 \times \text{total 1s} + 2)$

So, let's first figure out, what is the minimum amount of 0s we need to have for  $K = 0, 1, 2, \dots$  satisfying the condition.

if  $K = 0, i \geq 1$

$K = 1, i \geq 2$

$K = 2, i \geq 4$

$K = 3, i \geq 5$

$K = 4, i \geq 6$

$K = 5, i \geq 8$

$K = 6, i \geq 9$

$$3i \geq 4K + 2$$

$$\Rightarrow i \geq \left\lceil \frac{4K + 2}{3} \right\rceil$$

Now, can you find any pattern? Think in respect of  $K \% 3$ . [since,  $3i$ ]

$$K \% 3$$

$$\begin{aligned}
 K=0 &\Rightarrow K73=0, i \geq 1 \\
 K=1 &\Rightarrow K73=1, i \geq 2 \\
 K=2 &\Rightarrow K73=2, i \geq 4 \\
 K=3 &\Rightarrow K73=0, i \geq 5 \\
 K=4 &\Rightarrow K73=1, i \geq 6 \\
 K=5 &\Rightarrow K73=2, i \geq 8 \\
 K=6 &\Rightarrow K73=0, i \geq 9 \\
 K=7 &\Rightarrow K73=1, i \geq 10 \\
 K=8 &\Rightarrow K73=2, i \geq 12
 \end{aligned}$$

if we do  $K73$ , then we have three patterns.  
The minimum length of strings for each pattern

$$0^i 1^K$$

0

001

000011

$$K=0, K73=0 \quad K=1, K73=1 \quad K=2, K73=2$$

Now, see what amount of 0s and 1s  
you need to go the next pattern in same  $K73$

00000111

$$K=3, K73=0$$

0000001111

$$K=4, K73=1$$

00000000111111

$$K=5, K73=2$$

So, in each pattern we see, we can  
jump to the next string by adding 0000111

$$L = \{ w \in \{0,1\}^* : 0^i 1^k, \text{ where } i, k \geq 0 \}$$

So, now, if the condition was  $3i = 4k+2$  then,

$$S \rightarrow OA \mid 00A1 \mid 0000A11$$

$$A \rightarrow 0000A111 \mid \epsilon$$

Now, since, the condition is  $3i \geq 4k+2$ , hence, we will have some additional 0s as well. So,

$$S \rightarrow 0^{\overbrace{0}^{K73=0}} S \mid OA \mid 00A1 \mid 0000A11$$

$$A \rightarrow 0000A111 \mid \epsilon$$

Another Approach

Based on the increment of 0s  $\rightarrow i = 0, 1, 2, 4, 5, 6, 8$

$$\text{if } 3i = 4k+1$$

$$K = 0, 1, 2, 3, 4, 5, 6$$

$$S \rightarrow OA \mid \epsilon$$

$$A \rightarrow OB1 \mid \epsilon$$

$$B \rightarrow OOC1 \mid \epsilon$$

$$C \rightarrow OA1 \mid \epsilon$$

Now, for  $3i \geq 4k+1$

$$S \rightarrow 0S \mid OA \mid \epsilon$$

$$A \rightarrow OB1 \mid \epsilon$$

$$B \rightarrow OOC1 \mid \epsilon$$

$$C \rightarrow OA1 \mid \epsilon$$