# Non Linear Optimization Project

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### 1 Introduction

Implementation of Linear Programming Techniques for Solving Constrained Optimization Problems with a Simple User Interface.

The following report discuss the development of a Python program that can solve constrained optimization problems using Linear Programming (LP) techniques.

The program will allow the user to choose between:

- Graphical method
- Simplex method

The program has multiple functionalities

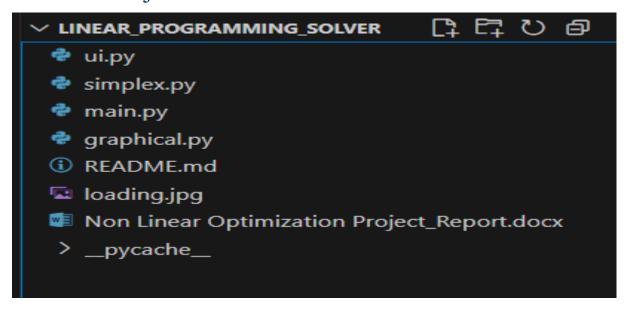
- Entering the problem data
- Determine the optimal solution
- Display results either visually or textually

Libraries used throughout the project:

- Numpy
- Matplotlib.pyplot
- Tkinter for User Interface

We will go through the implementation of the LP methods, User interface and finally a guide for how to run the code as well as using the program followed along with screenshots.

# 2 Project Structure



Ui.py: A simple user interface

Simplex.py: the simplex method algorithm

Graphical.py: the graphical method algorithm

Main.py: the entry point that launches the Tkinter GUI app.

README.md: provides a guide, explaining purpose, required libraries, project structure, how it works and features

## 3 Graphical Method

### 3.1 Logic

The graphical method for solving linear programming problems is a powerful visualization tool for problems with two variables. By plotting constraints and identifying the feasible region, one can find the optimal solution by evaluating the objective function at the corner points. This method not only provides insights into the problem but also helps in understanding the impact of each constraint on the solution. However, for problems with more than two variables, other techniques such as the Simplex method are required.

#### Needed inputs:

- Type of optimization (Max/Min)
- Objective functions and constraints

This program assumes that the non negativity constraint is always true.

### 3.2 Program Flow Analysis

```
import numpy as np
import matplotlib.pyplot as pit

def format_constraint(a, b, c, sense):
    """Formats constraint equation as a string"""
    parts = []
    if a != 0:
        parts.append(f"{a:.2f}x" if abs(a) != 1 else ("x" if a > 0 else "-x"))
    if b != 0:
        sign = "+" if (b > 0 and parts) else ("-" if b < 0 else "")
        coeff = abs(b) if abs(b) != 1 else ""
        parts.append(f" {sign} {coeff}y" if coeff else f" {sign} y")
    return f"{''.join(parts)} {sense} {c:.2f}"</pre>
```

A helper function used lately in code to plot a legend with the objective function and constraints.

```
def graphical_solver(opt_type, obj_coeffs, constraints):
```

The main function

```
def satisfies(x, y, a, b, c, sense):
    if sense -- "<-":
        return a * x + b * y <= c + 1e-9
    elif sense -- ">":
        return a * x + b * y >= c + 1e-9
    elif sense -- ">":
        return a * x + b * y >= c + 1e-9
    elif sense -- "-":
        return np.isclose(a * x + b * y, c, atol=1e-9)
    return False

cons_funcs = [lambda X, Y, a=a, b=b, c=c, s=sense: satisfies(X, Y, a, b, c, s) for a, b, c, sense in constraints]
```

Constraint checker function. (Note: tolerance of 1e-9 to handle floating point rounding errors)

#### 1. Find Intersection Points

Solve the system of equations for each pair to find intersection points.

We have a loop to access every unique point stored as tuples

Note: a1, b1, c1, = constraints[i]  $\rightarrow$  we ignored the sense variable as it is not needed here.

We check if lines aren't parallel then we solve for (x,y)

Before appending the intersection point we should ensure the non negativity constraint and that the point satisfies every single constraint.

#### 2. Find Axis Intercepts

Assume the coefficient of y=0 and find x then the coefficient of x=0 and find y finally append each point to points list after making sure it satisfies every constraint.

#### 3. Remove Duplicates

Start with empty list unique points and append points to it if no match found

```
opt_val - None
opt_pts = []
if not unique points:
   print ("Wo feasible area solution")
status = "infeasible"
# Check imbounded solution
elif status != "Infemsible"
    obj_dir = np.array(obj_coeffs, dtype=float)
    if opt_type — 'min':
obj_dir = obj_dir
     limiting constraints = 0
         (a, b, c, sense) in constraints:
         n = np.array([a, b], dtype=float)
         proj = np.dot(n, obj dir)
if sense == "c=" and proj > 1e-9:
              limiting constraints +- 1
          ell! sesse - ">-" und proj < -te-9:
             limiting constraints +- 1
             limiting constraints - 1
     if limiting constraints - 0:
```

1. Check infeasibility, if no unique\_points then there is no points that satisfies all the constraints and no solution is found

Check unbounded solution by examining constraint directions relative to objective

#### How this was done?

- A variable <u>limiting constraints</u> is set to zero. This counter will track the number of constraints that restrict movement in the objective function's improvement direction.
- Each constraint is represented as (a, b, c, sense), corresponding to: ax+by (sense) c

For each constraint:

#### 1. Convert to Normal Vector:

The coefficients (a, b) are stored in vector form (n).

#### 2. Project the Objective Direction:

The projection of the constraint's normal vector onto the objective direction vector is computed:  $proj = n \cdot obj \ dir$ 

This determines whether movement in the improvement direction will increase or decrease the left-hand side (LHS) of the constraint.

- Determine limiting behavior:
  - o For a <= constraint: If proj > 1e-9, movement in the objective direction increases the LHS, so this constraint **limits** progress.
  - o For a >= constraint: If proj < -1e-9, movement in the objective direction decreases the LHS, so this constraint **limits** progress.

- For an = constraint: This constraint always limits movement because the solution must remain exactly on the defined line.
- **Update Counter**, if the constraint limits movement, increment limiting\_constraints by 1.
- If, after checking all constraints, limiting\_constraints remains **zero**, it means there are no constraints restricting movement in the improvement direction. In this case, the solution is classified as **unbounded**.

```
# If optimal, compute optimum
if status == "optimal":
    values = [(p, obj_coeffs[0] * p[0] + obj_coeffs[1] * p[1]) for p in unique points]
if opt_type == 'max':
    opt_val = max(v for _, v in values)
else:
    opt_val = min(v for _, v in values)
opt_pts = [p for p, v in values if np.isclose(v, opt_val)]
    # multiple optimal solutions detection
if len(opt_pts) > 1:
    p1, p2 = opt_pts[0], opt_pts[-1]
    print(f"Multiple optimal solutions between [p1] and (p2)")
    print(f"All optimal points: P(λ) = λ*(p1) + (1 = λ)*(p2), 0 ≤ λ ≤ 1")
```

if optimal then compute optimum based on the optimization type (max/min) by calculating z among all feasible points and store the value in opt val

The list opt\_pts is created by selecting all feasible points whose objective value is numerically equal (within floating-point tolerance) to opt\_val.

If the list opt pts contains more than one point, it means multiple optimal solutions exist.

#### In such cases:

- The algorithm selects the first (p1) and last (p2) optimal points from opt pts.
- It reports that multiple optima exist between these two points.
- It also provides the **parametric form** of all optimal points:

$$Z = \lambda p1 + (1 - \lambda) p2$$
, where  $0 \le \lambda \le 1$ 

- 1. Sets plot boundaries based on solution points or defaults
- 2. Create grid for plotting

np.linspace(start,stop,num) -> Generates evenly spaced numbers over a specified interval.

np.meshgrid( $x_{vals,y_{vals}}$ ) Converts one-dimensional coordinate arrays into a full two-dimensional grid of points. Concept: The result is a grid of coordinates (x, y), where every x from x vals is paired with every y from y vals.

3. Compute feasible region by applying all constraints to the grid then plot it using plt.contourf (filled contour plots)

```
constraint_colors - plt.cm.tabl@(np.linspace(0, 1, len(constraints))) # Creates a color spectrum
legend handles = []
for i, (a, b, c, sense) in enumerate(constraints):
    label = f'C(i+1): (format_constraint(a, b, c, sense))'
        plt.plot(x_vals, (c - a * x_vals) / b,
                color-color,
                label=f"Constraint {i+1}: {a: 1f}x + (b: 1f}y {sense} (c: 1f)")
        plt.axvline(c / a,
               linestyle=
                color-color,
                label=f"Constraint (i+1): (a::1f)x (sense) (c::1f)")
    legend_handles.append(plt.Line2D([], [], color=color, linestyle= -- , linewidth=2,label=label))
obj_label - f"Obj: z - {obj_coeffs[0]:.2f}x + {obj_coeffs[1]:.2f}y ({opt_type})"
legend_handles.append(
    plt.tine20([], [], color='hlack', linestyle='-', linewidth=2,label=obj_label))
opt_color - 'greem' if opt_type -- 'max' else 'gold' opt_label = f"Optimal ({opt_type}))"
legend_handles.append(plt.Line2D([], [], color=opt_color, marker='o', linestyle='Mone',markersize=10, label=opt_label))
```

This code block is responsible for plotting the constraint lines, the objective function, and the optimal point.

- 1. Each constraint is drawn as a dashed line with a unique color for clarity.
- 2. The objective function is shown as a solid black line.
- 3. The optimal solution is highlighted with a colored point (green for maximization, yellow for minimization).
- 4. A legend is created to explain each line and the optimal point, making the plot easy to read and interpret. (legend\_handles)

```
# All Intersection points
for p in unique_points:
    plt.plot(p[0], p[1], 'ro')
    plt.text(p[0] + 0.1, p[1] + 0.1, f"({p[0]:.2f},{p[1]:.2f})", fontsize=8)

# Optimal points
for p in opt_pts:
    plt.plot(p[0], p[1], 'go' if opt_type == 'max' else 'yo', markersize=10)

plt.xlabel('x')
    plt.ylabel('y')
    plt.ylabel('y')
    plt.ylim(left=0) # Start x-axis at 0
    plt.ylim(bottom=0) # Start y-axis at 0
    plt.title(f"LP Graphical Solution - {status.capitalize()}")
    plt.tegend(handles=legend handles, loc='upper right', bbox_to_anchor=(1.3, 1))
    plt.tight_layout()
    plt.grid(True)
    plt.show()

return {'status': status, 'opt_value': opt_val, 'opt_points': opt_pts}
```

- 1. Marks all feasible corner points with red dots and coordinates and adds a text label with the point coordinates near the point
- 2. Highlights optimal points with green (max) or yellow (min) markers
- 3. Adds labels, title, grid, legend and displays the plot
- 4. Finally returns solution information as a dictionary

#### 3.3 Conclusion

The code successfully implements the graphical method for linear programming by calculating intersection points, checking feasibility, and determining the optimal solution. It also distinguishes between maximization and minimization problems, highlights optimal points on the graph, and even detects multiple optimal solutions when they occur. Overall, the code is efficient, clear, and provides both numerical and visual outputs that make the results easy to understand.

# 4 Simplex Method

### 4.1 Logic

The first step involved in the simplex method is to construct an auxiliary prob lem by introducing certain variables known as artificial variables into the standard form of the linear programming problem. The primary aim of adding the artificial variables is to bring the resulting auxiliary problem into a canonical form from which its basic feasible solution can be obtained immediately. Starting from this canonical form, the optimal solution of the original linear programming problem is sought in two phases. The first phase is intended to find a basic feasible solution to the original linear programming problem. It consists of a sequence of pivot operations that produces a succession of different canonical forms from which the optimal solution of the auxiliary problem can be found. This also enables us to find a basic feasible solution, if one exists, to the original linear programming problem. The second phase is intended to find the optimal solution to the original linear programming problem. It consists of a second sequence of pivot operations that enables us to move from one basic feasible solution to the next of the original linear programming problem. In this process, the optimal solution to the problem, if one exists, will be identified. ~reference: engineering-optimization-theory-and-practice

#### Needed inputs:

- Type of optimization (Max/Min)
- Objective functions and constraints

Note: the following program handles only the  $\leq$  constraints using the slack variables.

The surplus and artificial variables for  $[\geq, =]$  constraints aren't covere

### 4.2 Program Flow Analysis

A container to hold the result of the simplex algorithm.

#### Attributes:

- o **status:** String like "optimal", "unbounded", "infeasible", or "error".
- o x: decision variables
- o **z:** The optimal objective value.
- o **message:** Extra details or error messages.

```
def __repr__(self):
    return f"SimplexResult(status=(self.status!r), x=(self.x), z=(self.z), message=(self.message!r))"
```

string representation of the result when printed.

```
class SimplexSolver:

   def __init__(self, tol=1e-9, max_iter=10_000):
        self.tol = tol
        self.max_iter = max_iter

   def __nivot(self__T__row__col):
```

Initializes the solver

### **Utilities**

```
def _pivot(self, T, row, col):
    pivot = T[row, col] #pivot element
    T[row, :] /= pivot #Normalize the pivot row
    m, n = T.shape
    for r in range(m): #Gaussian elemination to update the table
        if r != row:
            T[r, :] -= T[r, col] * T[row, :]
```

#### 1.Pivot Operation:

- Normalize the pivot row
- o Eliminate the pivot column (Set all other entries to zero)
- Update the table

```
def _choose_entering(self, row):
    candidates = np.where(row < -self.tol)[0]
    if candidates.size -= 0:
        return None
    return candidates[np.argwin(row[candidates])]</pre>
```

#### 2. Choose the pivot column:

- o Pick the smallest negative coefficient at the objective row.
- o If ne negative coefficients then status is already optimal

```
def _choose_leaving(self, T, col):
    rhs = T[:-1, -1]
    col_vals = T[:-1, col]
    mask = col_vals > self.tol

if not np.any(mask):
    return None # Unbounded
    ratios = rhs[mask] / col_vals[mask]
    idx = np.argmin(ratios)

leaving_rows = np.where(mask)[0]
    return leaving_rows[idx]
```

#### 3. Choose the pivot row:

- Minimum ratio test (pick the smallest positive ratio) to determine which variable will leave the basis
- o If no positive entries then problem is unbounded

#### Constructs the initial simplex tableau:

- A: constraint matrix.
  - o n: number of constraints
  - o m: number of variables
- Add slack variables
  - o **np.eye(m)** creates an m×m identity matrix (slack variables).
  - o **np.hstack([...])** stacks A and the identity side by side.
- Objective row
  - o Since we're maximizing, we put -c (because simplex looks for negative coefficients to improve the solution).
  - $\circ$  Slack variables don't appear in the objective, so their coefficients = 0.
- Tableau layout:
  - o Rows = constraints + 1 objective row  $\rightarrow$  m+1.
  - Columns = variables  $(n+m) + RHS \rightarrow n+m+1$ .

#### Filling the tableau:

- 1. Constraint rows (T[:-1, :n+m] = full A)
  - Copies the constraint matrix with slacks.
- 2. **Right-hand side** (T[:-1, -1] = b)
  - Stores constraint bounds.
- 3. **Objective row** (T[-1, :n+m] = c full)
  - Stores the negated objective coefficients.
- 4. The last cell T[-1, -1] (bottom-right) is 0 at the start.

Conclusion; \_build\_tableau converts a human-readable LP into a matrix form where simplex operations (pivoting, entering/leaving, ratio test) can be applied.

```
def _optimize_tableau(self, T, basis):
    iters = 0
    while iters < self.max_iter:
        iters += 1
        obj_row = T[-1, :-1]
        col = self._choose_entering(obj_row)
        if col is None:
            return "optimal", basis

        row = self._choose_leaving(T, col)
        if row is None:
            return "unbounded", basis

        self._pivot(T, row, col)
        basis[row] = col

        return "iteration_limit", basis</pre>
```

Runs simplex iterations:

- Choose entering
- Choose leaving
- Pivot operation
- o Return optimal, unbounded or iteration limit

```
def solve(self, c, A, b, maximize=True):
    c = up.array(c, dtype=float).flatten()
    A = up.array(A, dtype=float)
    b = up.array(b, dtype=float).flatten()

if A.shape[e] != b.shape[e]:
    return SimplexResult("error", message="Number of constraints in A and b don't match.")

if A.shape[i] != c.shape[o]:
    return SimplexResult("error", message="Objective length must equal number of variables.")
```

Convert inputs to Numpy arrays and validates input dimensions

```
# Convert minimization to maximization
c_eff = c.copy()
if not maximize:
    c_eff = -c_eff
```

In case of minimization problem we negate the coefficient of the objective function and treat it like a maximization problem.

```
# Enter b >= 0
for i is range(len(b)):
    if b[i] < 0:
        return SimplexMesult("error", message="All b values must be non-negative for <= constraints.")

# Build and solve tabless
T, basis = self._build_tableau(A, b, c_eff)
    status, basis = self._optimize_tableau(T, basis)

If status == "unbounded":
    return SimplexMesult("unbounded", message="Objective is unbounded.")

if status != "optimal":
    return SimplexMesult("error", message=f"Simplex did not converge: (status)")</pre>
```

```
# Extract solution
m, n = A.shape[0], A.shape[1]
x = np.zeros(n)

# Find values of basic variables
for i in range(m):
    bj = basis[i]
    if bj < n:
        x[bj] = T[i, -1]

z = T[-1, -1]
if not maximize:
    z = -z

return SimplexResult("optimal", x=x, z=z)</pre>
```

Extracts the primal solution x from the tableau.

Note: Slack variables are ignored we only need the decision variables

Start with zeros for all decision variables.

For each row in the tableau:

If the basis variable is an original variable, copy its value from RHS into x.

If it's a slack variable, ignore it.

The result is the optimal solution vector x

Finally adjust the objective value for minimization and return the result.

#### 4.3 Conclusion

The code implements a **basic simplex algorithm solver** in a structured and modular way. The process begins with the *SimplexResult class*, which encapsulates the solver's outputs, ensuring that solutions, objective values, and error states are clearly reported. The *SimplexSolver* class then drives the solution process by:

- 1. **Building the initial tableau** (\_build\_tableau) with slack variables, ensuring constraints are properly represented in standard form.
- 2. **Iteratively optimizing** the tableau (\_optimize\_tableau) through pivot operations, where entering and leaving variables are systematically chosen to improve the objective.
- 3. **Extracting the final solution** from the tableau once optimality, unboundedness, or iteration limits are reached.

Each helper method—such as \_pivot, \_choose\_entering, and \_choose\_leaving performs a focused task, reflecting the mathematical operations of the simplex algorithm. The solve method integrates these components, handling input validation, converting minimization to maximization, and ensuring feasibility of constraints before optimization.

### 5 User Interface

### 5.1 Code Flow Analysis

```
import tkinter as tk
from tkinter import ttk, messagebox
from tkinter import *
import numpy as np
from PIL import Image, ImageTk
from graphical import graphical_solver
from simplex import SimplexSolver
```

#### **Importing Required Libraries**

- o tkinter: Python's built-in GUI library.
- o ttk: Themed Tkinter widgets (for a modern look).
- o PIL (Pillow): For loading and displaying images.
- o graphical\_solver : Custom module for solving LP problems graphically.
- o SimplexSolver: Custom class for solving LP problems using the Simplex method.

```
class LPSolverApp:
    def __init__(self, root):
        self.root = root
        self.root.title("Linear Programming Solver")
        self.root.geometry("900x900")
        image=Image.open('loading.jpg')
        self.icon=ImageTk.PhotoImage(image)
        self.root.iconphoto(True,self.icon)

# Create notebook for different solvers
        self.notebook = ttk.Notebook(root)
        self.notebook = ttk.Notebook(root)
        self.notebook,pack(fill='both' ,expand=True)

# Graphical Solver Tab
        self.create_graphical_tab()

# Simplex Solver Tab
        self.create_simplex_tab()
```

- -root is the main application window.
- -Sets the title and size of the window.
- -Loads an image (loading.jpg) to use as the app icon.
- -A tabbed interface with two tabs:
  - o Graphical
  - o Simplex

### Graphical Solver Tab Interface

```
tab = ttk.Frame(self.notebook)
self.notebook.add(tab, text="Graphical Solver")

# Problem Type
ttk.tabe(tab, toxt="problem Type:").grid(row=0, column=0, padx=5, pady=5, sticky=tk.W)
self.graphical_opt_type = tk.StringNar(value="max")
ttk.tabe(tab, toxt="rowblem Type:").grid(row=0, column=0, padx=5, pady=5, sticky=tk.W)
ttk.tabe(stom(tab, text="maximize", variable=self.graphical_opt_type, value="max").grid(row=0, column=1,sticky=tk.W)
ttk.tabe(tok, text="virinize", variable=self.graphical_opt_type, value="max").grid(row=0, column=2,sticky=tk.W)

# Objective Function

ttk.tabe(tab, text="objective Coefficients:").grid(row=1, column=0,padx=5, pady=5,sticky=tk.W)

# tkk.tabe(tab, text="y").grid(row=1, column=2,sticky=tk.W)

# self.obj x = tk.toublevar(value=1)
# self.obj x = tk.toublevar(value=1)
# self.obj y = tk.toublevar(value=1)
# ttk.tintry(tab, textvariable=self.obj x, width=5).grid(row=1, column=1,sticky=tk.W)

# Constraints Frame = ttk.tabe(frame(tab, text="Constraints")
# constraints frame = ttk.tabe(frow=2, column=0, columnpan=5, padx=10, pady=10, sticky=tk.W)

# Constraints frame, text="v").grid(row=0, column=2, padx=5, pady=5)
# ttk.tabe(constraints frame, text="v").grid(row=0, column=2, padx=5, pady=5)
# ttk.tabe(constraints frame, text="v").grid(row=0, column=0, padx=5, pady=5)
# ttk.tabe(constraints frame, text=
```

```
self.constraint_entries = []
for i in range(3):
    self.add_constraint_row(constraints_frame, i+1)

### Add/#conver constraint_button
button frame = tlk.frame(constraints_frame)
button_frame.grid(row-i0, column-0, columnspan-7, pady=5)
ttk.button(button_frame, text="Add Constraint", command=lumbda: self.add_constraint_row(constraints_frame, len(self.constraint_entries)+1)
ttk.Button(button_frame, text="Remove Constraint", command=self.remove_constraint_row).pack(side=tk.LEFT, padx=5)

### Solve Button
ttk.Button(tab, text="Solve Graphically", command=self.solve_graphical).grid(row=3, column=8, columnspan=5, pady=10)

### Button
### Button
tk.Button(tab, text="Solve Graphically", command=self.solve_graphical).grid(row=3, column=8, columnspan=5, pady=10)

### Button
tk.Button(tab, text="Solve Graphically", columnspan=5, padx=10, pady=10)
```

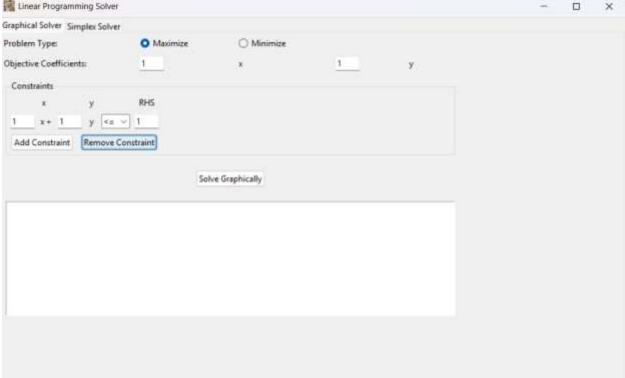
### **Backend Functions**

```
det add constraint_row(self, frame, row):
          coeff_x = tk.DoubleVar(value-1)
         coeff_y = tk.Doublevar(value=1)
rhs = tk.Doublevar(value=1)
          sense = tk.Stringvar(value="c=")
          x_entry = ttk.Entry(frame, textvariable=coeff_x, width=5)
          plus_label - ttk.tabel(frame, text-"x +")
         y entry - ttk.intry(frame, textvariable-coeff y, width-5)
y_label - ttk.tabel(frame, text-"y")
sense_combo - ttk.Combobox(frame, textvariable-sense,
                                     values=["<-", ">-", "-"],
width=1, state="readonly")
          rhs_entry = ttk.Entry(frame, textvariable=rhs, width=5)
          x_entry_grid(row-row, column=1, padx=2, pady=2)
          plus label.grid(row-row, column-2, padx-0, pady-2)
          y_entry_grid(row-row, column=3, padx=2, pady=2)
          y_label.grid(row-row, column=4, padx=8, pady=2)
         sense_combo.grid(row-row, column=5, padx=2, pady=2) rhs_entry.grid(row=row, column=6, padx=2, pady=2)
         self.constraint_entries.append((
              'coeff_x': coeff_x,
'coeff_y': coeff_y,
               'widgets': [x_entry, plus_label, y_entry, y_label, sense_combo, rhs_entry]
Linear Programming Solver
                                                                                                                                                       ×
Graphical Solver Simplex Solver
Problem Type:

    Maximize

                                                              O Minimize
Objective Coefficients:
                                                                                                            y.
  Constraints
                                    RHS
                      y <= v 1
                      y <= v
         x+ 1
                      y <= ~
        x+ 1
        x+ 1
                      y <= -
                     y <= v 1
        x+ 1
  Add Constraint
                    Remove Constraint
                                                    Solve Graphically
```





Calls *graphical\_solver()* and prints results.

### Simplex Solver Tab Interface

```
tab = ttk.Frame(self.notebook)
self.notebook.add(tab, text="Simplex Solver")

a Problem Type
ttk.Label(tab, text="Problem Type:").grid(row=0, column=0, padx=5, pady=5, sticky=tk.W)
self.simplex.opt type = tk.StrimpVar(value="nux")
ttk.Label(tab, text="Problem Type:").grid(row=0, column=0, padx=5, pady=5, sticky=tk.W)
self.simplex.opt type = tk.StrimpVar(value="nux")
ttk.Ladiobutton(tab, text="Minimize", variable=self.simplex.opt type, value="nux").grid(row=0, column=1, sticky=tk.W)
ttk.Ladiobutton(tab, text="Minimize", variable=self.simplex.opt type, value="nux").grid(row=0, column=2, sticky=tk.W)

a Variables Frame
vars frame = ttk.LabelFrame(tab, text="Variables")
vars frame.grid(row=1, column=0, columnspan=3, padx=10, pady=10, sticky=tk.W=tk.E)

ttk.Label(vars_frame, text="Number of Variables:").grid(row=0, column=0, padx=5, pady=5)
self.num_vars = tk.IntVar(value=2)
ttk.Spinbus(vars_frame, trow=1, to=10, textvariable=self.num_vars, width=5, command=self.update_variable_count).grid(row=0, column=1, pac

# Objective Function
self.objective entries = []
self.objective frame_grid(row=1, column=0, column=pan=2, padx=5, pady=5, sticky=tk.W=tk.E)
self.objective_entries = (column=0, column=0, column=0, pady=5, sticky=tk.W=tk.E)
self.objective_entries()
```

constraints frame - ttk.LabelFrame(tab, text="Constraints")	
constraints_frame.grid(row=2, column=0, columnspan=3, padx=10, pady=10, sticky=U	K.WELK.E)
ttk.Libel(constraints_frame, text="Coefficients").grid(row=0, column=0, columnsp	an=10, padx=5, pady=5)
# Constraint entries (start with 2 empty constraints)	
self.simplex constraint entries - []	
for i in range(2):	
self.add simplex constraint row(constraints frame, i+1)	
# Add/Remove constraint buttons	
button frame - ttk.Frume(constraints frame)	
button frame.grid(row=10, column=0, columnspan=10, pady=5)	
ttk.Button(button frame, text="Add Constraint", command=lambda: self.add_simplex	constraint row(constraints frame, len(self.simplex constr
ttk.Button(button_frame, text="Nameve Constraint", command-self.remove_simplex_c	onstraint_row).pack(side-tk.LEFT, padx-5)
# Salve Button	
<pre>ttk.Button(tab, text="Solve with Simplex", command=self.solve_simplex).grid(row=</pre>	3, column=0, columnspan=3, pady=10)
# Resulted	
self.simplex results -Text(tab, height-10, width-60)	
self.simplex results.grid(row-4, column-8, columnspan-3, padx-10, pady-10)	

raphical Solver [Simplex Solver]		Maximize	○ Minimize	
ariables		7-1000000000000000000000000000000000000		
lumber of Variables:	2 0			
Objective Coefficient				
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500 M				
onstraints				
Coefficie	ints			
+ 1	** 1			
+ 1	<= 1			
Add Constraint Re	move Constrair	W1		
add Constraint Re	move Constrair	it i		
		Solve with Signature		
		Solve with Simplex		

#### **Backend Functions**

```
def update_variable_count(self):
    self.update_objective_entries()
    self.update_constraints()
```

when the user changes the number of variables the entries for objective and constraints functions will change too.

The function is called at *line 174*:

ttk.Spinbox(vars\_frame, from\_=1, to=10, textvariable=self.num\_vars, width=5, command=self.update variable count).grid(row=0, column=1, padx=5, pady=5)

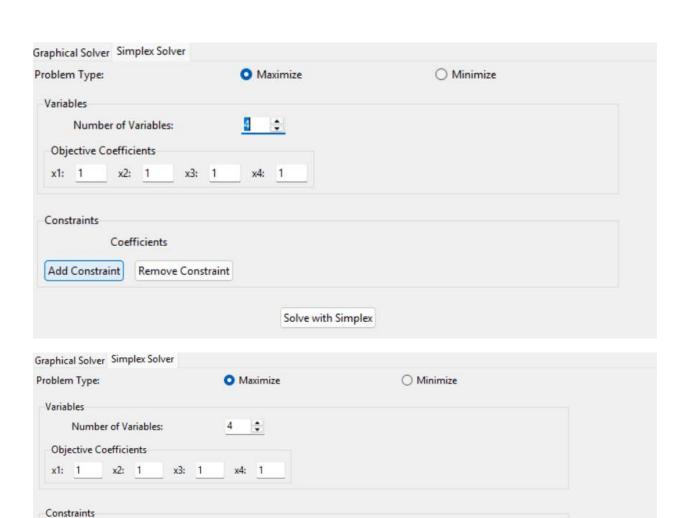
```
def update_objective_entries(self):

# Clear existing entries
for widget in self.objective_frame.winfo_children():
    widget.destroy()

self.objective_entries = []
    num_vars = self.num_vars.get()

for i in range(num_vars):
    ttx.Lahel(self.objective_frame, text=f"x(i+1):").grid(row=0, column=i*2, padx=5, pady=5)
    var = tk.DoubleVar(value=1)
    ttx.Entry(self.objective_frame, textvariable=var, width=5).grid(row=0, column=i*2+1, padx=5, pady=5)
    self.objective_entries.append(var)
```

```
update_constraints(self);
   current constraints = []
   for entry in self.simplex_constraint_entries:
      current constraints.append((
           'coeffs': [var.get() for var in entry['coeffs']],
           'sense': entry['sense'].get(),
'rhs': entry['rhs'].get()
   for entry in self.simplex_constraint_entries:
       for widget in entry['widgets']:
           widget.destroy()
   self.simplex constraint entries = []
       self.add_simplex_constraint_row(self.simplex_constraint_entries[0]['widgets'][0].master, row)
       # Set values from stored data (truncate or pad as needed)
       num_vars = self.num_vars.get()
       for j in range(min(num_vars, len(constraint['coeffs']))):
           self.simplex constraint entries[-1][ coeffs'][j].set(constraint[ coeffs'][j])
       self.simplex_constraint_entries[-1]['sense'].set(constraint['sense'])
       self.simplex_constraint_entries[-1]['rhm'].set(constraint['rhs'))
```



Coefficients

Add Constraint Remove Constraint

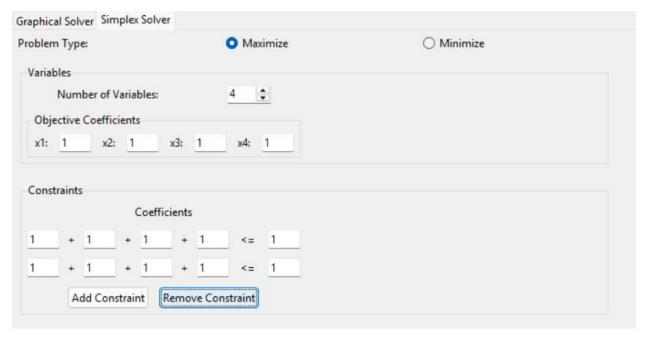
+ 1

<= <u>1</u> <= 1

<= 1

```
dof add simplex constraint row(self, frame, row):
   num_vars = self.num_vars.get()
   coeff vars -
   widgets = []
   sense = tk.StringVar(value="c=")
   rhs = tk.DoubleVar(value=1)
   for i in range(num_vars):
       var - tk.DoubleVar(value-1)
       entry = ttk.Entry(frame, textvariable=var, width=5)
       entry.grid(row-row, column-i*2, padx-5, pady-5)
       coeff_vars.append(var)
widgets.append(entry)
            lbl = ttk.tabel(frame, text="+")
            lbl.grid(row=row, column=i*2+1, padx=2)
widgets.append(lbl)
   col = num vars 2
   sense label = ttk.tabel(frame, text="<=", width=3)
sense label.grid(row=row, column=col, padx=5, pady=5)</pre>
   rhs_entry = ttk.Entry(frame, textvariable=rhs, width=5)
   rhs_entry.grid(row-row, column-col+1, padx-5, pady-5)
   widgets.extend([sense_label, rhs_entry])
    self.simplex_constraint_entries.append(|
         'coeffs': coeff vars,
'sense': sense,
         'widgets': widgets I have widgets for removal
```

Removing a constraint also removes all labels not only the entries (widget.destroy)



Calls *SimplexSolver().solve()* and prints results.

```
if __name__ == "__main__":
    root = tk,Tk()
    app = LPSolverApp(root)
    root.mainloop()
```

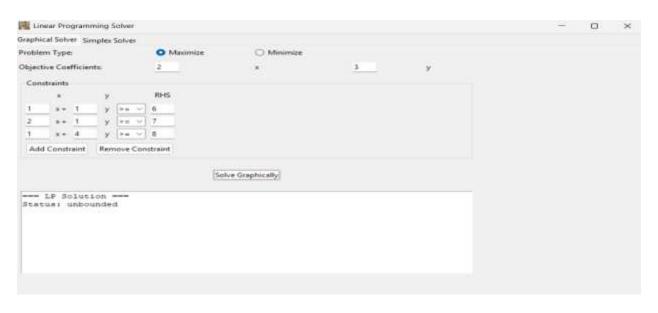
Standard Tkinter main loop.

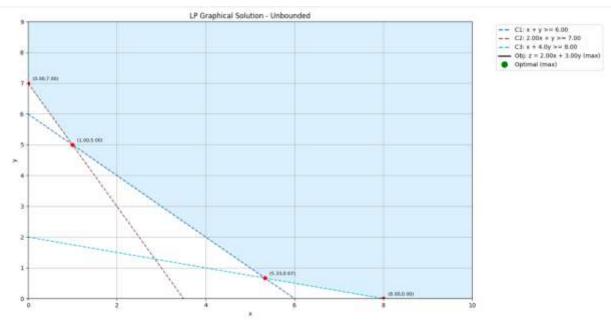
Runs the application until user closes it.

# 5.2 Example Usage

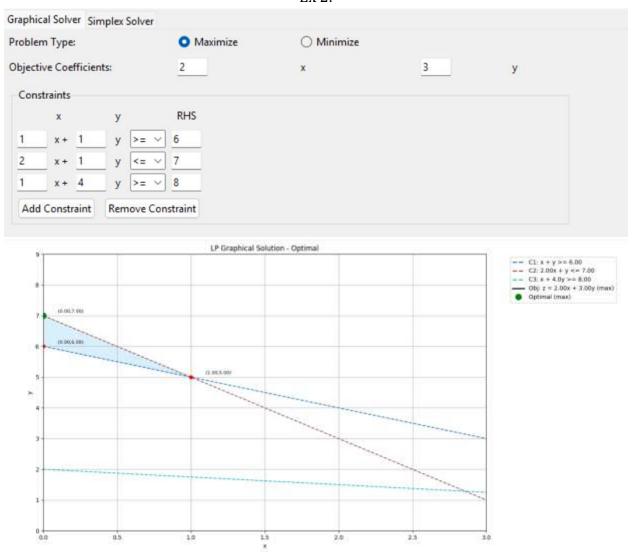
### Graphical problem

Min Z = 2x+3ySubject to:  $x+y \ge 6$ ,  $2x+y \ge 7$ ,  $x+4y \ge 8$ ,  $x \ge 0, y \ge 0$ 





Ex 2:



### Simplex Problem

Max 
$$z = 2x1+4x2+x3+x4$$
  
Subject to:  
 $x1 + 3x2 + x4 \le 4$ ,  
 $2x1 + x2 \le 3$ ,  
 $x2 + 4x3 + x4 \le 3$ ,  
 $x1,x2,x3,x4 \ge 0$ 

