

PERFORMANCE EVALUATION OF COMPUTER SYSTEMS AND NETWORKS

#### EPIDEMIC BROADCAST

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## Chapter 1

# Introduction

In this work a study on the broadcast of an epidemic message is carried out.

# Chapter 2

# Simulation

### Chapter 3

### **Appendices**

#### Appendix A

Given a scenario with N transmitter devices and a target device T in reach of all the transmitters, let us define the probabilities  $P_1(j, N)$  as the probability of j devices out of N transmitting at the same time during slot 1 and  $P_i(j)$  as the probability of j devices transmitting at the same time during slot i.

By specification, the successful reception of the message by device T happens if and only if **one** of the transmitters sends the message during the slot. Furthermore, the successful transmission of a device is "a Bernoullian RV with success probability p on every slot, until it achieves success"; therefore we can model  $P_1(j, N)$  as follows:

$$P_{1}(0,N) = (1-p)^{N}$$

$$P_{1}(1,N) = Np(1-p)^{N-1}$$

$$P_{1}(2,N) = \binom{N}{2}p^{2}(1-p)^{N-2}$$

$$P_{1}(3,N) = \binom{N}{3}p^{3}(1-p)^{N-3}$$
...
$$P_{1}(N-1,N) = \binom{N}{N-1}p^{N-1}(1-p)$$

$$P_{1}(N,N) = \binom{N}{N}p^{N}$$

As for  $P_i(j)$  we can model the system as if it was in the first slot, with N now being equal to N-t transmitters, where t is the number of devices that have transmitted in the (i-1)-th slot.

Therefore, for i = 2 we have:

$$\begin{split} P_2(0) &= P_1(0,N)P_1(0,N) + P_1(2,N)P_1(0,N-2) + \ldots + P_1(N-1,N)P_1(0,1) = \\ &= P_1(0,N)P_1(0,N) + \sum_{k=2}^{N-1} P_1(k,N)P_1(0,N-k) = \\ &= \sum_{k=0}^{N-1} P_1(k,N-k)P_1(0,N) - P_1(1,N)P_1(0,N-1) \\ P_2(1) &= P_1(0,N)P_1(1,N) + P_1(2,N)P_1(1,N-2) + \ldots + P_1(N-1,N)P_1(1,1) = \\ &= P_1(0,N)P_1(1,N) + \sum_{k=2}^{N-1} P_1(k,N)P_1(1,N-k) = \\ &= \sum_{k=0}^{N-1} P_1(k,N)P_1(1,N-k) - P_1(1,N)P_1(1,N-1) \\ P_2(2) &= P_1(0,N)P_1(2,N) + P_1(2,N)P_1(2,N-2) + \ldots + P_1(N-2,N)P_1(2,2) = \\ &= P_1(0,N)P_1(2,N) + \sum_{k=2}^{N-2} P_1(k,N)P_1(2,N-k) = \\ &= \sum_{k=0}^{N-2} P_1(k,N)P_1(2,N-k) - P_1(1,N)P_1(2,N-1) \\ P_2(3) &= \ldots = \sum_{k=0}^{N-3} P_1(k,N)P_1(3,N-k) - P_1(1,N)P_1(3,N-1) \\ \ldots \\ P_2(N-2) &= \sum_{k=0}^{2} P_1(k,N)P_1(N-2,N-k) - P_1(1,N)P_1(N-2,N-1) \\ P_2(N-1) &= \sum_{k=0}^{1} P_1(k,N)P_1(N-1,N-k) - P_1(1,N)P_1(N-1,N-1) \\ P_2(N) &= P_1(0,N)P_1(N,N) \end{split}$$

which has the general form:

$$P_{2}(j) = \begin{cases} \sum_{k=0}^{N-1} P_{1}(k,N) P_{1}(0,N-k) - P_{1}(1,N) P_{1}(0,N-1) & j = 0\\ \sum_{k=0}^{N-j} P_{1}(k,N) P_{1}(j,N-k) - P_{1}(1,N) P_{1}(j,N-1) & 0 < j < N\\ P_{1}(0,N) P_{1}(N,N) & j = N \end{cases}$$
(3.1)

where the term with the minus sign is due to the fact that, if only one device transmitted during slot i, the target device T will have correctly received the message and therefore,

starting from slot i + 1 onwards, it will not be listening for incoming messages any more but it will be itself transmitting instead.

For i = 3 we have:

$$\begin{split} P_3(0) &= \sum_{k=0}^{N-1} P_2(k) P_1(0,N-k) - P_2(1) P_1(0,N-1) \\ P_3(1) &= \sum_{k=0}^{N-1} P_2(k) P_1(1,N-k) - P_2(1) P_1(1,N-1) \\ P_3(2) &= \sum_{k=0}^{N-2} P_2(k) P_1(2,N-k) - P_2(1) P_1(2,N-1) \\ &\cdots \\ P_3(N-1) &= \sum_{k=0}^{1} P_2(k) P_1(N-1,N-k) - P_2(1) P_1(0,N-1) \\ P_3(N) &= P_2(0) P_1(N,N) \end{split}$$

which has a general form similar to (3.1)

$$P_3(j) = \begin{cases} \sum_{k=0}^{N-1} P_2(k) P_1(0, N-k) - P_1(1, N) P_1(0, N-1) & j = 0\\ \sum_{k=0}^{N-j} P_2(k) P_1(j, N-k) - P_1(1, N) P_1(j, N-1) & 0 < j < N\\ P_2(0) P_1(N, N) & j = N \end{cases}$$
(3.2)

We can further generalize formula 3.2 to obtain the probability  $P_i(j)$  we introduced at the beginning:

$$P_{i}(j) = \begin{cases} \sum_{k=0}^{N-1} P_{i-1}(k) P_{1}(0, N-k) - P_{i-1}(1, N) P_{1}(0, N-1) & j = 0\\ \sum_{k=0}^{N-j} P_{i-1}(k) P_{1}(j, N-k) - P_{i-1}(1, N) P_{1}(j, N-1) & 0 < j < N\\ P_{i-1}(0) P_{1}(N, N) & j = N \end{cases}$$
(3.3)