CSCI 247

Isaac Boaz

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Question 1

Given a floating-point format with a k-bit exponent and an n bit fraction, write formulas for the exponent E, the significand M, the fraction f, and the value V for the quantities that follow. In addition, describe the bit representation.

A) The number 7.0

Keeping in mind the general formula

$$V = (-1)^s \cdot M \cdot 2^E$$

where V is the value, s is the sign, M is the mantissa, and E is the exponent.

- 1. We'll first begin by converting 7.0 to its binary representation $7_{10}=111_2\,$
- 2. Normalizing this representation leaves us with

$$111_2 = 1.11 \cdot 2^2$$

- 3. Meaning the exponent Exp = 2
- 4. And the fractional part is $f = 111^*$ Note: IEEE implies a leading 1 at the beginning, so f = 110
- 5. Calculate the biasd E by adding $Bias = 2^{k-1} 1$

$$E = 2 + (2^{k-1} - 1) = 2^{k-1} + 1$$

6. Combing with an example of k = 4, M = 3

$$\begin{split} Exp &= 2 \\ E &= 2^{4-1} + 1 = 9_{10} = 1001_2 \\ M &= 1.11_2 \\ f &= 110_2 \\ V &= (-1)^s \cdot (1.) f \cdot 2^{Exp} \\ &= (1.)11_2 \cdot 2^2 = 7.0 \end{split}$$

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With a final bit representation of $\begin{array}{c|c} \operatorname{Sign} & \operatorname{Exp} & \operatorname{Mantissa} \\ 0 & 1001 & 110 \end{array}$

B) The largest odd integer that can be represented exactly

- ullet We know the sign must be 0 for positives.
- The mantissa should be all 1's to get the largest representable odd number.
- If our exponent goes beyond the mantissa's accuracy, we'll end up with a 0 as the least significant bit (IE an even number).
- Thus, our exponent is Exp = Min(k, M)

An example where k = 4, M = 3

$$Exp = 3$$

$$Bias = 2^{k-1} - 1 = 2^{4-1} - 1 = 7$$

$$E = Exp + Bias$$

$$= 3 + 7 = 10_{10} = 1010_{2}$$

$$f = 111_{2}$$

$$V = (1.)111_{2} \cdot 2^{3} = 15$$

With a final bit representation of $\begin{array}{c|c} \operatorname{Sign} & \operatorname{Exp} & \operatorname{Mantissa} \\ 0 & 1010 & 111 \end{array}$

C) The reciprocal of the smallest positive normalized value

In order to find the reciprocal of any number, we can simply "flip" the numerator and denominator. This can be achieved in floating point representation by multiplying the exponent by -1. The smallest possible value given k exponent bits and M fraction bits is when:

- \square M=0, s=0
- \square Exp is the biggest positive value multiplied by -1
- 1. Set M = 0, s = 0
- 2. Calculate E = Exp Bias with Exp = 1
- 3. One example where k = 5, M = 4

$$Bias = 2^{k-1} - 1 = 2^{5-1} - 1$$
 = 15
 $E = (1 - Bias) = (1 - 15) \cdot -1$ = 14
 $Exp = E + Bias = (14 + 15)$ = 29₁₀ = 11101₂
 $f = 0_2$

With a final bit representation of $\begin{array}{c|c} Sign & Exp & Mantissa \\ \hline 0 & 11101 & 000 \\ \end{array}$ Which equals 16384