# CSCI 305 Assignment 1

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## Problem 1

1	for i = 1 to n	cost	$_{ m time}$
2	i = 1	$C_1$	n+1
3	$\mathbf{while} \ \mathbf{j} \leq \mathbf{i} \ \mathbf{do}$	$C_2$	n
4	j = 2.5 * j	$C_3$	$(n+1)\log_{2.5}$
-	j = 2.0 · j	$C_4$	$n \log_{2.5} n$

Detailed Runtime:

$$C_1 \cdot (n+1) + C_2 \cdot n + C_3 \cdot \lceil (n+1) \log_{2.5} \rceil + C_4 \cdot \lceil n \log_{2.5} n \rceil$$

Asymptotic Runtime:

$$\Theta(n \log n)$$

#### Problem 2

0	def factorial(n):	cost	$_{ m time}$
1	x = 1	$C_1$	1
2	while $n > 1$ :	$C_2$	$n \\ n-1 \\ n-1$
3	x = x * n	$C_3$	n-1
4	n = n - 1	$C_4$	n-1 $n$
5	return x	$C_5$	n

Detailed Runtime:

$$C_1 \cdot 1 + C_2 \cdot n + C_3 \cdot (n-1) + C_4 \cdot (n-1) + C_5 \cdot n$$

Asymptotic Runtime:

O(n)

#### Problem 3

```
def Binary_InsertionSort(A):
 1
                for j = 2 to n:
 2
                        key = A[j]
 3
                         // insert A[j] into the sorted sequence A[1..j-1]
                        left = 1
 4
 5
                        right = j - 1
                        while right > left:
 6
                                  \begin{array}{l} \operatorname{mid} = \operatorname{floor} \left( \left( \operatorname{left} + \operatorname{right} \right) \; / \; 2 \right) \\ \mathbf{if} \; \operatorname{key} > \operatorname{A[mid]} : \end{array} 
 7
 8
9
                                         left = mid + 1
10
                                 {f else} {f right} = {f mid}
                        if \hspace{0.1cm} \text{key} \hspace{0.1cm} > \hspace{0.1cm} A \hspace{0.1cm} [\hspace{0.1cm} \hspace{0.1cm} \text{left} \hspace{0.1cm}] \hspace{0.1cm} \colon \hspace{0.1cm}
11
                        12
13
                                A[i] = A[i - 1]
14
                        A[left] = key
15
```

1. Note that regardless of input, the for loop will always run n times, with the contents of the loop running n-1 times.

Worst Case occurs when the array is sorted in reverse order, as the left and right bounds will need to be updated  $\log_2 j$  times.

Best Case occurs when the array is already sorted, as the left and right bounds will only need to be updated once.

2. Worst Case for the *while* block is the same in the best and worst case scenario, running  $\log_2 j$  times.