### Standard vs. Nested Form

### Methods

**Bisection Method** 

Standard Form:  $a_0 + a_1x + a_2x^2 + \cdots + a_nx^n$ Nested Form:  $a_0 + x(a_1 + a_2x + \cdots + a_nx^{n-1})$ 

# Sig Fig Rounding

4-Digit Example	
37.21429	37.21
0.002271828	0.002272
3000527.11059	3001000

#### Bit Preciscion

			$\operatorname{sign}$	exponent	mantissa	machine $\epsilon$
single		1	8	23	$2^{-23}$	
double (default)		1	11	52	$2^{-52}$	
long double		1	15	64	$2^{-64}$	
	sign	mantissa	expone	ent	$\rightarrow 1.b_1b_2\cdots b_n$	$b_n \times 2^P$

# **Error Calculation**

$$x_e = \text{Exact}, \ x_a = \text{Approximate}$$

$$\text{Abs err} = |x_e - x_a|$$

$$\text{Rel err} = \frac{|x_e - x_a|}{x_e}$$

Worst case error =  $0.5 \cdot 10^{-P+1}$ 

## Loss of Significance

$$a - b = \frac{a^2 - b^2}{a + b}$$

#### Theorems

#### Intermediate Value Theorem

If f(x) is continuous on [a, b] and  $f(a) \neq f(b)$ , then for any y in  $[f(a), f(b)] \exists c \in [a, b]$  such that f(c) = y

#### Mean Value Theorem

If 
$$f'(x)$$
 is continuous on  $[a, b]$ ,  
then  $\exists c \in [a, b]$  such that  $f'(c) = (f(b) - f(a))/(b - a)$ 

#### Rolle's Theorem

If 
$$f'(x)$$
 is continuous on  $[a, b]$  and  $f(a) = f(b)$ ,  
then  $\exists c \in [a, b]$  such that  $f'(c) = 0$ 

## Taylor Series

$$P_k(x) = \sum_{n=0}^k \frac{f^{(n)}(x_0)}{n!} (x - x_0)^n$$
$$E_{k+1} = \frac{f^{(k+1)}(c)}{(k+1)!} (x - x_0)^{k+1}$$

- 1. Choose a, b such that  $f(a) \cdot f(b) < 0$
- 2. Compute  $c = \frac{a+b}{2}$
- 3. If f(c) and f(a) have opposite signs, then b=c
- 4. If f(c) and f(b) have opposite signs, then a=c
- 5. If  $|b-a| < \epsilon$  or f(c) = 0 then c is the solution

$$\text{Error} = \frac{b-a}{2^{n+1}}$$
 
$$\text{Steps}(\epsilon) = \left\lceil \frac{\ln(b-a) - \ln(\epsilon)}{\ln(2)} - 1 \right\rceil$$

Newton's Method

- 1.  $x_{n+1} = x_n \frac{f(x_n)}{f'(x_n)}$
- 2. If  $|x_n x_{n-1}| < \epsilon$  or  $f(x_n) = 0$  then  $x_n$  is the solution

Multiplicity

If 
$$f^{(m-1)}(r_0) = 0$$
 and  $f^{(m)}(r_0) \neq 0$ ,  
then  $f^{(m)}(r_0)$  is a root of multiplicity  $m$ 

Convergence 
$$m=1$$
  $m>1$   $M=\frac{1}{2}\left|\frac{f''(r_0)}{f'(r_0)}\right|$   $e_{i+1}\approx \frac{m-1}{m}e_i$   $e_{i+1}\approx Me_i^2$ 

Secant Method

$$x_{n+1} = x_n - f(x_n) \frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})}$$

Convergence

$$e_{i+1} \approx M e_i^p$$
 
$$p = \frac{1+\sqrt{5}}{2} \text{ (Golden Ratio)}$$

False Position Method

Same as Bisection Method, but

$$c = a - \frac{f(a)(b-a)}{f(b) - f(a)}$$

Convergence - Generally same as Bisection **Fixed Point Iteration** 

1. Solve f(x) for x = g(x)

$$g(x) = 2.8x - x^{2} = x$$
$$1.8x - x^{2} = 0$$
$$x(1.8 - x) = 0$$
$$x = 0, 1.8$$

$$f(x) = 2x^3 - 6x - 1$$
$$g(x) = \frac{2x^3 - 1}{6}$$

2. If  $|g'(r_n)| < 1$ , fixed point will converge to  $r_n$ 

#### Matrix

# Gaussian Elimination $O(\frac{2n^3}{3})$

Back Substitution 
$$O(n^2)$$

# **Summations**

$$\sum_{i=1}^{n} i = \frac{n(n+1)}{2} \qquad \sum_{i=1}^{n+5} (k-1) = \sum_{i=1}^{n+4} j = \frac{(n+4)(n+5)}{2} \qquad \frac{s_b^2}{2 \cdot s_e^3/3} = \frac{t_b}{t_e}$$

$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

$$\frac{s_b^2}{2 \cdot s_e^3/3} = \frac{t_b}{t_e}$$

Complexity Ratio

### LU Factorization

Goal:  $A = LU \to A\overrightarrow{x} = \overrightarrow{b}$ 

#### Slow Method

- 1. Row reduce A, recording operations as elementary matrices
- 2. The row reduced A is now U
- 3. You now have  $E_n E_{n-1} \cdots E_2 E_1 A = LU$
- 4. Invert the elementary matricies to find L
- 5.  $L = (E_n E_{n-1} \cdots E_2 E_1)^{-1}$

$$A = \begin{bmatrix} 1 & 2 & -1 \\ 2 & 1 & -2 \\ -3 & 1 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -\frac{7}{3} & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 3 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$e_3 \qquad e_2 \qquad e_1$$

Example

$$E_{3}E_{2}E_{1}A = U$$

$$A = E_{1}^{-1}E_{2}^{-1}E_{3}^{-1}U$$

$$= (E_{3}E_{2}E_{1})^{-1}U$$

$$A = LU = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -3 & -\frac{7}{3} & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 & -1 \\ 0 & -3 & 0 \\ 0 & 0 & -2 \end{bmatrix}$$

#### Fast Method

- 1. Apply pivots internally within A
- 2. Note that you will subtract rather than add with these operations

$$\begin{bmatrix} 1 & 2 & -1 \\ 2 & 1 & -2 \\ -3 & 1 & 1 \end{bmatrix}$$
 (1) 
$$\begin{bmatrix} 1 & 2 & -1 \\ \frac{2}{1} & -3 & 0 \\ -\frac{3}{1} & -7 & -2 \end{bmatrix}$$
 (2)

3. Don't forget to recurse through the inner matrix (row 2 and beyond)

$$\begin{bmatrix} 1 & 2 & -1 \\ \frac{2}{1} & -3 & 0 \\ -\frac{3}{1} & \frac{7}{3} & -2 \end{bmatrix}$$
 (3)

- 4. The circled (lower triangular) elements are now L (plus the identity matrix)
- 5. The remaining upper triangular elements are U

#### **Derivatives**

$$\frac{df}{dx}f(g(x)) = f'(g(x))g'(x) \quad \frac{d}{dx}ax^n = nax^{n-1} \quad \frac{d}{dx}ln_a(x) = \frac{1}{xln(a)}$$
$$\frac{d}{dx}a^x = \ln(a)a^x \quad \frac{d}{dx}sin(x) = cos(x) \quad \frac{d}{dx}cos(x) = sin(x)$$

#### PA = LU Factorization

Goal:  $PA = LU \to A\overrightarrow{x} = \overrightarrow{b}$ 

- 1. Same logic as LU, but move the |biggest| element to the top of the column
- 2. Make sure you do the same for inner matrices
- 3. You'll have multiple P matrices, such that

$$PA = LU$$

4. Make sure you apply P when solving for  $\overrightarrow{x}$ 

#### Solving

- 1. Solve  $L\overrightarrow{y} = \overrightarrow{b}$  for  $\overrightarrow{y}$
- 2. Solve  $U\overrightarrow{x} = \overrightarrow{y}$  for  $\overrightarrow{x}$

Example 1: A = LU

Let 
$$\overrightarrow{b} = \begin{bmatrix} 2\\1\\-1 \end{bmatrix}$$

1.  $L\overrightarrow{y} = \overrightarrow{b}$ 

$$\begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 & 1 & 0 & 1 \\ -3 & -\frac{7}{3} & 1 & -1 \end{bmatrix} \rightarrow \overrightarrow{y} = \begin{pmatrix} 2 \\ -3 \\ -2 \end{pmatrix}$$

2.  $U\overrightarrow{x} = \overrightarrow{y}$ 

$$\begin{bmatrix} 1 & 2 & -1 & 2 \\ 0 & -3 & 0 & -3 \\ 0 & 0 & -2 & -2 \end{bmatrix} \rightarrow \overrightarrow{x} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$$

Example 2: - PA = LU

Let 
$$\overrightarrow{b} = \begin{bmatrix} 1\\2\\3 \end{bmatrix}$$
 and  $A = \begin{bmatrix} 2 & 1 & 5\\4 & 4 & -4\\1 & 3 & 1 \end{bmatrix}$ 

 Find P such that A's pivots are the largest in magnitude per column.
 In this instance:

$$P = P_2 P_1 \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

 $2. \ L\overrightarrow{y} = P\overrightarrow{b}$ 

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ \frac{1}{4} & 1 & 0 & 6 \\ \frac{1}{2} & -\frac{1}{2} & 1 & 5 \end{bmatrix} \rightarrow \overrightarrow{y} = \begin{pmatrix} 0 \\ 6 \\ 8 \end{pmatrix}$$

3.  $U\overrightarrow{x} = \overrightarrow{y}$ 

$$\begin{bmatrix} 4 & 4 & -4 & 0 \\ 0 & 2 & 2 & 6 \\ 0 & 0 & 8 & 8 \end{bmatrix} \rightarrow \overrightarrow{x} = \begin{pmatrix} -1 \\ 2 \\ 1 \end{pmatrix}$$