

## 4200 - Formal Languages: Homework #5

Due on Oct 17, 2019 at 2:00pm

*Instructor: Dr. Anh Nguyen*

## Problem 1

### 30 points

Exercise 2.6. Give context-free grammars (CFGs) generating the following languages.

1. The set of strings over the alphabet  $\Sigma = \{a, b\}$  with more a's than b's

**Answer:**

$$\begin{aligned} S &\rightarrow Aa \mid MS \mid SMA \\ A &\rightarrow Aa \mid \epsilon \\ M &\rightarrow \epsilon \mid MM \mid bMA \mid aMb \end{aligned}$$

2. The complement of the language  $\{a^n b^n \mid n \geq 0\}$

**Answer:**

$A$  is the complement of the language  $\{a^n b^n \mid n \geq 0\}$ .

That is,  $A$  is the language:  $A = \{a^n b^m \mid n \neq m\} \cup \{(a \cup b)^* ba(a \cup b)^*\}$

Let's call the leftmost language  $A_1$  and rightmost language  $A_2$  i.e.  $A = A_1 \cup A_2$ .

The CFG that generates  $A_1$  is

$$\begin{aligned} S_1 &\rightarrow aS_1b \mid T \mid U \\ T &\rightarrow aT \mid a \\ U &\rightarrow Ub \mid b \end{aligned}$$

The CFG that generates  $A_2$  is:

$$\begin{aligned} S_2 &\rightarrow RbaR \\ R &\rightarrow RR \mid a \mid b \mid \epsilon \end{aligned}$$

Therefore, the CFG that generates language  $A$  when  $A = A_1 \cup A_2$  is:

$$\begin{aligned} S &\rightarrow S_1 \mid S_2 \\ S_1 &\rightarrow aS_1b \mid T \mid U \\ S_2 &\rightarrow RbaR \\ T &\rightarrow aT \mid a \\ U &\rightarrow Ub \mid b \\ R &\rightarrow RR \mid a \mid b \mid \epsilon \end{aligned}$$

## Problem 2

### 15 points

Exercise 2.9. Give context-free grammars (CFGs) generating the following language:

$$A = \{a^i b^j c^k \mid i = j \text{ or } j = k \text{ where } i, j, k \geq 0\}$$

Is your grammar ambiguous? Why or why not?

### Answer

Note that  $A = A_1 \cup A_2$  where  $A_1 = \{a^i b^j c^k \mid i = j\}$  and  $A_2 = \{a^i b^j c^k \mid j = k\}$ .

The CFG that generates  $A_1$  is:

$$S_1 \rightarrow aS_2bS_3 \mid \epsilon$$

$$S_2 \rightarrow aS_2b \mid \epsilon$$

$$S_3 \rightarrow cS_3 \mid \epsilon$$

The CFG that generates  $A_2$  is:

$$S_4 \rightarrow S_5bS_6c \mid \epsilon$$

$$S_5 \rightarrow aS_5 \mid \epsilon$$

$$S_6 \rightarrow bS_6c \mid \epsilon$$

The CFG that generates  $A$  would contain all the above 6 rules and starts with the following rule:

$$S \rightarrow S_1 \mid S_4$$

The above CFG is ambiguous because you can derive two leftmost derivations each has a different parse tree.

## Problem 3

### 15 points

Exercise 2.14. Convert the following CFG into an equivalent CFG in Chomsky normal form, using the procedure given in Theorem 2.9.

$$A \rightarrow BAB \mid B \mid \epsilon$$

$$B \rightarrow 00 \mid \epsilon$$

**Answer:** Below is the final result after following the procedure in the book. *Note:* For the Midterm 2, for this type of question, you will be asked to provide all intermediate steps, not just the final one.

$$S_0 \rightarrow BA_1 \mid BS \mid SB \mid UU \mid BB \mid \epsilon$$

$$S \rightarrow BA_2 \mid BS \mid SB \mid UU \mid BB$$

$$B \rightarrow UU$$

$$U \rightarrow 0$$

$$A_1 \rightarrow SB$$

$$A_2 \rightarrow SB$$