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Solving the maximum edge disjoint path problem using a modified Lagrangian particle swarm optimization hybrid

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Maximum Edge Disjoint Path (MEDP)

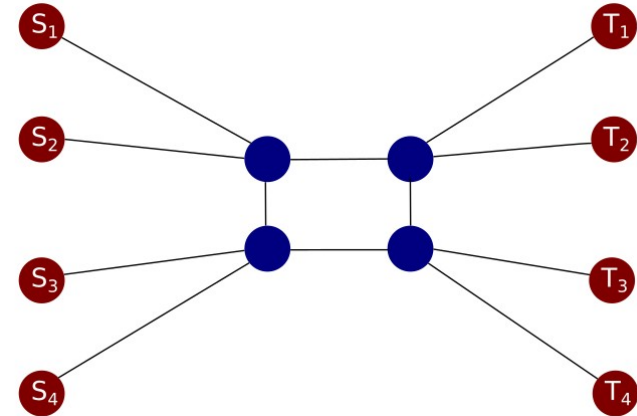


Maximum Edge Disjoint Path:

- Set of source-terminal node pairs has to be connected on a graph.
- Edges can only be used by at most one source-destination pair.

Objective: Connect the most amount of node pairs.

Applications: Optical and energy-efficient networks
(Raghavan & Upfal – 1994)
(Gerez – 1998)



Integer Problem formulation



$$\max \sum_{k \in K} \sum_{j: (s_k, j) \in E} x_{s_k j k}$$

Vars:

x = edge use for each commodity pair

Parameters:

k = source-terminal pair number

i, j = node number

s_k = source node of pair k

Assumptions:

- All edges have a cost of 0, therefore all paths are equally good as long as they satisfy the constraints

$$\text{Node flow: } \sum_{j: (i, j) \in E} x_{ijk} - \sum_{j: (j, i) \in E} x_{jik} = 0$$

$$\forall k \in K, i \in V \setminus \{s_k, t_k\}$$

$$\text{Source flow: } \sum_{j: (s_k, j) \in E} x_{s_k j k} \leq 1$$

$$\forall k \in K$$

$$\text{Capacity constraint: } \sum_{k \in K} (x_{ijk} + x_{jik}) \leq 1$$

$$\forall (i, j) \in E : i < j$$

$$x_{ijk} = \{0, 1\}$$

$$(i, j) \in E, k \in K$$

Multi-start Simple Greedy Algorithm (MSGGA)



MSGGA:

- One of the first approaches (*Kleinberg – 1996, Blesa & Blum – 2007*)
- Tries to connect the commodity pairs using Dijkstra in random orders, then removes the edges used from the graph if a path was found.

Advantages:

- Very easy to implement.
- Very fast.
- Works very well in small and well-connected graphs.

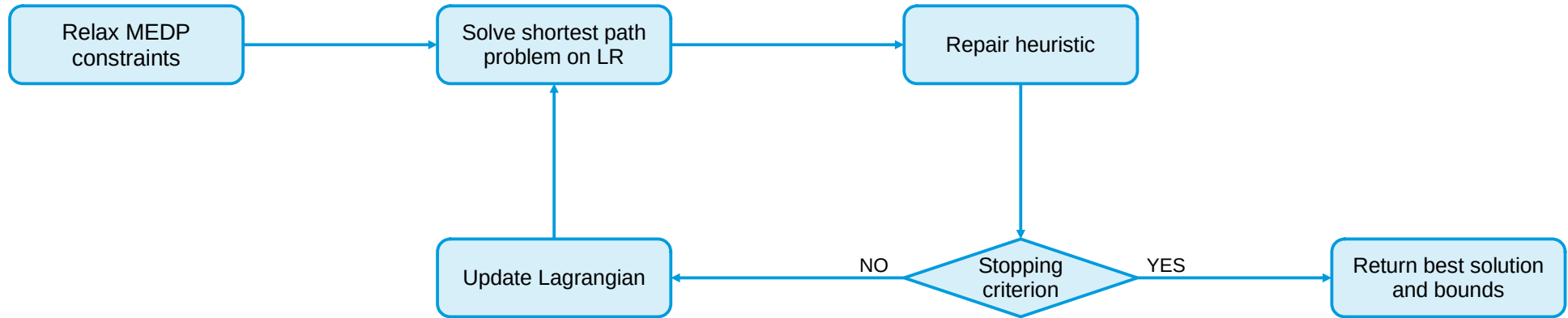
Disadvantages:

- Poor solutions on large or poorly connected graphs.
- Does not provide an upper bound solution.
- Does not explore the entire solution space.

```
Sbest = 0
while n < max_iter:
    reset Graph
    shuffle pairs
    Scurrent = 0
    for pair in pairs:
        path = dijkstra(Graph, pair)
        if path exists:
            Scurrent += 1
            remove(Graph, path)
    if Scurrent > Sbest:
        Sbest = Scurrent
```

LaPSO (2021):

- Short for Lagrangian Relaxation and Particle Swarm Optimisation.
- Consists in relaxing the capacity constraints and reformulating the problem as a minimization problem.
- The MEDP turns in a series of shortest path problems with variable weights.
- If solutions are infeasible, apply a repair heuristic.
- Given the solutions and infractions, update the Lagrangian through PSO framework.
- If solution hasn't improved, max time or iterations have been reached, return the best solution found.



Lagrangian Relaxation



Relaxing the capacity constraint:

$$\begin{aligned}
 \min_{\lambda \geq 0} LR(\lambda) &= \max_{x \in \mathcal{F}} \sum_{k \in K} \sum_{(s_k, j) \in E} x_{s_k j k} + \sum_{(i, j) \in \bar{E}} \lambda_{ij} \left(1 - \sum_{k \in K} (x_{ijk} + x_{jik}) \right) \\
 &= \sum_{(i, j) \in \bar{E}} \lambda_{ij} + \max_{x \in \mathcal{F}} \left\{ \sum_{(s_k, j) \in E} x_{s_k j k} (1 - \lambda_{s_k j}) - \sum_{k \in K} \sum_{(i, j) \in \bar{E}, i \neq s_k} \lambda_{ij} x_{ijk} \right\} \\
 &= \sum_{(i, j) \in \bar{E}} \lambda_{ij} - \sum_{k \in K} \left\{ \min_{x_k \in \mathcal{F}_k} \sum_{(s_k, j) \in E} (\lambda_{s_k j} - 1) x_{s_k j k} + \sum_{(i, j) \in \bar{E}, i \neq s_k} \lambda_{ij} x_{ijk} \right\}
 \end{aligned}$$

Parameters:

$$\bar{E} = \{(i, j) \in E : i < j\}$$

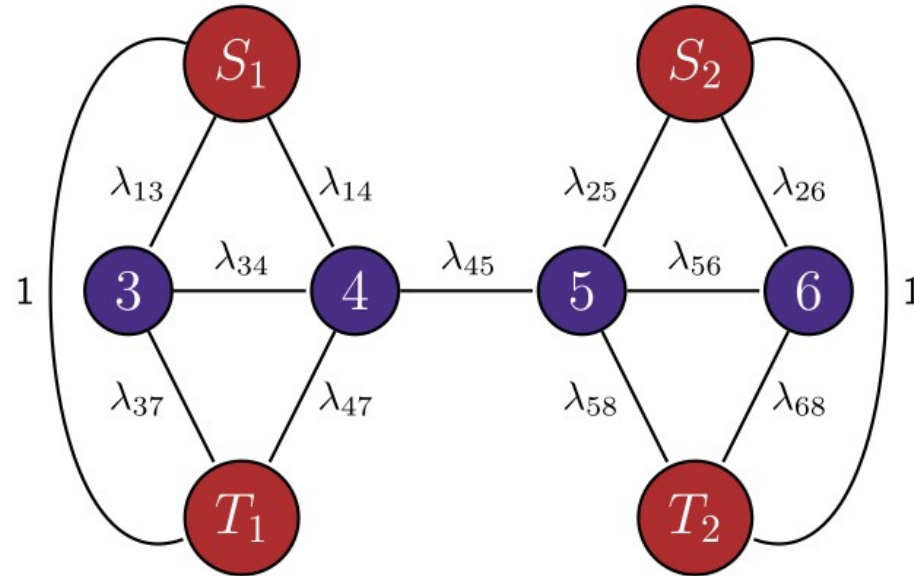
\mathcal{F} set of feasible solutions to all constraints except capacity

$$\min_{\lambda \geq 0} LR(\lambda) = \sum_{(i, j) \in E} \lambda_{ij} + |K| - \sum_{k \in K} \min \left\{ 1, \min_{s_k - t_k \text{ path } P_k} \sum_{(i, j) \in P_k} \lambda_{ij} \right\}$$

Graph representation



Graph representation of the relaxed
Lagrangian minimization problem:

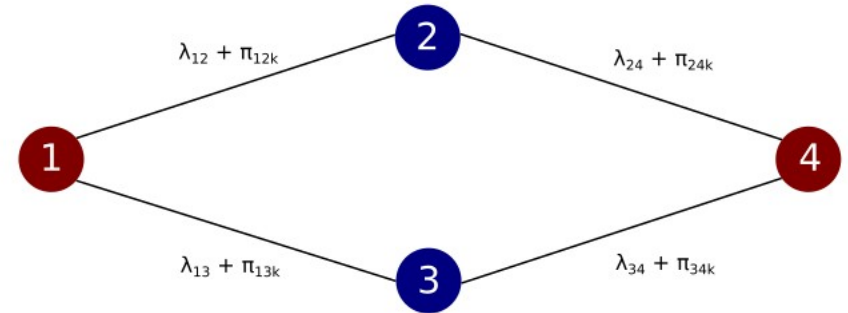


New problem:

If the edge weights are all the same for all commodities, then they will most likely all be connected through the same lowest cost edges, leading to poor diversity and high violation count.

Solution: Perturbations

By adding small perturbations to the graph's edges, commodity pairs will diversify their paths, leading to overall better results. These perturbations are learned through the PSO framework to penalize edges already used by other commodity pairs.



Repair heuristics



Since the MEDP problem has been relaxed, generated solutions may (and very often are) infeasible.

Solution: Repair heuristic

A repair heuristic is applied to the relaxed solution in order to make it feasible.
It takes the infeasible paths and tries to reroute them onto unused edges using Dijkstra.

There are three possible repair heuristics:

- **Rand**: Randomly picks an infeasible path and tries to reroute it.
- **LVA** (Largest Violation Arbitrary): Picks the path with the largest amount of violations and tries to reroute it.
- **LVP** (Largest Violation Perturbation): Picks the path with the largest amount of violations and tries to reroute it using the perturbation values as edge weights.

Particle Swarm Optimization



$$\mathbf{v}_i \leftarrow \omega \mathbf{v}_i + r^L \alpha_i \frac{UB - LB_i}{||\mathbf{g}_i||^2} \mathbf{g}_i + r^G \phi(\lambda_i^G - \lambda_i)$$

$$\mathbf{w}_i \leftarrow \omega \mathbf{w}_i + r^G \phi(\pi_i^G - \pi_i) + \delta r^G \phi(1 - 2\mathbf{x}_i^G)$$

$$\lambda_i \leftarrow \lambda_i + \mathbf{v}_i$$

$$\pi_i \leftarrow \theta \pi_i + \mathbf{w}_i$$

Parameters:

$\mathbf{v}_i, \mathbf{w}_i$ = Particle speeds

ω = Velocity factor

r^L, r^G = random uniform variables – [0,1]

\mathbf{g}_i = subgradient vector

UB, LB_i = Upper and lower bounds

α_i = Subgradient factor

ϕ = Global factor

δ = Perturbation factor

θ = Perturbation decay

Algorithm 1 MEDP LaPSO Algorithm.

Require: Number of particles n

Require: CPU Runtime Limit T

Require: Velocity parameters $\omega, \phi, \delta > 0$

Require: α^0, β the initial value and update factor of α

```

1: Define global best  $\lambda^G = 0, p^G = 0$ 
2: Set initial perturbation and Lagrangian Multipliers for  $p_i(\lambda_i, \pi_i) \in n : \lambda_i \in [0, 0.1],$ 
    $\pi_i = 0$ 
3: Set initial velocities  $v_i = 0, w_i = 0$ 
4: Set maxIter and LBC  $maxIter = 30000, LBC = 10$ 
5: for  $iter = 0 \dots maxIter$  and while  $LB \leq UB - \epsilon$  do and while  $CPU() \leq T$ 
6:   for  $p_i(\lambda_i, \pi_i) \in n$  do
7:     if  $iter \bmod LBC = 0$  then                                // Lower Bound check
8:       Solve subproblem  $L(\lambda_i, 0)$                           // using Dijkstra
9:     else
10:      Solve subproblem  $L(\lambda_i, \pi_i)$                         // using Dijkstra
11:    end if
```

```

12:    Let  $x'$  be the optimal solution to the minimisation subproblem
13:     $g \leftarrow b - Ax'$                                          // update subgradient g
14:    if  $L(\lambda_i, \pi_i) > LB$  then
15:       $LB \leftarrow L(\lambda_i, \pi_i)$ 
16:    else if  $LB_i$  not improved for 10 iterations of particle  $p_i$  then
17:       $\alpha_i \leftarrow \beta \alpha_i$                              // reduce step size
18:    end if
19:    if  $x'$  is not feasible then
20:       $x' \leftarrow \text{REPAIRHEURISTIC}(x')$                        // repair solution
21:    end if
22:    if  $x'$  is feasible and  $(\bar{c}x' < UB)$  then
23:       $UB \leftarrow \bar{c}x'$                                          // update best Upper Bound
24:       $p^G \leftarrow p_i$                                          // update best particle
25:    end if
26:    Randomly generate numbers  $r^G, r^L \in [0, 1]$ 
27:    Update particle velocities  $v_i, w_i$  using Eq. (13) and (14)
28:    Update the particle positions using Eq. (15) and (16)
29:  end for
30: end for
```

Experimental setup



Parameter choices (PSO):

Parameter	Tuning Domain	Value Chosen
Velocity Factor - ω	{0.01,0.05,0.1,0.5}	0.1
Global Factor - ϕ	{0.01,0.05,0.1,0.2,0.3}	0.05
Subgradient Factor - α_i	{1.5,2,3,3.5,5}	2
Perturbation Factor - δ	{0.01,0.05,0.1,0.5}	0.5
Swarm Size - p	{1,8,16,32,64,96}	8

Maximum iterations: 1000
Maximum time: 20 minutes

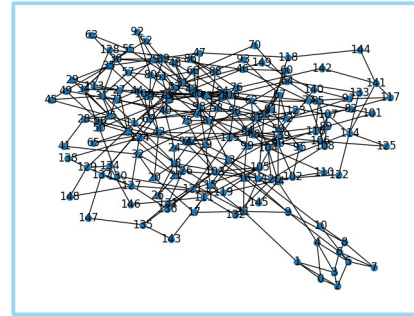
3 runs per graph combination
66% nodes used for pairs

Computational:

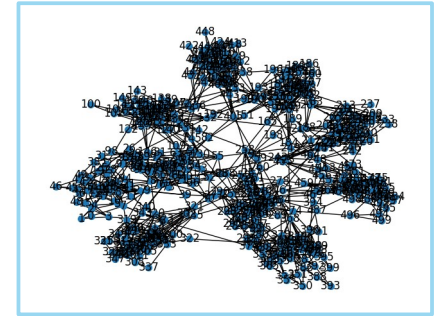
- Python codebase, with RustworkX for fast graph operations
- Ryzen 7 7800X3D 8-cores – 32 GBs of RAM
- MIP solver: Gurobi

Graph choices:

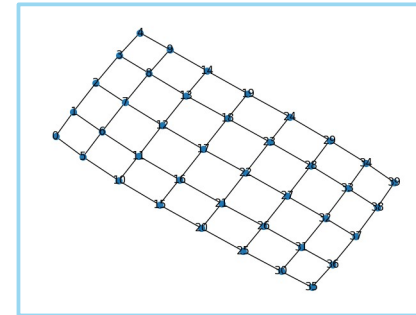
Ordinary_N:



Collated_{C,N}:



Grid_{w,h}:

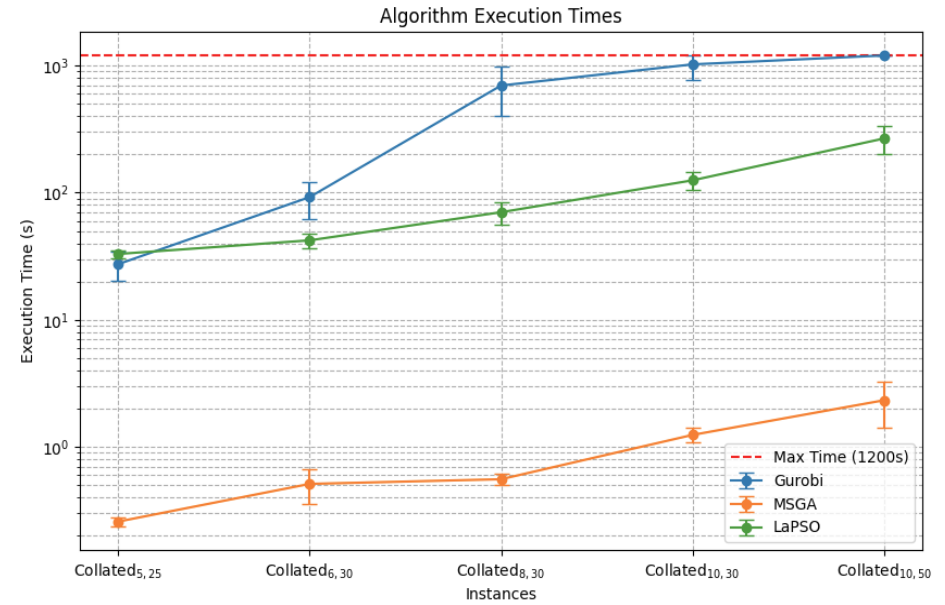
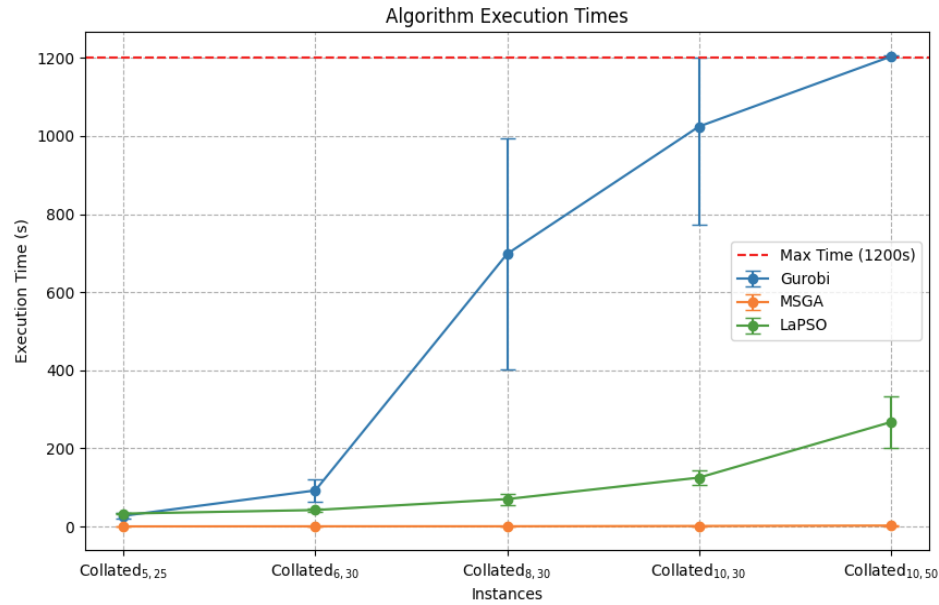


Results

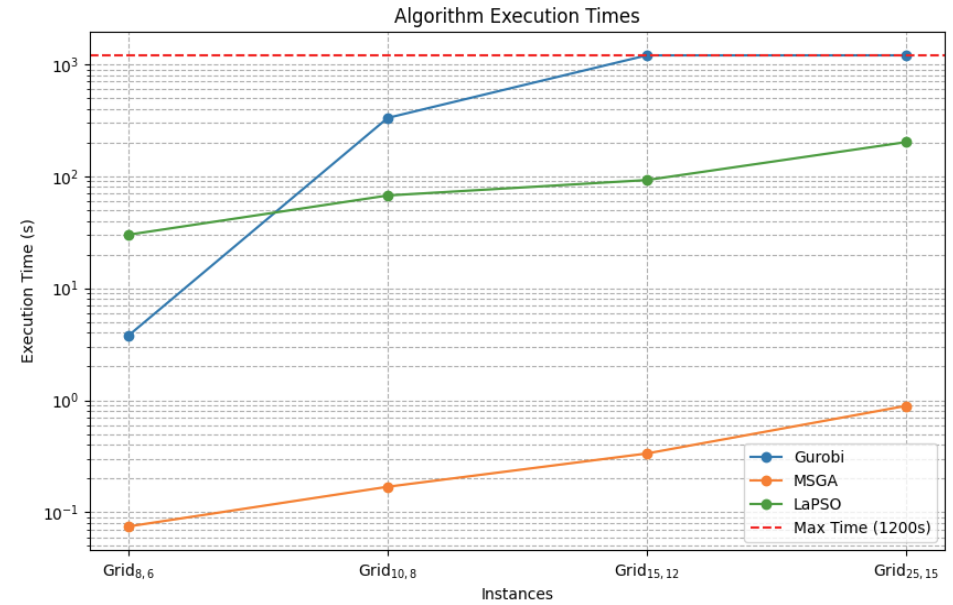
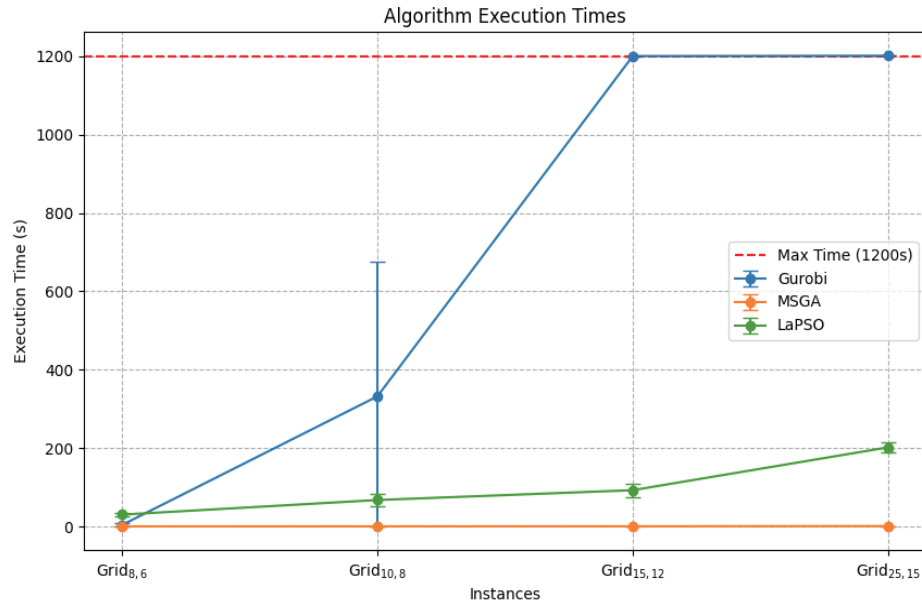


CASE:	Gurobi - Avg	MSGA - Avg	LaPSO - Avg	Gurobi - Gap	MSGA - Gap	LaPSO - Gap	Gurobi - Time	MSGA - Time	LaPSO - Time
Collated _{5,25}	35.6666	35.6666	35.6666	0.0000	0.2866	0.1160	27.3711	0.2577	32.9486
Collated _{6,30}	49.6666	49.3333	49.6666	0.0000	0.2850	0.1240	92.3655	0.5105	42.1783
Collated _{8,30}	55.0000	55.0000	55.0000	0.0000	0.2361	0.1628	697.9784	0.5555	70.2271
Collated _{10,30}	75.3333	72.6666	75.3333	0.0492	0.1925	0.0848	1024.0420	1.2402	125.1619
Collated _{10,50}	124.3333	117.6666	124.3333	0.1907	0.4260	0.1077	1203.2094	2.3249	267.2233
Grid _{8,6}	11.3333	11.3333	11.3333	0.0000	0.2444	0.1288	3.7978	0.0748	30.0885
Grid _{10,8}	17.6666	16.0000	17.6666	0.0000	0.3846	0.1180	331.8981	0.1690	67.3477
Grid _{15,12}	32.3333	26.0000	28.6666	0.1101	0.5593	0.2252	1200.2497	0.3349	92.4551
Grid _{25,15}	54.0000	39.3333	47.0000	0.3799	0.6802	0.2703	1201.1051	0.8911	201.6951
Regular ₅₀	16.0000	16.0000	15.6666	0.0000	0.0000	0.0208	0.3434	0.0870	4.0694
Regular ₁₀₀	31.3333	28.3333	28.0000	0.0000	0.1414	0.1515	255.6080	0.2212	31.5226
Regular ₁₅₀	43.0000	38.0000	36.6666	0.0720	0.2245	0.2517	1200.2062	0.3314	23.6007
Regular ₂₅₀	65.6666	57.6666	55.0000	0.1244	0.2967	0.3293	1200.5760	0.9000	70.6878

Graph plots (Collated)



Graph plots (Grid)



LaPSO vs:

Gurobi:

- Faster computation
- Better scalability
- Better gap estimation for larger instances

MSGGA:

- Better gap estimation
- Better solutions
- Comparable time complexity

Thank you for your attention.