	recursion tree, each node represents the cost of a single newhere in the set of recursive function invocations. *
0	Problem
<b>()</b>	sub problem
0	instruction
0	node
0	Other:

Solve the following recurrence using Master's theorem. \*

$$T(n) = T(n/2) + 2^n$$

- T(n) = O(n\*n)
- $T(n) = O(n*n \log n)$
- $T(n) = O(2^n)$
- cannot be solved

running time of a		algorithm. *
0	Dynamic	
0	Greedy	
•	Divide-and-conquer	
0	Backtracking	

Solve 
$$T(n) = 8T(n/2) + n^2$$

- O (n^2)
- () O(n)
- O(n^3)
- O(n^2 log n)

Solve the recurrence by master method

$$T(n) = T(\sqrt{n}) + 1$$

(a) 
$$T(n) = \Theta(loglogn)$$

(b) 
$$T(n) = \Theta(log n)$$

(c) 
$$T(n) = \Theta(\sqrt{n})$$

(d) None of these.

- a
- O 6
- O c
- O d

Under what case of Masters theorem will the recurrence relation of merge sort fall?

- (a) 1st case
- (b)  $2^{nd}$  case
- (c) 3<sup>rd</sup> case
- (d) It cannot be solved using masters theorem
- O :
- b
- 0 0
- $\bigcirc$  a

If  $T(n) = T(n/4) + T(n/2) + n^2$ , then using recursion tree method

- (a)  $T(n) = \theta(n)$ .
- (b)  $T(n) = \theta(n^2)$ .
- (c) T(N) = θ(n³).
- (d) None of the above
- 0 :
- b
- 0 0
- $\bigcirc$  d

Recurrence relation for binary search

(a) 
$$T(n) = 2T(n/2) + \theta(1)$$

$$T(n) = T(n/2) + \theta(1)$$

(c) 
$$T(n) = 2T(n/2) + \theta(n)$$

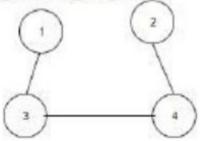
(d) 
$$T(n) = 2T(n/2) + \theta(n^2)$$

Assume that a merge sort algorithm in the worst case takes 30 seconds for an input of size 64. Which of the following most closely approximates the maximum input size of a problem that can be solved in 6 minutes?

- (a) 256
- (b) 512
- (c) 1024
- (d) 2048
- () E
- b
- $\bigcirc$   $\circ$
- $\bigcirc$   $\circ$

\*

What would be the number of zeros in the adjacency matrix of the given graph?



- (a) 10
- (b) 6
- (c) 16
- (d) 0
- a
- $\bigcirc$  b
- 0 0
- $\bigcirc$  d