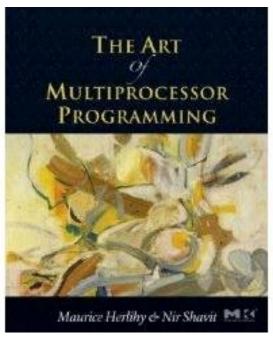


COS 226

Chapter 10
Concurrent Queues and the ABA Problem

Acknowledgement



Some of the slides are taken from the companion slides for "The Art of Multiprocessor Programming" by Maurice Herlihy & Nir Shavit



- Similar to Set class of Chapter 9 except:
 - Pool object does not necessarily have a contains() method
 - Same item can appear more than once



Bounded vs Unbounded

- Bounded
 - Fixed capacity
 - Good when resources an issue
- Unbounded
 - Holds any number of objects

Total, partial or synchronous

- Total:
 - A method is total if calls do not wait for conditions to become true
 - ☐ For example:
 - An attempt to dequeue from an empty queue throws an exception or failure code

Total, partial or synchronous

- Partial:
 - A method may wait for conditions to hold
 - ☐ For example:
 - Thread that tries to dequeue from an empty queue blocks (suspends) until item becomes available

Total, partial or synchronous

- Synchronous:
 - A method is synchronous if it waits for another method to overlap its call interval
 - ☐ For example:
 - A method call that adds an item to the pool is blocked until that item is removed by another method call
 - Used by CSP and Ada for threads to rendevous and exchange values

Blocking vs Non-Blocking

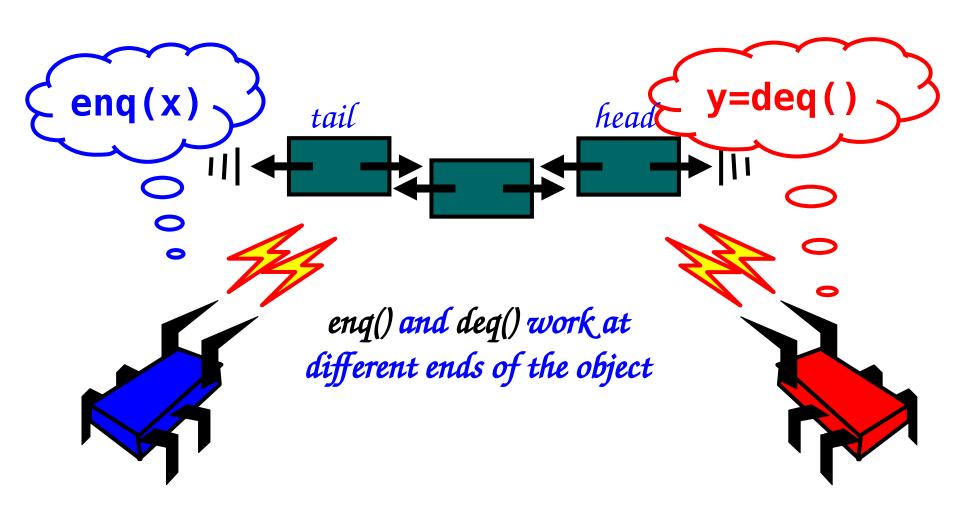
- Problem cases:
 - Removing from empty pool
 - Adding to full (bounded) pool
- Blocking
 - Caller waits until state changes
- Non-Blocking
 - Method throws exception



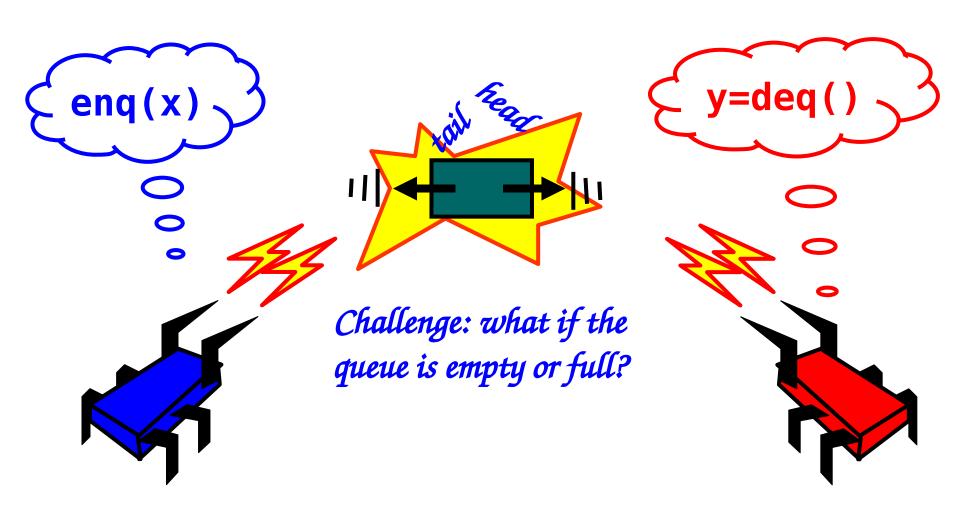
How much concurrency can we expect a bounded queue implementation with multiple concurrent enqueuers and dequeuers to have?

- Informally:
 - enq() and deq() methods operate on opposite ends of the queue
 - So an enq() call and a deq() call should be able to proceed concurrently without interference
 - Concurrent deq() calls and concurrent enq() calls will however interfere

Queue: Concurrency



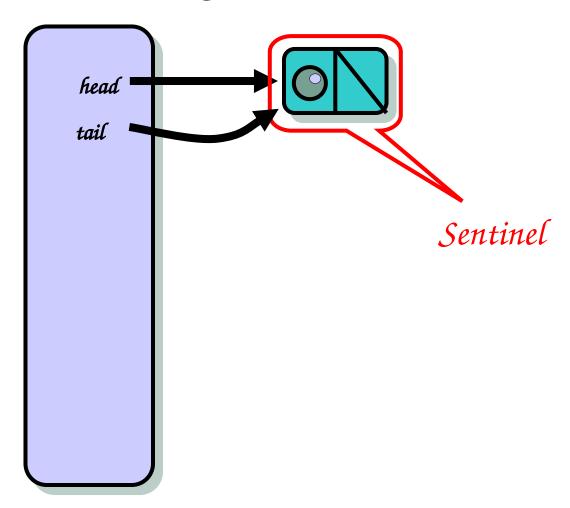
Concurrency

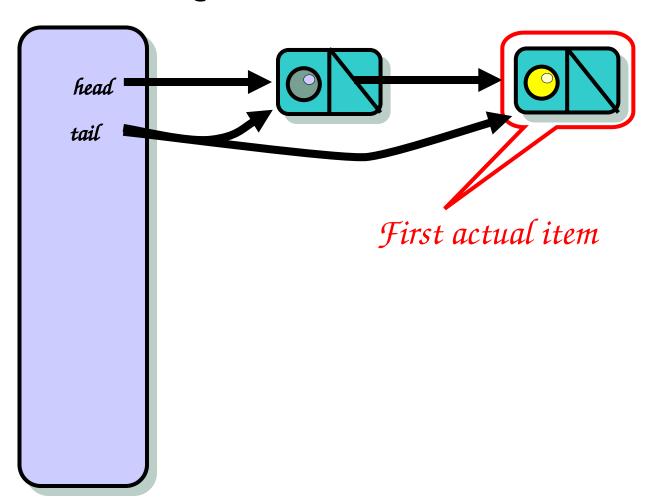


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Bounded lock-based queue

- Implement queue as linked list
- Queue has head and tail fields that refer to first and last nodes in list
- The queue contains a sentinel node a place-holder

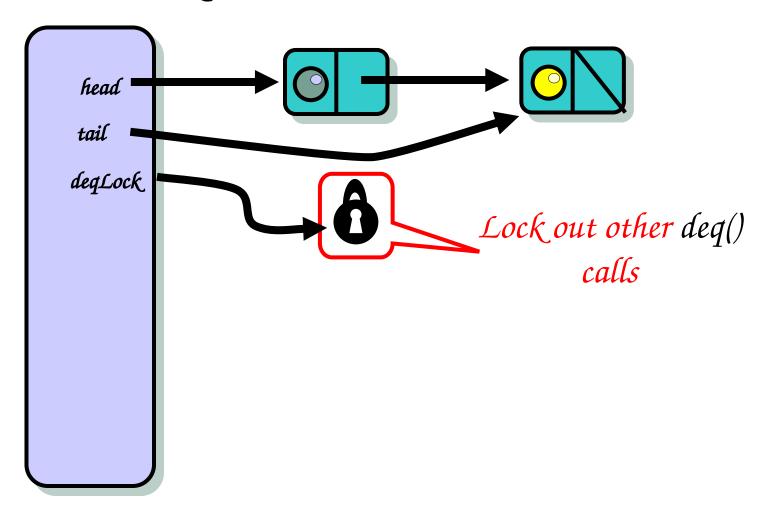


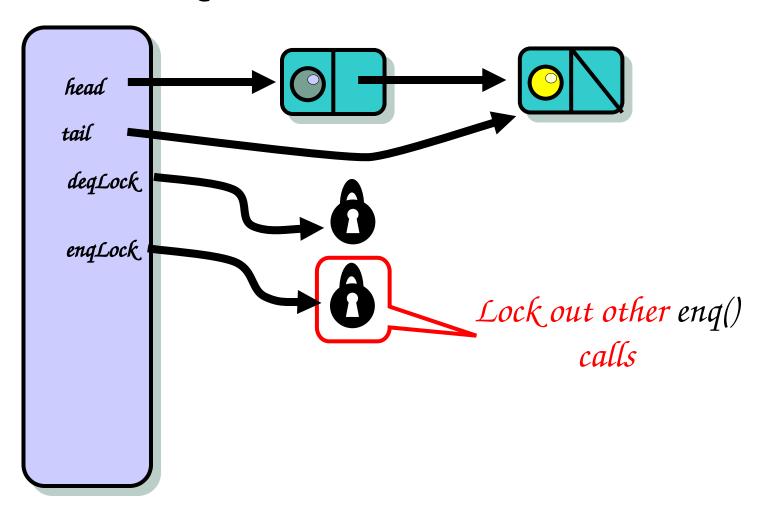


- Two distinct locks are used:
 - deqlock to lock the front of the queue to ensure that there is not more than one dequeuer at a time
 - enqlock to lock the end of the queue to ensure that there is not more than one enqueuer at a time

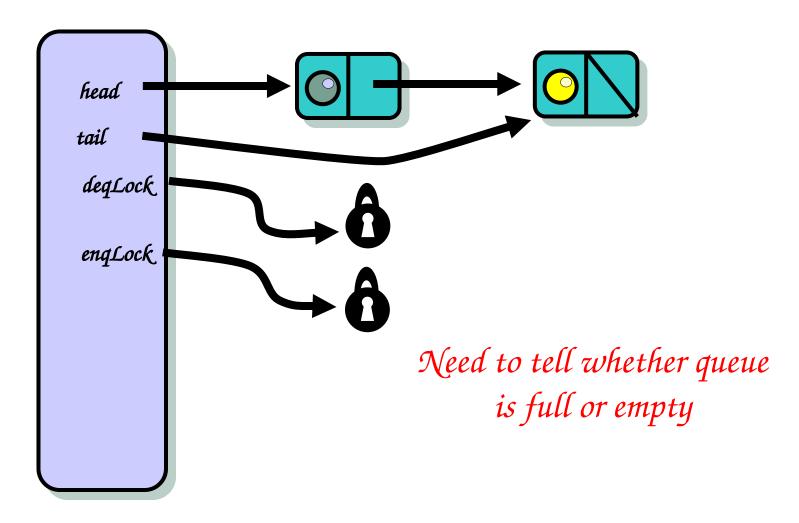


Using two locks instead of one means that enqueuers and dequeuers can operate independently and does not lock one another out unnecessarily





Not Done Yet



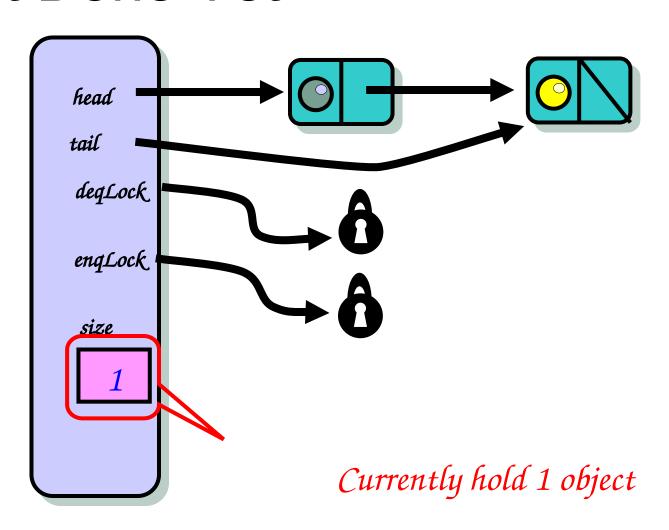
- enq() calls cannot continue while list is full
- deq() calls cannot continue while list is empty

- Each lock has an associated condition field:
 - enqlock is associated with the notFullCondition
 - Notifies enqueuers when list is not full anymore
 - deqlock is associated with the notEmptyCondition
 - Notifies dequeuers when list is not empty anymore

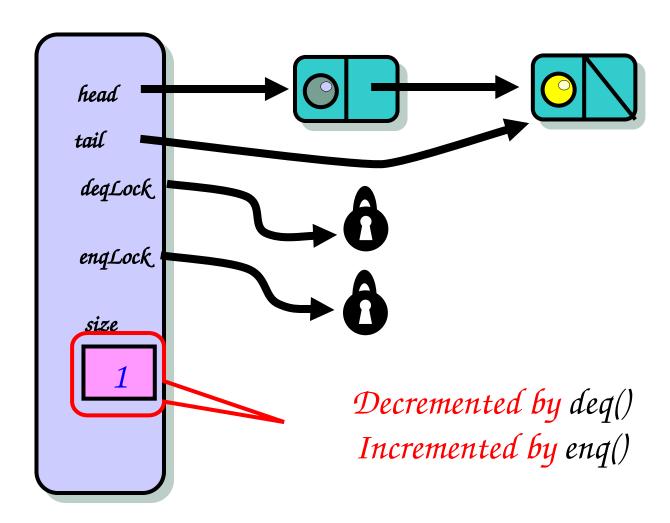


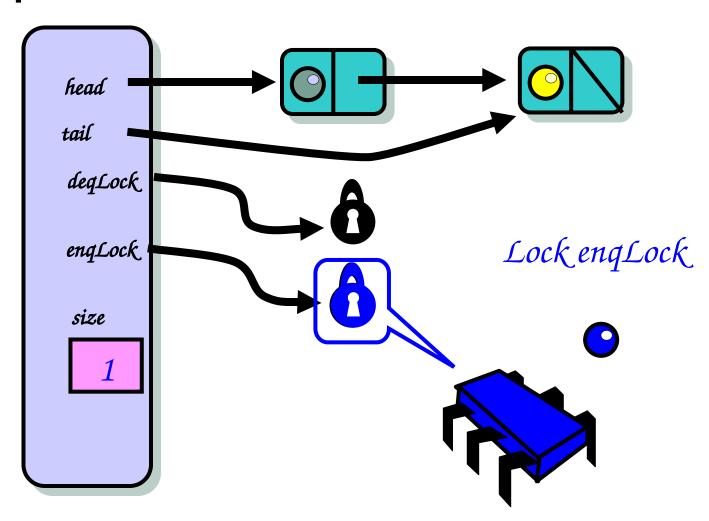
- Since the queue is bounded we must also keep track of the number of empty/available slots
- The size field keeps track of the number of objects currently in the queue
- The size field is an AtomicInteger

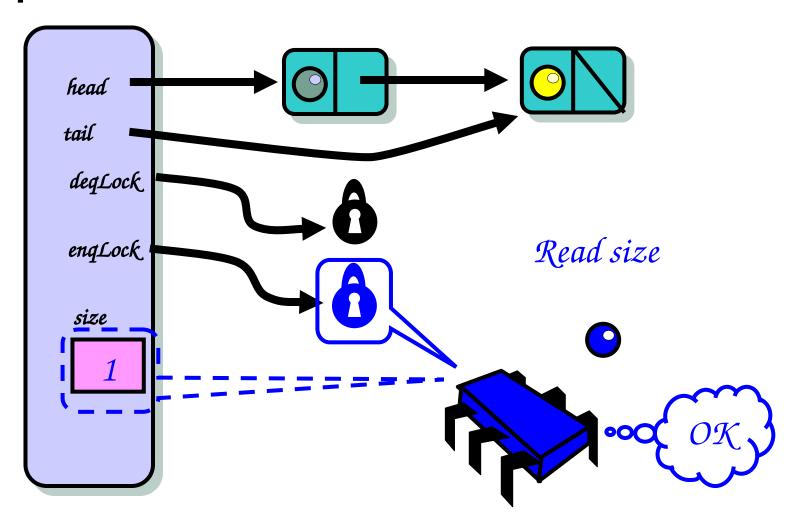
Not Done Yet

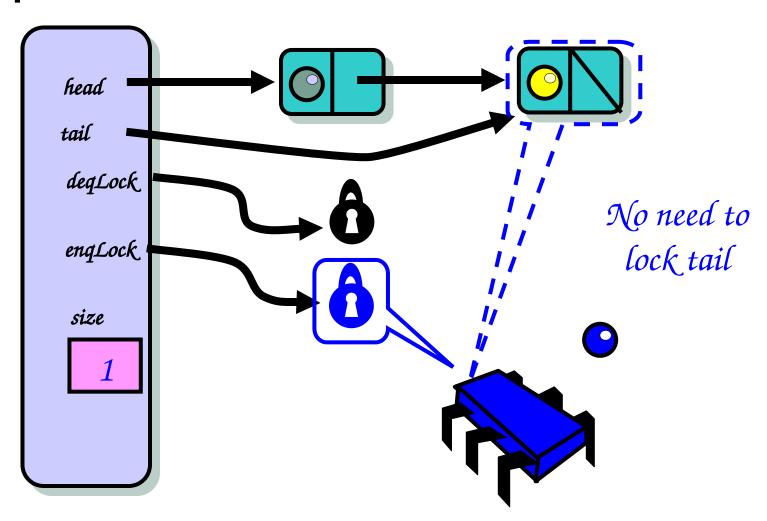


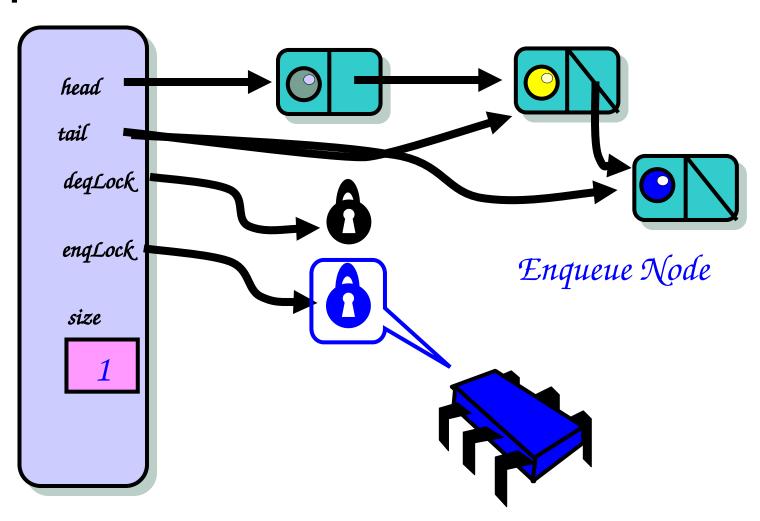
Not Done Yet

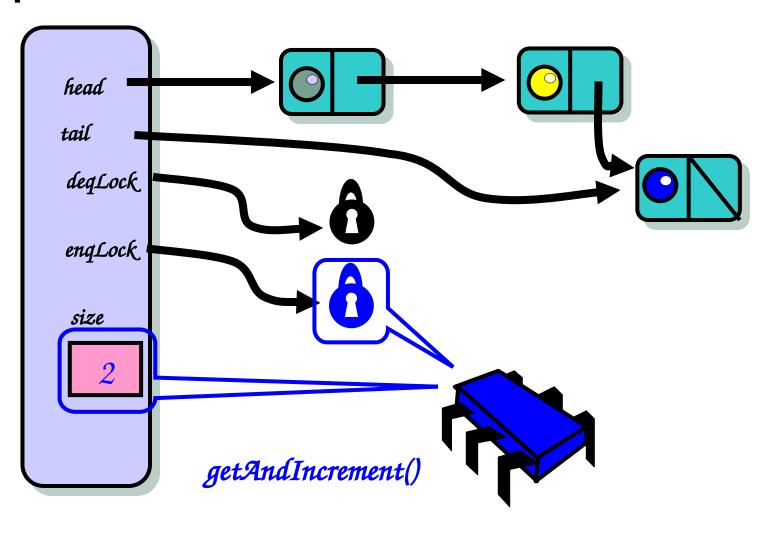


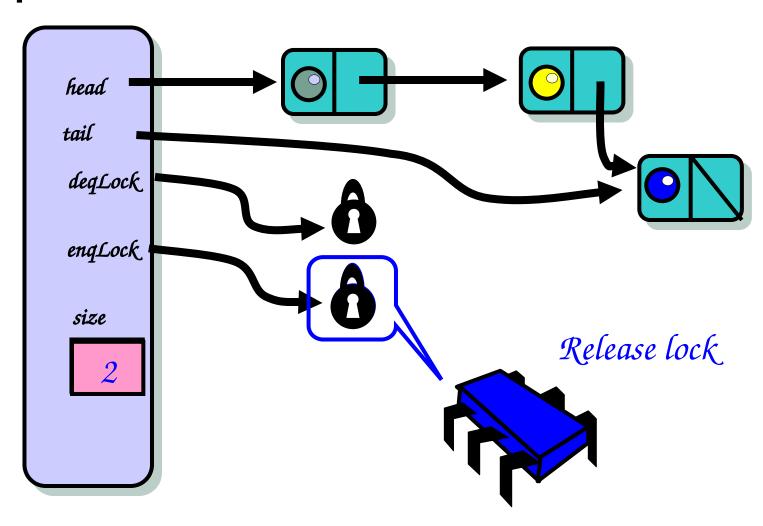


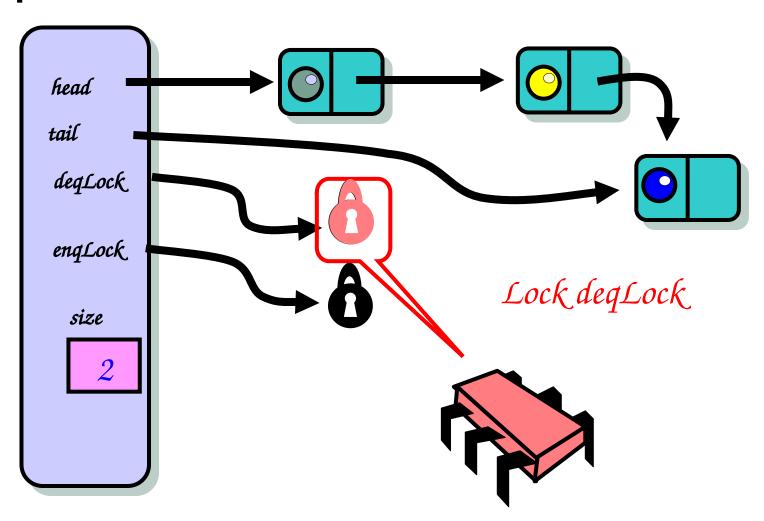


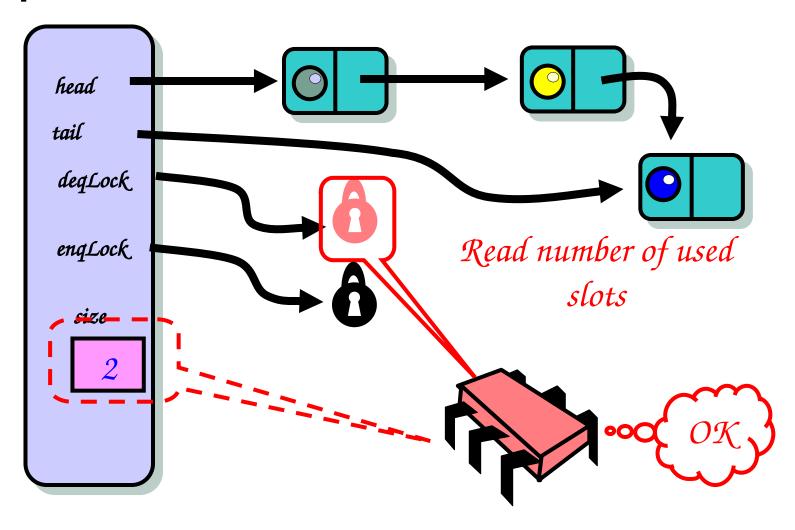


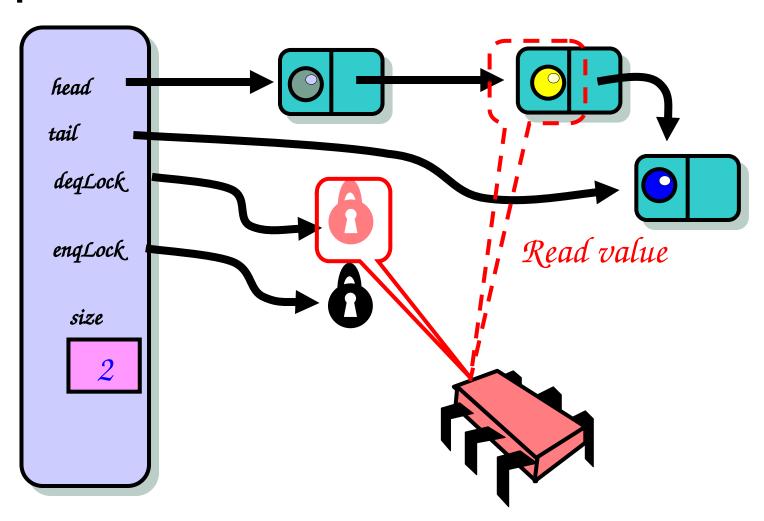




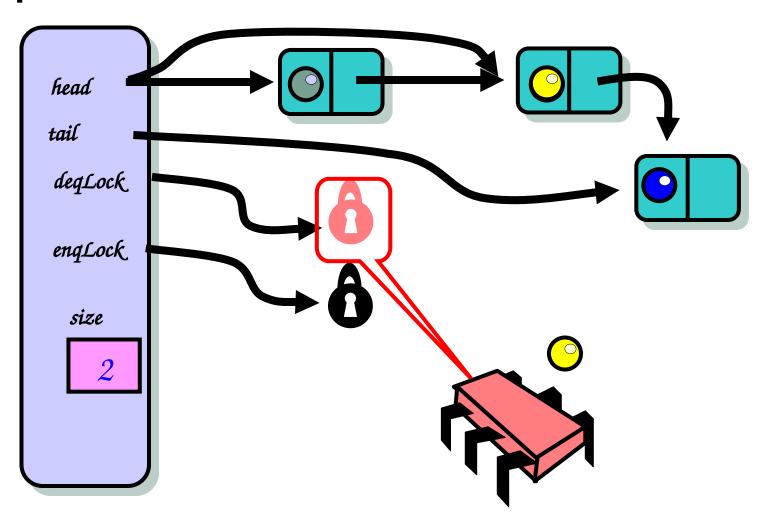


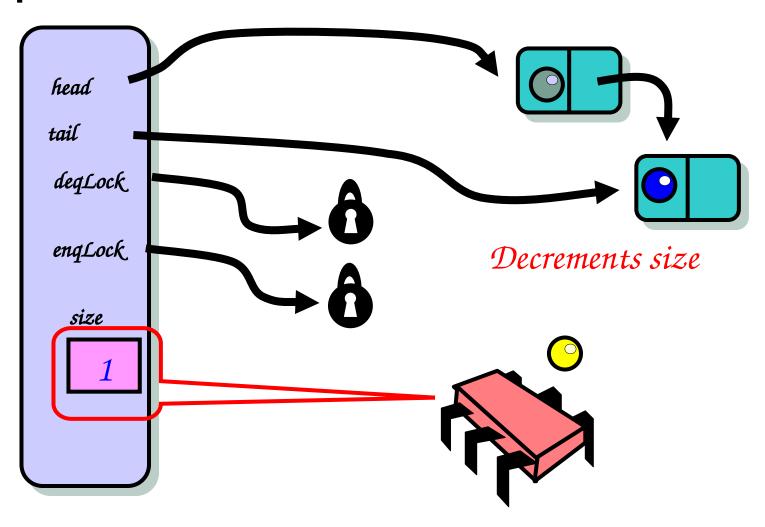




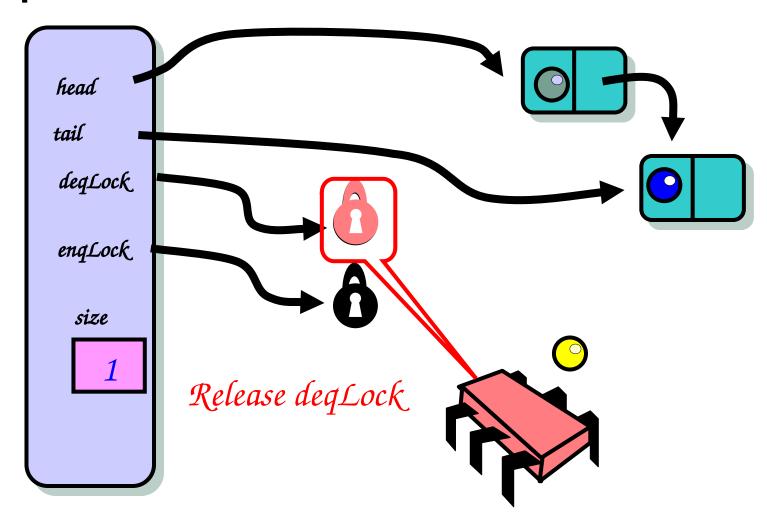


Make first Node new sentinel





Dequeuer





```
public class BoundedQueue<T> {
  ReentrantLock enqLock, deqLock;
  Condition notEmptyCondition, notFullCondition;
 AtomicInteger size;
  Node head;
 Node tail;
  int capacity;
  enqLock = new ReentrantLock();
  notFullCondition = enqLock.newCondition();
  deqLock = new ReentrantLock();
  notEmptyCondition = deqLock.newCondition();
```



```
public class BoundedQueue<T> {
 ReentrantLock enqLock, deqLock;
  Condition notEmptyCondition, notFullCondition;
 AtomicInteger size;
  Node head;
 Node tail;
  int capacity;
                               Eng & deq locks
 enqLock = new ReentrantLock();
  notFullCondition = enqLock.newCondition();
  deqLock = new ReentrantLock();
  notEmptyCondition = deqLock.newCondition();
```



```
public class BoundedQueue<T> {
  ReentrantLock engLock, degLock;
 Condition notEmptyCondition, notFullCondition;
 AtomicInteger size;
  Node head;
 Node tail;
  int capacity;
                                  Conditions
 enqLock = new ReentrantLock();
  notFullCondition = enqLock.newCondition();
  deqLock = new ReentrantLock();
  notEmptyCondition = deqLock.newCondition();
```



```
public class BoundedQueue<T> {
  ReentrantLock enqLock, deqLock;
  Condition notEmptyCondition, notFullCondition;
 AtomicInteger size;
  Node head;
 Node tail;
  int capacity;
  enqLock = new ReentrantLock();
  notFullCondition = enqLock.newCondition();
 deqLock = new ReentrantLock();
  notEmptyCondition = deqLock.newCondition();
```



```
public void enq(T x) {
  boolean mustWakeDequeuers = false;
  enqLock.lock();
  try {
    while (size.get() == capacity)
      notFullCondition.await();
    Node e = new Node(x);
    tail.next = e;
    tail = tail.next;
    if (size.getAndIncrement() == 0)
     mustWakeDequeuers = true;
  } finally {
    enqLock.unlock();
```



```
public void enq(T x) {
  boolean mustWakeDequeuers = false;
 enqLock.lock();
                                Lock and unlock enq
  try {
    while (size.get() == capacity)
      notFullCondition.await()
    Node e = new Node(x);
    tail.next = e;
    tail = tail.next
    if (size.getAndIncrement() == 0)
     mustWakeDequevers = true;
    finally {
    enqLock.unlock();
```

```
public void enq(T x) {
  boolean mustWakeDequeuers = false;
  enqLock.lock();
  try {
   while (size.get() == capacity)
      notFullCondition.await();
    Node e = new Node(x);
    tail.next = e;
    tail = tail.next;
    if (size.getAndIncrement() == 0)
      mustWakeDequeuers = \true;
  } finally {
    enqLock.unlock();
                  If queue is full, patiently await further
                             instructions ...
```

```
public void enq(T x) {
  boolean mustWakeDequeuers = false;
  enqLock.lock();
  try {
    while (size.get() == capacity)
      notFullCondition.await();
    Node e = new Node(x);
    tail.next = e;
    tail = tail.next;
    if (size.getAndIncrement() == 0)
      mustWakeDequeuers = true;
  } finally {
    enqLock.unlock();
                              Add new node
```

```
public void enq(T x) {
  boolean mustWakeDequeuers = false;
  enqLock.lock();
  try {
    while (size.get() == capacity)
      notFullCondition.await();
    Node e = new Node(x);
    tail.next = e;
    tail = tail next:
    if (size.getAndIncrement() == 0)
      mustWakeDequeuers = true;
   finally {
    enqLock.unlock()
               If queue was empty, allows dequeuers to be
```



Enq Method Part Deux

```
public void enq(T x) {
    ...
    if (mustWakeDequeuers)
notEmptyCondition.signalAll();
    }
}
```

Enq Method Part Deux

```
public void enq(T x) {
         if (mustWakeDequeuers) {
             notEmptyCondition.signalAll();
         }
     }
```

Are there dequeuers to be signaled?



```
public void enq(T x) {
         if (mustWakeDequeuers) {
             notEmptyCondition.signalAll();
         }
}
```

Signal dequeuers that queue no longer empty



Deq Method Part One

```
public T deq() {
  T result;
  boolean mustWakeEnqueuers = false;
  deqLock.lock();
  try {
    while (size.get() == 0)
      notEmptyCondition.await();
    result = head.next.value;
    head = head.next;
    if (size.getAndDecrement() == capacity)
     mustWakeEnqueuers = true;
  } finally {
    deqLock.unlock();
```



Deq Method Part Deux

```
public T deq() {
    ...
    if (mustWakeEnqueuers) {
        notFullCondition.signalAll();
    }
    return result;
    }
```

The Enq() & Deq() Methods

- Share no locks
 - ☐That's good
- But do share an atomic counter
 - Accessed on every method call
 - ☐ That's not so good
- Can we alleviate this bottleneck?

Split the Counter into two counters

- enqSideSize
 - Decremented by deq()
- deqSideSize
 - Incremented by enq()

Two counters

- The enq() method:
 - Checks enqSideSize and proceeds as long as less than capacity
 - Cares only if value is capacity
- The deq() method:
 - Checks deqSideSize and proceeds as long as greater than zero
 - Cares only if value is zero

Unbounded Queues

- Queue can hold unbounded number of items:
 - enq() always enqueues its item
 - deq() throws an EmptyException if there is no item to dequeue
- Same representation as bounded queue, only no need to count number of items



Unbounded queue enq()

```
public void enq(T x) {
   enqLock.lock();
   try {
     Node e = new Node(x);
     tail.next = e;
     tail = tail.next;
   } finally {
     enqLock.unlock();
   }
}
```

Unbounded queue deq()

```
public T deq() throws EmptyException {
  T result;
  deqLock.lock();
  try {
   if (head.next == null)
      throw new EmptyException();
    result = head.next.value;
    head = head.next;
  } finally {
    deqLock.unlock(); Only need to check that queue
                             is not empty
  return result;
```



Now an unbounded lock-free queue

- Natural extension of the unbounded queue
- However it prevents method calls from starving by having quicker threads help slower threads

Unbounded Lock-free Queue

- Again represented as a Linked List
- However each node's next field is now an AtomicReference (to facilitate lock-free operations)

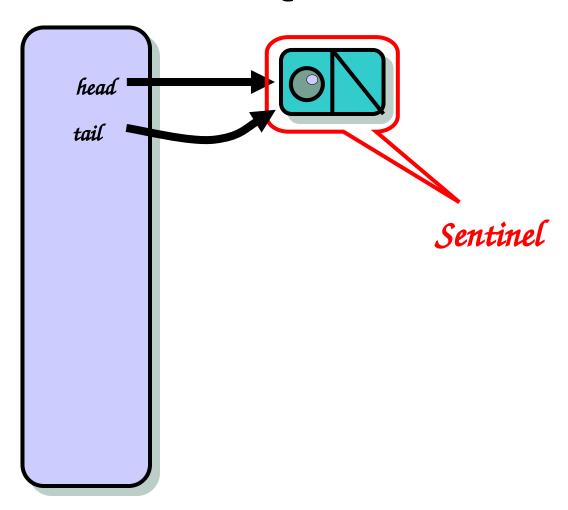


Lock-free queue Node class

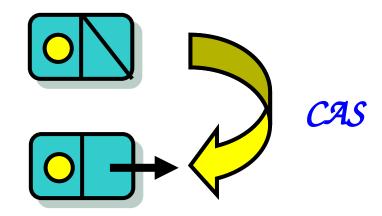
```
public class Node {
  public T value;
  public AtomicReference<Node> next;

public Node(T value) {
    value = value;
    next = new AtomicReference<Node>(null);
  }
}
```

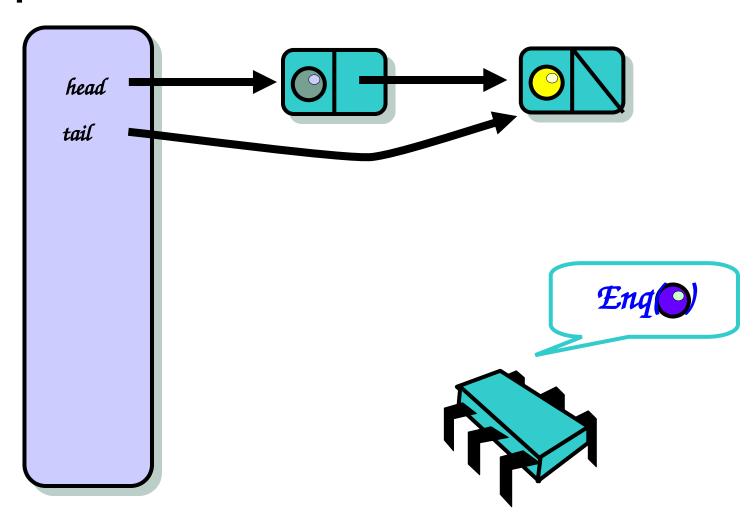
A Lock-Free Queue



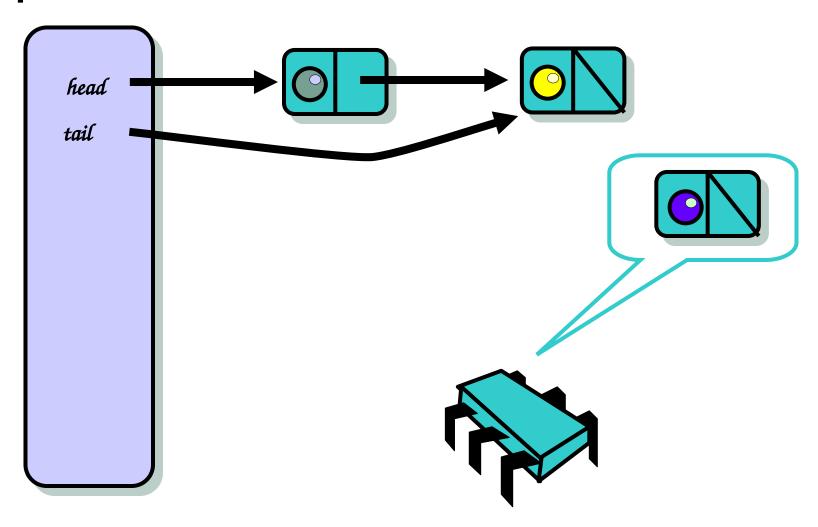
Compare and Set



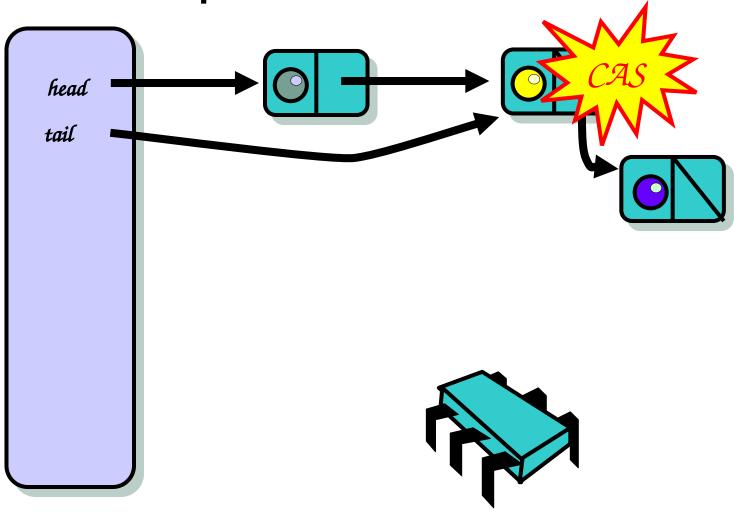
Enqueue



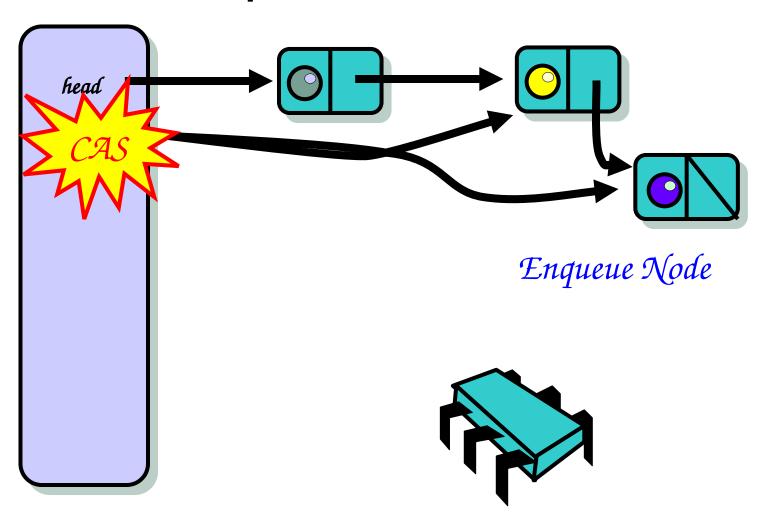
Enqueue



Logical Enqueue



Physical Enqueue



Enqueue

- The enq() method is lazy
- Enqueue consists of two steps:
 - Logical enqueue appends the node to the linked list
 - Physical enqueue changes the tail field to point to the new node
- These two steps are not atomic



- It is thus possible that a new item is reachable by other threads before it actually becomes the tail
- The tail field can thus refer to either
 - ☐ The actual last node or
 - ☐ The penultimate node



Enqueue

- When enqueueing you have to make sure that you are adding an item to the actual end of the queue
- So check whether the tail you are working with has a successor



- What do you do if you find
 - ☐ Tail pointing to the wrong value?
- Stop and help fix it
 - □ If **tail** node has non-*null* next field
 - CAS the queue's tail field to tail.next



Lock-free enqueue method

```
public void enq(T value) {
  Node node = new Node(value);
 while (true) {
    Node last = tail.get();
    Node next = tail.next.get();
    if (last == tail.get()) {
      if (next == null) {
        if (last.next.compareAndSet(next, node)) {
          tail.compareAndSet(last, node);
          return;
        } else
          tail.compareAndSet(last, last.next);
```

Lock-free enqueue method

```
public void enq(T value) {
 Node node = new Node(value); What does tail point to
 while (true) {
   Node last = tail.get();
    Node next = tail.next.get();
    if (last == tail.get()) {
      if (next == null) {
        if (last.next.compareAndSet(next, node)) {
          tail.compareAndSet(last, node);
          return;
        } else
          tail.compareAndSet(last, last.next);
```

Lock-free enqueue method

```
public void enq(T value) {
  Node node = new Node(value);
                                What does tail's next
 while (true) {
    Node last = tail.get();
   Node next = tail.next.get();
    if (last == tail.get()) {
      if (next == null) {
        if (last.next.compareAndSet(next, node)) {
          tail.compareAndSet(last, node);
          return;
        } else
          tail.compareAndSet(last, last.next);
```



```
public void enq(T value) {
                                  If the value pointed to
  Node node = new Node(value);
  while (true) {
                                      by tail has not changed...
    Node last = tail.get();
    Node next = tail.next_get();
   if (last == tail.get()) {
      if (next == null) {
        if (last.next.compareAndSet(next, node)) {
          tail.compareAndSet(last, node);
          return;
        } else
          tail.compareAndSet(last, last.next);
```



Lock-free enqueue method

```
public void enq(T value) {
  Node node = new Node(value); ... and tail's next points
  while (true) {
                                       to null...
    Node last = tail.get();
    Node next = tail.next.get();
    if (last == tail.get()) {
     if (next == null) {
        if (last.next.compareAndSet(next, node)) {
          tail.compareAndSet(last, node);
          return;
        } else
          tail.compareAndSet(last, last.next);
```



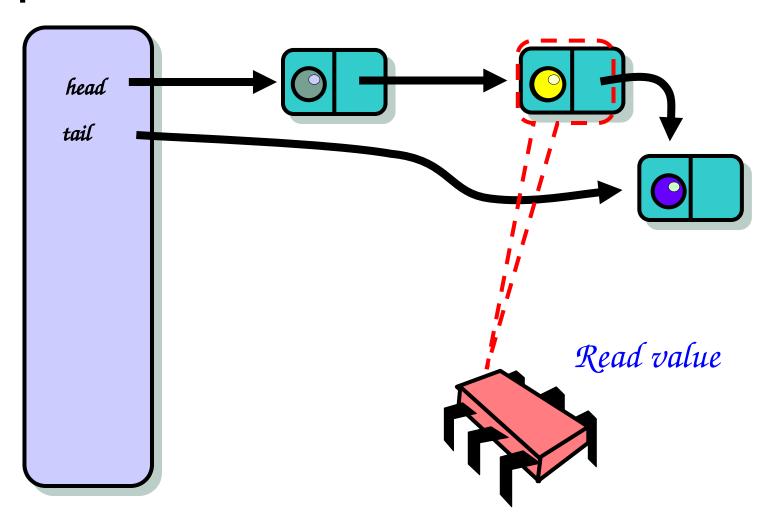
```
public void enq(T value) {
  Node node = new Node(value);
 while (true) {
    Node last = tail.get();
    Node next = tail.next.get();
    if (last == tail.get()) {
      if (next == null) {
       if (last.next.compareAndSet(next, node)) {
          tail.compareAndSet(last, node);
          return;
        } else
          tail.compareAndSet(last, last.next);
                      Logical enqueue
```



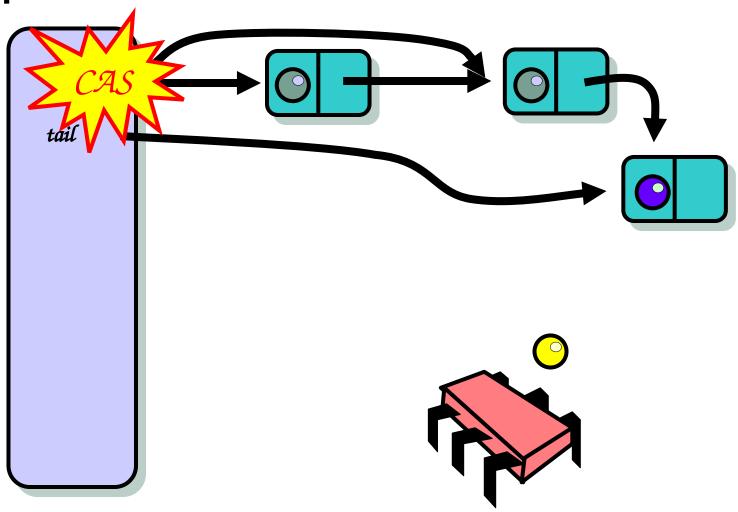
```
public void enq(T value) {
  Node node = new Node(value);
 while (true) {
   Node last = tail.get();
   Node next = tail.next.get();
    if (last == tail.get()) {
      if (next == null) {
        if (last.next.compareAndSet(next, node)) {
         tail.compareAndSet(last, node);
          return;
        } else
          tail.compareAndSet(\last.next);
```



```
public void enq(T value) {
  Node node = new Node(value);
  while (true) {
    Node last = tail.get();
    Node next = tail.next.get();
    if (last == tail.get()) {
      if (next == null) {
        if (last.next.compareAndSet(next, node)) {
          tail.compareAndSet(last, node);
          return;
          else
          tail.compareAndSet(last, last.next);
          If tail has a successor it 'helps' other nodes by
             advancing tail to refer to the successor
```



Make first Node new sentinel



×

- Although queue is unbounded, you always have to check that queue is not empty
- Without counter, queue is empty when head = tail = sentinel

Lock-free dequeue method

```
public T deq() throws EmptyException {
 while (true) {
    Node first = head.get();
    Node last = tail.get();
    Node next = first.next.get();
    if (first == head.get()) {
      if (first == last) {
        if (next == null) {
          throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        T value = next.value;
        if (head.compareAndSet(first, next))
          return value;
```



```
public T deq() throws EmptyException {
 while (true) {
                                What does head point
   Node first = head.get()
    Node last = tail.get();
    Node next = first.next.get();
    if (first == head.get()) {
      if (first == last) {
        if (next == null) {
          throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        T value = next.value;
        if (head.compareAndSet(first, next))
          return value;
```



```
public T deq() throws EmptyException {
  while (true) {
                                What does tail point to
    Node first - head.get():
   Node last = tail.get();
    Node next = first.next.get();
    if (first == head.get()) {
      if (first == last) {
        if (next == null) {
          throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        T value = next.value;
        if (head.compareAndSet(first, next))
          return value;
```



```
public T deq() throws EmptyException {
 while (true) {
                                 Is there a node after
   Node first = head.get();
    Node last = tail.get();
   Node next = first.next.get();
    if (first == head.get()) {
      if (first == last) {
        if (next == null) {
          throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        T value = next.value;
        if (head.compareAndSet(first, next))
          return value;
```



```
public T deq() throws EmptyException {
  while (true) {
                                Has the value of head
    Node first = head.get();
    Node last = tail.get();
    Node next = first_next.get();
   if (first == head.get()) {
      if (first == last) {
        if (next == null) {
          throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        T value = next.value;
        if (head.compareAndSet(first, next))
          return value;
```



```
public T deq() throws EmptyException {
 while (true) {
                                 Is the queue empty?
    Node first = head.get();
    Node last = tail.get();
    Node next = first.next.get():
    if (first == head.get()) {
     if (first == last) {
        if (next == null) {
          throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        T value = next.value;
        if (head.compareAndSet(first, next))
          return value;
```



```
public T deq() throws EmptyException {
 while (true) {
                               Is the queue still empty?
    Node first = head.get();
    Node last = tail.get();
    Node next = first.next.get():
    if (first == head.get())
      if (first == last)
       if (next == null) {
          throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        T value = next.value;
        if (head.compareAndSet(first, next))
          return value;
```



```
public T deq() throws EmptyException {
                                       If a new value node have arrived, 'help'
  while (true) {
    Node first = head.get();
    Node last = tail.get();
    Node next = first.next.get(); others by advancing tail
    if (first == head.get()) {
      if (first == last) {
        if (next == null) {
           throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        T value = next.value;
        if (head.compareAndSet(first, next))
           return value;
```



```
public T deq() throws EmptyException {
  while (true) {
    Node first = head.get();
    Node last = tail.get();
    Node next = first.next.get();
    if (first == head.get()) {
                                     If the queue is not
      if (first == last) {
        if (next == null) {
          throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
        <del>T value</del> = next.value;
        if (head.compareAndSet(first, next))
          return value;
```

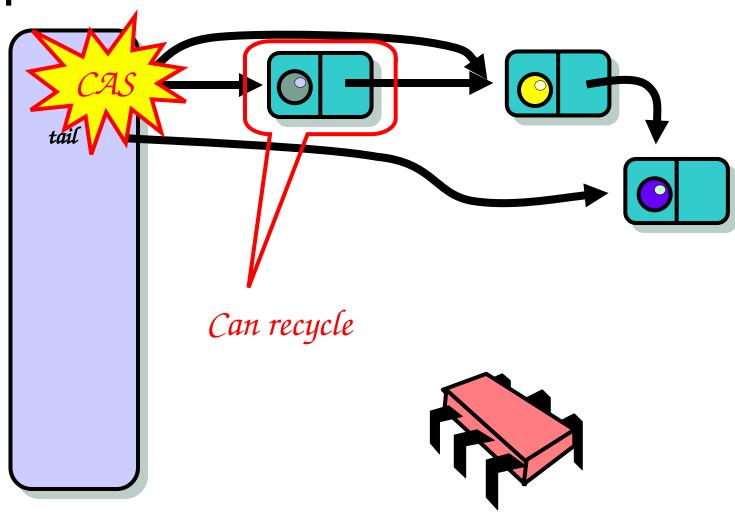


```
public T deq() throws EmptyException {
  while (true) {
    Node first = head.get();
    Node last = tail.get();
                                      Try to remove the first value from the queue
    Node next = first.next.get();
    if (first == head.get()) {
      if (first == last) {
        if (next == null) {
           throw new EmptyException();
        tail.compareAndSet(last, next);
      } else {
         value = next.value;
        if (head.compareAndSet(first, next))
           return value;
```



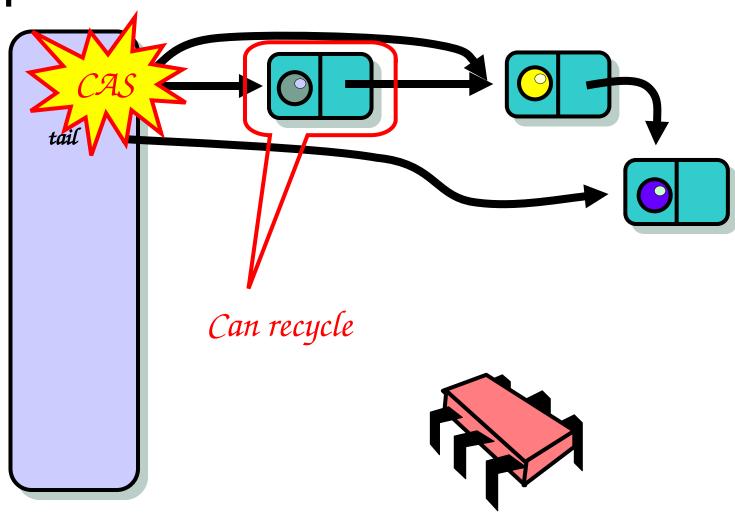
Memory reuse

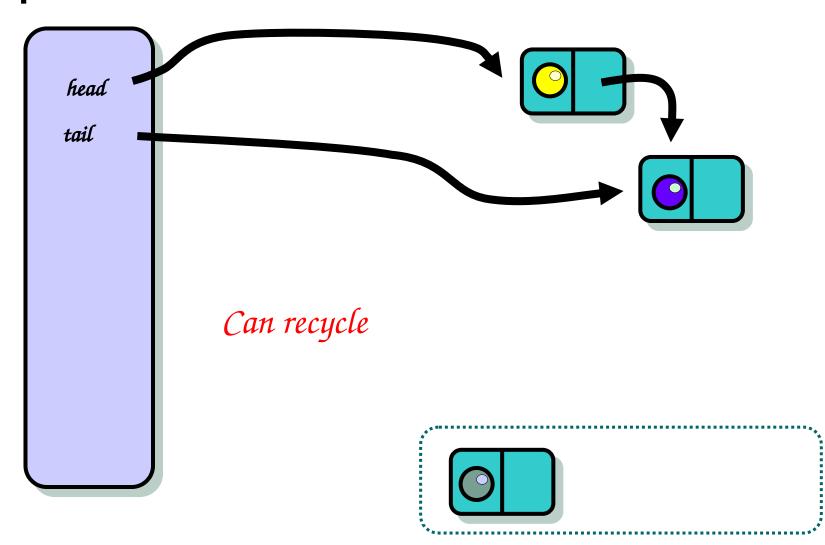
- What do we do with nodes after we dequeue them?
- With current implementation nodes are unlinked from linked list, but we depend on built-in garbage collection to remove them from memory
- Suppose there is no garbage collector?



Simple solution

- Each thread maintains its own private freelist of unused queue entries
 - □ When enqueue get a node from the free-list
 - ■When dequeue add removed node to free-list
- If free-list is empty, simply allocate new node through dynamic memory allocation

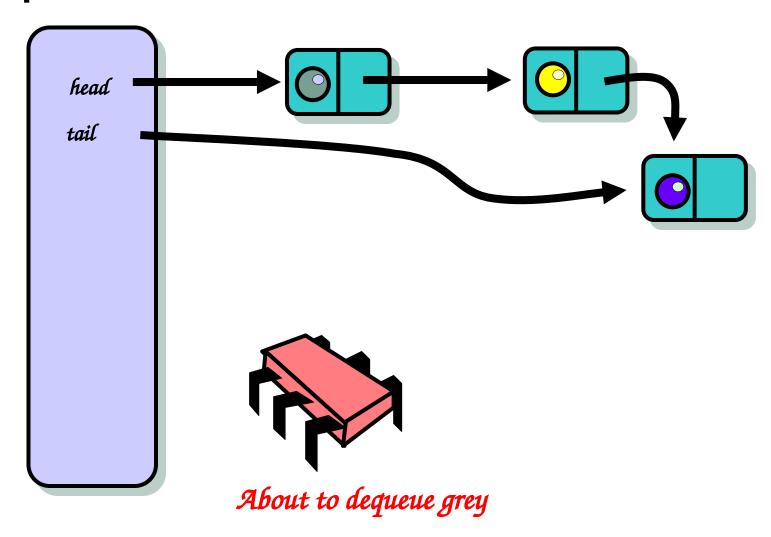


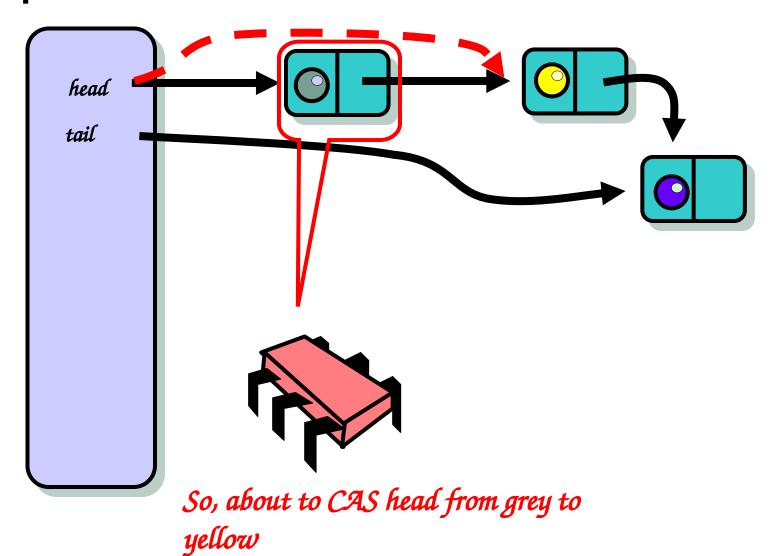


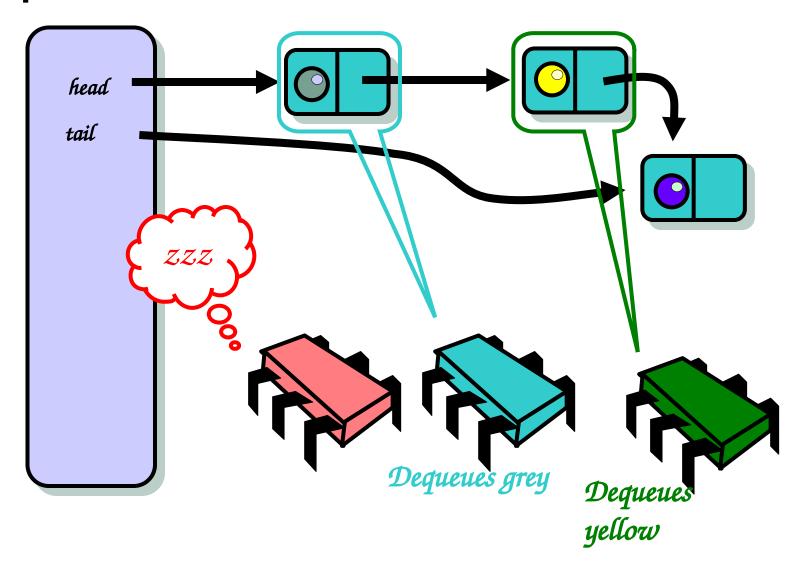


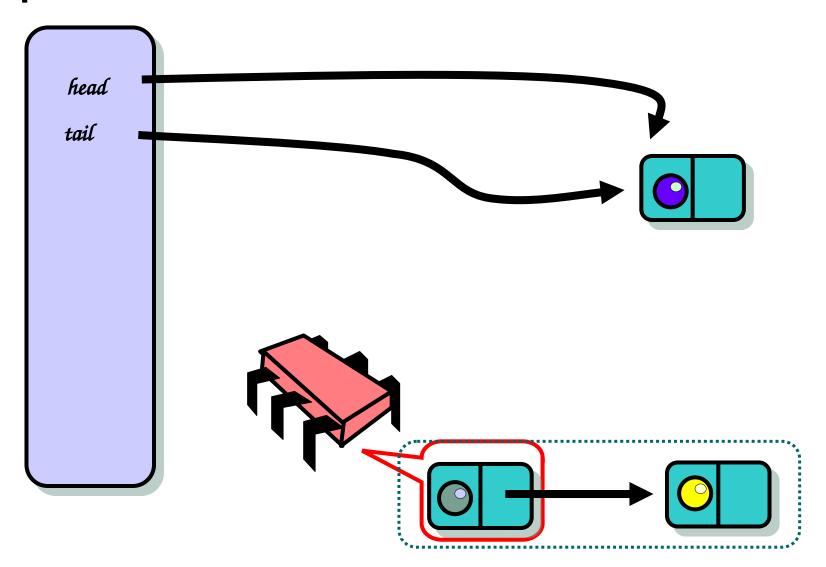
Potential problem

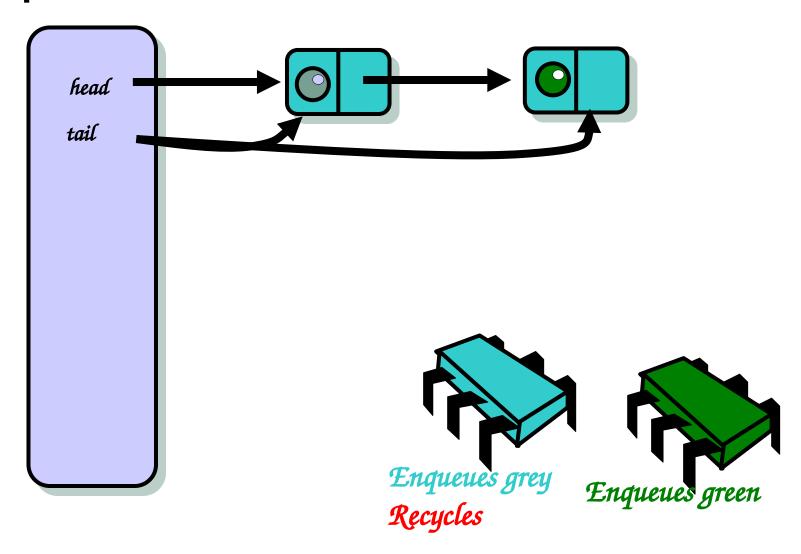
- When nodes are recycled removed from the queue and later added again
 - Dreaded ABA problem



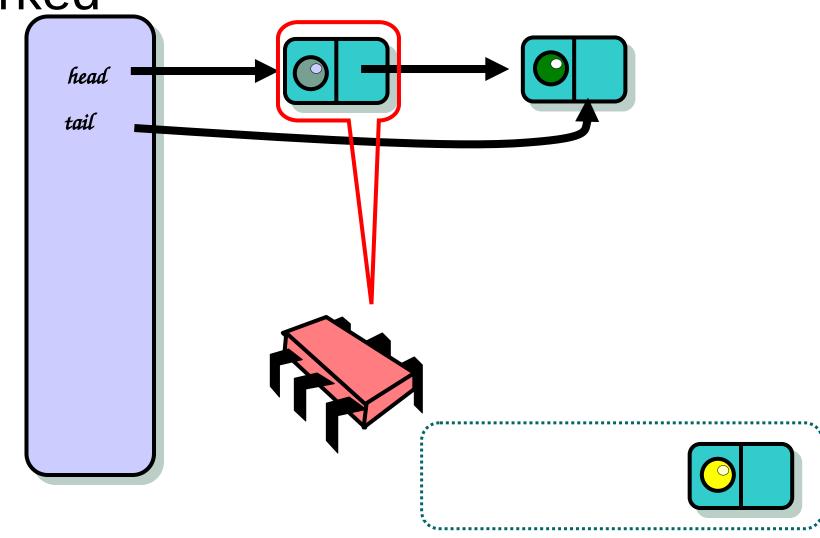


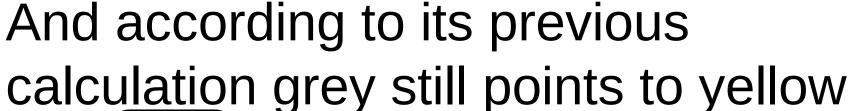


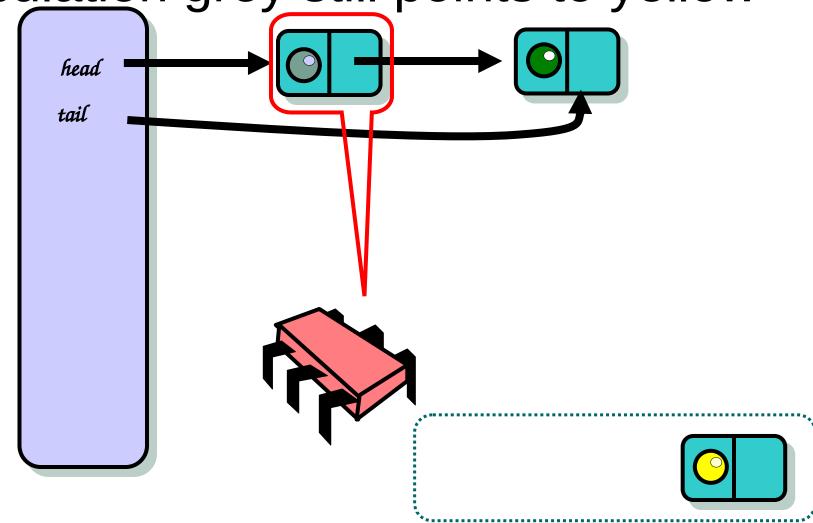


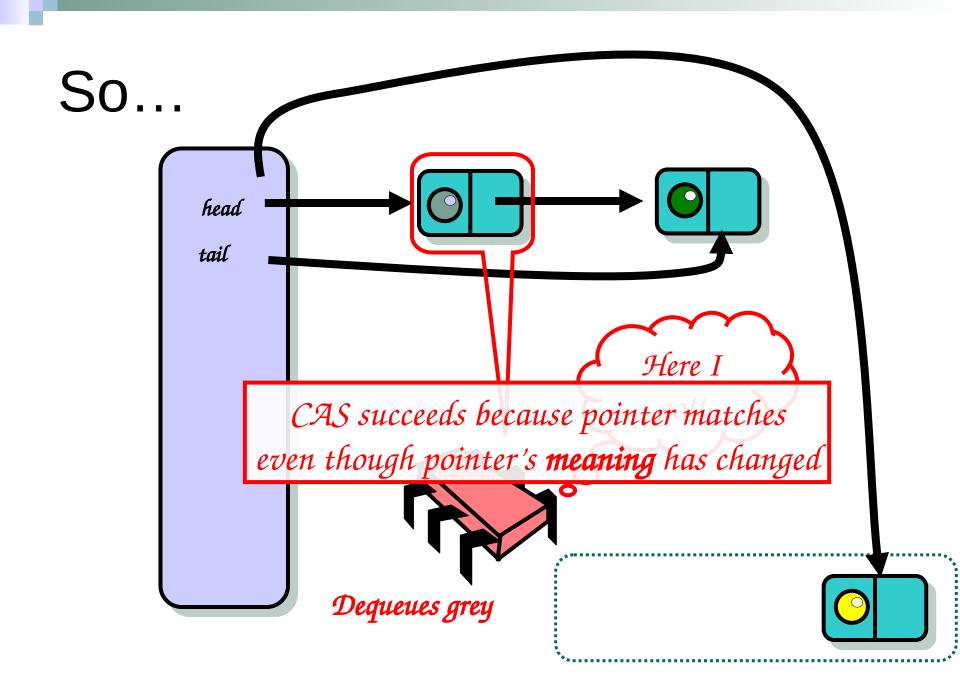


But the red thread still had grey marked











ABA problem

Head ends up pointing to a node that is not in the queue anymore, but has been added to the free-list

Dreaded ABA - A Solution

- Tag each pointer with a counter
- Unique over lifetime of node
- Pointer size vs word size issues
- Overflow?
 - □ Don't worry be happy?
 - □ Bounded tags?
- AtomicStampedReference class



Dual Data Structures

To reduce the synchronization overhead of the synchronous queue we split enq() and deq() into two steps

deq() method

- If a dequeuer tries to remove an item from an empty queue:
 - It puts a reservation object in the queue to indicate that the deq() is waiting for an enq() to rendezvous with
 - \Box deq() spins on flag in reservation

enq() method

- When enq() sees reservation:
 - ☐ Fulfills the reservation by depositing an item
 - Notifies deq() by setting object's flag



Dual Data Structure

- An enq() can also add its own reservation to the queue and spin on the flag waiting for a deq()
- At any time the queue contains either enq() reservations, deq() reservations or is empty

Dual Data Structure

- Methods take effect in two stages:
 - Reservation and fulfillment



Dual Data structures

- Good things:
 - Waiting threads spin on a locally cached flag
 - Ensures fairness
 - Linearizable



Node class

```
private enum NodeType {ITEM, RESERVATION};
private class Node {
      volatile NodeType type;
      volatile AtomicReference<T> item;
      volatile AtomicReference<Node> next;
      Node (T myItem, NodeType myType) {
            item = new AtomicReference<T> (myItem);
            next = new AtomicReference<Node>(null);
            type = myType;
```

```
public void enq(T e) {
      Node offer = new Node (e, ITEM);
      while (true) {
        Node t = tail.get(), h = head.get();
        if (h == t || t.type == ITEM) {
          Node n = t.next.get();
          if (t == tail.get()) {
            if (n != null) {
              tail.compareAndSet(t, n);
            } else if (t.next.CAS(n, offer)) {
              tail.compareAndSet(t, offer);
              while (offer.item.get() == e) {}
              h = head.get();
              if (offer == h.next.get())
                head.compareAndSet(h, offer);
              return; }}
```



Create new node

```
public void eng(T e) {
      Node offer = new Node (e, ITEM);
      while (true) {
        Node t = tail.get(), h = head.get();
        if (h == t || t.type == ITEM) {
          Node n = t.next.get();
          if (t == tail.get()) {
            if (n != null) {
              tail.compareAndSet(t, n);
            } else if (t.next.CAS(n, offer)) {
              tail.compareAndSet(t, offer);
              while (offer.item.get() == e) {}
              h = head.get();
              if (offer == h.next.get())
                head.compareAndSet(h, offer);
              return; }}
```



If the queue is empty OR if it contains only ITEMS...

```
public void eng(T e) {
      Node offer = new Node (e/I)
      while (true) {
        Node t = tail.get(), h = head.get();
       if (h == t || t.type == ITEM) {
          Node n = t.next.get();
          if (t == tail.get()) {
            if (n != null) {
              tail.compareAndSet(t, n);
            } else if (t.next.CAS(n, offer)) {
              tail.compareAndSet(t, offer);
              while (offer.item.get() == e) {}
              h = head.get();
              if (offer == h.next.get())
                head.compareAndSet(h, offer);
              return; }}
```



```
public void enq(T e) {
      Node offer = new Node (e ITEM);
     while (true) {
        Node t = tail.get(), h = head.get();
       if (h == t || t.type == ITEM) {
          Node n = t.next.get();
          if (t == tail.get()) {
            if (n != null) {
              tail.compareAndSet(t, n);
            } else if (t.next.CAS(n, offer)) {
              tail.compareAndSet(t, offer);
              while (offer.item.get() == e) {}
              h = head.get();
              if (offer == h.next.get())
                head.compareAndSet(h, offer);
              return; }}
```

```
public void enq(T e) {
     Node offer = new Node (e//ITEM);
     while (true) {
        Node t = tail.get(), h/= head.get();
        if (h == t || t.t/pe = ITEM) {
          Node n = t.nex(t.get/();
         if (t == tail.get()) {
           if (n != null) {
              tail.compareAndSet(t, n);
            } else if (t.next.CAS(n, offer)) {
              tail.compareAndSet(t, offer);
              while (offer.item.get() == e) {}
              h = head.get();
              if (offer == h.next.get())
                head.compareAndSet(h, offer);
              return; }}
```

```
public void enq(T e) {
      Node offer = new Node (e, ITEM);
      while (true) {
        Node t = tail.get(), h = head, get();
        if (h == t || t.type == ITEM)
          Node n = t.next.get();
          if (t == tail.get()) {
            if (n != null) {
              tail.compareAndSet/(t, n);
            } else if (t.next.CAS(n, offer)) {
              tail.compareAndSet(t, offer);
              while (offer.item.get() == e) {}
              h = head.get();
              if (offer == h.next.get())
                head.compareAndSet(h, offer);
              return; }}
```



```
public void enq(T e) {
      Node offer = new Node (e, ITEM);
      while (true) {
        Node t = tail.get(), h = head.get();
        if (h == t || t.type == ITEM) {
          Node n = t.next.get();
          if (t == tail.get()) {
            if (n != null) {
              tail.compareAndSet(t, p/);
            } else if (t.next.CAS(p, offer)) {
              tail.compareAndSet(t, offer);
             while (offer.item.get() == e) {}
              h = head.get();
              if (offer == h.next.get())
                head.compareAndSet(h, offer);
              return; }}
```



After it is dequeued remove node from queue

```
public void eng(T e) {
      Node offer = new Node (e, ITEM);
     while (true) {
        Node t = tail.get(), h = head.get();
        if (h == t || t.type == ITEM) {
          Node n = t.next.get();
          if (t == tail.get()) {
            if (n != null) {
              tail.compareAndSet(t, n)
            } else if (t.next.CAS(n, offer)) {
              tail.compareAndSet(t/offer);
              while (offer.item.get() == e) {}
              h = head.get();
             if (offer == h.next.get())
                head.compareAndSet(h, offer);
              return; }}
```



```
} else {
    Node n = h.next.get();
    if (t != tail.get() || h != head.get() || n ==
      null) {
      continue;
    boolean success = n.item.compareAndSet(null,e);
    head.compareAndSet(h, n);
    if (success)
      return;
}}}
```



Otherwise fulfill reservation

```
else {
    Node n = h.next.get();
    if (t != tail.get() || h != head.get() ||
      null) {
      continue;
    boolean success = n.item.compareAndSet(null,e);
    head.compareAndSet(h, n);
    if (success)
      return;
}}}
```



Make sure nothing has changed

```
} else {
    Node n = h.next.get():
    if (t != tail.get() || h != head.get()
      null) {
      continue;
    boolean success = n.item.compareAndSet(null,e);
    head.compareAndSet(h, n);
    if (success)
      return;
}}}
```



Change item in

```
reservation node to enq
  } else {
                                         item
    Node n = h.next.get();
    if (t != tail.get() || h != head.get() ||
      null) {
      continue;
    boolean success = n.item.compareAndSet(null,e);
    head.compareAndSet(h, n);
    if (success)
      return;
}}}
```