

## The Prim - Tarnik's Algo: -

1. Start with a single cluster  $C$  having some "root" vertex  $v$ .  $C = \{v\}$
2. Choose a minimum-weight edge  $e = \{u, v\}$  connecting  $v$  in the cloud  $C$  and  $u$  outside the cloud  $C$
3. Add vertex  $u$  to the cloud  $C = \{u, v\}$
4. Repeat until a MST is formed.

For efficiency we have a label  $D[u]$  for each vertex  $u$  outside the cloud  $C$ .

- stores the weight of the best current edge for joining  $u$  to the cloud  $C$ . [the smallest wt of an edge connecting  $u$  to a vertex in cloud]
- ∴ this reduces the no. of edges we need to consider to decide which vertex should join the cloud next.

## Algorithm Prim's Min MST (G):

Input: A weighted connected graph  $G$  with  $n$  vertices and  $m$  edges.

Output: A minimum spanning tree  $T$  for  $G$ .

Pick any vertex  $u \notin V$  do

$$D[u] \leftarrow \infty$$

for each vertex  $u \notin V$  do

$$D[u] \leftarrow \infty$$

Initialize  $T \leftarrow \emptyset$ .

Initialize a priority queue  $Q$  with an item  $((u, \text{null}), D[u])$  for each vertex  $u$ ,

where  $(u, \text{null})$  is the element and

$D[u]$  is the key.

while  $Q$  is not empty do

$$(u, e) \leftarrow Q.\text{removeMin}$$

Add vertex  $u$  and edge  $e$  to  $T$ .

for each vertex  $z$  adjacent to  $u$  such  
that  $z$  is in  $Q$  do

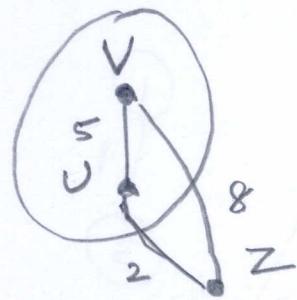
if  $w(u, z) < D[z]$  then

$$D[z] \leftarrow w(u, z)$$

Change the element of

vertex  $z$  in  $\pi$  to

$$(z, (u, z))$$



Here  $D[z]$  is  
changed to 7.

change the key of vertex  $z$  in  $\pi$  to  
 $D[z]$

return the tree  $T$ .

Running Time :-

$O(m \log n)$

Source 1 ..

$$D[1] = 0$$

$$\begin{array}{|c|c|} \hline V & D[V] \\ \hline 1 & 0 \\ 2 & \infty \\ 3 & \infty \\ 4 & \infty \\ 5 & \infty \\ 6 & \infty \\ \hline \end{array}$$

$$T = \emptyset$$

$T = \{1\}$

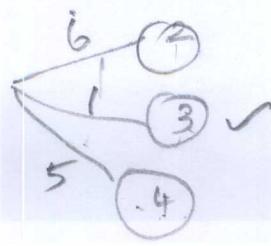
$m = 1$

$(3)$

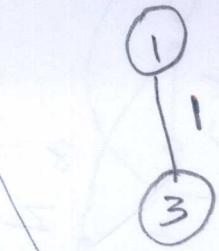
Step 1 :

$$T = \{1\}$$

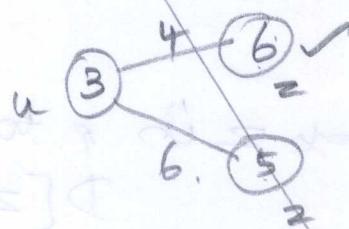
$$m = 1$$



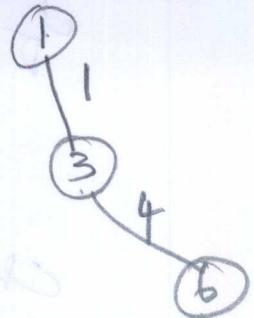
(3)



$$u = 3$$

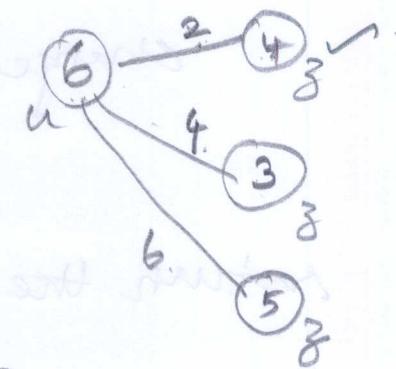


V	D[V]
1	0
2	(6)
3	1 ✓
4	5
5	0 ←
6	0 ← ✓

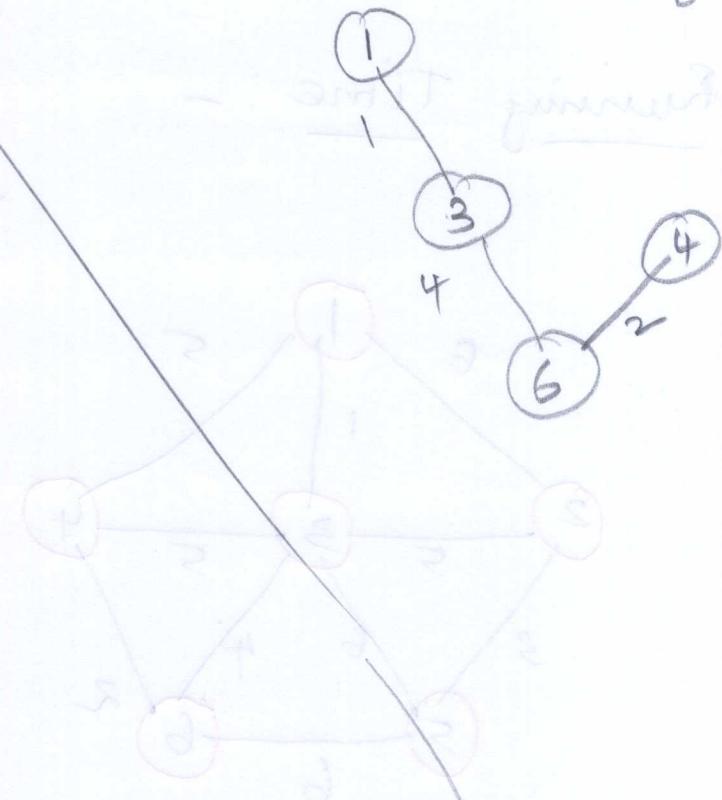


$$u = 6.$$

$$T = \{1, 3, 6\}$$

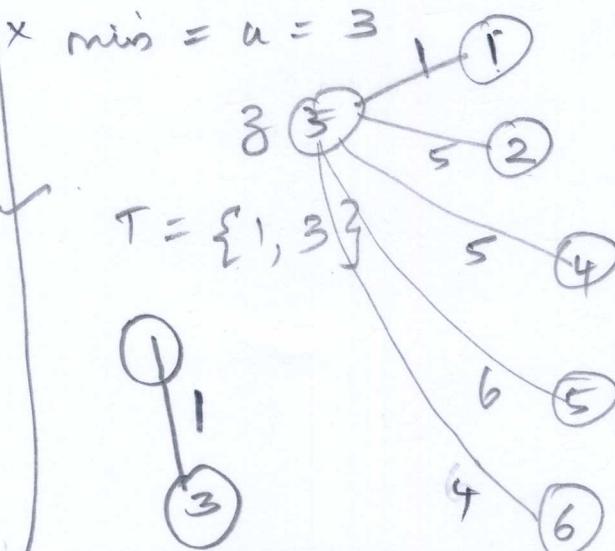


$v$	$D[r]$
1	0
2	6
3	1
4	2
5	6



V	D[V]
1	0
2	6
3	1
4	5
5	2
6	1

STEP 2:

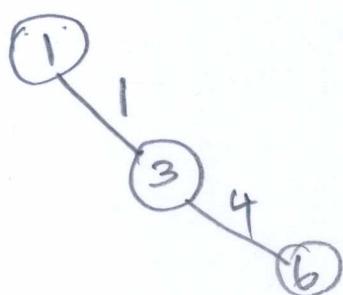


V	D[V]
1	0
2	6
3	1
4	5
5	2
6	1

STEP 3:

$$\min = u = 6$$

$$T = \{1, 3, 6\}$$

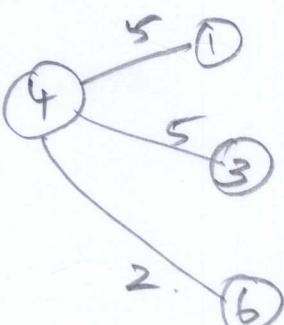
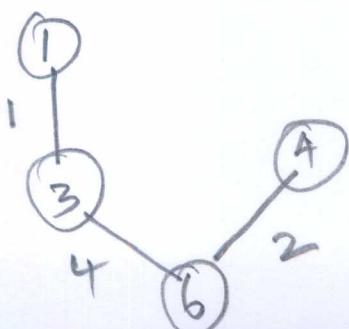


V	D[V] <sup>old</sup>	D[V] <sup>new</sup>
1	0	0
2	5	5
3	1	1
4	2	2
5	6	6
6	4	4

STEP 4:

$$\min = u = 4$$

$$T = \{1, 3, 6, 4\}$$



V	D[V] <sup>old</sup>	D[V] <sup>new</sup>
1	0	x
2	5	5
3	1	x
4	2	2
5	6	6
6	4	2

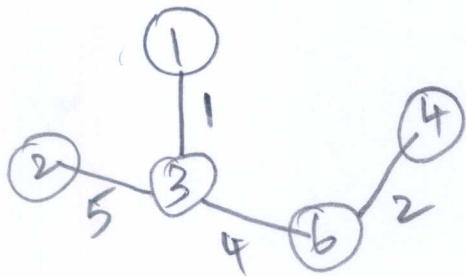
(4)



STEP 5:

$$mio = u = 2$$

$$T = \{1, 3, 6, 4, 2\}$$

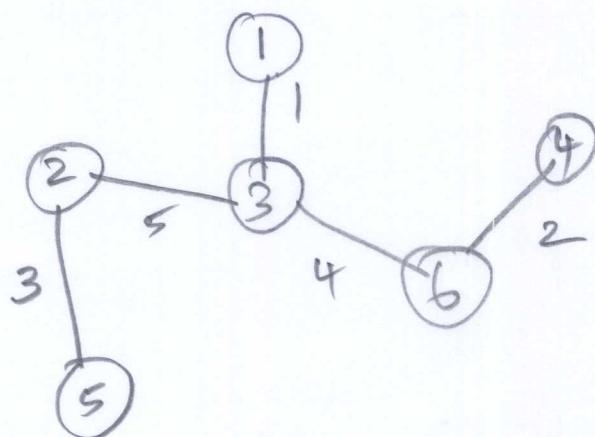


STEP 6:

$$T = \{1, 3, 6, 4, 2\} \cup \{5\}$$

v	old D[v]	new
+	0	*
2	5	5
3	1	*
4	2	2
5	3	{with edge 2}
6	4	4

$\therefore HST =$



(5)

