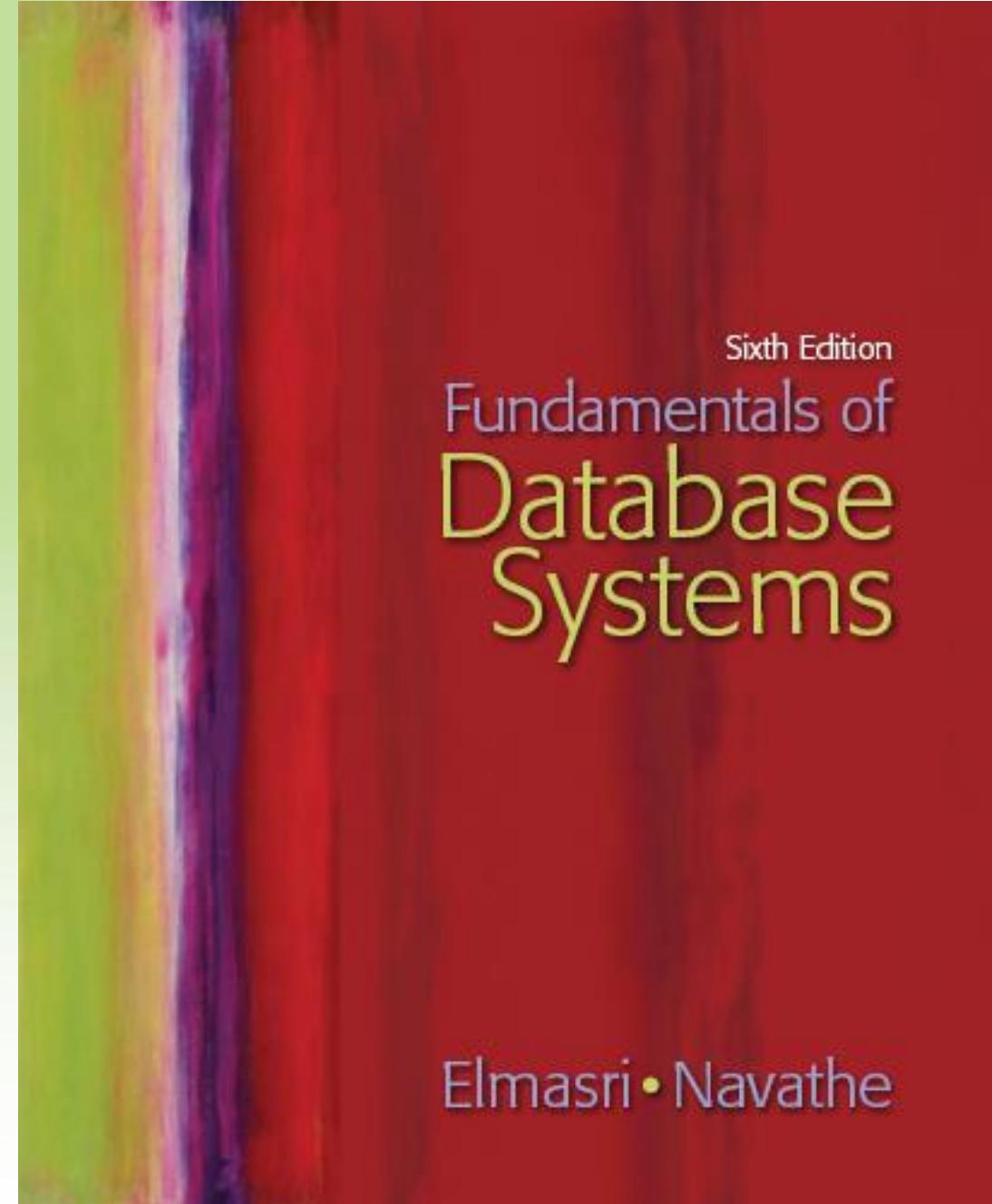


Chapter 3

The Relational Data Model and Relational Database Constraints



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Chapter 3 Outline

- The Relational Data Model and Relational Database Constraints
- Relational Model Constraints and Relational Database Schemas
- Update Operations, Transactions, and Dealing with Constraint Violations

The Relational Data Model and Relational Database Constraints

- Relational model
 - First commercial implementations available in early 1980s
 - Has been implemented in a large number of commercial systems
- Hierarchical and network models
 - Preceded the relational model

Relational Model Concepts

- Represents data as a collection of relations
- **Table** of values
 - Row
 - Represents a collection of related data values
 - Fact that typically corresponds to a real-world entity or relationship
 - *Tuple*
 - Table name and column names
 - Interpret the meaning of the values in each row *attribute*

Relational Model Concepts (cont'd.)

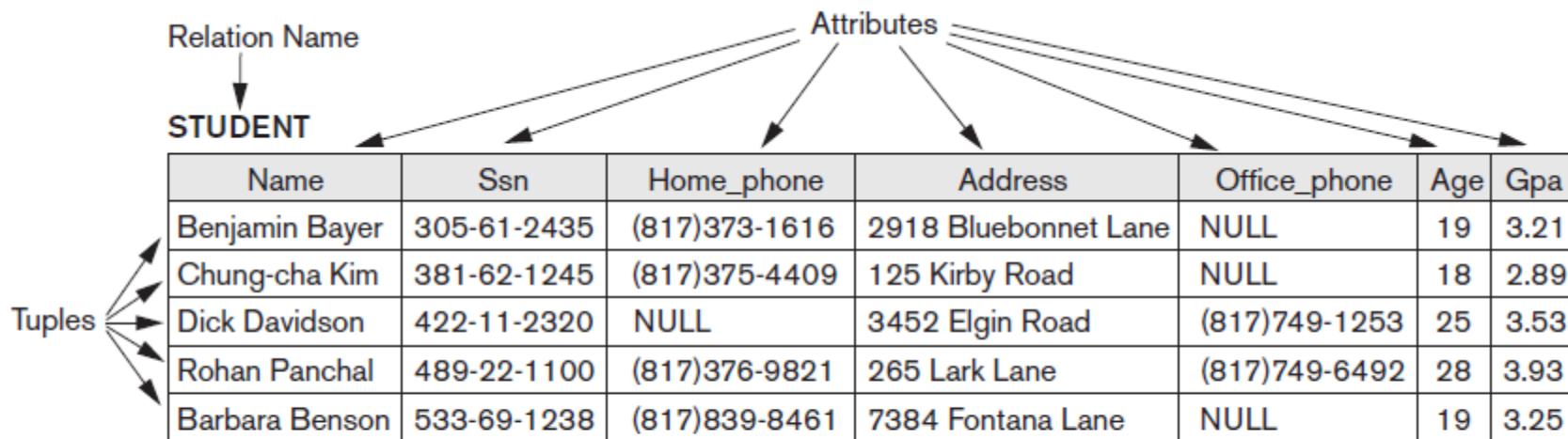


Figure 3.1

The attributes and tuples of a relation STUDENT.

Domains, Attributes, Tuples, and Relations

- **Domain D**
 - Set of atomic values
- **Atomic**
 - Each value indivisible
- Specifying a domain
 - **Data type** specified for each domain

Domains, Attributes, Tuples, and Relations (cont'd.)

- **Relation schema R**
 - Denoted by $R(A_1, A_2, \dots, A_n)$
 - Made up of a relation name R and a list of attributes, A_1, A_2, \dots, A_n
- **Attribute A_i**
 - Name of a role played by some domain D in the relation schema R
- **Degree (or arity) of a relation**
 - Number of attributes n of its relation schema

Domains, Attributes, Tuples, and Relations (cont'd.)

■ Relation (or relation state)

- Set of n -tuples $r = \{t_1, t_2, \dots, t_m\}$
- Each n -tuple t
 - Ordered list of n values $t = \langle v_1, v_2, \dots, v_n \rangle$
 - Each value v_i , $1 \leq i \leq n$, is an element of $\text{dom}(A_i)$ or is a special NULL value

Domains, Attributes, Tuples, and Relations (cont'd.)

- Relation (or relation state) $r(R)$
 - **Mathematical relation** of degree n on the domains $\text{dom}(A_1), \text{dom}(A_2), \dots, \text{dom}(A_n)$
 - **Subset** of the **Cartesian product** of the domains that define R:
 - $r(R) \subseteq (\text{dom}(A_1) \times \text{dom}(A_2) \times \dots \times \text{dom}(A_n))$

Domains, Attributes, Tuples, and Relations (cont'd.)

- **Cardinality**
 - Total number of values in domain
- **Current relation state**
 - Relation state at a given time
 - Reflects only the valid tuples that represent a particular state of the real world
- **Attribute names**
 - Indicate different **roles**, or interpretations, for the domain

Characteristics of Relations

- Ordering of tuples in a relation
 - Relation defined as a set of tuples
 - Elements have no order among them
- Ordering of values within a tuple and an alternative definition of a relation
 - Order of attributes and values is not that important
 - As long as correspondence between attributes and values maintained

Characteristics of Relations (cont'd.)

- Alternative definition of a relation
 - Tuple considered as a set of ($<\text{attribute}>$, $<\text{value}>$) pairs
 - Each pair gives the value of the mapping from an attribute A_i to a value v_i from $\text{dom}(A_i)$
- Use the first definition of relation
 - Attributes and the values within tuples are ordered
 - Simpler notation

Characteristics of Relations (cont'd.)

Figure 3.2

The relation STUDENT from Figure 3.1 with a different order of tuples.

STUDENT

Name	Ssn	Home_phone	Address	Office_phone	Age	Gpa
Dick Davidson	422-11-2320	NULL	3452 Elgin Road	(817)749-1253	25	3.53
Barbara Benson	533-69-1238	(817)839-8461	7384 Fontana Lane	NULL	19	3.25
Rohan Panchal	489-22-1100	(817)376-9821	265 Lark Lane	(817)749-6492	28	3.93
Chung-cha Kim	381-62-1245	(817)375-4409	125 Kirby Road	NULL	18	2.89
Benjamin Bayer	305-61-2435	(817)373-1616	2918 Bluebonnet Lane	NULL	19	3.21

Characteristics of Relations (cont'd.)

- Values and NULLs in tuples
 - Each value in a tuple is atomic
 - **Flat relational model**
 - Composite and multivalued attributes not allowed
 - **First normal form** assumption
 - Multivalued attributes
 - Must be represented by separate relations
 - Composite attributes
 - Represented only by simple component attributes in basic relational model

Characteristics of Relations (cont'd.)

- **NULL values**
 - Represent the values of attributes that may be unknown or may not apply to a tuple
 - Meanings for NULL values
 - *Value unknown*
 - *Value exists but is not available*
 - *Attribute does not apply to this tuple (also known as value undefined)*

Characteristics of Relations (cont'd.)

- Interpretation (meaning) of a relation
 - **Assertion**
 - Each tuple in the relation is a **fact** or a particular instance of the assertion
 - **Predicate**
 - Values in each tuple interpreted as values that satisfy predicate

Relational Model Notation

- Relation schema R of degree n
 - Denoted by $R(A_1, A_2, \dots, A_n)$
- Uppercase letters Q, R, S
 - Denote relation names
- Lowercase letters q, r, s
 - Denote relation states
- Letters t, u, v
 - Denote tuples

Relational Model Notation

- Name of a relation schema: STUDENT
 - Indicates the current set of tuples in that relation
- Notation: STUDENT(Name, Ssn, ...)
 - Refers only to relation schema
- Attribute A can be qualified with the relation name R to which it belongs
 - Using the dot notation $R.A$

Relational Model Notation

- *n-tuple t in a relation r(R)*
 - Denoted by $t = \langle v_1, v_2, \dots, v_n \rangle$
 - v_i is the value corresponding to attribute A_i
- Component values of tuples:
 - $t[A_i]$ and $t.A_i$ refer to the value v_i in t for attribute A_i
 - $t[A_u, A_w, \dots, A_z]$ and $t.(A_u, A_w, \dots, A_z)$ refer to the subtuple of values $\langle v_u, v_w, \dots, v_z \rangle$ from t corresponding to the attributes specified in the list

Relational Model Constraints

- Constraints
 - Restrictions on the actual values in a database state
 - Derived from the rules in the miniworld that the database represents
- **Inherent model-based constraints or implicit constraints**
 - Inherent in the data model

Relational Model Constraints (cont'd.)

- **Schema-based constraints or explicit constraints**
 - Can be directly expressed in schemas of the data model
- **Application-based or semantic constraints or business rules**
 - Cannot be directly expressed in schemas
 - Expressed and enforced by application program

Domain Constraints

- Typically include:
 - Numeric data types for integers and real numbers
 - Characters
 - Booleans
 - Fixed-length strings
 - Variable-length strings
 - Date, time, timestamp
 - Money
 - Other special data types

Key Constraints and Constraints on NULL Values

- No two tuples can have the same combination of values for all their attributes.
- **Superkey**
 - No two distinct tuples in any state r of R can have the same value for SK
- **Key**
 - Superkey of R
 - Removing any attribute A from K leaves a set of attributes K' that is not a superkey of R any more

Key Constraints and Constraints on NULL Values (cont'd.)

- Key satisfies two properties:
 - Two distinct tuples in any state of relation cannot have identical values for (all) attributes in key
 - Minimal superkey
 - Cannot remove any attributes and still have uniqueness constraint in above condition hold

Key Constraints and Constraints on NULL Values (cont'd.)

- **Candidate key**
 - Relation schema may have more than one key
- **Primary key of the relation**
 - Designated among candidate keys
 - Underline attribute
- Other candidate keys are designated as **unique keys**

Key Constraints and Constraints on NULL Values (cont'd.)

Figure 3.4

The CAR relation, with two candidate keys: License_number and Engine_serial_number.

CAR

License_number	Engine_serial_number	Make	Model	Year
Texas ABC-739	A69352	Ford	Mustang	02
Florida TVP-347	B43696	Oldsmobile	Cutlass	05
New York MPO-22	X83554	Oldsmobile	Delta	01
California 432-TFY	C43742	Mercedes	190-D	99
California RSK-629	Y82935	Toyota	Camry	04
Texas RSK-629	U028365	Jaguar	XJS	04

Relational Databases and Relational Database Schemas

- **Relational database schema S**
 - Set of relation schemas $S = \{R_1, R_2, \dots, R_m\}$
 - Set of integrity constraints IC
- **Relational database state**
 - Set of relation states $DB = \{r_1, r_2, \dots, r_m\}$
 - Each r_i is a state of R_i and such that the r_i relation states satisfy integrity constraints specified in IC

Relational Databases and Relational Database Schemas (cont'd.)

- **Invalid state**
 - Does not obey all the integrity constraints
- **Valid state**
 - Satisfies all the constraints in the defined set of integrity constraints IC

Integrity, Referential Integrity, and Foreign Keys

- **Entity integrity constraint**
 - No primary key value can be NULL
- **Referential integrity constraint**
 - Specified between two relations
 - Maintains consistency among tuples in two relations

Integrity, Referential Integrity, and Foreign Keys (cont'd.)

- **Foreign key rules:**

- The attributes in FK have the same domain(s) as the primary key attributes PK
- Value of FK in a tuple t_1 of the current state $r_1(R_1)$ either occurs as a value of PK for some tuple t_2 in the current state $r_2(R_2)$ or is NULL

Integrity, Referential Integrity, and Foreign Keys (cont'd.)

- Diagrammatically display referential integrity constraints
 - Directed arc from each foreign key to the relation it references
- All integrity constraints should be specified on relational database schema

Other Types of Constraints

- Semantic integrity constraints
 - May have to be specified and enforced on a relational database
 - Use **triggers** and **assertions**
 - More common to check for these types of constraints within the application programs

Other Types of Constraints (cont'd.)

- **Functional dependency constraint**
 - Establishes a functional relationship among two sets of attributes X and Y
 - Value of X determines a unique value of Y
- **State constraints**
 - Define the constraints that a valid state of the database must satisfy
- **Transition constraints**
 - Define to deal with state changes in the database

Update Operations, Transactions, and Dealing with Constraint Violations

- Operations of the relational model can be categorized into retrievals and updates
- Basic operations that change the states of relations in the database:
 - Insert
 - Delete
 - Update (or Modify)

Figure 3.6

One possible database state for the COMPANY relational database schema.

EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	NULL	1

DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

DEPT_LOCATIONS

Dnumber	Dlocation
1	Houston
4	Stafford
5	Bellaire
5	Sugarland
5	Houston

Figure 3.6

One possible database state for the COMPANY relational database schema.

WORKS_ON

<u>Essn</u>	<u>Pno</u>	Hours
123456789	1	32.5
123456789	2	7.5
666884444	3	40.0
453453453	1	20.0
453453453	2	20.0
333445555	2	10.0
333445555	3	10.0
333445555	10	10.0
333445555	20	10.0
999887777	30	30.0
999887777	10	10.0
987987987	10	35.0
987987987	30	5.0
987654321	30	20.0
987654321	20	15.0
888665555	20	NULL

PROJECT

<u>Pname</u>	<u>Pnumber</u>	<u>Plocation</u>	<u>Dnum</u>
ProductX	1	Bellaire	5
ProductY	2	Sugarland	5
ProductZ	3	Houston	5
Computerization	10	Stafford	4
Reorganization	20	Houston	1
Newbenefits	30	Stafford	4

DEPENDENT

<u>Essn</u>	<u>Dependent_name</u>	<u>Sex</u>	<u>Bdate</u>	<u>Relationship</u>
333445555	Alice	F	1986-04-05	Daughter
333445555	Theodore	M	1983-10-25	Son
333445555	Joy	F	1958-05-03	Spouse
987654321	Abner	M	1942-02-28	Spouse
123456789	Michael	M	1988-01-04	Son
123456789	Alice	F	1988-12-30	Daughter
123456789	Elizabeth	F	1967-05-05	Spouse

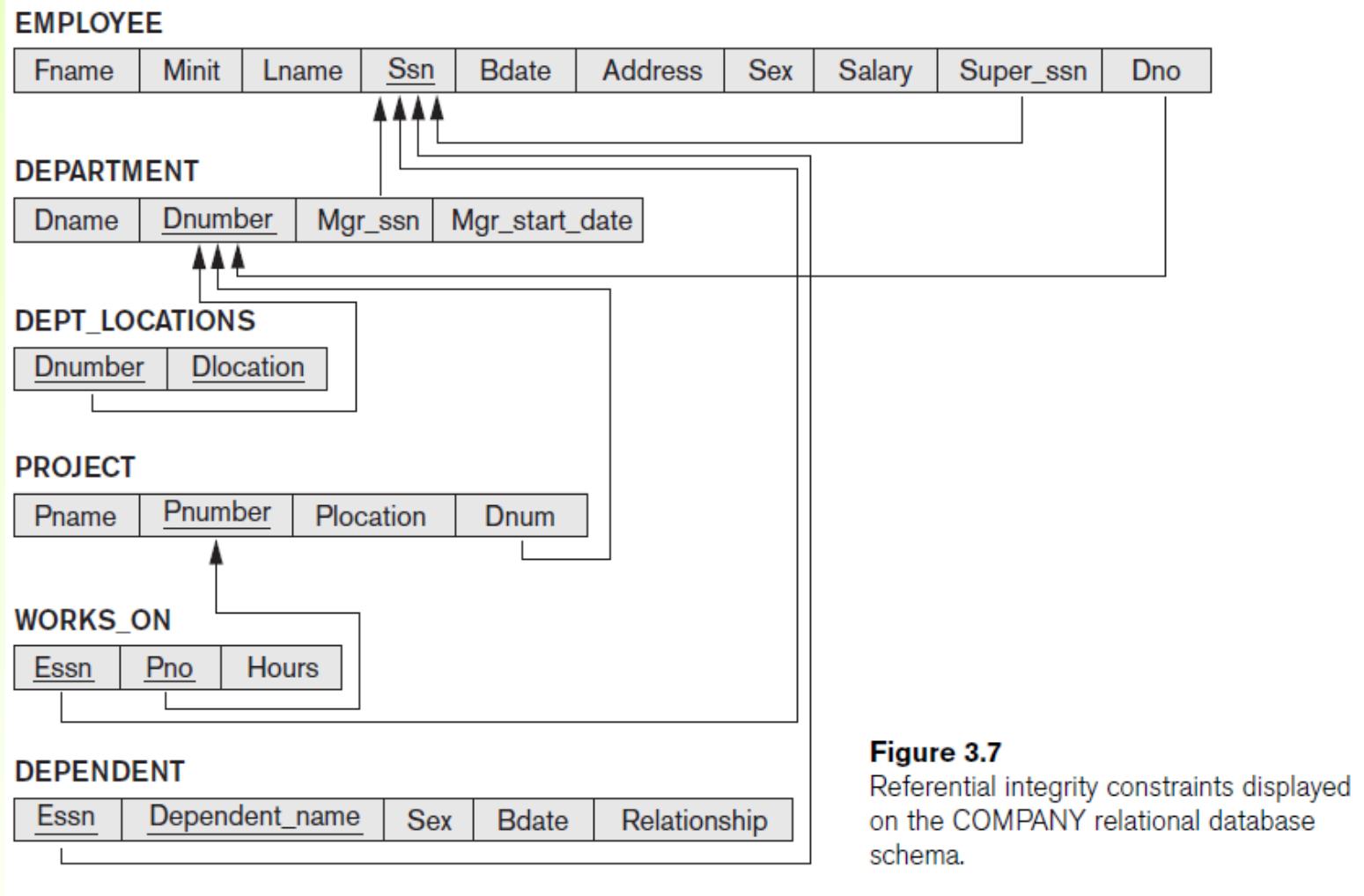


Figure 3.7

Referential integrity constraints displayed on the COMPANY relational database schema.

The Insert Operation

- Provides a list of attribute values for a new tuple t that is to be inserted into a relation R
- Can violate any of the four types of constraints
- If an insertion violates one or more constraints
 - Default option is to reject the insertion

The Delete Operation

- Can violate only referential integrity
 - If tuple being deleted is referenced by foreign keys from other tuples
 - **Restrict**
 - Reject the deletion
 - **Cascade**
 - Propagate the deletion by deleting tuples that reference the tuple that is being deleted
 - **Set null or set default**
 - Modify the referencing attribute values that cause the violation

The Update Operation

- Necessary to specify a condition on attributes of relation
 - Select the tuple (or tuples) to be modified
- If attribute not part of a primary key nor of a foreign key
 - Usually causes no problems
- Updating a primary/foreign key
 - Similar issues as with Insert/Delete

The Transaction Concept

- **Transaction**
 - Executing program
 - Includes some database operations
 - Must leave the database in a valid or consistent state
- **Online transaction processing (OLTP) systems**
 - Execute transactions at rates that reach several hundred per second

Summary

- Characteristics differentiate relations from ordinary tables or files
- Classify database constraints into:
 - Inherent model-based constraints, explicit schema-based constraints, and application-based constraints
- Modification operations on the relational model:
 - Insert, Delete, and Update