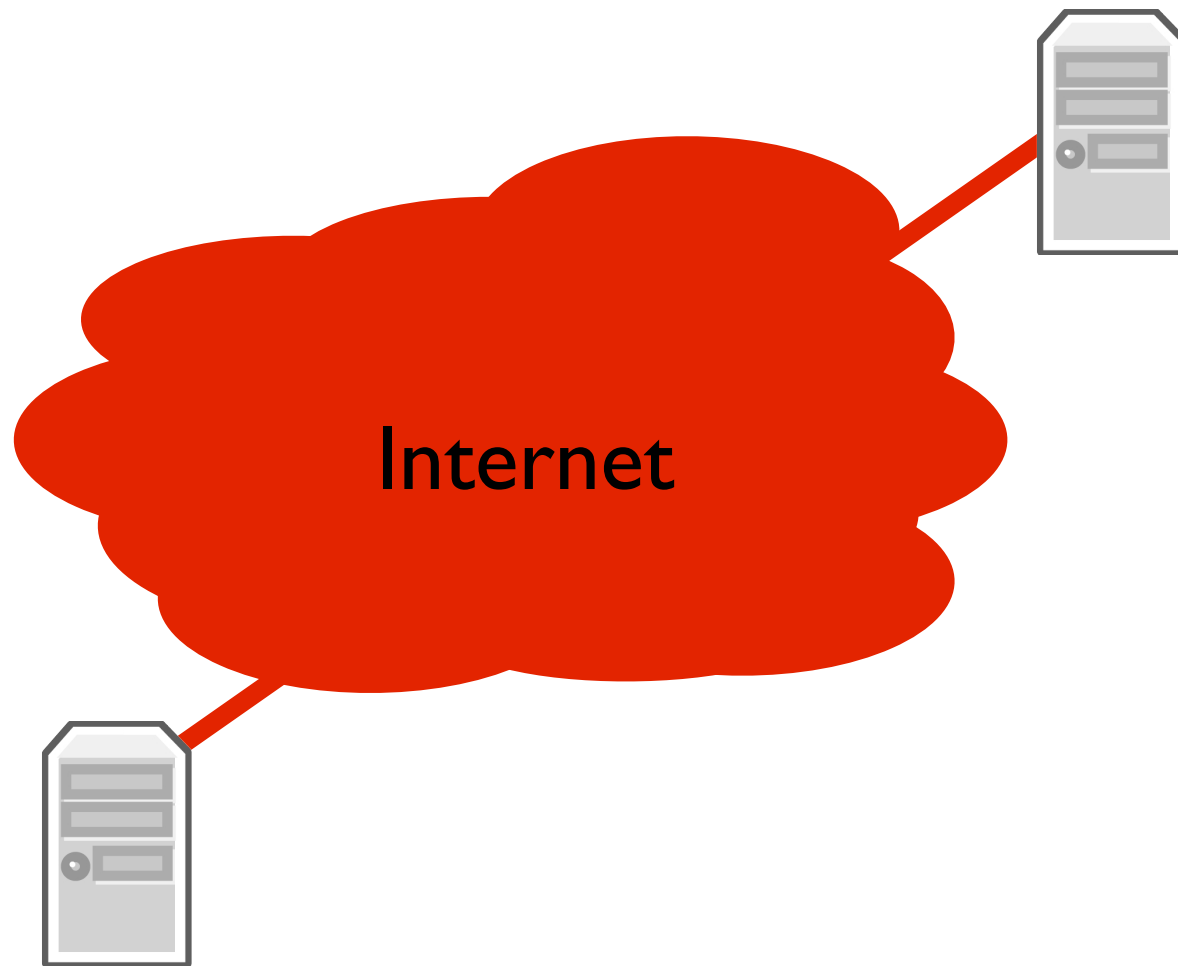


Security Principles

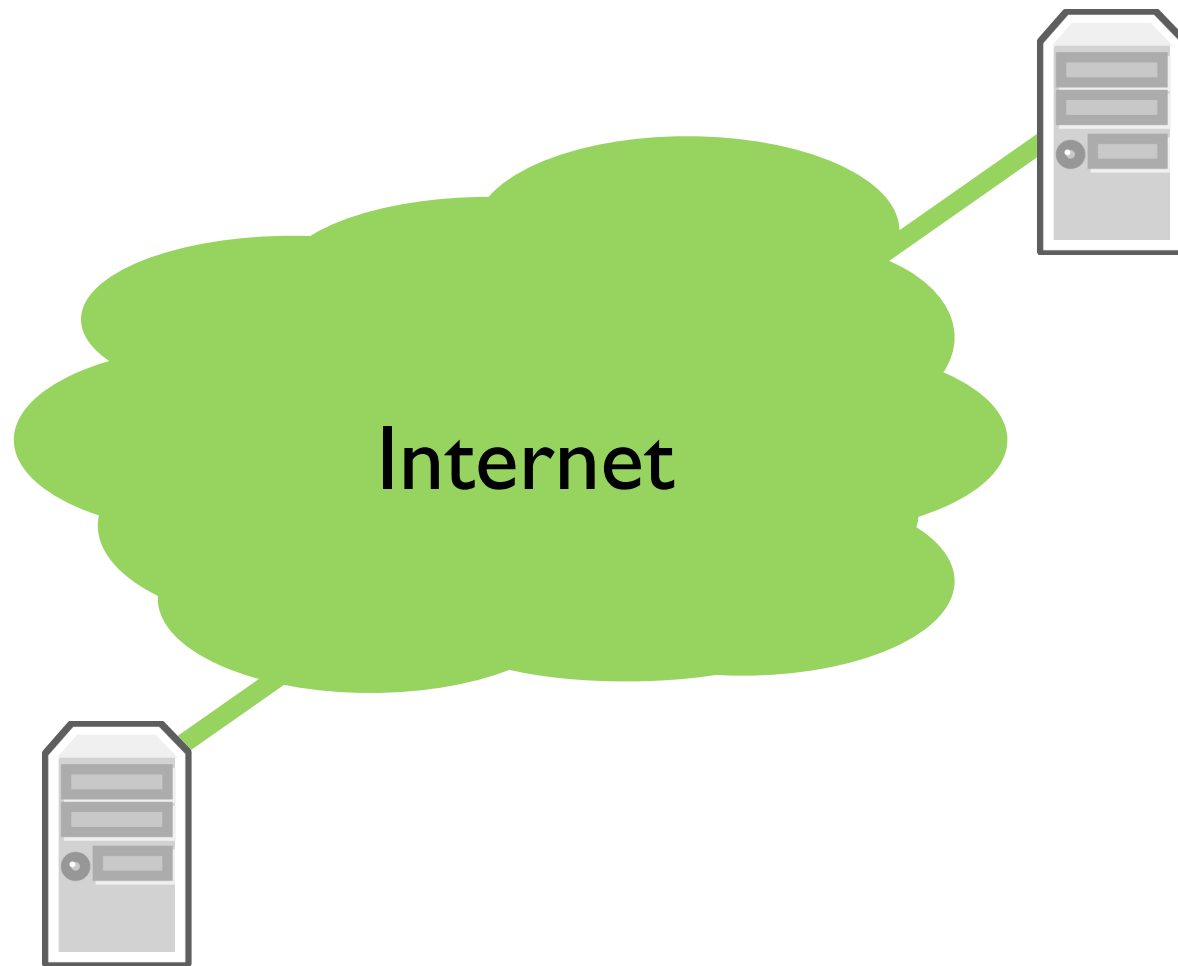
Confidentiality, Integrity, and Availability

Basic Problem



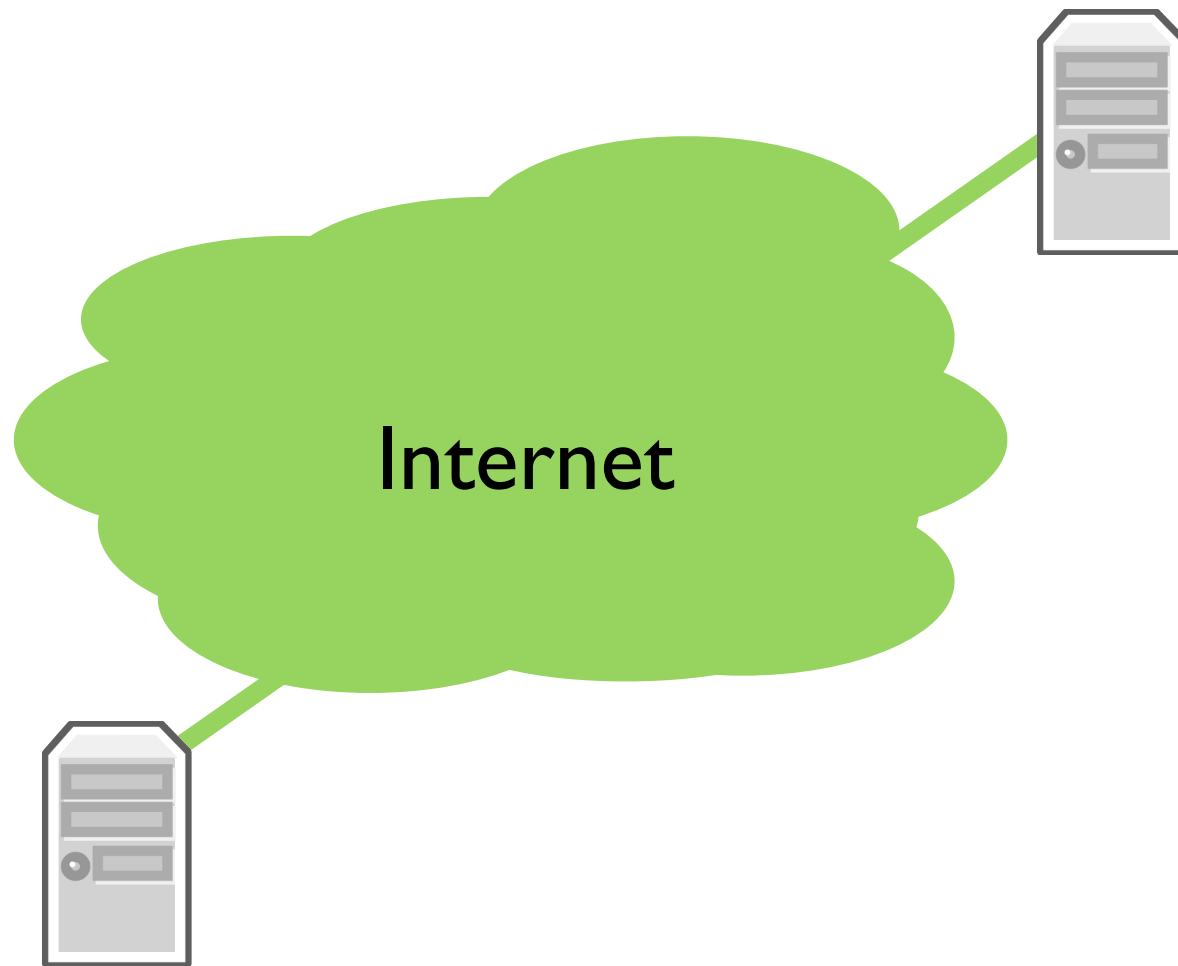
- To first approximation, attackers control the network
 - Can snoop, replay, suppress, send
- How do we defend against this?
 - Communicate securely despite insecure networks -- *cryptography*
 - Secure small parts of network despite insecurity of wider network
 - Design systems to scale well in response to attacks
- Two approaches: cryptography and scalable system design

Cryptography



- A set of mathematical principles and ideas for securing communication
- Be careful: easy to mess up!
 - ▶ Often misused!
 - ▶ Need to understand what it guarantees and what it doesn't
- How cryptography helps
 - ▶ Confidentiality: we can communicate privately (encryption)
 - ▶ Integrity: protect from tampering (hashes, signatures, MACs)
 - ▶ Authenticity: you are whom you say you are (certificates, MACs, signatures)

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Confidentiality

- I want to tell you something secretly, so no-one else knows what I said
 - ▶ “My credit card number is....”
- Perfect confidentiality: one-time pad
 - ▶ You and I share a perfectly random *key* of zeroes and ones, K , nobody else has it
 - ▶ I XOR my message M with K , producing C , send C to you ($C = M \oplus K$)
 - ▶ You XOR C with K , you have reconstructed M ($M = C \oplus K$)
- Advantages: informationally theoretic secure and fast
 - ▶ Given any C , any M is equally likely
- Disadvantage: need a K as long as all data I might ever send
- Ways to exchange a small K such that there are 2^k possible C s

Integrity

- I want to make sure you received my message unchanged/untampered
 - ▶ “I recalculated his grade and it is a C...”, program to run on your computer
- I want to make sure you sent the message
 - ▶ Nick says: “I said it was alright if he handed in the assignment late”
- Cryptographic hash: $H(M)$
 - ▶ Turn arbitrary length input into fixed length hash
 - ▶ Collision-resistance: given $x \neq y$, intractable to find $H(x) = H(y)$
- Message authentication code: $MAC(M,K)$
 - ▶ Use key K to generate $MAC(M,K)$, use K to check $MAC(M,K)$
 - ▶ Intractable to generate $MAC(M,K)$ unless you have K

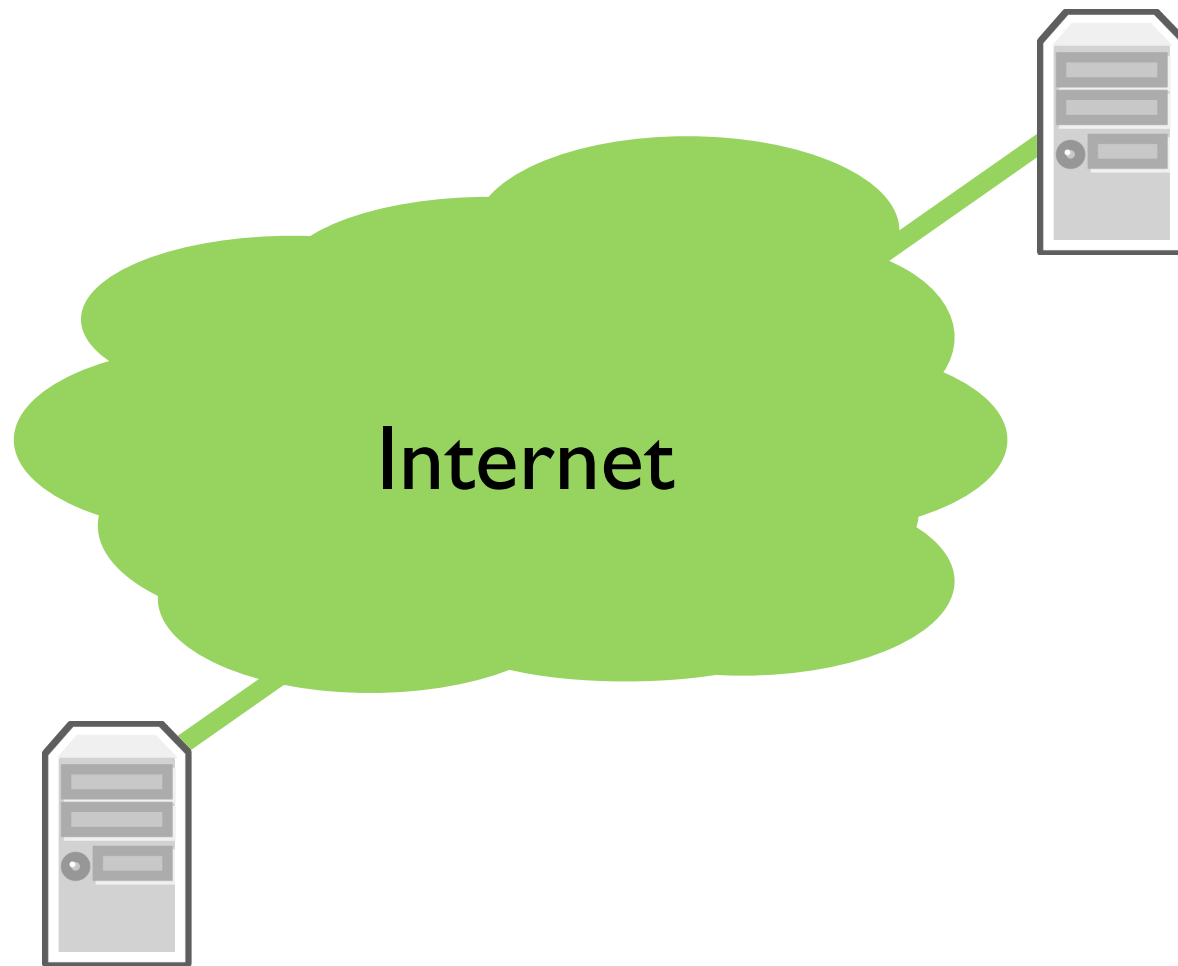
Authenticity

- I want to be sure you are whom you say you are
 - ▶ “This is the provost... I need your credit card number.”
 - ▶ “You should send all of my traffic through this third party”
- We’ve exchanged a secret K beforehand: $\text{MAC}(\text{“This is the...”}, K)$
- If we haven’t: chain of of trust
 - ▶ We can trust Verisign by design (root of trust)
 - ▶ Verisign says “here’s a secret for Stanford”
 - ▶ Stanford says “here’s a secret for the provost”

High Availability Design

- Denial of service (DoS), Distributed Denial of Service (DDoS)
- Many kinds of attacks, many defenses
 - ▶ Replication (scale-out)
 - ▶ Keeping costs symmetric
 - ▶ Upstream filtering (stop letting those pings through, router!)
- Continual arms race: not going to talk much more about it

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