

# CS144

## An Introduction to Computer Networks

### Routing

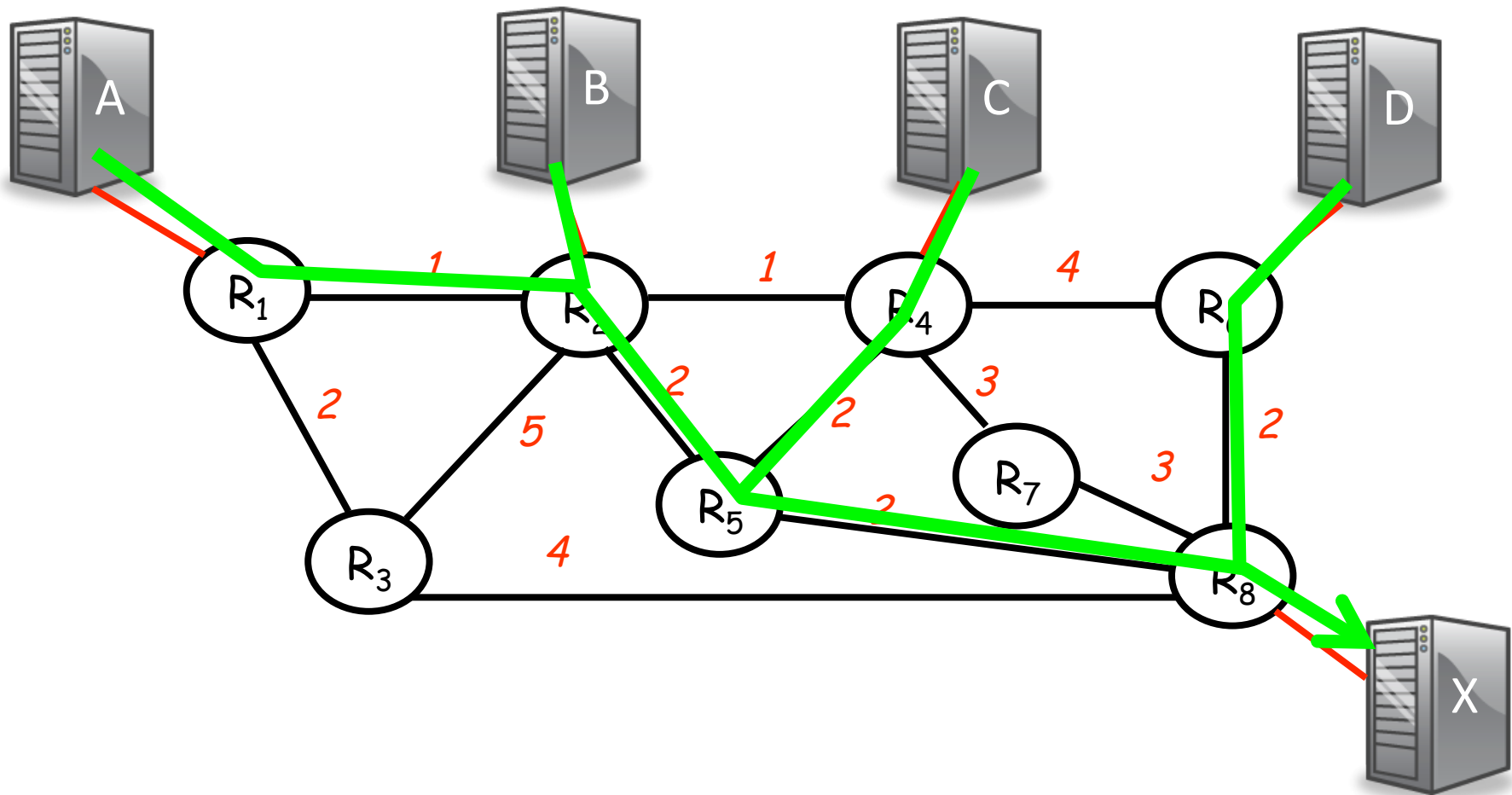
*Distance Vector Protocol:  
Bellman Ford algorithm*



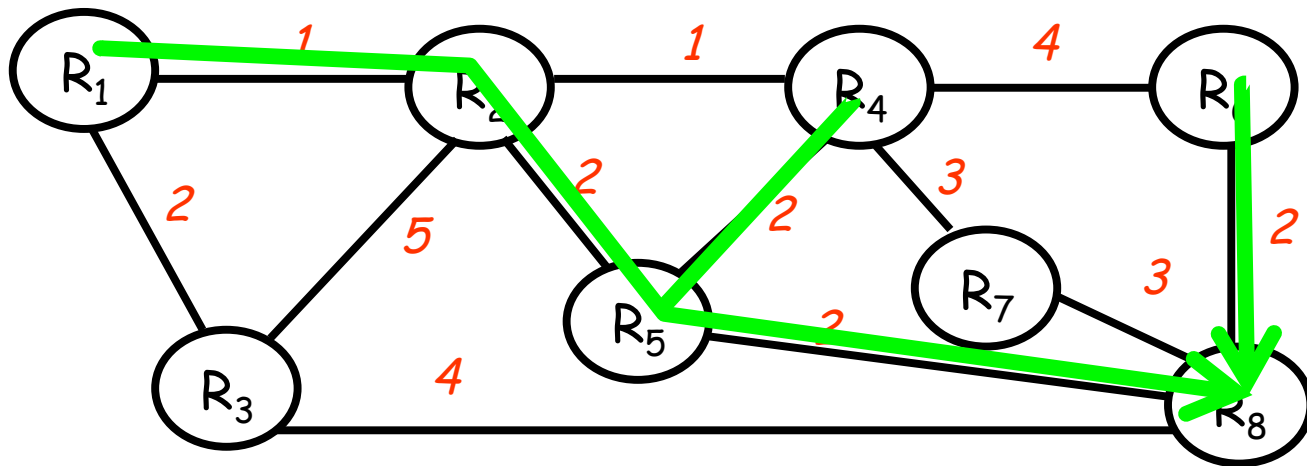
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Problem: How can routers work together to find minimum cost spanning tree?



Equivalent to finding minimum cost spanning tree among routers only



# The Distributed Bellman-Ford Algorithm

Example: Find min-cost spanning tree to  $R_8$

Assume routers know cost of link to each neighbor.

Router  $R_i$  maintains value cost  $C_i$  to reach  $R_8$ .

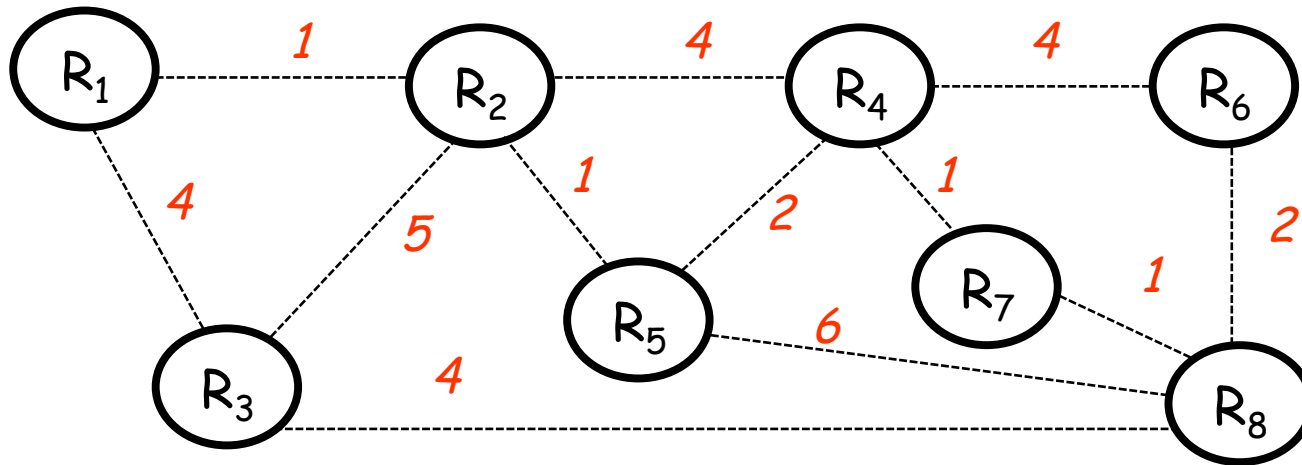
Vector  $\underline{C}=(C_1, C_2, \dots, C_7)$  is the *distance vector* to  $R_8$ .

Initially, set  $\underline{C} = (\infty, \infty, \dots, \infty)$ .

1. After  $T$  seconds,  $R_i$  sends  $C_i$  to its neighbors.
2. If  $R_i$  learns lower cost path, update  $C_i$ .
3. Repeat.

Natural extension to calculate tree for  $R_1 - R_7$ .

# An example



R <sub>1</sub>	∞	R <sub>1</sub>	∞	R <sub>1</sub>	8, R <sub>3</sub>	R <sub>1</sub>	8, R <sub>3</sub>	R <sub>1</sub>	7, R <sub>2</sub>	R <sub>1</sub>	6, R <sub>4</sub>
R <sub>2</sub>	∞	R <sub>2</sub>	∞	R <sub>2</sub>	7, R <sub>5</sub>	R <sub>2</sub>	6, R <sub>4</sub>	R <sub>2</sub>	5, R <sub>5</sub>	R <sub>2</sub>	5, R <sub>5</sub>
R <sub>3</sub>	∞	R <sub>3</sub>	4	R <sub>3</sub>	4	R <sub>3</sub>	4	R <sub>3</sub>	4	R <sub>3</sub>	4
R <sub>4</sub>	∞	R <sub>4</sub>	∞	R <sub>4</sub>	2, R <sub>7</sub>	R <sub>4</sub>	2, R <sub>7</sub>	R <sub>4</sub>	2, R <sub>7</sub>	R <sub>4</sub>	2, R <sub>7</sub>
R <sub>5</sub>	∞	R <sub>5</sub>	6	R <sub>5</sub>	6	R <sub>5</sub>	4, R <sub>4</sub>	R <sub>5</sub>	4, R <sub>4</sub>	R <sub>5</sub>	4, R <sub>4</sub>
R <sub>6</sub>	∞	R <sub>6</sub>	2	R <sub>6</sub>	2	R <sub>6</sub>	2	R <sub>6</sub>	2	R <sub>6</sub>	2
R <sub>7</sub>	∞	R <sub>7</sub>	1	R <sub>7</sub>	1	R <sub>7</sub>	1	R <sub>7</sub>	1	R <sub>7</sub>	1

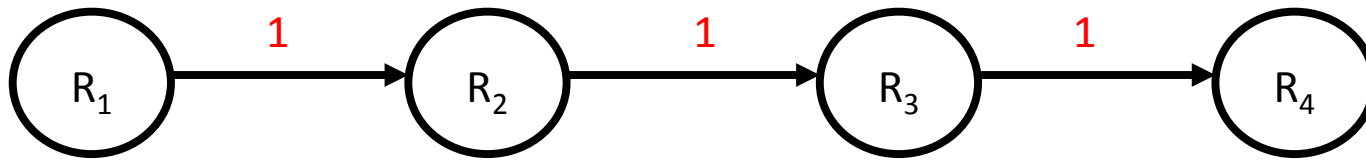
# Distributed Bellman-Ford Algorithm

## Questions:

1. What is the maximum run time of the algorithm?
2. Will the algorithm always converge?
3. What happens when link costs change, or when routers/links fail?

# A Problem with Bellman-Ford

“Bad news travels slowly”



Consider the calculation of distances to  $R_4$ :

Time	$R_1$	$R_2$	$R_3$
0	3, $R_2$	2, $R_3$	1, $R_4$
1	3, $R_2$	2, $R_3$	3, $R_2$
2	3, $R_2$	4, $R_3$	3, $R_2$
3	5, $R_2$	4, $R_3$	5, $R_2$
...	“Counting to infinity”		...

← Link  $R_3 \rightarrow R_4$  fails

# Counting to Infinity Problem

## *Solutions*

1. Set infinity = “some small integer” (e.g. 16).  
Stop when count = 16.
2. Split Horizon: Because  $R_2$  received lowest cost path from  $R_3$ , it does not advertise cost to  $R_3$ .
3. Split-horizon with poison reverse:  $R_2$  advertises infinity to  $R_3$ .
4. There are many problems with (and fixes for) the Bellman-Ford algorithm.



# Bellman Ford in practice

Bellman-Ford algorithm is an example of a Distance Vector algorithm.

It was used in the first Internet routing protocol, called Routing Information Protocol (RIP).

It requires very little computation on the routers, is distributed, and will eventually converged.

Over time it was replaced by algorithms that calculate the entire spanning tree at each router.