

Intel X86 Processors

only 14 assembly instructions make up to 90 percent of the code!

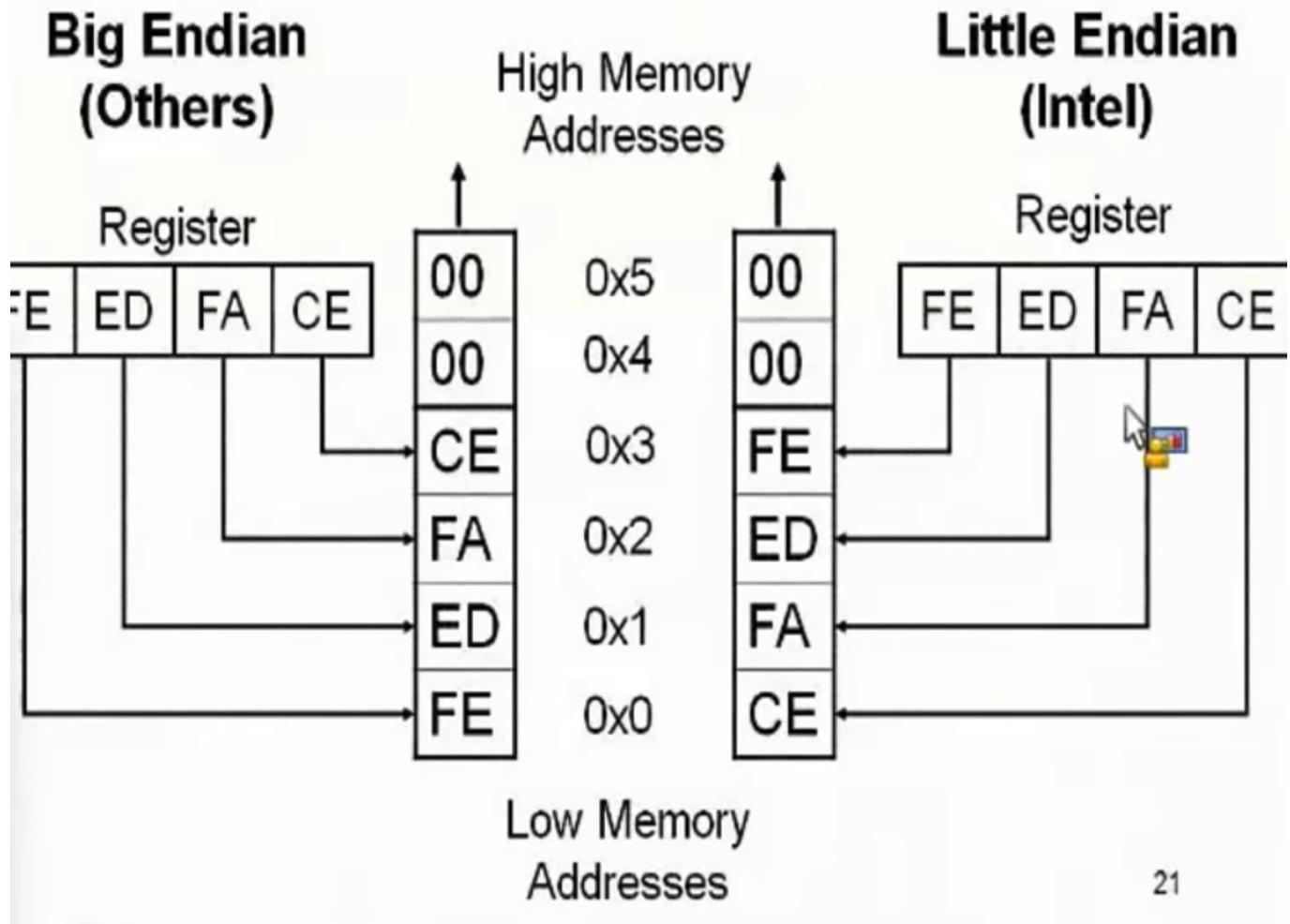
in intel terminology : a byte is a char(8) , a word is a short(16), int/long is double word(32), double long(64)

If you don't know this, you must memorize tonight

Decimal (base 10)	Binary (base 2)	Hex (base 16)
00	0000b	0x00
01	0001b	0x01
02	0010b	0x02
03	0011b	0x03
04	0100b	0x04
05	0101b	0x05
06	0110b	0x06
07	0111b	0x07
08	1000b	0x08
09	1001b	0x09
10	1010b	0x0A
11	1011b	0x0B
12	1100b	0x0C
13	1101b	0x0D
14	1110b	0x0E
15	1111b	0x0F

Intel is little endian. 0x12345678, least significant bit is registered first, then more significant and more significant.

Endianess pictures



21

Endians are only meaningful in byte level not in bit level!

Register Conventions!

EAX ==> stores function return value **EBX** ==> Base pointer to the data section **ECX**==>counter for string and loop ops **EDX** ==> I/O pointer **ESI**==>source pointer for string operations **EDI**==>Destination pointer for string operations **ESP**==>stack pointer **EBP**==>Stack frame base pointer **EIP**==>pointer to next instruction to execute(instruction pointer)

There are also other type of registers.

EAX, EDX , ECX == > Caller-save registers

I am a function and I will call another function. And when I call I am sure I am going to destroy the current registers(eax edx ecx) so we save the copy of the registers before calling the function so we do not lose the register

EBP, EBX, ESI, EDI ==> Callee-save registers

If I call a function, it wont change and modify the info in the register of ebp, ebx etc.

this E at the head of registers mean Extended. Since these are written originally for 16 bit , AX became Extended AX with 32 bits. etc.

EFLAGS

EFLAGS register holds many single bit flags. Will only ask you to remember the following for now :

Zero Flag(ZF)- Set if the result of some instruction is zero, cleared otherwise Sign Flah(SF) - Set equal to the most significant bit of the result.

INSTRUCTIONS

NOP No Operation. No registers no values nothing. Just there to pad bytes or delay time. Bad guys use it to make simple exploits more reliable. But that's another class.

NOP actually takes a register and re-registers it on itself 😊

PUSH ==> push word, double word or quadword onto the STACK. can either be immediate(a numeric constant) or the value in a register.

For our purposes, it will always be a DWORD(for this course) the push instruction automatically decrements the stack pointer, esp, by 4. Why decrements? Because its stack, its FILO.

POP ==> Pop a value from the Stack

take a Dword off the stack, put it in a register and increment esp by 4 (reverse of Push)

CALL ==> Call's job is to transfer control to a different function in a way that control can later be resumed where it left off.

First it pushes the address of the next instruction onto the stack. for use by RET(return from procedure) for when the procedure is done.

Because idea behind a function is , do this and I will keep executing from the next line. So call also pushes the address of the next instruction set, so that once the called function is completed, the program keeps continuing from where it left off. Its a reminder to where to return to after executing the function.

Then it changes eip to the address given in the instruction.

Destination address can be specified in a multiple ways.

- Absolute address ==> address 0x0030434
- Relative address(relative to the end of the instruction) ==> some address that hex 50 bytes ta ta ta.

RET==> Return from Procedure.

Two forms :

1. pop the top of the stack into the eip

in this form, the instruction is just written as ret.

typically used by cdecl functions

2. pop the top of the stack into eip and add a constant number of bytes to esp.

in this form, the instruction is written as ret 0x8 or 0x20 etc. typically used by stdcall functions

MOV ==> Move.

Can move:

register to register

memory to register, register to memory, immediate to register immediate to memory.

immediate is hardcoded value.

NEVER MEMORY TO MEMORY

Memory addresses are given in r/m32 form.

ADD ==> a = a + b for example add eax,ebx means eax = eax + ebx. so it sums them and writes on what is on the left side

SUB ==> subtraction. sub eax, [ebx*2]

LEA ==> Load Effective Address. ==> frequently used with pointer arithmetic sometimes for just arithmetic in general. uses the r/m32 form but **is the exception to the rule**

Lea makes the calculation and appends to the destination without disturbing the memory. Does the calc on the spot.

that the [] syntax means dereference("value at")

example: ebx = 0X2, edx = 0x1000 --- lea eax,[edx+ebx*2] **eax = 0X1004, not the value at 0x1004!!**

JMP ==> Jump. Changes the eip unconditionally to the given address. Main forms of the address.

short relative : 1 byte displacement from the end of the instruction.

01151015 jmp 1155000h doesn't have the number **1155000h** anywhere in it, it is really **jmp 0x0E** bytes forward.

near relative: 4 bytes displacement from the current eip

absolute: hardcoded address in instruction. like **01151013 jmp 01151023**

absolute indirect : indirect address calculated with r/m32

jmp -2 is infinite loop for short relative jmp 😊

Conditional Jumps: **je, jge, jne, jle**

JCC ==> Jump if Condition is met.

JNE == JNZ jump if not equal means jump if not zero. What does it mean ? jne a b for example. the operation is a-b. so flag is either 0 or <0 or 0>

so 0 means equal, if not equal, its not zero. So it makes zero flag check

THE STACK

stack is the conceptual area of main memory which is designated by OS when program is started. Stack is LIFO/FILO data structure where data is pushed on to top of the stack and popped off the top

Stack grows over lower memory addresses. Adding something to the stack means the top of the stack is now at a lower memory address.

ESP points to the top the stack. The lowest address which is being used(since it is LIFO)

Stack keeps track of which functions were called before the current one. It holds LOCAL VARIABLES and is used to pass arguments to the next function to be called

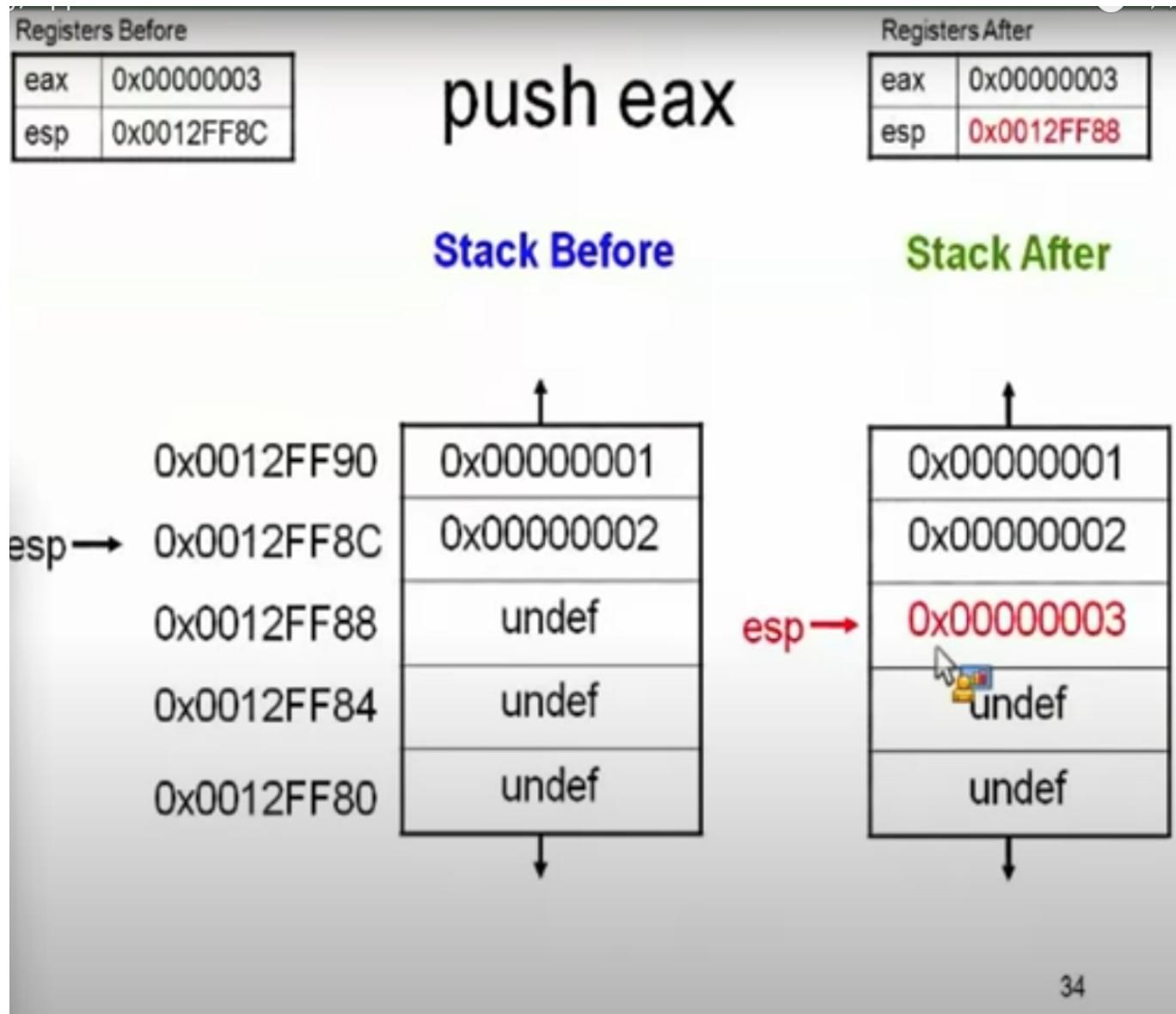
Register in time 1 :

eax : 0x00000003 esp : 0x0012FF8C

So when we call lets say , **push eax**, we expect the eax to be written 4 bytes after(decrement) the last esp register in the stack.

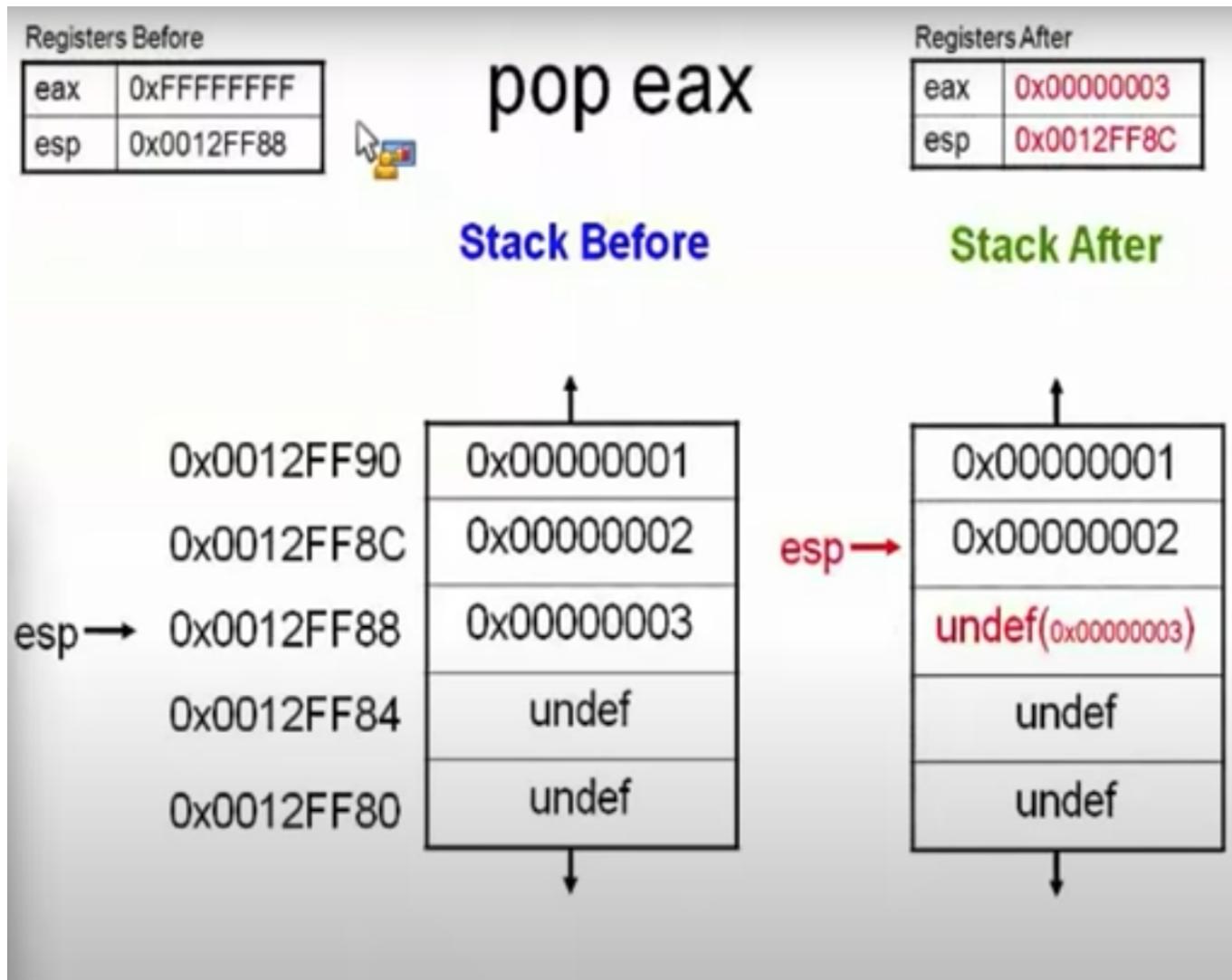
Register after push eax:

eax: 0x00000003 esp:0x0012FF88 ==> here, we jumped 4 bytes from FF8C to FF88. (where FF8C>FF88)



to reverse it , we use **POP**

in order to reverse the action above, we **pop eax** and voila, esp jumped by 4 bytes up (incremented) and top shelf of the stack is emptied.



Everything numerically under `esp` is considered to be undefined. Although there is a data over there, it is not of our business.

once we call `pop eax`, the the value at the top of the stack (`eax`) is taken to the register and that memory spot is not undefined.

Calling Convention

calling convention is about how to pass parameters and how to get parameters back.

how code calls a subroutine is compiler-dependent and configurable. But there are few conventions. We will deal with cdecl and stdcall conventions.

cdecl

C Declaration. The most common calling convention. Function parameters pushed (using push) onto stack RIGHT TO LEFT. Saves the old stace frame pointer and sets up a new stack frame.

`eax` or `edx:eax` returns the result for primitive data types.

Most significant 32 bits go into `edx`, least significant bits go into `eax`

Also values are always big endian in registers, in memory they are little endian.

so for example : `printf(%d\n, myVariable)` ==> from right to left: push my variable to the stack then push the pointer and then call the function.

function here, saves the old frame pointer, saves the address to the stack.

Caller is responsible for cleaning up the stack. ==>

so we have 2 parameters and 1 function call. push push and call. so whatever register is calling the function is also responsible to clean up the stack

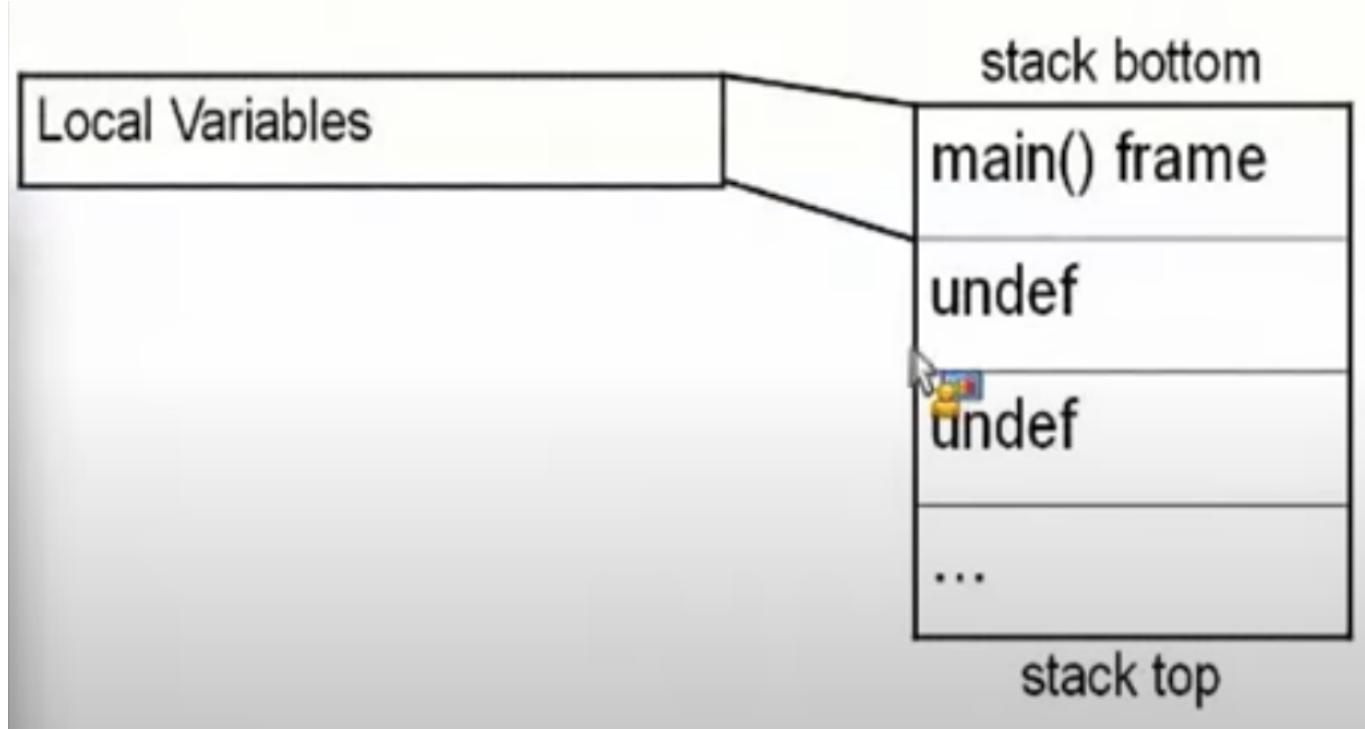
stdcall

Microsoft C++ code e.g. Win32API parameters pushed onto stack right to left saves the old stack frame pointer and sets up a new stack frame pointer

HERE, CALLEE is responsible for cleaning up any stack parameters it takes not the CALLER!

General Stack Frame Operation

We are going to assume that the `main()` is the very first function being executed in a program. This is what its stack looks like to start with.



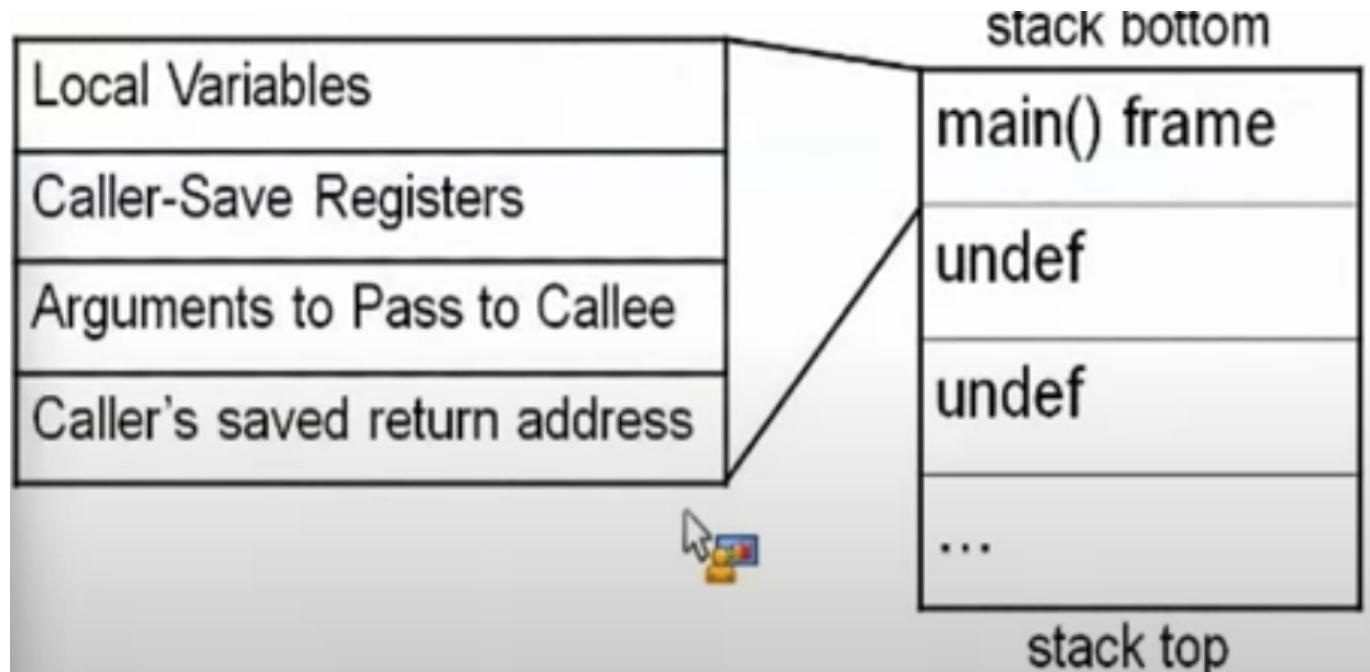
Low addresses at the bottom, high addresses at the top.

The more I add to the stack, the more I expect it to grow to the bottom.

So first thing `main` does is to preserve space on the stack for the local variables. When `main()` decides to call a subroutine, `main()` becomes the caller. We will assume `main()` has some registers it would like to remain the same, so it will save them. We will also assume that the callee function stakes some input args.

so , local variables, caller-save registers, and Args to pass to callee are in `main()` stack frame at the bottom of the stack.

so right to left, local variables, caller-save registers, args to pass to callee and then function call. git when we execute the call instruction, the return address(RET) gets save onto the stack and because the next instruction after the call will be the beginning of the called function, we consider the frame to have changed to the callee.



next, when subroutine starts, the frame pointer(ebp) still points to the main()'s frame. so the first thing it does is to save the old frame pointer on the stack and set the new value to point to its own frame.

so at the main() function, ebp is pointing at the top of the frame. when a subroutine is called, it will save this address to stack , take the ebp and continue its function, once the function is finished, it will restore the value of ebp so that the next subroutines can take up from original ebp register.

STACK FRAMES ARE A LINKED LIST!

The ebp in the current frame points at the saved ebp of the previous frame

Example1.c

The stack frames in this example will be very simple.
Only saved frame pointer (ebp) and saved return addresses (eip).

```
//Example1 - using the stack
//to call subroutines
//New instructions:
//push, pop, call, ret, mov
int sub(){
    return 0xbeef;
}
int main(){
    sub();
    return 0xf00d;
}

sub:
00401000 push    ebp
00401001 mov     ebp,esp
00401003 mov     eax,0BEEFh
00401008 pop    ebp
00401009 ret

main:
00401010 push    ebp
00401011 mov     ebp,esp
00401013 call    sub (401000h)
00401018 mov     eax,0F00Dh
0040101D pop    ebp
0040101E ret
```

```
int sub(){
    return 0xbeef;
}

int main(){
    sub();
    return 0xf00d;
}
```

NOTE : ON INTEL REGISTERS AND CALLS, DESTINATION IS ON THE LEFT AND SOURCE IS ON THE RIGHT

SO

MOV EBP,ESP MEANS

Move esp TO EBP!!!

IN AT&T syntax, this is the opposite tho.

sub:

```
00401000 push    ebp
00401001 mov     ebp,esp
```

```

00401003 mov eax,0BEEFh
00401008 pop ebp
00401009 ret

main:
00401010 push ebp
00401011 mov ebp,esp
00401013 call sub(401000h)
00401018 mov eax,0F00Dh
0040101D pop ebp
0040101E ret

```

note : push ebp mov ebp,esp codes are generated by compiler automatically

```

main:
00401010 push ebp

```

when code started to executing, there was some value in ebp. so initially, main pushes ebp(base pointer). so when we push ebp, since esp always points to the top of the stack, we also see that esp is also changed.

stack is decremented by 4 bytes with this push since it pushes stack to down(LIFO)

mov ebp,esp means take whatever value in the esp and put it onto ebp. because intel x86 syntax is **instruction destination,source**

so now ebp and esp are the same. so whatever is at the address of ebp and low is the new stack frame.

00401013 call sub(401000h) in the main function is very self-explanatory. call the function called sub at the address of 401000h which is the beginning of the address of sub.

so what is at the stack now ?

```

00401013 call sub(401000h)
00401018 mov eax,0F00Dh

```

subroutine(function sub) needs to register the address of the next register so that after the execution of the subroutine, main function picks up from where it left off. hence, the address of mov eax,0F00Dh is registered to the stack.

now the stack is hosting 00401018 right after ebp.

next step is to calling sub routine. when subroutine is called, it will save the ebp to the register (push ebp) this ebp is pointing at the main function normally. then it will move esp to ebp like main function(mov ebp,esp). now ebp and esp are pointing at the same place. why **mov ebp,esp** is important? **by setting esp to ebp, sub is creating its own stack!!** next step, sub can start functioning, which is writing 0XBEEF to eax.

why? because eax holds the return values and the sub() function was only returning 0xbeef;. thus it makes sense!

thus , in the stack, eax is changed to 0x0000BEEF address since this is our return value.

now sub() function is executed and we do not need the base pointer anymore. next step is pop ebp since this is the end of the stack.

end we will exit the sub() function with **ret**.

what **ret** does is that whatever is on the top iof the stack, go ahead and take it from there and I will put there the insturction pointer and so this will be our next destination.

remember the code :

```
sub:  
00401000 push ebp  
00401001 mov ebp,esp  
00401003 mov eax,0BEEFh  
00401008 pop ebp  
00401009 ret  
  
main:  
00401010 push ebp  
00401011 mov ebp,esp  
00401013 call sub(401000h)  
00401018 mov eax,0F00Dh  
0040101D pop ebp  
0040101E ret
```

so when **ret** (0x000401009) is called, the top of the stack will have the value of will be replaced to what?

it will be replaced with **0x000401018** since this is the first address after the sub() function call. thus main() function will keep goint from here. so next step is **00401018 mov eax,0F00Dh**

whici moves the value 0x00F00D to the eax because is the return value. if we did not put this register, the eax would be stil under the effect of sub() function so would keep pointing at 0x00BEEF. but main() wants a return of 0xFOOD so we immediately overwrite it.

next, we tear down the stack with **pop ebp** and returning to the function.

with pop ebp and ret, we cleared the last 2 remaining registers of stack. now stack is all undefined.

Some notes:

since sub() is a deadcode , meaning its return value 0xFOOD, is not used for anything, compiler with optimization option on would delete sub() function.

because there are not input parameters to sub(), there is no diff whether we compile cdecl or stdcall calling convention.

in the debugger when we jump to the assembly debugger here is the C code and the corresponding assembly code :

```

int main(){
011E1010 push      ebp
011E1011 mov       ebp,esp
    sub();
011E1013 call      sub (11E1000h)
    return 0xf00d;
011E1018 mov       eax,0F00Dh
}
011E101D pop      ebp
011E101E ret

```

Lets do another example.

```

#include <stdlib.h>
int sub(int x, int y){
    return 2*x+y;
}

int main(int argc, char ** argv){
    int a;
    a = atoi(argv[1]);
    return sub(a,a);
}

.text:00000000 _sub:    push  ebp
.text:00000001          mov   ebp,esp
.text:00000003          mov   eax,[ebp+8]
.text:00000006          mov   ecx,[ebp+0Ch]
.text:00000009          lea   eax,[ecx+eax*2]
.text:0000000C          pop   ebp
.text:0000000D          retn 
.text:00000010 _main:   push  ebp
.text:00000011          mov   ebp,esp
.text:00000013          push  ecx
.text:00000014          mov   eax,[ebp+0Ch]
.text:00000017          mov   ecx,[eax+4]
.text:0000001A          push  ecx
.text:0000001B          call  dword ptr ds:_imp_atoi
.text:00000021          add   esp,4
.text:00000024          mov   [ebp-4],eax
.text:00000027          mov   edx,[ebp-4]
.text:0000002A          push  edx
.text:0000002B          mov   eax,[ebp+8]
.text:0000002E          push  eax
.text:0000002F          call  _sub
.text:00000034          add   esp,8
.text:00000037          mov   esp,ebp
.text:00000039          pop   ebp
.text:0000003A          retn 

```

88

r/m32 Addressing Forms

Anywhere you see an r/m32 it means it could be taking a value either from a register or a memory address. (r/m == register/memory) in **Intel Syntax**, most of the time square brackets([]) means to treat the value within as a memory address and fetch the value at the address (like dereferencing pointer)

- mov eax,ebx
- mov eax, [ebx]
- mov eax, [ebx+ecx*X] => X can be 1,2,4 or 8.

so for example here, it means, take the ecx(lets say 0x00344F) and multiply it with X, go to the address in this multiplication, add it to ebx, and write the cumulation to eax.

the most complicated form is : [base + index * scale + disp]

base +(index * scale) + displacement.

mov eax ,[ebx+ecx*X+Y] where Y is one byte(0-255) or 4 bytes(0-2^32-1)

Scenario : data type we have is an array. Its consequential data type. So ebx points to the base of the Array, beginning point of the array. Ecx is the index number(so A[0]) in this case what should be the X? it could be a 4 byte element, 2 bytes elements. so to jump to correct next position, X is the length of the element.

base + 0x4

base + 1x4

base + 2x4

so 4 is the elements length. hence, ecx * X is jumping to the next element in each iteration.

so what is displacement then ? Why is is used?

When the given array is **multidimensional** array. so to give an idea:

Scenario : in the C code I have 2 arrays of 10 bytes. I want to iterate over each array. using **base +(index * scale)**, I started at the base0 of the array1 and went through all the elements. Now I want to jump to the second array right? So in order to move the ebx to the base of the second array, I use displacement constant.

[0,1,2,3,4,5,6,7,8,9] [B0,B1,B2,B3,B4,B5,B6,B7,B8,B9]

when ebp is pointing to **0**, **base +(index * scale)** will iterate through the array. now I eax is pointing to **9** but I want to jump to **B0**. Addinng the **displacement** constant to ebp will help pointer to jump to **B0**

what does this mean ? ==> – **mov eax, [ebx]** go to ebx, read what is in that memory address, read 4 bytes from memory and write it on eax.

Note that brackets can be on one side or on the other side but never on both sides because there is not memory to memory operation!

just before exiting a function, generally the main function, we generally see the register mnemonic **mov ebp, esp** meaning level the base and the top pointers of the stack.

In practice this means destroy all the local variables in the stack. We are closing down the shop, nothing should be left in the stack.

Yani we are carrying the bottom level pointer(epb) to the top level (esp) and thus whatever downwards of the esp is garbage, undefined. No more in the stack.

The next move, of course, is to destroy the pointers from the stack which is `pop ebp` and finally `retn`.

`retn == ret`

NOTE: In VM Visual C++ Express, standard C libraries are not auto-recognized and yields debug error. In order to fix it, Under the project properties we need to add the address manually. In my case this was `C:\Program Files (x86)\Microsoft Visual Studio 10.0\VC\include`.

Example2.c

```
#include <stdlib.h>
int sub(int x, int y){
    return 2*x+y;
}

int main(int argc, char ** argv){
    int a;
    a = atoi(argv[1]);
    return sub(argc,a);
}
```

Takes arguments, since arguments (argv) are string, it converts it to int using atoi. then uses this variable in the sub() function, makes the arithmetic.

then returns the return of sub function. So main returns the return of sub.

here is the disassembled version of the code alongside the C code :

```
#include <stdlib.h>
int sub(int x, int y){
009E1020 push      ebp
009E1021 mov       ebp,esp
    return 2*x+y;
009E1023 mov       eax,dword ptr [ebp+8]
009E1026 mov       ecx,dword ptr [ebp+0Ch]
009E1029 lea       eax,[ecx+eax*2]
}
009E102C pop      ebp
009E102D ret
--- No source file -----
-----
009E102E int      3
009E102F int      3
--- c:\users\mehmetyavuzyagis\Desktop\intotoasm\example2\example2.c -----
-----

int main(int argc, char ** argv){
009E1030 push      ebp
009E1031 mov       ebp,esp
```

```

009E1033 push      ecx
    int a;
    a = atoi(argv[1]);
009E1034 mov       eax,dword ptr [ebp+0Ch]
009E1037 mov       ecx,dword ptr [eax+4]
009E103A push     ecx
009E103B call    dword ptr ds:[009E6238h]
009E1041 add      esp,4
009E1044 mov      dword ptr [ebp-4],eax
    return sub(argc,a);
009E1047 mov      edx,dword ptr [ebp-4]
009E104A push     edx
009E104B mov      eax,dword ptr [ebp+8]
009E104E push     eax
009E104F call    009E1005
009E1054 add      esp,8
}
009E1057 mov      esp,ebp
009E1059 pop      ebp
009E105A ret

```

NOTE if you see `ebp + something` this something is a passed argument that is trying to reach to memory

NOTE if you see `ebp - something` this passing a local variable.

`011A1033 push ecx` actually smart way to allocate a place for `int a;` 4 byte integer allocation. made by the compiler automatically. Thus, pushing ecx, 4 bytes, between ebp and esp , not stack has one 4 byte variable `a`.

Also checking the stack currently we can see that there is a 4 bytes difference between them:

eax	0x005f1cd8
ebx	0x7efde000
ecx	0x005f1480
argv[1]	0x005f14cb "256"
eip	0x009e1034
esp	0x0033f8a4
ebp	0x0033f8a8

`009E1034 mov eax,dword ptr [ebp+0Ch]` here, `[ebp +0Ch]` is the argv[1] specified in r/m32 Addressing Forms. So it is going to ebp, skipping the argv[0] which is the name of the program(example2) and reaching to argv[1]

`01001047 mov edx,dword ptr [ebp-4]` here, `[ebp-4]` is the first parameter we want in the subroutine and

we save the value to the edx and next, we push the `edx` to the stack.

`0100104B mov eax,dword ptr [ebp+8]` here `[ebp+8]` is the second parameter we want in the subroutine.

we save the value to the `eax` and we push the value `eax` to the stack.

Now both of the values of the parameters are defined and in the stack.

next `0100104F call 01001005` this calls the subroutine sub() at the address 0x01001005. that is the initial address of the sub()

> NOTE if you see `ebp + something` this something is a passed argument that is trying to reach to memory

> NOTE if you see `ebp - something` this passing a local variable.

`01001054 add esp,8` I had passed two parameters above, now the function is executed so I need to clean them up. for 2 integers from 4 bytes, I add 8 to esp.

then moving esp to ebp to destroy the local variables in the stack and popping the ebp and retting. exiting the program

Until now we used the following instructions:

```
NOP ==> copies eax to eax. so doing nothing
PUSH ==> pushes from register to memory
POP ==> pops from memory to register
CALL ==> pushes the address of the next instruction to the stack, changes eip to the next instruction.
RET-RETN
MOV ==> register to register, register to memory , immediate(hardcoded)
NEVER Memory to Memory!
LEA ==> instead of using r/m32 form, It makes its operations right on the spot
ADD ==> takes second parameter, adds to the first parameter and saves the result to the first parameter.
SUB
```

```
#include<stdio.h>
int main(){
    printf("Hello World!\n");
    return 0x1234;
}
```

This is a basic C code for hello world.

the assembly version is basic:

```
push ebp
mov ebp, esp
push offset HelloWorld;"Hello world\n"
call ds:_imp__printf
add esp,4
mov eax,1234h
```

```
pop ebp
retn
```

```
push ebp
mov ebp, esp
```

creating the stack

```
push offset aHelloWorld;"Hello world\n"
call ds:_imp_printf
```

pushing the hello World string from the pointer into the stack (from Right to left remember?) calling the printf function.

add esp,4

adding 4 to esp to destroy the data in the stack.

mov eax,1234h

per code, returing 0x1234 since eax holds the return values

```
pop ebp
retn
```

popping off the ebp to clear the stackframe and exiting the function.

```
#include<stdio.h>
int main(){
01091000 push      ebp
01091001 mov       ebp,esp
    printf("Hello World!\n");
01091003 push      offset __globalallocstatus-10h (1093000h)
01091008 call      dword ptr [_imp_printf (10920A8h)]
0109100E add       esp,4
    return 0x1234;
01091011 mov       eax,1234h
}
01091016 pop      ebp
01091017 ret
```

Here is the step by step resolution.

Control Flow in Assembly

There are 2 kinds of control flows.

- conditional : go somewhere if a condition is met. Think "if"s, switches, loops
- unconditional : go somewhere no matter what. Procedure calls, goto, exceptions, interrupts.

some of the procedure calls manifest themselves as push/call/ret

Now its time to see **goto** in assembly.

here is the source code for the current example to learn GOTO

```
//Goto example
#include <stdio.h>
int main(){
    goto mylabel;
    printf("skipped\n");
mylabel:
    printf("goto ftw!\n");
    return 0xf00d;
}
```

here is the step by step resolution.

```
#include <stdio.h>
int main(){
01151010  push      ebp
01151011  mov       ebp,esp
    goto mylabel;
01151013  jmp       01151023
    printf("skipped\n");
01151015  push      115500h
0115101A  call      dword ptr ds:[01156238h]
01151020  add       esp,4
mylabel:
    printf("goto ftw!\n");
01151023  push      115500Ch
01151028  call      dword ptr ds:[01156238h]
0115102E  add       esp,4
    return 0xf00d;
01151031  mov       eax,0F00Dh
}
01151036  pop       ebp
01151037  ret
```

so here the interesting code piece is **01151013 jmp 01151023** part. It jumps to the given eip and code keeps executing from there on.

`printf("skipped\n");`Keep in mind that, whatever under the GOTO was skipped due to the unconditional jump of the jmp. So we literally skipped the part `printf("skipped\n");`

Lets see the new C code and the resolution alongside assembly for jmp comparisons:

here is the C code:

```
int main(){
    int a=1, b=2;
    if(a == b){
        return 1;
    }
    if(a > b){
        return 2;
    }
    if(a < b){
        return 3;
    }
    return 0xdefea7;
}
```

int main(){	
00831010 push	ebp
00831011 mov	ebp,esp
00831013 sub	esp,8
int a=1, b=2;	
00831016 mov	dword ptr [ebp-4],1
0083101D mov	dword ptr [ebp-8],2
if(a == b){	
00831024 mov	eax,dword ptr [ebp-4]
00831027 cmp	eax,dword ptr [ebp-8]
0083102A jne	00831033
return 1;	
0083102C mov	eax,1
00831031 jmp	00831056
}	
if(a > b){	
00831033 mov	ecx,dword ptr [ebp-4]
00831036 cmp	ecx,dword ptr [ebp-8]
00831039 jle	00831042
return 2;	
0083103B mov	eax,2
00831040 jmp	00831056
}	
if(a < b){	
00831042 mov	edx,dword ptr [ebp-4]
00831045 cmp	edx,dword ptr [ebp-8]
00831048 jge	00831051
return 3;	
0083104A mov	eax,3

```

0083104F jmp      00831056
}
return 0xdefea7;
00831051 mov      eax, 0DEFEA7h
}
00831056 mov      esp, ebp
00831058 pop      ebp
00831059 ret

```

Now be careful in this assembly, we have some important mnemonics: **cmp, jne, jle, jge**

In order to create space for two local variables, the compiler pushed the esp by 8 bytes : **00831013 sub esp, 8**

after creating space in the stackframe, values 1 and 2 are assigned. :

```

00831016 mov      dword ptr [ebp-4], 1
0083101D mov      dword ptr [ebp-8], 2

```

ebp-4 ==1 and ebp-8 ==2

```

00831024 mov      eax, dword ptr [ebp-4]
00831027 cmp      eax, dword ptr [ebp-8]
0083102A jne      00831033

```

it takes the dereferenced pointer value of [ebp-4] to the eax. (which is 1)

then it compares the eax to dereferenced pointer value of [ebp-8] (which is 2)

if not equal (jne) jumps to the given address ahead : **00831033**

Then How do we set flags?

before making a conditional jump, we need something to set the condition flags for me.

this is typically done with CMP(compare) TEST, or some other flag-setting mnemonics.

so for example ZeroFlag is set if some operation is zero. In this codebase, this was the comparison of two integers by cmp.

although cmp is same as sub, sub saves the result but cmp just uses it to set flags, never touches the original values.

so also there should be another jump instruction after the cmp to follow if the comparison is 0.

in our example:

```
if(a == b){  
00831024  mov      eax,dword ptr [ebp-4]  
00831027  cmp      eax,dword ptr [ebp-8]  
0083102A  jne      00831033  
        return 1;  
0083102C  mov      eax,1  
00831031  jmp      00831056
```

00831031 jmp 00831056 this one.

this line jumps to thje end where stackframe is destroyed and functrion is terminated.