

Intel X86 Processors

Part 1(Intro)

Windows(Intel Syntax)

only 14 assembly instructions make up to 90 percent of the code!

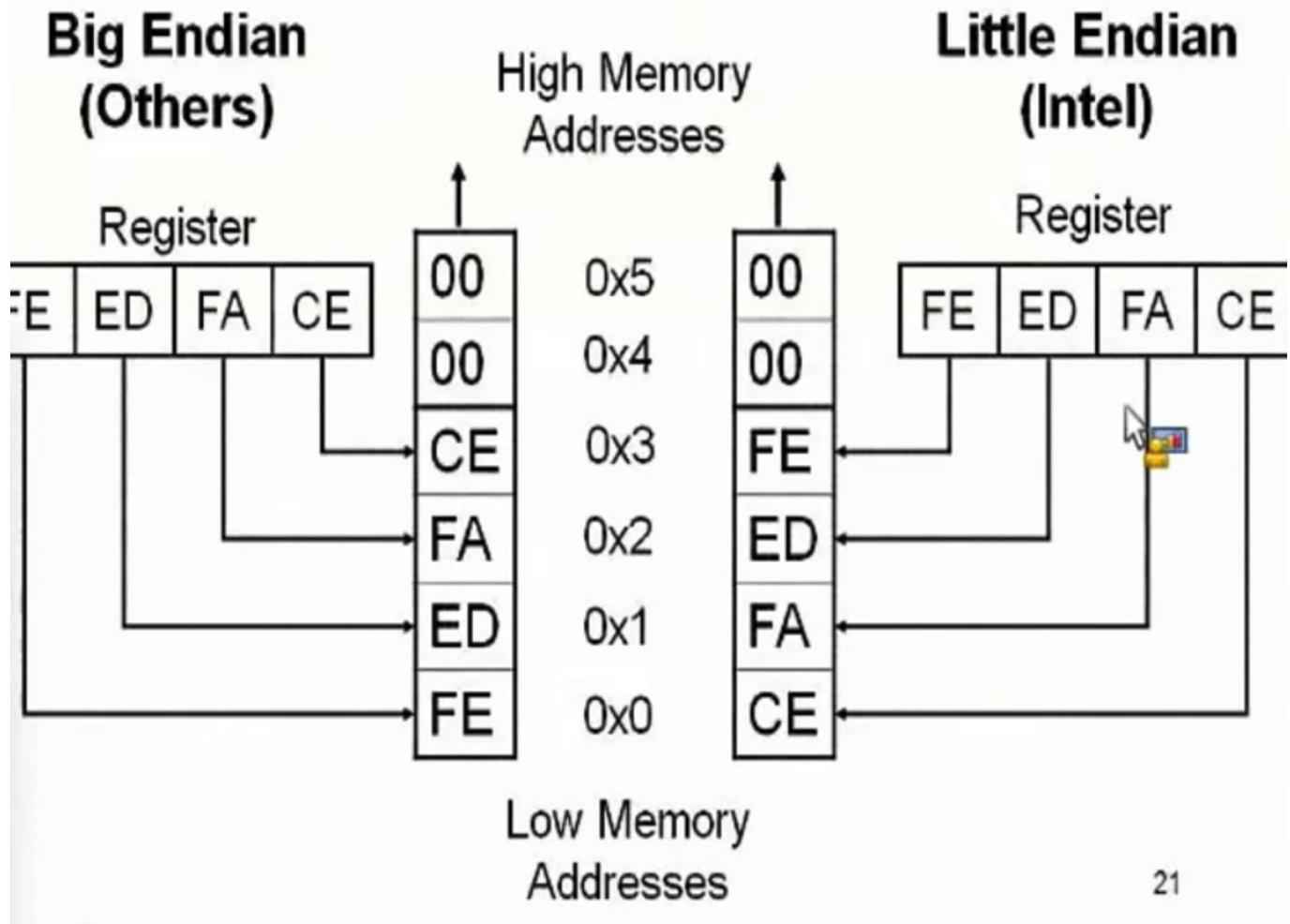
in intel terminology : a byte is a char(8) , a word is a short(16), int/long is double word(32), double long(64)

If you don't know this, you must memorize tonight

Decimal (base 10)	Binary (base 2)	Hex (base 16)
00	0000b	0x00
01	0001b	0x01
02	0010b	0x02
03	0011b	0x03
04	0100b	0x04
05	0101b	0x05
06	0110b	0x06
07	0111b	0x07
08	1000b	0x08
09	1001b	0x09
10	1010b	0x0A
11	1011b	0x0B
12	1100b	0x0C
13	1101b	0x0D
14	1110b	0x0E
15	1111b	0x0F

Intel is little endian. 0x12345678, least significant bit is registered first, then more significant and more significant.

Endianess pictures



Endians are only meaningful in byte level not in bit level!

Register Conventions!

EAX ==> stores function return value **EBX** ==> Base pointer to the data section **ECX**==>counter for string and loop ops **EDX** ==> I/O pointer **ESI**==>source pointer for string operations **EDI**==>Destination pointer for string operations **ESP**==>stack pointer **EBP**==>Stack frame base pointer **EIP**==>pointer to next instruction to execute(instruction pointer)

There are also other type of registers.

EAX, EDX , ECX == > Caller-save registers

I am a function and I will call another function. And when I call I am sure I am going to destroy the current registers(eax edx ecx) so we save the copy of the registers before calling the function so we do not lose the register

EBP, EBX, ESI, EDI ==> Callee-save registers

If I call a function, it wont change and modify the info in the register of ebp, ebx etc.

this **E** at the head of registers mean Extended. Since these are written originally for 16 bit , AX became Extended AX with 32 bits. etc.

EFLAGS

EFLAGS register holds many single bit flags. Will only ask you to remember the following for now :

Zero Flag(ZF)- Set if the result of some instruction is zero, cleared otherwise **Sign Flah(SF)** - Set equal to the most significant bit of the result.

INSTURCTIONS

NOP No Operation. No registers no values nothing. Just there to pad bytes or delay time. Bad guys use it to make simple exploits more reliable. But that's another class.

NOP actually takes a register and re-registers it on itself 😊

PUSH ==> push word, double word or quadword onto the STACK. can either be immediate(a numeric constant) of the value in a register.

For our purposes, it will always be a DWORD(for this course) the push instruction automatically decrements the stack pointer, esp, by 4. Why decrements? Because its stack, its FILO.

POP ==> Pop a value from the Stack

take a Dword off the stack, put it in a register and increment esp by 4 (reverse of Push)

CALL ==> Call's job is to transfer control to a different function in a way that control can later be resumed where it left off.

First it pushes the address of the next instruction onto the stack. for use by RET(return from procedure) for when the procedure is done.

Because idea behind a function is , do this and I will keep executing from the next line. So call also pushes the address of the next instruction set, so that once the called function is completed, the program keeps continuing from where it left off. Its a reminder to where to return to after executing the function.

Then it changes eip to the address given in the instruction.

Destination address can be specified in a multiple ways.

- Absolute address ==> address 0x0030434
- Relative address(relative to the end of the instruction) ==> some address that hex 50 bytes ta ta ta.

RET==> Return from Procedure.

Two forms :

1. pop the top of the stack into the eip

in this form, the instruction is just written as ret.

typically used by cdecl functions

2. pop the top of the stack into eip and add a constant number of bytes to esp.

in this form, the instruction is written as ret 0x8 or 0x20 etc. typically used by stdcall functions

MOV ==> Move.

Can move:

register to register

memory to register, register to memory, immediate to register immediate to memory.

immediate is hardcoded value.

NEVER MEMORY TO MEMORY

Memory addresses are given in r/m32 form.

ADD ==> a = a + b for example add eax,ebx means eax = eax + ebx. so it sums them and writes on what is on the left side

SUB ==> subtraction. sub eax, [ebx*2]

LEA ==> Load Effective Address. ==> frequently used with pointer arithmetic sometimes for just arithmetic in general. uses the r/m32 form but **is the exception to the rule**

Lea makes the calculation and appends to the destination without disturbing the memory. Does the calc on the spot.

that the [] syntax means dereference("value at")

example: ebx = 0X2, edx = 0x1000 --- lea eax,[edx+ebx*2] **eax = 0X1004, not the value at 0x1004!!**

JMP ==> Jump. Changes the eip unconditionally to the given address. Main forms of the address.

short relative : 1 byte displacement from the end of the instruction.

01151015 jmp 1155000h doesn't have the number **1155000h** anywhere in it, it is really **jmp 0x0E** bytes forward.

near relative: 4 bytes displacement from the current eip

absolute: hardcoded address in instruction. like **01151013 jmp 01151023**

absolute indirect : indirect address calculated with r/m32

jmp -2 is infinite loop for short relative jmp 😊

Conditional Jumps: **je, jge, jne, jle**

JCC ==> Jump if Condition is met.

JNE == JNZ jump if not equal means jump if not zero. What does it mean ? jne a b for example. the operation is a-b. so flag is either 0 or <0 or 0>

so 0 means equal, if not equal, its not zero. So it makes zero flag check

TEST ==> computes the bit-wise logical AND of first operand(source 1 operand) and teh second operand(source 2 operand) and sets the SF, ZFand PF staus flags accordingly.

like CMP, sets the flags and does not save the results

AND ==> destination operand can be r/m32 or register. source operand can be r/m32 or register or immediate!

`and al, bl`

`and eax, 0x42`

OR==> destination operand can be r/m32 or register. source operand can be r/m32 or register or immediate!

XOR ==> A xor B = If A , then A ; else then B.

there is a trick with XOR, when we xor something with itself, the result will always be 0. so ZF will be set.

`xor eax eax` sets the Zero Flag in that sense

thus, it is commonly used to zero a register by XORing it with itself because it is faster than `mov`.

NOT ==> single source/destination operand can be r/m32. One's complement negation.

`not al` or `not [al+bl]`.

example :

```
NOT 0x00110011b(al-0x33)
result 0x11001100b((al-0xCC))
```

SHL ==> Shift Logical Left. Can actually be used with the C and C++ << operator.

first operand(source and destination) operand is an r/m32

second operand is either a cl(lowest byte of ecx) or or a 1 byte immediate(hardcoded).

the 2nd operand is the number of places to shift

it multiplies the register by 2 for each placeteh value is shifted. More efficient than a multiply insturction.

bits shifted off to the left hand side are `shifted into` (set) the carry flag.(CF)

SHR ==> Shift Logical Right

Same thing with SHL, but this time each bit registers divide by power of that bit.

like `shr eax, 5 ==> divide eax with 2**5(32) and write it on eax.`

Still much effective than division and faster. because just shifting bits around.

when shifting bits to right so much so that it is now outbound of 8 bits, Carry flag(CF) is set to 1 to show that it is not actually 0 but it is carried.

IMUL ==> signed Multiplier. 3 forms:

1. imul r/m32 ==> edx:eax=eax*r/m32

for example : **imul ecx** take single parameter.

what does this mean : multiply ecx with eax and write the return to edx:eax

2. imul reg, r/m32 ==> reg=reg*r/m32

for example : **imul eax,ecx**

so eax times ecx and puts it back into eax.

if there is a chance for overflow, we make some modifications so it fits into the first form.

3. imul reg,r/m32, immediate ==> reg=r/m32*immediate

Multiplication and division have some possible side effects. For example, multiplication of two 32 bits may exceed the bounds of 32 bits of data.

Easily.

Thus, it is saved to edx:eax, so that it fits onto 64 bits.

note that **edx:eax** means concatenation. When eax is saved onto edx:eax it means the result of the operation is outbound of 32 bit limits!.

in 32 bit architecture, in this kind of a scenario, half of the result is written on edx and half of it is written on eax and compiler figures out how to output it.

DIV ==> Unsigned Divide.

2 Forms:

1. divide ax(16 bit of eax) by r/m8 , al=quotient, ah=remainder

2) edx:eax by r/m32, eax=quotient,edx=remainder

if the divisor is 0, a **divide by 0 exception** is raised.

REP STOS ==> Repeat Store String

One of a family of "rep" operations which repeat a single instruction multiple times(i.e. stos is also a standalone instruction)

Rep is not a standalone instruction, it's more of a prefix to stos.

all rep operations use **ecx** register as counter to determine how many times to loop through the instruction. Each time it executes, it decrements ecx. Once ecx=0, it continues to the next instruction.

either moves one byte at a time or one dword at a time.

either fill btye at [edi] with al or fill dword at [edi] with ax.

moves teh edi register forward one byte or one dword at at time, so that the repeated store operation is **storing into consecutive locations**.

thus, there are 4 pieces which must happen before the actual **rep stos** occurs

1. set **edi** to start the destination
2. (e)ax to the value to store
3. ecx to the number of times to store

NEG==> 2s Compliment negation. take a positive number make it negative for 2s.

so takes a value and outputs the negative value.

MUL ==> unsigned multiplication.

DEC ==> Decrement a register (substract one)

LEAVE==> tearing stackframe(mov esp, ebp pop ebp ret)

REP MOV==> rep repeats the following string operation ecx times. movs copies data from ds:esi to es:edi and increments or decrements the pointers based on the setting of the direction flag. As such, repeating it will move a range of memory to somewhere else

THE STACK

stack is the conceptual area of main memory which is designated by OS when program is started. Stack is LIFO/FILO data structure where data is pushed on to top of the stack and popped off the top

Stack grows over lower memory addresses. Adding something to the stack means the top of the stack is now at a lower memory address.

ESP points to the top the stack. The lowest address which is being used(since it is LIFO)

Stack keeps track of which functions were called before the current one. It holds LOCAL VARIABLES and is used to pass arguments to the next function to be called

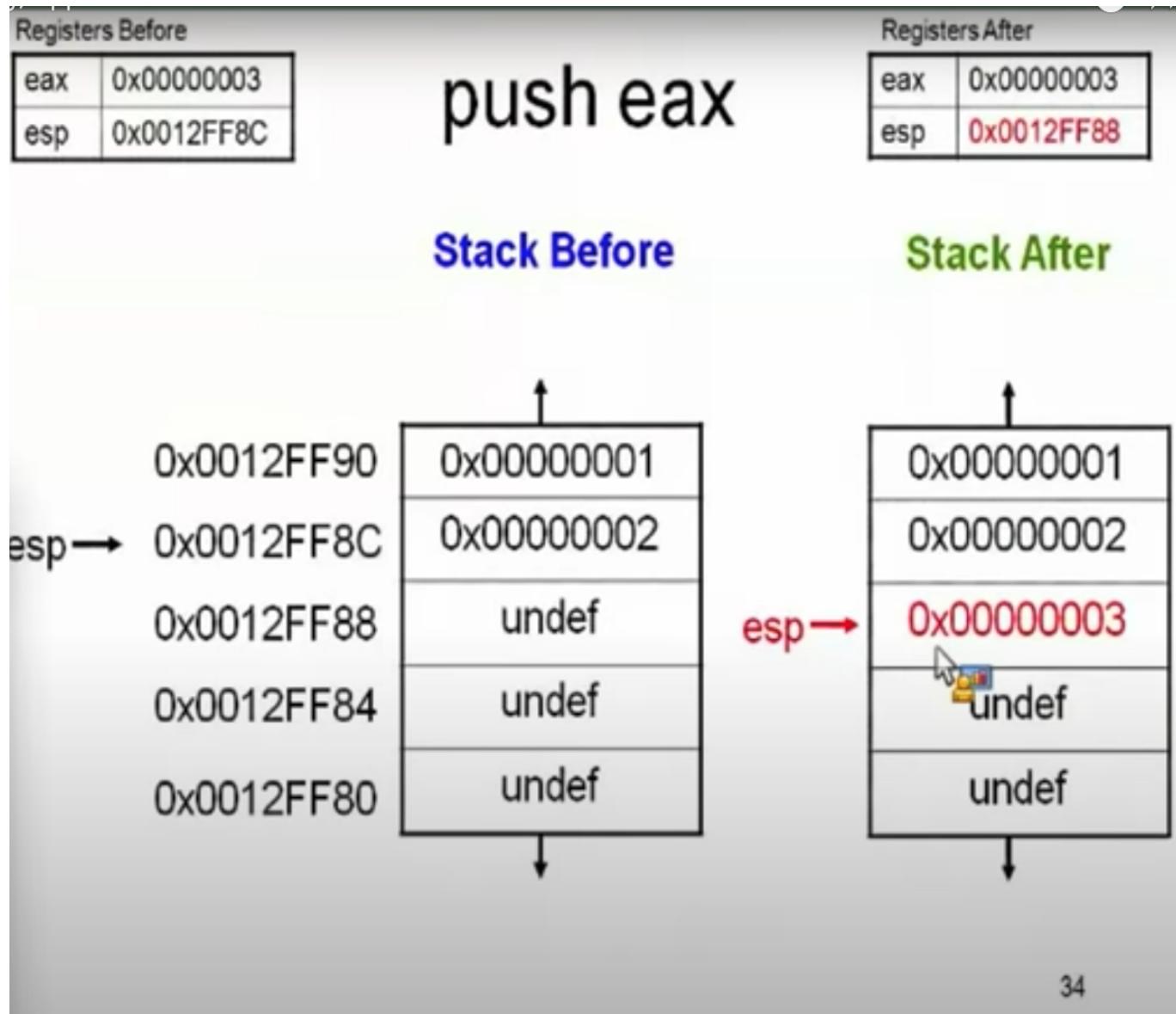
Register in time 1 :

eax : 0x00000003 esp : 0x0012FF8C

So when we call lets say , **push eax**, we expect the eax to be written 4 bytes after(decrement) the last esp register in the stack.

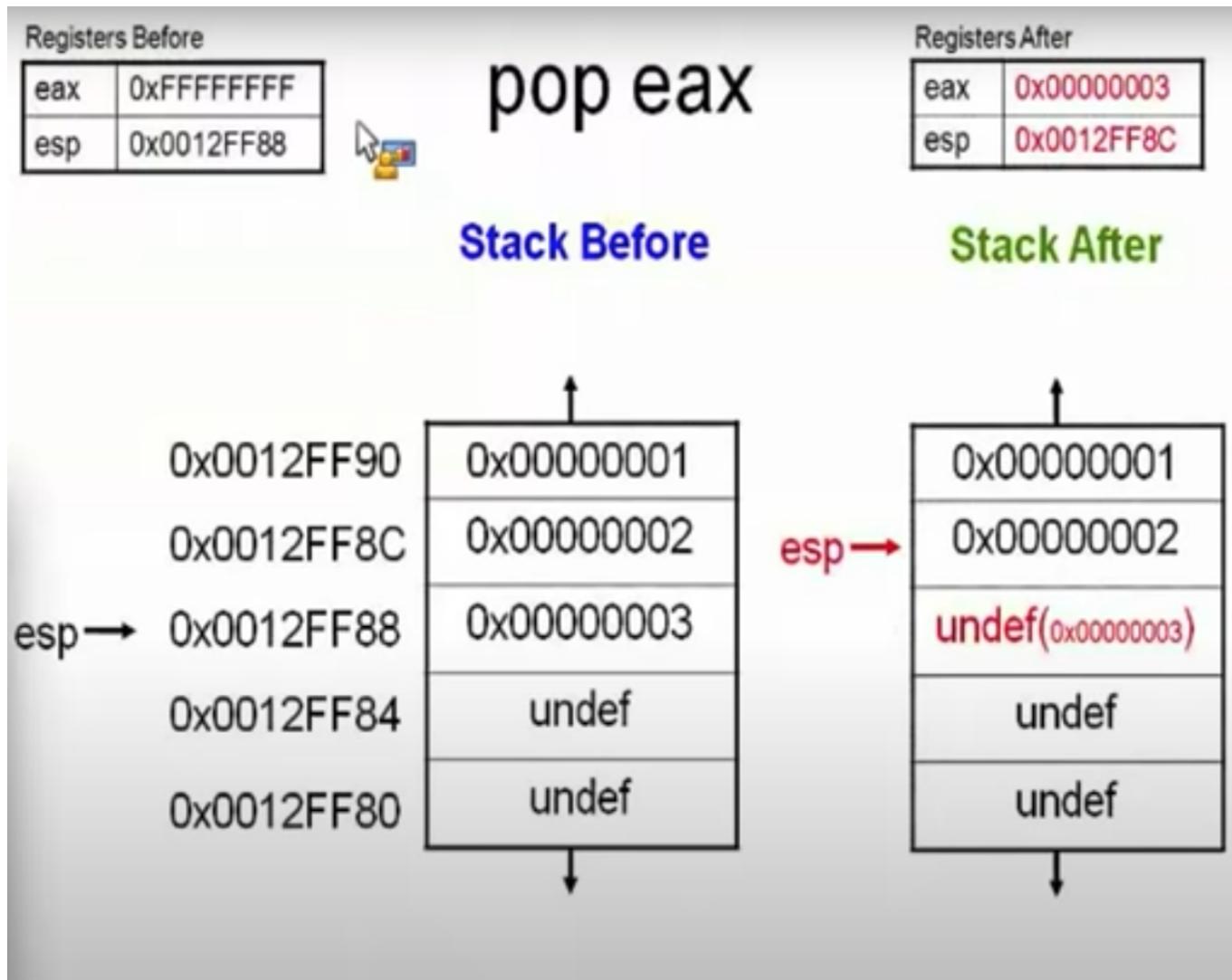
Register after push eax:

eax: 0x00000003 esp:0x0012FF88 ==> here, we jumped 4 bytes from FF8C to FF88. (where FF8C>FF88)



to reverse it , we use **POP**

in order to reverse the action above, we **pop eax** and voila, esp jumped by 4 bytes up (incremented) and top shelf of the stack is emptied.



Everything numerically under `esp` is considered to be undefined. Although there is a data over there, it is not of our business.

once we call `pop eax`, the the value at the top of the stack (`eax`) is taken to the register and that memory spot is not undefined.

Calling Convention

calling convention is about how to pass parameters and how to get parameters back.

how code calls a subroutine is compiler-dependent and configurable. But there are few conventions. We will deal with cdecl and stdcall conventions.

cdecl

C Declaration. The most common calling convention. Function parameters pushed (using push) onto stack RIGHT TO LEFT. Saves the old stack frame pointer and sets up a new stack frame.

`eax` or `edx:eax` returns the result for primitive data types.

Most significant 32 bits go into `edx`, least significant bits go into `eax`

Also values are always big endian in registers, in memory they are little endian.

so for example : `printf(%d\n, myVariable)` ==> from right to left: push my variable to the stack then push the pointer and then call the function.

function here, saves the old frame pointer, saves the address to the stack.

Caller is responsible for cleaning up the stack. ==>

so we have 2 parameters and 1 function call. push push and call. so whatever register is calling the function is also responsible to clean up the stack

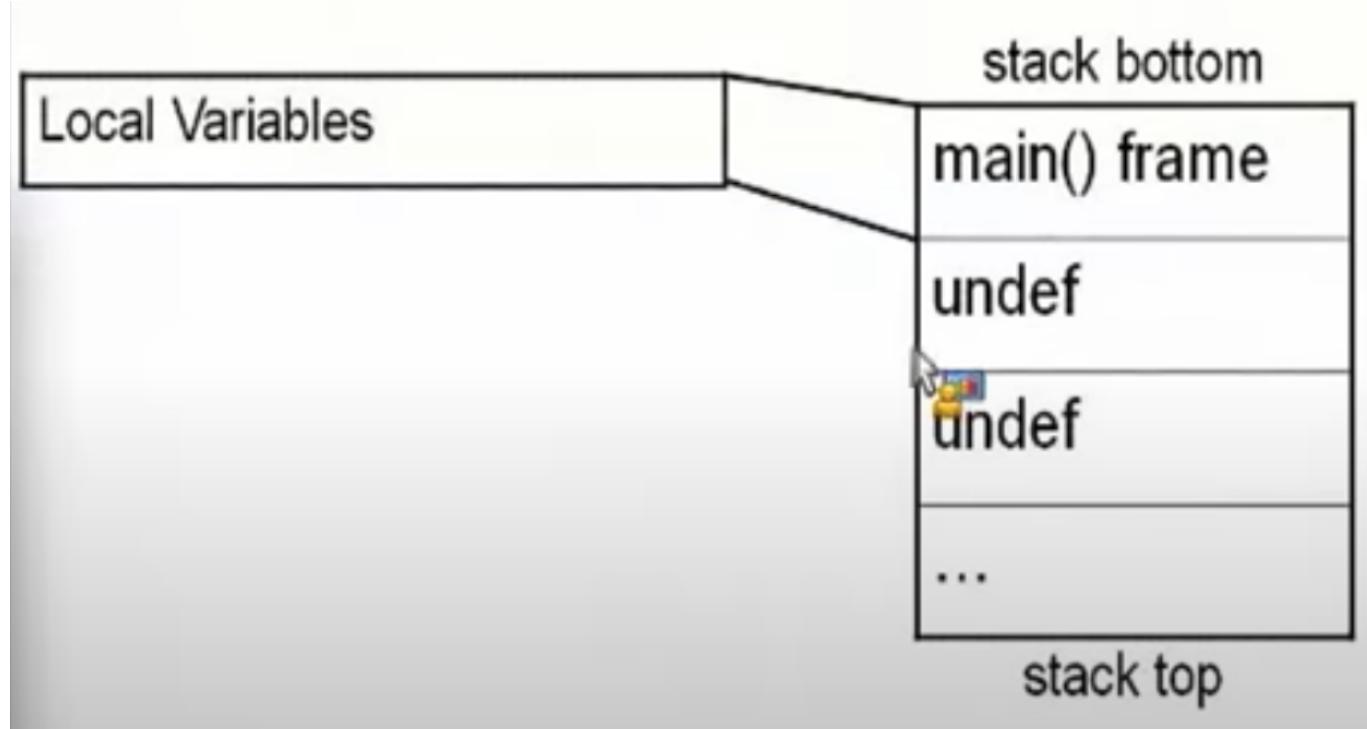
stdcall

Microsoft C++ code e.g. Win32API parameters pushed onto stack right to left saves the old stack frame pointer and sets up a new stack frame pointer

HERE, CALLEE is responsible for cleaning up any stack parameters it takes not the CALLER!

General Stack Frame Operation

We are going to assume that the `main()` is the very first function being executed in a program. This is what its stack looks like to start with.



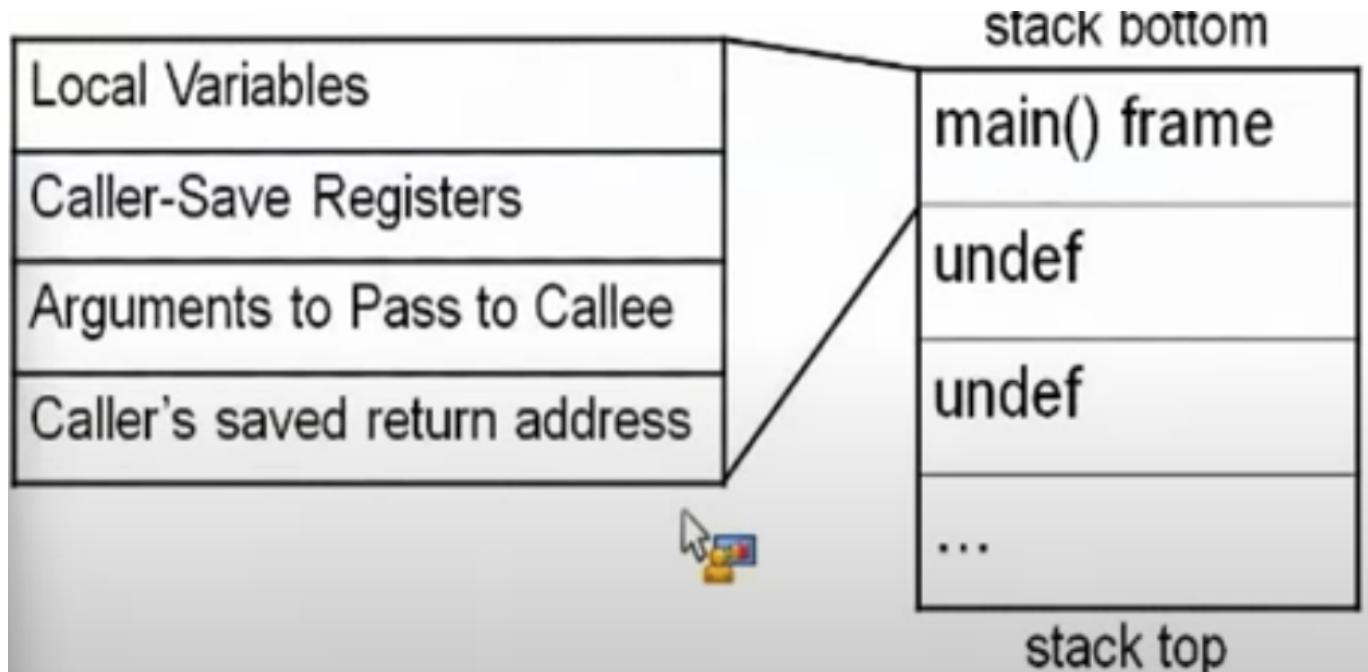
Low addresses at the bottom, high addresses at the top.

The more I add to the stack, the more I expect it to grow to the bottom.

So first thing `main` does is to preserve space on the stack for the local variables. When `main()` decides to call a subroutine, `main()` becomes the caller. We will assume `main()` has some registers it would like to remain the same, so it will save them. We will also assume that the callee function stakes some input args.

so , local variables, caller-save registers, and Args to pass to callee are in `main()` stack frame at the bottom of the stack.

so right to left, local variables, caller-save registers, args to pass to callee and then function call. git when we execute the call instruction, the return address(RET) gets save onto the stackm and becuse the next instruction after the call will be the beginning of the called function, wer consiter the frame to have changed to the callee.



next, when subroutine starts, the frame pointer(ebp) still points to the main()'s frame. so tthe first thing it does is to save the old frame pointer on the starck and set the new value to point to its own frame.

so at the main() fnction, ebp is pointing at the top of the frame. when a subroutine is called, it will save this adress to stack , take the ebp and contibnue its function, once the function is finished, it will restore the value of ebp so that the next subroutines can take up to from original ebp register.

STACK FRAMES AR A LINKED LIST!

The ebp in the current frame points at the saved ebp of the previous frame

Example1.c

The stack frames in this example will be very simple.
Only saved frame pointer (ebp) and saved return addresses (eip).

```
//Example1 - using the stack
//to call subroutines
//New instructions:
//push, pop, call, ret, mov
int sub(){
    return 0xbeef;
}
int main(){
    sub();
    return 0xf00d;
}

sub:
00401000 push    ebp
00401001 mov     ebp,esp
00401003 mov     eax,0BEEFh
00401008 pop    ebp
00401009 ret

main:
00401010 push    ebp
00401011 mov     ebp,esp
00401013 call    sub (401000h)
00401018 mov     eax,0F00Dh
0040101D pop    ebp
0040101E ret
```

```
int sub(){
    return 0xbeef;
}

int main(){
    sub();
    return 0xf00d;
}
```

NOTE : ON INTEL REGISTERS AND CALLS, DESTINATION IS ON THE LEFT AND SOURCE IS ON THE RIGHT

SO

MOV EBP,ESP MEANS

Move esp TO EBP!!!

IN AT&T syntax, this is the opposite tho.

sub:

```
00401000 push    ebp
00401001 mov     ebp,esp
```

```

00401003 mov eax,0BEEFh
00401008 pop ebp
00401009 ret

main:
00401010 push ebp
00401011 mov ebp,esp
00401013 call sub(401000h)
00401018 mov eax,0F00Dh
0040101D pop ebp
0040101E ret

```

note : push ebp mov ebp,esp codes are generated by compiler automatically

```

main:
00401010 push ebp

```

when code started to executing, there was some value in ebp. so initially, main pushes ebp(base pointer). so when we push ebp, since esp always points to the top of the stack, we also see that esp is also changed.

stack is decremented by 4 bytes with this push since it pushes stack to down(LIFO)

mov ebp,esp means take whatever value in the esp and put it onto ebp. because intel x86 syntax is **instruction destination,source**

so now ebp and esp are the same. so whatever is at the address of ebp and low is the new stack frame.

00401013 call sub(401000h) in the main function is very self-explanatory. call the function called sub at the address of 401000h which is the beginning of the address of sub.

so what is at the stack now ?

```

00401013 call sub(401000h)
00401018 mov eax,0F00Dh

```

subroutine(function sub) needs to register the address of the next register so that after the execution of the subroutine, main function picks up from where it left off. hence, the address of mov eax,0F00Dh is registered to the stack.

now the stack is hosting 00401018 right after ebp.

next step is to calling sub routine. when subroutine is called, it will save the ebp to the register (push ebp) this ebp is pointing at the main function normally. then it will move esp to ebp like main function(mov ebp,esp). now ebp and esp are pointing at the same place. why **mov ebp,esp** is important? **by setting esp to ebp, sub is creating its own stack!!** next step, sub can start functioning, which is writing 0XBEEF to eax.

why? because eax holds the return values and the sub() function was only returning 0xbeef;. thus it makes sense!

thus , in the stack, eax is changed to 0x0000BEEF address since this is our return value.

now sub() function is executed and we do not need the base pointer anymore. next step is pop ebp since this is the end of the stack.

end we will exit the sub() function with **ret**.

what **ret** does is that whatever is on the top iof the stack, go ahead and take it from there and I will put there the insturction pointer and so this will be our next destination.

remember the code :

```
sub:
00401000 push ebp
00401001 mov ebp,esp
00401003 mov eax,0BEEFh
00401008 pop ebp
00401009 ret

main:
00401010 push ebp
00401011 mov ebp,esp
00401013 call sub(401000h)
00401018 mov eax,0F00Dh
0040101D pop ebp
0040101E ret
```

so when **ret** (0x000401009) is called, the top of the stack will have the value of will be replaced to what?

it will be replaced with **0x000401018** since this is the first address after the sub() function call. thus main() function will keep goint from here. so next step is **00401018 mov eax,0F00Dh**

whici moves the value 0x00F00D to the eax because is the return value. if we did not put this register, the eax would be stil under the effect of sub() function so would keep pointing at 0x00BEEF. but main() wants a return of 0xFOOD so we immediately overwrite it.

next, we tear down the stack with **pop ebp** and returning to the function.

with pop ebp and ret, we cleared the last 2 remaining registers of stack. now stack is all undefined.

Some notes:

since sub() is a deadcode , meaning its return value 0xFOOD, is not used for anything, compiler with optimization option on would delete sub() function.

because there are not input parameters to sub(), there is no diff whether we compile cdecl or stdcall calling convention.

in the debugger when we jump to the assembly debugger here is the C code and the corresponding assembly code :

```

int main(){
011E1010 push      ebp
011E1011 mov       ebp,esp
    sub();
011E1013 call      sub (11E1000h)
    return 0xf00d;
011E1018 mov       eax,0F00Dh
}
011E101D pop      ebp
011E101E ret

```

Lets do another example.

```

#include <stdlib.h>
int sub(int x, int y){
    return 2*x+y;
}

int main(int argc, char ** argv){
    int a;
    a = atoi(argv[1]);
    return sub(a,a);
}

.text:00000000 _sub:    push  ebp
.text:00000001          mov   ebp,esp
.text:00000003          mov   eax,[ebp+8]
.text:00000006          mov   ecx,[ebp+0Ch]
.text:00000009          lea   eax,[ecx+eax*2]
.text:0000000C          pop   ebp
.text:0000000D          retn 
.text:00000010 _main:   push  ebp
.text:00000011          mov   ebp,esp
.text:00000013          push  ecx
.text:00000014          mov   eax,[ebp+0Ch]
.text:00000017          mov   ecx,[eax+4]
.text:0000001A          push  ecx
.text:0000001B          call  dword ptr ds:_imp_atoi
.text:00000021          add   esp,4
.text:00000024          mov   [ebp-4],eax
.text:00000027          mov   edx,[ebp-4]
.text:0000002A          push  edx
.text:0000002B          mov   eax,[ebp+8]
.text:0000002E          push  eax
.text:0000002F          call  _sub
.text:00000034          add   esp,8
.text:00000037          mov   esp,ebp
.text:00000039          pop   ebp
.text:0000003A          retn 

```

88

r/m32 Addressing Forms

Anywhere you see an r/m32 it means it could be taking a value either from a register or a memory address. (r/m == register/memory) in **Intel Syntax**, most iof the time square brackets([]) means to treat the value within as a memory address and fetch the value at the address (like dereferencing pointer)

- mov eax,ebx
- mov eax, [ebx]
- mov eax, [ebx+ecx*X] => X can be 1,2,4 or 8.

so for example here, it means, take the ecx(lets say 0x00344F) and multiply it with X, go to the address in this multiplication, add it to ebx, and write the cumulation to eax.

the most complicated form is : [base + index * scale + disp]

base +(index * scale) + displacement.

mov eax ,[ebx+ecx*X+Y] where Y is one byte(0-255) or 4 bytes(0-2^32-1)

Scenario : data type we have is an array. Its consequential data type. So ebx points to the base of the Array, beginning point of the array. Ecx is the index number(so A[0]) in this case what should be the X? it could be a 4 byte element, 2 bytes elements. so to jump to correct next position, X is the length of the element.

base + 0x4

base + 1x4

base + 2x4

so 4 is the elements length. hence, ecx * X is jumping to the next element in each iteration.

so what is displacement then ? Why is is used?

When the given array is **multidimensional** array. so to give an idea:

Scenario : in the C code I have 2 arrays of 10 bytes. I want to iterate over each array. using **base +(index * scale)**, I started at the base0 of the array1 and went through all the elements. Now I want to jump to the second array right? So in order to move the ebx to the base of the second array, I use displacement constant.

[0,1,2,3,4,5,6,7,8,9] [B0,B1,B2,B3,B4,B5,B6,B7,B8,B9]

when ebp is pointing to **0**, **base +(index * scale)** will iterate through the array. now I eax is pointing to **9** but I want to jump to **B0**. Addinng the **displacement** constant to ebp will help pointer to jump to **B0**

what does this mean ? ==> – **mov eax, [ebx]** go to ebx, read what is in that memory address, read 4 bytes from memory and write it on eax.

Note that brackets can be on one side or on the other side but never on both sides because there is not memory to memory operation!

just before exiting a function, generally the main function, we generally see the register mnemonic **mov ebp, esp** meaning level the base and the top pointers of the stack.

In practice this means destroy all the local variables in the stack. We are closing down the shop, nothing should be left in the stack.

Yani we are carrying the bottom level pointer(epb) to the top level (esp) and thus whatever downwards of the esp is garbage, undefined. No more in the stack.

The next move, of course, is to destroy the pointers from the stack which is `pop ebp` and finally `retn`.

`retn == ret`

NOTE: In VM Visual C++ Express, standard C libraries are not auto-recognized and yields debug error. In order to fix it, Under the project properties we need to add the address manually. In my case this was `C:\Program Files (x86)\Microsoft Visual Studio 10.0\VC\include`.

Example2.c

```
#include <stdlib.h>
int sub(int x, int y){
    return 2*x+y;
}

int main(int argc, char ** argv){
    int a;
    a = atoi(argv[1]);
    return sub(argc,a);
}
```

Takes arguments, since arguments (argv) are string, it converts it to int using atoi. then uses this variable in the sub() function, makes the arithmetic.

then returns the return of sub function. So main returns the return of sub.

here is the disassembled version of the code alongside the C code :

```
#include <stdlib.h>
int sub(int x, int y){
009E1020  push      ebp
009E1021  mov       ebp,esp
    return 2*x+y;
009E1023  mov       eax,dword ptr [ebp+8]
009E1026  mov       ecx,dword ptr [ebp+0Ch]
009E1029  lea       eax,[ecx+eax*2]
}
009E102C  pop      ebp
009E102D  ret
--- No source file -----
-----
009E102E  int      3
009E102F  int      3
--- c:\users\mehmetyavuzyagis\Desktop\intotoasm\example2\example2.c -----
-----

int main(int argc, char ** argv){
009E1030  push      ebp
009E1031  mov       ebp,esp
```

```

009E1033 push      ecx
    int a;
    a = atoi(argv[1]);
009E1034 mov       eax,dword ptr [ebp+0Ch]
009E1037 mov       ecx,dword ptr [eax+4]
009E103A push     ecx
009E103B call    dword ptr ds:[009E6238h]
009E1041 add      esp,4
009E1044 mov      dword ptr [ebp-4],eax
    return sub(argc,a);
009E1047 mov      edx,dword ptr [ebp-4]
009E104A push     edx
009E104B mov      eax,dword ptr [ebp+8]
009E104E push     eax
009E104F call    009E1005
009E1054 add      esp,8
}
009E1057 mov      esp,ebp
009E1059 pop      ebp
009E105A ret

```

NOTE if you see `ebp + something` this something is a passed argument that is trying to reach to memory

NOTE if you see `ebp - something` this passing a local variable.

`011A1033 push ecx` actually smart way to allocate a place for `int a;` 4 byte integer allocation. made by the compiler automatically. Thus, pushing ecx, 4 bytes, between ebp and esp , not stack has one 4 byte variable `a`.

Also checking the stack currently we can see that there is a 4 bytes difference between them:

eax	0x005f1cd8
ebx	0x7efde000
ecx	0x005f1480
argv[1]	0x005f14cb "256"
eip	0x009e1034
esp	0x0033f8a4
ebp	0x0033f8a8

`009E1034 mov eax,dword ptr [ebp+0Ch]` here, `[ebp +0Ch]` is the argv[1] specified in r/m32 Addressing Forms. So it is going to ebp, skipping the argv[0] which is the name of the program(example2) and reaching to argv[1]

`01001047 mov edx,dword ptr [ebp-4]` here, `[ebp-4]` is the first parameter we want in the subroutine and

we save the value to the edx and next, we push the `edx` to the stack.

`0100104B mov eax,dword ptr [ebp+8]` here `[ebp+8]` is the second parameter we want in the subroutine.

we save the value to the `eax` and we push the value `eax` to the stack.

Now both of the values of the parameters are defined and in the stack.

next **0100104F call 01001005** this calls the subroutine sub() at the address 0x01001005. that is the initial address of the sub()

> NOTE if you see ebp + something this something is a passed argument that is trying to reach to memory

> NOTE if you see ebp - something this passing a local variable.

01001054 add esp,8 I had passed two parameters above, now the function is executed so I need to clean them up. for 2 integers from 4 bytes, I add 8 to esp.

then moving esp to ebp to destroy the local variables in the stack and popping the ebp and retting. exiting the program

Until now we used the following instructions:

```
NOP ==> copies eax to eax. so doing nothing
PUSH ==> pushes from register to memory
POP ==> pops from memory to register
CALL ==> pushes the address of the next instruction to the stack, changes eip to the next instruction.
RET-RETN
MOV ==> register to register, register to memory , immediate(hardcoded)
NEVER Memory to Memory!
LEA ==> instead of using r/m32 form, It makes its operations right on the spot
ADD ==> takes second parameter, adds to the first parameter and saves the result to the first parameter.
SUB
```

```
#include<stdio.h>
int main(){
    printf("Hello World!\n");
    return 0x1234;
}
```

This is a basic C code for hello world.

the assembly version is basic:

```
push ebp
mov ebp, esp
push offsetHelloWorld;"Hello world\n"
call ds:_imp__printf
add esp,4
mov eax,1234h
```

```
pop ebp
ret
```

```
push ebp
mov ebp, esp
```

creating the stack

```
push offset aHelloWorld;"Hello world\n"
call ds:_imp_printf
```

pushing the hello World string from the pointer into the stack (from Right to left remember?) calling the printf function.

add esp,4

adding 4 to esp to destroy the data in the stack.

mov eax,1234h

per code, returing 0x1234 since eax holds the return values

```
pop ebp
ret
```

popping off the ebp to clear the stackframe and exiting the function.

```
#include<stdio.h>
int main(){
01091000 push      ebp
01091001 mov       ebp,esp
    printf("Hello World!\n");
01091003 push      offset __globalallocstatus-10h (1093000h)
01091008 call      dword ptr [_imp_printf (10920A8h)]
0109100E add       esp,4
    return 0x1234;
01091011 mov       eax,1234h
}
01091016 pop      ebp
01091017 ret
```

Here is the step by step resolution.

Control Flow in Assembly

There are 2 kinds of control flows.

- conditional : go somewhere if a condition is met. Think "if"s, switches, loops
- unconditional : go somewhere no matter what. Procedure calls, goto, exceptions, interrupts.

some of the procedure calls manifest themselves as push/call/ret

Now its time to see **goto** in assembly.

here is the source code for the current example to learn GOTO

```
//Goto example
#include <stdio.h>
int main(){
    goto mylabel;
    printf("skipped\n");
mylabel:
    printf("goto ftw!\n");
    return 0xf00d;
}
```

here is the step by step resolution.

```
#include <stdio.h>
int main(){
01151010  push      ebp
01151011  mov       ebp,esp
    goto mylabel;
01151013  jmp       01151023
    printf("skipped\n");
01151015  push      115500h
0115101A  call      dword ptr ds:[01156238h]
01151020  add       esp,4
mylabel:
    printf("goto ftw!\n");
01151023  push      115500Ch
01151028  call      dword ptr ds:[01156238h]
0115102E  add       esp,4
    return 0xf00d;
01151031  mov       eax,0F00Dh
}
01151036  pop       ebp
01151037  ret
```

so here the interesting code piece is **01151013 jmp 01151023** part. It jumps to the given eip and code keeps executing from there on.

`printf("skipped\n");`Keep in mind that, whatever under the GOTO was skipped due to the unconditional jump of the jmp. So we literally skipped the part `printf("skipped\n");`

Lets see the new C code and the resolution alongside assembly for jmp comparisons:

here is the C code:

```
int main(){
    int a=1, b=2;
    if(a == b){
        return 1;
    }
    if(a > b){
        return 2;
    }
    if(a < b){
        return 3;
    }
    return 0xdefea7;
}
```

int main(){	
00831010 push	ebp
00831011 mov	ebp,esp
00831013 sub	esp,8
int a=1, b=2;	
00831016 mov	dword ptr [ebp-4],1
0083101D mov	dword ptr [ebp-8],2
if(a == b){	
00831024 mov	eax,dword ptr [ebp-4]
00831027 cmp	eax,dword ptr [ebp-8]
0083102A jne	00831033
return 1;	
0083102C mov	eax,1
00831031 jmp	00831056
}	
if(a > b){	
00831033 mov	ecx,dword ptr [ebp-4]
00831036 cmp	ecx,dword ptr [ebp-8]
00831039 jle	00831042
return 2;	
0083103B mov	eax,2
00831040 jmp	00831056
}	
if(a < b){	
00831042 mov	edx,dword ptr [ebp-4]
00831045 cmp	edx,dword ptr [ebp-8]
00831048 jge	00831051
return 3;	
0083104A mov	eax,3

```

0083104F jmp      00831056
}
return 0xdefea7;
00831051 mov      eax, 0DEFEA7h
}
00831056 mov      esp, ebp
00831058 pop      ebp
00831059 ret

```

Now be careful in this assembly, we have some important mnemonics: **cmp, jne, jle, jge**

In order to create space for two local variables, the compiler pushed the esp by 8 bytes : **00831013 sub esp, 8**

after creating space in the stackframe, values 1 and 2 are assigned. :

```

00831016 mov      dword ptr [ebp-4], 1
0083101D mov      dword ptr [ebp-8], 2

```

ebp-4 ==1 and ebp-8 ==2

```

00831024 mov      eax, dword ptr [ebp-4]
00831027 cmp      eax, dword ptr [ebp-8]
0083102A jne      00831033

```

it takes the dereferenced pointer value of [ebp-4] to the eax. (which is 1)

then it compares the eax to dereferenced pointer value of [ebp-8] (which is 2)

if not equal (jne) jumps to the given address ahead : **00831033**

Then How do we set flags?

before making a conditional jump, we need something to set the condition flags for me.

this is typically done with CMP(compare) TEST, or some other flag-setting mnemonics.

so for example ZeroFlag is set if some operation is zero. In this codebase, this was the comparison of two integers by cmp.

although cmp is same as sub, sub saves the result but cmp just uses it to set flags, never touches the original values.

so also there should be another jump instruction after the cmp to follow if the comparison is 0.

in our example:

```

    if(a == b){
00831024  mov      eax,dword ptr [ebp-4]
00831027  cmp      eax,dword ptr [ebp-8]
0083102A  jne      00831033
            return 1;
0083102C  mov      eax,1
00831031  jmp      00831056

```

00831031 jmp 00831056 this one.

this line jumps to the end where stackframe is destroyed and function is terminated.

je is the same as **jz** because zero flag is still essentially an equality check. In assembly code, je and jz does exactly the same thing.

jle is the same as **jng** because jump less than or equal to exactly means jump if not greater than. so many of the jump instructions have multiple ways to put.

Important note regarding the difference between less than and below!

JBE ==> jump below or equal to. Less than has to do with signed numbers, below has to do with unsigned numbers.

Compiler has to know that -1 is not equal to 1 for example. FFFFF is -1 and 1 is 1. In chart, this is confusing for unsigned numbers. So to make things easier, we use below, since FFFFFFF is below 1.

For Loop in X86 ASM

Here is the C code:

```

//Simple loop with single variable
#include <stdio.h>

int main(){
    int i;
    for(i = 0; i < 10; i++){
        printf("i = %d\n", i);
    }
}

```

Now let's deconstruct this code in the 32-bit assembly.

```

2: #include <stdio.h>
3:
4: int main(){
00951010  push      ebp
00951011  mov       ebp,esp
00951013  push      ecx

```

```

5:    int i;
6:    for(i = 0; i < 10; i++){
00951014  mov      dword ptr [ebp-4],0
0095101B  jmp      00951026
0095101D  mov      eax,dword ptr [ebp-4]
00951020  add      eax,1
00951023  mov      dword ptr [ebp-4],eax
00951026  cmp      dword ptr [ebp-4],0Ah
0095102A  jge      00951040
7:        printf("i = %d\n", i);
0095102C  mov      ecx,dword ptr [ebp-4]
0095102F  push     ecx
00951030  push     955000h
00951035  call     dword ptr ds:[00956238h]
0095103B  add      esp,8
8:    }
0095103E  jmp      0095101D
9: }
00951040  xor      eax,eax
00951042  mov      esp,ebp
00951044  pop      ebp
00951045  ret

```

veeery neat and beautiful!

Step 1)

```

00951010  push     ebp
00951011  mov      ebp,esp
00951013  push     ecx

```

In order to initialize the stackframe, pushed base pointer. then created the top of the stack with esp then pushed **ecx** in order to open up space for **i**

Step 2)

```

00951014  mov      dword ptr [ebp-4],0
0095101B  jmp      00951026
0095101D  mov      eax,dword ptr [ebp-4]
00951020  add      eax,1
00951023  mov      dword ptr [ebp-4],eax
00951026  cmp      dword ptr [ebp-4],0Ah
0095102A  jge      00951040

```

what is [ebp-4]? it is ecx right? so adding 0 to ecx to realize **i=0**. then jumping to address **00951026** which is **cmp dword ptr [ebp-4],0Ah**.

this cmp comparison, compares the dword pointer [ebp-4] with 0Ah. 0A is 10 in hexadecimal. so comparing the current value of ecx to 10.

since it is not great or equal to 10, we come to the point **00951020 add eax,1** and incrementing i by 1. then pushing eax (now 1) to ecx(epb-4) and repeating the same steps.

step 3)

```
00951040 xor      eax,eax
00951042 mov      esp,ebp
00951044 pop     ebp
00951045 ret
```

once we have the equality, there comes **00951040 xor eax, eax**

as we know, xor a , a is the syntax of zeroing the register instead of moving zero to that place. then regular esp ebp and popping ebp and exiting the function.

0095103B add esp,8 and here, esp + 8 means clearing the stack. but why 8 ?

4 is for i

4 is for printf function.

nice to notice:

```
0095101D mov      eax,dword ptr [ebp-4]
00951020 add      eax,1
00951023 mov      dword ptr [ebp-4],eax
```

this part corresponds to **i++**

take i, add to eax,

add 1 to eax,

take eax and put it into i. 😊

Working with Multiply / Divide (shift left/ right)

to Demonstrate the multiplication/ division operations using **shift**, we will be using the following C code:

```
//Multiply and divide transformations
//New instructions:
//shl - Shift Left, shr - Shift Right

int main(){
    unsigned int a, b, c;
```

```

    a = 0x40;
    b = a * 8;
    c = b / 32;
    return c;
}

```

And this is the corresponding resolution with assembly:

```

5: int main(){
01001010 push      ebp
01001011 mov       ebp,esp
01001013 sub       esp,0Ch
6:     unsigned int a, b, c;
7:     a = 0x40;
01001016 mov       dword ptr [ebp-4],40h
8:     b = a * 8;
0100101D mov       eax,dword ptr [ebp-4]
01001020 shl       eax,3
01001023 mov       dword ptr [ebp-8],eax
9:     c = b / 32;
01001026 mov       ecx,dword ptr [ebp-8]
01001029 shr       ecx,5
0100102C mov       dword ptr [ebp-0Ch],ecx
10:    return c;
0100102F mov       eax,dword ptr [ebp-0Ch]
11: }
01001032 mov       esp,ebp
01001034 pop       ebp
01001035 ret

```

as can be seen, there are shift right and shift left mnemonics. These correspond to the mult/div operations because there is no mult/ div operations in assembly like there is add/sub. but how do they work?

SHL, the first instruction is the source and destination. the second operand is the number of the multiplication you want source to be multiplied by and you take it and re-assign it onto the first parameter which is the destination.

Do not forget that each time there is a shift iteration, it multiplies by two.

`b = a * 8;` corresponds to the asm code of `01001020 shl eax,3`

shift left eax, argument is 3 so $2^{**3} = 8$. means. $\text{eax} = \text{eax} * 8$.

one bit movement is multiplication by two.

why? because in binary notation, each time you step to another bit, it is 1 level up of base of two.

010100000

(20 * 0) + (21 *1) + (22 * 0)** ikiler basamagi mantigi.

each bit to the left is multiplication of base 2 in binary.

When there is a shift left, last bit of the binary value changes to 0.

like from **0101001** to **1010010**, shifted 1 bit to left for multiplication. the last bit is now zero, yea? in order cpu not to raise 0 flag, this SHL sets carry flag(CF).

Its like I have zeroes all over but the real value is not 0.

sometimes we have 8 bit of **10000000**. when shifted it becomes **1|00000000** 1 is out of the bound of 8 bits. Machine should not read it as value **0** CF helps machine to say that 1 is outbound but valid so number is not 0.

When all things are inbound, CF is set to 0; when after shifting there is an outbound, then CF is set to 1.

What is the difference between logical and arithmetic shift?

Logical shift is unsigned integers, arithmetic shift signed integers.

SHR is the same with SHL.

for example **shr eax, 5 ==> divide eax with 2**5(32) and write it on eax.**

when shifting to right, the 8 bits can return to zero and the rightmost bit can be outbound of 8 bit limits.(it could be carried to 10th bit) in that case it would look like

0x00000000|001 see, 1 is outbound of 8 bit limits. But the number is not actually 0!.

In order to prevent such confusion, in this situation a Carry Flag is set to 1 to note that although first 8 bits are 0, actual value is not 0.

turning back to the assembly code:

```
01001013 sub esp, 0Ch
6:    unsigned int a, b, c;
```

3 integers are initialized. thus, esp should be 12 bits lower than initial phase. here,

sub esp, 0Ch means (c ==12 in hex) open space for 12 so I fit 3 integers.

a == ebp-4 b==ebp-8 c== ebp-12

then **01001016 mov dword ptr [ebp-4],40h**

moving the value of 40h to ebp-4, which is a.

```
9:    c = b / 32;
01001026 mov ecx, dword ptr [ebp-8]
01001029 shr ecx, 5
0100102C mov dword ptr [ebp-0Ch], ecx
```

registering b, to ebp-8

shifting right by 5, dividing by 32, and writing the ecx on ebp-0c(12)

now after the calculations, it is time to return **c**.

return is assigned to **eax**

```
10:      return c;
0100102F  mov         eax,dword ptr [ebp-0Ch]
```

finally tearing down the stackframe

```
01001032  mov         esp,ebp
01001034  pop         ebp
01001035  ret
```

First time so far, we have more than 1 local variables, a,b, and c.

here comes the power of mov ebp,esp. in one move, it tears down the stackframe.

once ebp is brought up to esp, whatever under the esp-ebp line is garbage and stackframe is cleared at once! perfect

Then , when multiplication and division is to be made by power of 2, shr and shl are used.

Multiply and Divide when operand is not power of 2

Sometimes a division is made by 3, multiplication is made by 17 etc. not powers of 2.

In that case, we use **div** and **imul**

the C code we will examine is the following:

```
//Multiply and divide operations
//when the operand is not a
//power of two
//New instructions: imul, div
int main(){
    unsigned int a = 1;
    a = a * 6;
    a = a / 3;
    return 0x2bad;
}
```

The disassembled version of the code is the following:

```

5: int main(){
012A1010 push      ebp
012A1011 mov       ebp,esp
012A1013 push      ecx
    6:     unsigned int a = 1;
012A1014 mov       dword ptr [ebp-4],1
    7:     a = a * 6;
012A101B mov       eax,dword ptr [ebp-4]
012A101E imul     eax,eax,6
012A1021 mov       dword ptr [ebp-4],eax
    8:     a = a / 3;
012A1024 mov       eax,dword ptr [ebp-4]
012A1027 xor       edx,edx
012A1029 mov       ecx,3
012A102E div       eax,ecx
012A1030 mov       dword ptr [ebp-4],eax
    9:     return 0x2bad;
012A1033 mov       eax,2BADh
  10: }
012A1038 mov       esp,ebp
012A103A pop

```

so in this example multiplier and divisor are not powers of 2. Hence, we use **imul** and **div**

so after the generic memory allocation, we need to multiply a with 6.

012A101E imul eax,eax,6 so take eax, multiply it with 6 and write it over 6.

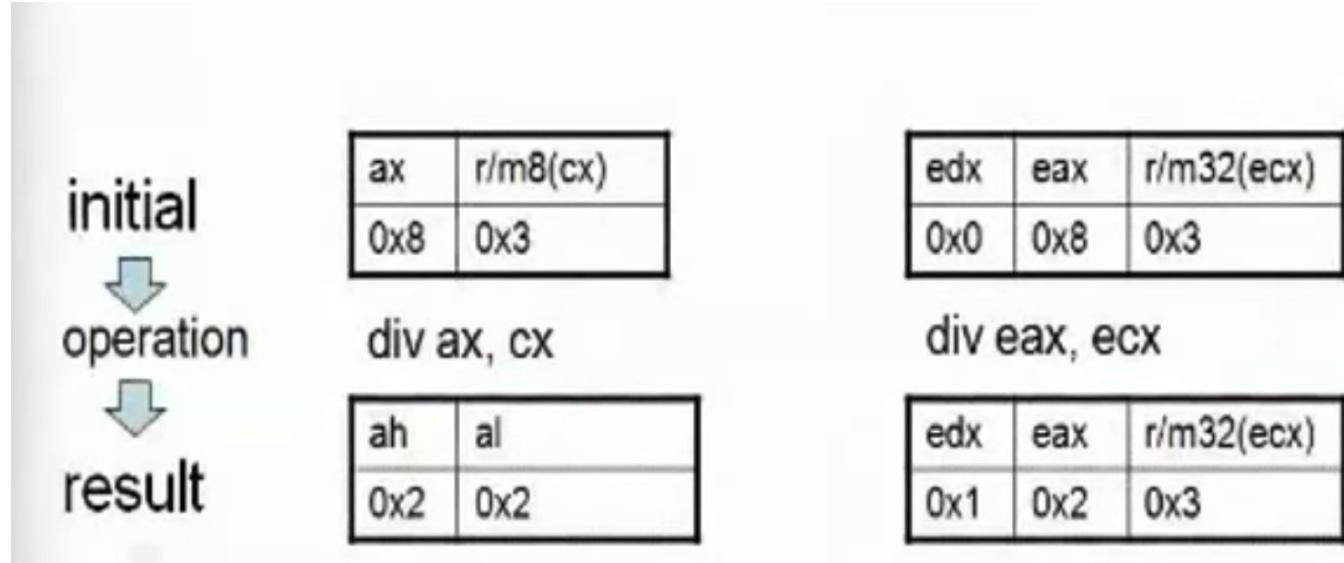
Since the multiplication is still inbounds of 32 bit limits, we used the **r/m32 * immediate** syntax, not the other **imul** syntaxes.

in the scenario of **imul ecx**

so for example, lets say we have 0x44000000 and we want to multiply it with 0x4(r/m32)

since the result is greater than what 32 bit can handle, edx:eax will share the result and while edx holds the 1 at the most left side, eax will hold the rest of the 8 digits.

012A102E div eax,ecx here we have div mnemonic, that is unsigned division operator.



so for example:

```
divide 0x8 by 0x3.
0x8==ax
r/m8==cx
ah==0x2 (remainder)
al==0x2 (quotient)
```

Buffer initialization

The C code we will examine in the part is :

```
//VisualStudio runtime check
//buffer initialization
//auto-generated code
//New instruction: rep stos

int main(){
    char buf[40];
    buf[39] = 42;
    return 0xb100d;
}
```

what does it do is basically allocating a memory space for 40 bytes and allocating the 42 to the 39th index place of the buffer.

then returning 0xb100d

here is the disassembled code :

```
int main()
008E1010  push        ebp
```

```

008E1011    mov        ebp,esp
008E1013    sub        esp,30h
008E1016    push       edi
008E1017    lea        edi,[ebp-30h]
008E101A    mov        ecx,0Ch
008E101F    mov        eax,0CCCCCCCCCh
008E1024    rep stos   dword ptr es:[edi]
    char buf[40];
    buf[39] = 42;
008E1026    mov        byte ptr [ebp-5],2Ah
    return 0xb100d;
008E102A    mov        eax,0B100Dh
}
008E102F    push       edx
008E1030    mov        ecx,ebp
008E1032    push       eax
008E1033    lea        edx,[ (8E1048h)]
008E1039    call      _RTC_CheckStackVars (8E10B0h)
008E103E    pop        eax
008E103F    pop        edx
008E1040    pop        edi
008E1041    mov        esp,ebp
008E1043    pop        ebp
008E1044    ret

```

but wait. what is **01331024 rep stos dword ptr es:[edi]** ?

what is rep stos here?

lets remember:

all rep operations use **ecx** register as counter to determine how many times to loop through the instruction. Each time it executes, it decrements ecx. once ecx=0, it continues to the next instruction.

moves the edi register forward one byte or one dword at a time, so that the repeated store operation is **storing into consecutive locations**.

thus, there are 4 pieces which must happen before the actual **rep stos** occurs

1. set **edi** to start the destination
2. (e)ax to the value to store
3. ecx to the number of times to store

rep stos setup

004113AC lea edi,[ebp-0F0h]

Set edi - the destination

004113B2 mov ecx,3Ch

Set ecx - the count

004113B7 mov eax,0CCCCCCCCh

Set eax - the value

004113BC rep stos dword ptr es:[edi]

Start the repeated store

- So what's this going to do? Store 0x3C copies of the dword 0xCCCCCCCC starting at ebp-0xF0
- And that just happens to be 0xF0 bytes of 0xCC!

So stos takes a byte or a dword, and writes it. Rep prefix tells it how many times to write. For counter to work, **ecx** is set, and in each iteration it decrements by one until it reaches to 0.

```
008E101F    mov      eax,0CCCCCCCCh  
008E1039    call     _RTC_CheckStackVars (8E10B0h)
```

These two lines are protection against buffer overflow at the end of two ends of stack. Do not get confused. These are auto-generated breakpoints to guard up.

008E1016 push edi

so why we pushed edi? also we popped it later on? we will use rep mnemonics , compiler knows it, and rep always uses edi for its operations .

we allocated 40 bytes of buffer in the C code but in assembly ,`00E6101A mov ecx,0Ch`

we moved 0c(12) to ecx, and $12 \times 4 = 48$ bytes of stackframe created

there are 4 on top 4 on bottom byte extra space created and put CCCs below and above.(`00E6101F mov eax,0CCCCCCCCCh`)

Journey To the Center of MEMCPY

So far we could not transfer data memory to memory, we used mov etc to use registers. Instead, memcpy helps direct data-transfer from mem to mem.

```
//Journey to the center of memcpy
#include <stdio.h>
#include <string.h>

typedef struct mystruct{
    int var1;
    char var2[4];
} mystruct_t;

int main(){
    mystruct_t a, b;
    a.var1 = 0xFF;
    memcpy(&b, &a, sizeof(mystruct_t));
    return 0xAce0Ba5e;
}
```

Now what the code does?

It defines a custom data type: mystruct_t.

This data type has int(var1) and an array of char(var2) of 4 bytes.

when the main function is called, the function declares 2 mystruct_t(a,b)

a's var1(int) is set to 0xFF which is 255

then there is a `memcpy` function call. What is memcpy?

The C library function `void *memcpy(void *dest, const void *src, size_t n)`
copies n characters from memory area src to memory area dest.

for example the following code returns:

```
#include <stdio.h>
#include <string.h>
```

```

int main () {
    const char src[50] = "http://www.tutorialspoint.com";
    char dest[50];
    strcpy(dest,"Heloooo!!");
    printf("Before memcpy dest = %s\n", dest);
    memcpy(dest, src, strlen(src)+1);
    printf("After memcpy dest = %s\n", dest);

    return(0);
}

```

Before memcpy dest = Heloooo!!
After memcpy dest = http://www.tutorialspoint.com

thus , memcpy takes 3 arguments: destination to fill, source to fill from, and size. `void *memcpy(void *dest, const void * src, size_t n)`

Hence, in our example,

`memcpy(&b, &a, sizeof(mystruct_t));`

value in memory address b is destination, value of memory address a is source and the size is mystruct_t

here is the disassembled version of the code :

```

//Journey to the center of memcpy
#include <stdio.h>
#include <string.h>
typedef struct mystruct{
    int var1;
    char var2[4];
} mystruct_t;

int main(){
01281010  push        ebp
01281011  mov         ebp,esp
01281013  sub         esp,10h  ==> decimal 16
    mystruct_t a, b;
    a.var1 = 0xFF;
01281016  mov         dword ptr [ebp-8],0FFh
    memcpy(&b, &a, sizeof(mystruct_t));
0128101D  push        8
0128101F  lea         eax,[ebp-8]
01281022  push        eax
01281023  lea         ecx,[ebp-10h]
01281026  push        ecx
01281027  call        01281042
0128102C  add         esp,0Ch
    return 0xAce0Ba5e;
}

```

```
0128102F    mov          eax, 0ACE0BA5Eh  
}  
01281034    mov          esp, ebp  
01281036    pop          ebp  
01281037    ret
```

nice easter eggs in the return values 😊

return 0xAce2Ba5e; return ace to base 😊

01281013 sub esp,10h ==> decimal 16

creating place for 16 bits. why? because mystruct has an int and 4-bytes array.

we declare two of them(a,b) in total we initially need 16.

Read this very very carefully to understand what is happening in memcpy!

- create stackframe
- SUB ESP, 10H

WHY? our typedef is 4 byte array and an integer in total 8 bytes.

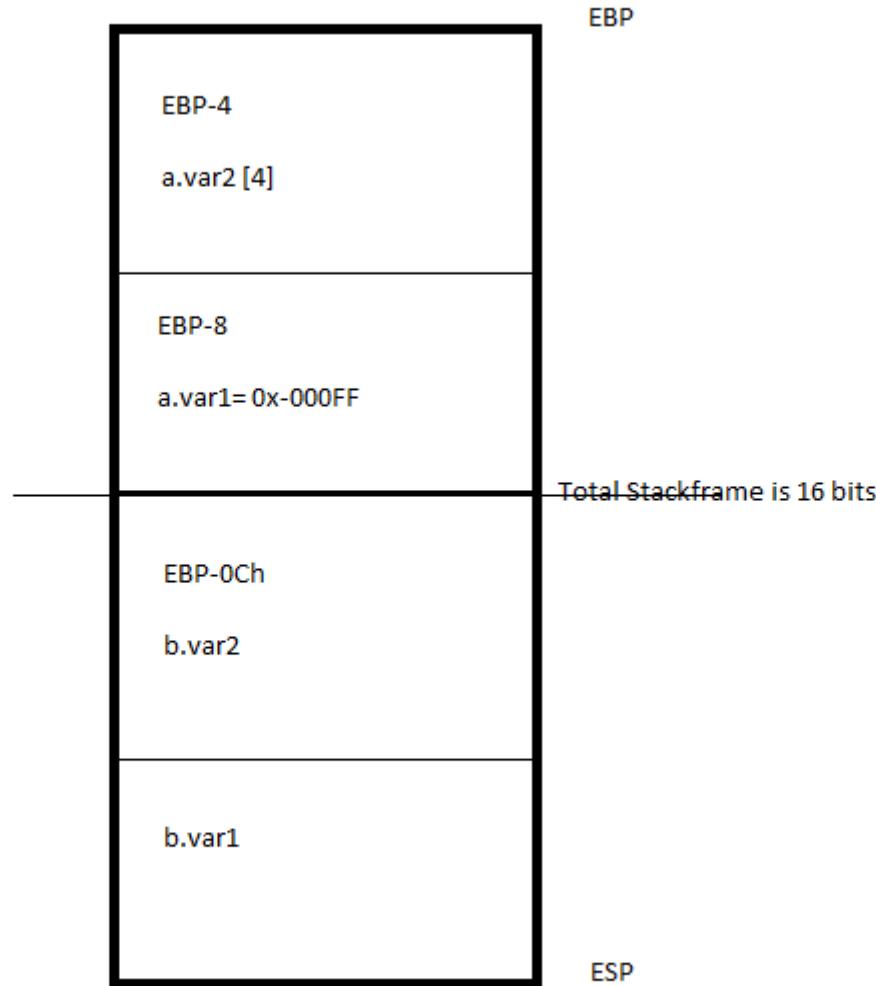
we have two of them

so hex 10 == dec 16. 16 bytes place is created.

- THEN, 01281016 mov dword ptr [ebp-8], 0FFh
-

DWORD == 4 bits so adding 0xFF to ebp-8

now we know the structure of our stackframe



So the stackframe is set for 16 bits and divided into 4 equal pieces per `mystruct_t` structure.

- third step : `memcpy(&b, &a, sizeof(mystruct_t));`

this code will take `&a`, copy into `&b` with the size of `mystruct_t` which is 8 bits.

Hence, the next step is to allocate place for it in the stackframe.

there comes the line `push 8`

- Esp goes down by 8.

Now we opened a space by 8, thich is the size of `mystruct_t`.

next, we will copy the values `&a` into the `&b`, where are the values of `&a`?

`ebp-8` and `ebp-10h`

each time we calculate these and add to `eax`, we then push the `eax` to the stackframe

```

0128101F  lea      eax, [ebp-8] ; b.var2  ebp-8
01281022  push    eax
01281023  lea      ecx, [ebp-10h] ; b.var1  ebp-16(10==16 in dec)
01281026  push    ecx

```

- now stackframe is ready for the operation, we can call the memcpy function. **01281027 call 01281042**
- after memcpy function is called, data swap is done, but we have extra space left in the stackframe. we need to destroy it with **0128102C add esp,0Ch** so removing all the 12 bits from it.

then regular tear down operations and exiting the program.

If we step into the memcpy call, we see that it also creating its own stackframe alongside our routine's stackframe.

Part 2

Linux Counterpart(AT&T Syntax)

IN AT&T Syntax EVERYTHING IS BACKWARDS!!! MNEMONIC SOURCE, DESTINATION.

REGISTERS GET % AND IMMEDIATES GET \$ IN FRONT OF THEM

mov %esp, %ebp ==> move esp to ebp, not the reverse.

add \$0x14, %esp ==> same, and immediate got \$ prefix

Also difference in m/r32 values:

Intel Syntax ==> [base+ index* scale + displacement]

Unix Syntax ==> displacement(base,index,scale)

Example :

call DWORD PTR[ebx+esi*4-0xe8] ==> Intel

call *-0xe8(%ebx,%esi,4) ==> Unix Syntax ==> * here is absolute indirect call syntax

Another example:

lea eax, [ebx-0xe8] ==> Intel

lea -x08(%ebx), %eax ==> Unix Syntax

In AT&T Syntax, for instructions which can operate on different sizes, **the mnemonic will have an indicator of the size**

movb operates on bytes mov/movw operates on word(2bytes) movl operatos on long (dword - 4bytes)

Intel does indicate size with things like `mov dword ptr[ebp-4]` but its just not in the actual mnemonic of the instruction.

In unix systems, we have GCC. It has many features but we mostly care about `-ggdb` command line argument to build debug symbols

gcc basic usage :

```
gcc -o <output filename> < input filename>
```

our case we will be using the following syntax:

```
gcc -ggdb -o<filename> <filename>.c
```

```
gcc -ggdb -o Example Example.c
```

objdump basic usage :

Where object file can be an intermediate file created during compilation but before linking, or a fully linked exec.

for our purposes, means any ELF file- the exec format standard for linux

the main usage for us is `objdump -d file`

can override the output syntax with `‐M intel`

so we can go back and forth between intel and at&t

```
gcc -ggdb -M intel -o example2 Example2.c
```

once we have the binary, now we can see the dump

```
objdump -d example2 | less
```

in order to compile program in 32 bit, after installing `gcc-multilib`, we use `-m32` flag

```
gcc -m32 Example2.c -o example_32
```

here is the At&t syntax assembly code of the same example2.c we had used previously

```
0804840b <sub>:
```

804840b:	55	push %ebp
804840c:	89 e5	mov %esp,%ebp
804840e:	8b 45 08	mov 0x8(%ebp),%eax
8048411:	8d 14 00	lea (%eax,%eax,1),%edx
8048414:	8b 45 0c	mov 0xc(%ebp),%eax
8048417:	01 d0	add %edx,%eax

8048419:	5d	pop %ebp
804841a:	c3	ret
 0804841b <main>:		
804841b:	8d 4c 24 04	lea 0x4(%esp),%ecx
804841f:	83 e4 f0	and \$0xffffffff0,%esp
8048422:	ff 71 fc	pushl -0x4(%ecx)
8048425:	55	push %ebp
8048426:	89 e5	mov %esp,%ebp
8048428:	53	push %ebx
8048429:	51	push %ecx
804842a:	83 ec 10	sub \$0x10,%esp
804842d:	89 cb	mov %ecx,%ebx
804842f:	8b 43 04	mov 0x4(%ebx),%eax
8048432:	83 c0 04	add \$0x4,%eax
8048435:	8b 00	mov (%eax),%eax
8048437:	83 ec 0c	sub \$0xc,%esp
804843a:	50	push %eax
804843b:	e8 b0 fe ff ff	call 80482f0 <atoi@plt>
8048440:	83 c4 10	add \$0x10,%esp
8048443:	89 45 f4	mov %eax,-0xc(%ebp)
8048446:	83 ec 08	sub \$0x8,%esp
8048449:	ff 75 f4	pushl -0xc(%ebp)
804844c:	ff 33	pushl (%ebx)
804844e:	e8 b8 ff ff ff	call 804840b <sub>
8048453:	83 c4 10	add \$0x10,%esp
8048456:	8d 65 f8	lea -0x8(%ebp),%esp
8048459:	59	pop %ecx
804845a:	5b	pop %ebx
804845b:	5d	pop %ebp
804845c:	8d 61 fc	lea -0x4(%ecx),%esp
804845f:	c3	

Hexdump & xxd

sometimes useful to look at a hexdump to see opcodes or raw file format. [hexdump -C](#) for canonical hex & ASCII view

xxd -make a hexdump or do the reverse

- use as quick hex editor
- xxd hello > hello.dump
- edit hello.dump
- xxd -r hello.dump>hello

So take a file, dump the hex, edit the hex of that dump, reverse compile that dump into binary again. cool and dirty!

GDB - The GNU Debugger

command line debugger

useful syntax for gdb : ==

`gdb <program name> -x <command file> ==> gdb Example-1 -x myCmds`

gdb can be pain in the neck since it is very barebones unless I specifically ask for something. Thus, using the `-x` flag, we will give a plain text file that encompasses all the required fields so that we don't bear the pain of keep typing commands all the time.

```
display/10i $eip
display/x $eax
display/x $ebx
display/x $ecx
display/x $edx
display/x $edi
display/x $esi
display/x $ebp
display/16xw $esp
break *main
run
```

basically this is the content of the file we will feed into gdb.

so let's say we named the file `commands`, `gdb filename -x commands`

the output will be like

```
Reading symbols from hello...done.
Breakpoint 1 at 0x40052a: file Hello.c, line 3.
```

that works! cool

some useful commands in gdb :

`run ==>` runs the program `r <argv>` ==> run the program with the given arguments. `start ==>` set a breakpoint on main and run ``

given the commands up there are 32 bit commands, they won't work with 64 bit programs. Thus, a recompile (using `-m32` flag) the file and running the gnu debugger does the trick!

```
Starting program: /home/myy/Desktop/asm/Example2/example_32

Breakpoint 1, 0x0804842a in main ()
1: x/10i $eip
0x08048425 <+10>: push    %ebp
0x08048426 <+11>: mov     %esp,%ebp
0x08048428 <+13>: push    %ebx
```

```

0x08048429 <+14>:    push   %ecx
=> 0x0804842a <+15>:    sub    $0x10,%esp
  0x0804842d <+18>:    mov    %ecx,%ebx
  0x0804842f <+20>:    mov    0x4(%ebx),%eax
  0x08048432 <+23>:    add    $0x4,%eax
  0x08048435 <+26>:    mov    (%eax),%eax
  0x08048437 <+28>:    sub    $0xc,%esp
  0x0804843a <+31>:    push   %eax
  0x0804843b <+32>:    call   0x80482f0 <atoi@plt>
  0x08048440 <+37>:    add    $0x10,%esp
  0x08048443 <+40>:    mov    %eax,-0xc(%ebp)
  0x08048446 <+43>:    sub    $0x8,%esp
  0x08048449 <+46>:    pushl  -0xc(%ebp)
  0x0804844c <+49>:    pushl  (%ebx)
  0x0804844e <+51>:    call   0x804840b <sub>
  0x08048453 <+56>:    add    $0x10,%esp
  0x08048456 <+59>:    lea    -0x8(%ebp),%esp
  0x08048459 <+62>:    pop    %ecx
  0x0804845a <+63>:    pop    %ebx
  0x0804845b <+64>:    pop    %ebp
  0x0804845c <+65>:    lea    -0x4(%ecx),%esp
  0x0804845f <+68>:    ret

```

```

2: /x $eax = 0xf7fb8dbc
3: /x $ebx = 0x0
4: /x $ecx = 0xfffffcfb0
5: /x $edx = 0xfffffcfd4
6: /x $edi = 0xf7fb7000
7: /x $esi = 0xf7fb7000
8: /x $ebp = 0xfffffcf98
9: x/16xw $esp
0xfffffcf90: 0xfffffcfb0 0x00000000 0x00000000 0xf7e1f647
0xfffffcfa0: 0xf7fb7000 0xf7fb7000 0x00000000 0xf7e1f647
0xfffffcfb0: 0x00000001 0xfffffd044 0xfffffd04c 0x00000000
0xfffffcfc0: 0x00000000 0x00000000 0xf7fb7000 0xf7fffdc04

```

another command is **x** to examine the memory

Always assumes the given value is a memory address and it **dereferences** it to look at the value at that memory address.

print/FMT EXP print the value of an expression

example : **x/x \$ebp** 0xfffffcf98: 0x00000000

example : **print/x %ebp** \$1 = 0xfffffcf98

break or **b==>** sets breakpoint. **b *<address>** or **b main**

to know about current breakpoints, **info b**

```
(gdb) info b
Num      Type            Disp Enb Address     What
1        breakpoint      keep y  0x0804842a <main+15>
breakpoint already hit 1 time
```

delete <breakpoint number>

sometimes we need to go back in the stack here are the commands to use then:

reverse-step, rs ==> step back until reaching the beginning of a previous source line

reverse-stepi==> step 1 instruction

reverse-continue ==> continue program being debugged but run it in reverse

reverse-finish==> execute backwards until just before the selected stack frame is called

reverse-next, rn ==> step backwardsm proceeding through subroutine calls

reverse-nexti==> step backwards one instruction

set exec-direction==> set direction of exzecution.

stepi, si ==> steps one asm instruction at at time. **will always step into subroutines**

step , s ==> steps one source line at at time. If no source is available then acts like stepi

until, u ==> steps until the next source line, not stepping into subroutines.

nexti, n ==> next instruction, **step over**

gdb -x command example_32

inside, **break *main** then **run**

each time we **stepi**, it moves one instruction ahead and pops the one on the top.

one of the breakpoints in the file is x/10 i , that gives first 10 instructions starting from **main** function.

```
Breakpoint 2, 0x0804841b in main ()
1: x/20i $eip
=> 0x804841b <main>:    lea     0x4(%esp),%ecx
  0x804841f <main+4>:   and    $0xffffffff0,%esp
  0x8048422 <main+7>:  pushl   -0x4(%ecx)
  0x8048425 <main+10>: push    %ebp
  0x8048426 <main+11>: mov     %esp,%ebp
  0x8048428 <main+13>: push    %ebx
  0x8048429 <main+14>: push    %ecx
  0x804842a <main+15>: sub     $0x10,%esp
  0x804842d <main+18>: mov     %ecx,%ebx
  0x804842f <main+20>: mov     0x4(%ebx),%eax
  0x8048432 <main+23>: add     $0x4,%eax
```

```

0x8048435 <main+26>: mov    (%eax),%eax
0x8048437 <main+28>: sub    $0xc,%esp
0x804843a <main+31>: push   %eax
0x804843b <main+32>: call   0x80482f0 <atoi@plt>
0x8048440 <main+37>: add    $0x10,%esp
0x8048443 <main+40>: mov    %eax,-0xc(%ebp)
0x8048446 <main+43>: sub    $0x8,%esp

2: /x $eax = 0xf7fb7dbc
3: /x $ebx = 0x0
4: /x $ecx = 0xfffffcfb0
5: /x $edx = 0xfffffcfd4
6: /x $edi = 0xf7fb6000
7: /x $esi = 0xf7fb6000
8: /x $ebp = 0x0
9: x/16xw $esp
0xfffffcfac: 0xf7e1e647 0x00000001 0xfffffd044 0xfffffd04c
0xfffffcfbc: 0x00000000 0x00000000 0x00000000 0xf7fb6000
0xfffffcfcc: 0xf7ffdc04 0xf7ffd000 0x00000000 0xf7fb6000
0xfffffcfdc: 0xf7fb6000 0x00000000 0xeb8040f 0xc2eb6a1f

```

in gbd, stack is configured a little bit differently.

when main called, at the top of the stack there is **eip** of what called **main()** to run. So address of xyz, which called the main to run, is saved to eip.

then pushed the ebp and esp.

when we **stepi** and reach to the next instruction after **main()**, which is **0x804841f <main+4>: and \$0xffffffff0,%esp**,

our \$esp is **0xfffffcfac**, ebp is 0. when we execute the next instruction, which is the **and** instruction above,

it will be **and'd** with 0xffffffff0

in this line **0x804842f <main+20>: mov 0x4(%ebx),%eax**, this is moving our **argc** (which we did not give when running actually)

to give an argument, (in this case our argument will be the integer 256) , basically **run 256** this will restart the program

now I want to go 4 steps further from beginning, **stepi 4** voila

so we are here **0x804842f <main+20>: mov 0x4(%ebx),%eax** (ebp + 0x4) into the eax register.

What is eax +4 ? it is actually the saved eip 😐

I'm actually confused here, compiler should have given us **mov 0xc(%ebp),%eax** 😐/

But we know the structure of the stack, how linux kernel puts them

```
argv ==> parameters
argc ==> parameters
saved eip
saved ebp

aligned space
```

btw: how to read this `mov 0x4(%ebx),%eax`? add 0x4 and % ebx, go to that memory address, take whatever the value is and stick it into eax.

so here `0x8048432 <main+23>: add $0x4,%eax`, we are going into the array of argvs and passing the first one.

after executing this, our eax is `0xfffffd048`

lets dig in what is happening there. to dig and see the string in that memory space is done with the syntax `x/s *address` in our case `x/s *0xfffffd044`

you wanna see a magic trick?

we had passed an argv with `run 256` right? now once we stepped into the argv array up there, we now pulled the string version of it with `x/s *0xfffffd028` this address is the %eax currently.

this is the outcome:

```
(gdb) x/s *0xfffffd028
0xfffffd247: "256" ;this was the argv we had passed!
```

argv is a pointer to the pointer. So argv itself points to multiple pointers, and each of these pointers points to a ASCII string values somewhere in the memory.

* before the address is dereferenced. also, If I do not want to take it as a string but I want the character, I use `x/c *<address>` that gives the first value

```
(gdb) x/c *0xfffffd028
0xfffffd247: 50 '2' ; 2 of 256
```

since the next instruction is a pointer to a pointer, we need to take this into the eax. thus `0x8048435 <main+26>: mov (%eax),%eax`

move the memory pointing to the eax, dereference it and write it into eax. Now the value of %eax is not an address but a value that resides in this mem address.

so eax is pointing at `0xfffffd028`, whatever in this memory address should be taken as value and to be fed into eax.

this value is, of course 256.

now eax holds the value of 256 not the pointer.

next step (stepi) is `0x8048435 <main+26>: mov (%eax),%eax` moving what is into the eax into eax.

if we want to see what is in the eax now, we should **not dereference it since it is not a pointer anymore**, so instead, take the address of eax and x/s it.

```
(gdb) x/s 0xfffffd247 ; NOTICE THERE IS NO * HERE.
0xfffffd247: "256"
```

to see char by char, we move bit by bit using x/c.

```
(gdb) x/c 0xfffffd247
0xfffffd247: 50 '2'
(gdb) x/c 0xfffffd248
0xfffffd248: 53 '5'
(gdb) x/c 0xfffffd249
0xfffffd249: 54 '6'
```

right before the `atoi` function is called, the next address, `0x8048440` is saved to eip so that after the function, it will follow up from there.

we do not want to step into `atoi` function's internal mechanism so we do not stepi. instead we used `nexti` its like step over.

what is the outcome of eax after `atoi` is implemented??

it is **2**: `/x $eax = 0x100` 0x100 corresponds to 256 in decimal 😊 so we converted the string into integer directly and successfully!!!

if I am in a subroutine and want to leave, the command is `finish`

`ebp+4, ebp + 8` these are taking argv and using them! ALWAYS!

Inline Assembly , gdb & visual studio

inline assembly is a way to include assembly directly into C/C++ file. However, syntax will differ between compilers and assemblers. There are times when you actually have to code asm in order to do something from a C/C++ file.

for example, in OS or driver design inline asm is used. because there are many special aspects of hardware which can only be accessed with a special instruction. For

in crypto, you might want to access to the `rol/ror` rotate right/ left instructions which don't have corresponding C syntax.

different assemblers require different syntax.

Visual Studio Inline Assembly:

we use intel syntax

`__asm{code here};` ==> instructions are separated by \n

two underscores at the beginning, actually does not require ; at the beginning but it helps of auto-indentation

```
__asm{
    mov eax, [esp+0x4]
    cmp eax, 0xdeadbeef
    je mylabel
    xor eax, eax
mylabel:   mov bl, al
};
```

syntax using C variables is the same. just put the variable in place of a register name for instance.

the assembler will substitute the correct address for the variable.

```
// value into C variable from register
__asm{mov myVar eax};      where myVar is a C variable.

// value into register from C variable

__asm{mov myVar, eax};
```

so address of `myVar` will be inferred by the assembler!

we cannot use like `__asm{mov var1, var2};` since that would be a mem-mem transfer which is not allowed!!

so reg-reg, reg-mem, mem-reg operations are allowed just like regular assembly

GCC Inline Assembly:

uses AT&T Syntax, obviously.

syntax is :

`asm("instructions separated by \n");`

NEEDS SEMICOLON AFTER THE PARENTHESES!! THIS ONE IS NOT OPTIONAL!

we do not need `__` but we need \n after every line!

```
asm(
    "mov 0x4(%esp), %eax\n"
    "cmp $0xdeedbeef, %eax\n"
    "je mylabel\n"
    "xpr %eax,%eax\n"
    "myLabel: movw %bx, %ax"
);
```

syntax using C variables are little bit nasty in GCC format!

```
asm(assembler template
:output operands /optional/
:input operands /optional/
:list of clobbered registers /optional/
);
```

for example:

```
int myVar;
//value into C variable from register

asm("movl %eax %%eax, %0" : "=r"(myVar));

//value into register from C variable

asm("movl %0, %%eax"::"r"(myVar));
```

Seriously, go to hell Richard Stallman! what is this douchbag syntax!

`%%` double percent sign means 'this is a regular eax, do not treat it like something special!' whereas `%` single percent sign means, take the `number`th value from output/input operand and put it into the right. `%0` means take the 0th argument, write it onto myVar

`=` means to, `:` means from. we can infer the direction of the operation from these

`_emit & .byte`

once you learn about opcodes later on, you can even specify exactly the instructions you want to use by using `_emit` or `.byte` keywords to place specific bytes into the code.

those bytes can then be interpreted as instructions or data.

this sometimes useful if you cannot figure out the inline asm syntax for the instruction you want to use, but you know its opcodes(either from seeing them elsewhere or reading the manual)

examples

- `__asm{_emit0x55}` is `__asm{push ebp}`
- `__asm{_emit0x89}; __asm{emit0xE5}` is `__asm{mov ebp,esp}`
- `asm{".byte 0x55"};` is `asm("push %ebp")`
- `asm{".byte0x89; .byte0xE5"};` is `asm("mov %esp,%ebp");`