

ALGORITHMS TASK

SPRING 2023

TASK NUMBER 5 K-TH ELEMENT OF TWO SORTED ARRAYS

NAMES

20210887

20210867

20210900

20210926

20210859

20210836

1. محي الدين ماهر

2. محمود عبد الدايم أبو المجد

3. مروة عمر محمد محمود

4. مريم محمود سيد محمود

5. محمود حسام الدين محمود

6. محمد مصطفى محمد عبد ربه



MERGE SORT APPROACH

A. PSEUDOCODE

```
mergeSort( arr1 , size1 ,arr2 ,size2 , finalArr ,i , j, k) :  
  
    if i<size1 && j<size2  
        if arr1[i]<arr2[j]  
            finalArr[k] = arr1[i]  
            i++  
  
        else  
            finalArr[k] = arr2[j]  
            j++  
  
        k++  
        mergeSort( arr1 , size1 ,arr2 ,size2 , finalArr ,i , j, k)  
  
    else  
        while i<size1  
            finalArr[k] = arr1[i]  
            i++  
            k++  
  
        while j<size2  
            finalArr[k] = arr2[ij]  
            j++  
            k++  
  
    end
```

B. SOURCE CODE

- [GITHUB LINK](#)

C. CODE ANALYSIS

The function takes in two arrays `arr1`, `arr2`, their sizes `s1`, `s2`, and an array `finalArray` to hold the merged and sorted result. It also takes pointers `i`, `j`, and `k` to keep track of the current index being compared in both arrays and merged array respectively. The function recursively merges the two arrays until one of the arrays has been fully traversed using the following: If `arr1[i]` is less than `arr2[j]`, `arr1[i]` is added to the final array, `i` and `k` are incremented. If `arr1[i]` is greater than or equal to `arr2[j]`, `arr2[j]` is added to `finalArray`, `j` and `k` are incremented. In the main function, two sorted arrays `ar1` and `ar2`, and an integer `k` are initialized. The `sizeof` operator is used to find the lengths of the arrays `ar1` and `ar2` respectively, and the sum of these lengths is stored in the variable `tot`. An integer array `finalar` of size `s1+s2` is initialized to hold the concatenated arrays and the `mergeSort` function is called with pointers `i`, `j`, and `k` initialised to 0. The `k`-th element of the sorted concatenated array, `finalar`, is printed to the console. Note that `k-1` is used since `finalar` is

D. TIME COMPLEXITY:

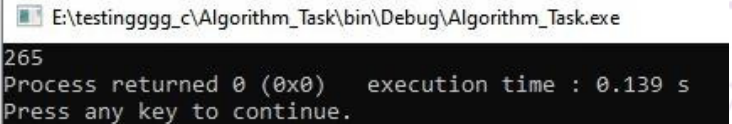
$O(n)$

E. OUTPUT SCREENSHOT

```
int main()
{
    int ar1[]={100,112,256,349,770};
    int ar2[]={72,86,119,265,445,892};
    int s1 = sizeof(ar1)/sizeof(ar1[0]);
    int s2 = sizeof(ar2)/sizeof(ar2[0]);
    int finalar[s1+s2];

    int k = 7;
    mergeSort(ar1,s1,ar2,s2,finalar,0,0,0);

    cout<<finalar[k-1]; // because the array is 0 indexed
    return 0;
}
```



NON-RECURSIVE APPROACH

A. PSEUDOCODE

```
nonRecursive( arr1 , size1 ,arr2 ,size2 , finalArr) :  
  
int i=0,j=0,k=0  
while i<size1 && j<size2  
    if arr1[i]<arr2[j]  
        finalArr[k] = arr1[i]  
        i++  
        k++  
  
    else  
        finalArr[k] = arr2[j]  
        j++  
        k++  
  
while i<size1  
    finalArr[k] = arr1[i]  
    i++  
    k++  
  
while j<size2  
    finalArr[k] = arr2[j]  
    j++  
    k++  
  
end
```

B. SOURCE CODE

- [GITHUB LINK](#)

C. CODE ANALYSIS

The `nonRecusize` function takes in two arrays `arr1`, `arr2`, their respective lengths `s1`, `s2`, and an integer array `finalArr` to hold the merged and sorted result. It then uses three pointers `i`, `j`, and `k` to keep track of the current index being compared in both arrays and the merged array respectively. The function merges the two arrays in a while loop using the following as a guideline: If `arr1[i]` is less than `arr2[j]`, `arr1[i]` is added to `finalArr`, `i` is incremented, and `k` is incremented. If `arr1[i]` is greater than or equal to `arr2[j]`, `arr2[j]` is added to `finalArr`, `j` is incremented, and `k` is incremented. Once one of the arrays has been fully traversed, the remaining elements in the other array are copied to `finalArr`. In the main function, two sorted arrays `a` and `b`, and an integer `k` are initialized. The `sizeof` operator is used to find the lengths of the arrays `a` and `b` respectively, and the sum of these lengths is stored in the variable `s1+s2`. An integer array `f` of size `s1+s2` is initialized to hold the concatenated arrays and the `nonRecusize` function is called with pointers `i`, `j`, and `k` initialised to 0. The `k`-th element of the sorted concatenated array, `f`, is printed to the console. Note that `k-1` is used since `f` is zero-indexed

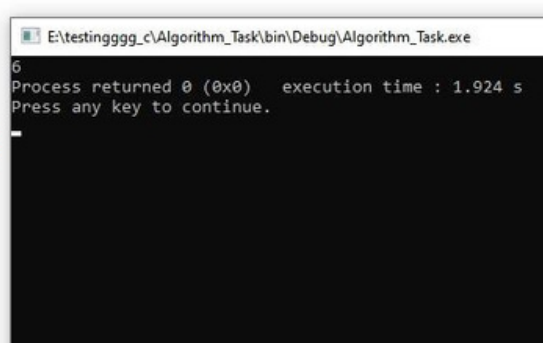
D. TIME COMPLEXITY:

$O(n)$

E.OUTPUT SCREENSHOT

```
int main()
{
    fast();
    int a[] = {2,3,6,7,9};
    int b[] = {1,4,8,10};
    int s1 = sizeof(a)/sizeof(a[0]);
    int s2 = sizeof(b)/sizeof(b[0]);
    int f[s1+s2];
    nonRecusize(a,s1,b,s2,f);

    int k=5;
    cout<<f[k-1];
    return 0;
}
```



```
E:\testingggg_c\Algorithm_Task\bin\Debug\Algorithm_Task.exe
6
Process returned 0 (0x0)   execution time : 1.924 s
Press any key to continue.
```

DIVIDE AND CONQUER APPROACH

A. PSEUDOCODE

```
Algorithm kth_element(arr1,arr2,last1,last2, k)

    if arr1 = last1
        return arr2[k]
    if arr2 = last2
        return arr1[k]

    mid1 ← (last1 - arr1)/2
    mid2 ← (last2 - arr2)/2

    if mid1 + mid2 < k
        if arr1[mid1] > arr2[mid2]
            return kth_element(arr1, arr2+mid2+1, last1, last2, k-mid2-1)
        else
            return kth_element(arr1+mid1+1, arr2, last1, last2, k-mid1-1)
    else
        if arr1[mid1] > arr2[mid2]
            return kth_element(arr1, arr2, arr1 + mid1, last2, k)
        else
            return kth_element(arr1, arr2, last1 ,arr2 + mid2, k)
    end
```

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B. SOURCE CODE

- [GITHUB LINK](#)

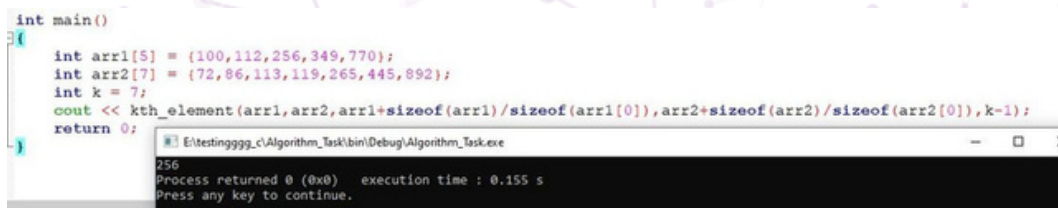
C. CODE ANALYSIS

The `kth_element` function takes 5 parameters: two sorted arrays (`arr1` and `arr2`), two pointers to the last element of each array (`last1` and `last2`), and an `int k` representing the index of the desired element. If the length of one of the arrays is 0, then the function returns the `k`-th element of the other array. Two variables, `mid1` and `mid2`, are set to the middle index of each array. If the sum of `mid1` and `mid2` is less than `k`, the function recursively calls itself with either `arr2+mid2+1` (if `arr1[mid1]` is greater than `arr2[mid2]`) or `arr1+ mid1+1` (if `arr1[mid1]` is less than or equal to `arr2[mid2]`), and an updated value of `k`. If `k` is less than the sum of `mid1` and `mid2`, the function recursively calls itself with either `arr1+mid1` and `last2` (if `arr1[mid1]` is greater than `arr2[mid2]`) or `last1` and `arr2+mid2` (if `arr1[mid1]` is less than or equal to `arr2[mid2]`). The main function initializes two sorted arrays (`arr1` and `arr2`) and an integer `k`. It then calls the `kth_element` function with the two arrays, their last elements, and `k-1` (since arrays are 0-indexed). Finally, the `kth_element` function returns the `k`th element of the combined arrays.

D. TIME COMPLEXITY:

$O(\log m + \log n)$

E.OUTPUT SCREENSHOT



The screenshot shows a C++ program in a code editor and its execution output in a terminal window. The code defines two sorted arrays, `arr1` and `arr2`, and a variable `k`. It calls the `kth_element` function to find the `k`-th element of the combined arrays. The output shows the value 256, which is the 7th element (index 6) of the combined sorted array.

```
int main()
{
    int arr1[5] = {100,112,256,349,770};
    int arr2[7] = {72,86,113,119,265,445,892};
    int k = 7;
    cout << kth_element(arr1,arr2,arr1+sizeof(arr1)/sizeof(arr1[0]),arr2+sizeof(arr2)/sizeof(arr2[0]),k-1);
    return 0;
}
```

Output:

```
256
Process returned 0 (0x0)   execution time : 0.155 s
Press any key to continue.
```

BUBBLE SORT APPROACH

A. PSEUDOCODE

```
bubbleSort(arr1 , n )  
flag := true  
for i := 0 to n-2 do  
    for j := 0 to n-i-2 do  
        if a[j] > a[j+1] then  
            swap(a[j+1], a[j])  
            flag := false  
        end if  
    end for  
    if flag = true then  
        break  
    end if  
end for  
End
```

B. SOURCE CODE

- [GITHUB LINK](#)

C. CODE ANALYSIS

The bubbleSort function takes in an array a and its length n, and sorts it using bubble sort. It sets flag to true initially, which acts as a checker for whether the array is sorted. The outer loop goes from 0 to $i < n-1$ and the inner loop, from $j=0$ to $j < n-i-1$. If $a[j]$ is greater than $a[j+1]$, the function swaps the two elements and sets flag to false. If flag is still true after the inner loop finishes, it means that the array is already sorted, so the function breaks out of the outer loop. In the main function, two sorted arrays a and b, and an integer k are initialized. The sizeof operator is used to find the lengths of the arrays a and b, and the sum of these lengths is stored in the variable tot. An integer array finalarray of size tot is initialized, and the first s1 elements of finalarray are filled with the elements of a. The remaining elements of finalarray are filled with the elements of b, starting from finalarray[s1]. The bubbleSort function is called with finalarray and tot, to ensure that the concatenated array is sorted. The k-th element of the sorted concatenated array, finalarray, is printed to the console. Note that k-1 is used since finalarray is 0-indexed.

D. TIME COMPLEXITY:

$O(n^2)$

E. OUTPUT SCREENSHOT

```
int main()
{
    int a[]={2,3,6,7,9};
    int b[]={1,4,8,10};
    int k=5;
    int s1 = sizeof(a)/sizeof(a[0]);
    int s2 = sizeof(b)/sizeof(b[0]);
    int tot = s1 + s2; //the size of the final array
    int finalarray[tot];
    for(int i=0;i<s1;i++){
        finalarray[i]=a[i];
    }
    for(int i=s1;i<tot;i++){
        finalarray[i]=b[i-s1];
    }
    bubbleSort(finalarray,tot);
    cout<<finalarray[k-1] ; //because the array is 0 indexed
    return 0;
}
```

```
E:\testingggg_c\Algorithm_Task\bin\Debug\Algorithm_Task.exe
6
Process returned 0 (0x0) execution time : 1.620 s
Press any key to continue.
```

INSERTION SORT APPROACH

A. PSEUDOCODE

```
insertionSort(ar , n)
  for i := 1 to n-1 do
    key := ar[i]
    j := i-1
    while j ≥ 0 and ar[j] > key do
      ar[j+1] := ar[j]
      j := j-1
    end while
    ar[j+1] := key
  end for
End
```

B. SOURCE CODE

- [GITHUB LINK](#)

C. CODE ANALYSIS

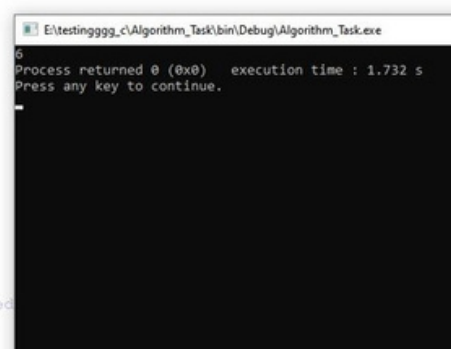
The `insertionSort` function takes an array `ar` of length `n`. It iterates through each element of the array, starting from index 1. For each element, it compares it with the previous elements to find the correct position to insert it into the sorted subarray that precedes it. The while loop moves each element that is greater than key one position to the right until the correct position for key is found. Once the correct position is found, key is inserted into the array. In the main function, two sorted arrays `a` and `b` are initialized, along with an integer `k` to find the `k`-th element of the sorted concatenated array. The `sizeof` operator is used to find the lengths of the arrays `a` and `b` respectively, and the sum of these lengths is stored in `tot`. An integer array `finalarray` of size `tot` is declared to hold the concatenated arrays. The elements of `a` are copied into `finalarray` using a for loop, and the elements of `b` are copied into `finalarray` using another for loop that starts from index `s1`, where `s1` is the length of `a`. The `insertionSort` function is called with `finalarray` and `tot` as arguments to sort the concatenated array. The `k`-th element of the sorted concatenated array, `finalarray`, is printed to the console. Note that `k-1` is used since `finalarray` is zero-indexed.

D. TIME COMPLEXITY:

$O(n^2)$

E. OUTPUT SCREENSHOT

```
int main()
{
    int a[]={2,3,6,7,9};
    int b[]={1,4,8,10};
    int k=5;
    int s1 = sizeof(a)/sizeof(a[0]);
    int s2 = sizeof(b)/sizeof(b[0]);
    int tot = s1 + s2; // the tot size of the final array
    int finalarray[tot];
    for(int i=0;i<s1;i++){
        finalarray[i]=a[i];
    }
    for(int i=s1;i<tot;i++){
        finalarray[i]=b[i-s1];
    }
    insertionSort(finalarray,tot);
    cout<<finalarray[k-1] ; // because the array is 0 indexed
    return 0;
}
```



Comparing between approaches from the complexity

Merge Sort	Non-Recursive	Divide and Conquer	Bubble Sort	Insertion Sort
$O(n)$	$O(n)$	$O(\log m + \log n)$	$O(N^2)$	$O(N^2)$