

SageMath Toolbox for Rank Error-Correcting Codes

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Introduction

- Quantum computing breaks
 - Integer factorization problem \rightarrow RSA
 - Discrete logarithm problem \rightarrow ElGamal, ECC
- (General) Syndrome Decoding Problem (NP-hard)
- Rank Syndrome Decoding Problem!



Background

Error-Correcting Codes

An $[n, k]$ -code $C \subseteq \mathbb{F}_{q^m}$ has

- Generator matrix $G \in \mathbb{F}_{q^m}^{k \times n}$
- Parity check matrix $H \in \mathbb{F}_{q^m}^{(n-k) \times n}$
- Messages $m \in \mathbb{F}_{q^m}^k$
- Codewords $c \in \mathbb{F}_{q^m}^n$

Properties: $GH^T = 0$, $Hc^T = 0$

Encoding: $c = mG$

Syndrome decoding: $s(y) = Hy^T = H(c + e)^T = He^T$

Background

Rank

$c = (z_2, z_2 + 1)$ in $[2, 1]$ -code $C \subseteq \mathbb{F}_{2^2}$ in matrix form: $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}$

$$\rightarrow \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \rightarrow |c|_R = 2$$

Support

$\text{Supp}(c) = \langle z_2, z_2 + 1 \rangle_2$ is the subspace with basis $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$

Syndrome Decoding Problems

Rank Syndrome Decoding (RSD) Problem

$$\begin{cases} He^T = s^T \\ |e|_R = w \end{cases}$$

Knowing s and $E = \text{Supp}(e)$ gives e !

bit security	128	192	256
RSA	3 072	7 680	15 360
Goppa-McEliece	$2 \cdot 10^6$	$4 \cdot 10^6$	$6 \cdot 10^6$
Gabidulin (DRANKULA)	62 000	118 160	216 000
QC LRPC (LOCKER)	5 893	8 383	9 523

IF
GSD
RSD

Toolbox

- In SageMath
- Basics:
 - Support
 - Finding codewords
 - Rewriting elements in certain basis
- Attack: GRS algorithm
- Family of codes: LRPC codes



GRS Algorithm

- 1 Pick $F \subseteq \mathbb{F}_{q^m}$ of dimension r
- 2 Rewrite $He^T = s$, where e is rewritten in F via $e_i = \sum_{j=0}^{r-1} \lambda_{ij} F_j$, nr unknowns
- 3 Rewrite system, where H , F and s are rewritten in basis, gives $(n - k)m$ equations
- 4 Find solution, check rank

Low Rank Parity Check (LRPC) Codes

- Construction: pick d random vectors for F , find H
- Encoding: generator matrix
- Decoding:
 - ① Find the support of s
 - ② Recover $E = S_1 \cap \dots \cap S_d$, where $S_i = F_i^{-1} S$
 - ③ Rewrite $He^T = s$ in terms of $P = \langle E, F \rangle$
 - ④ Solve system, then $c = y - e$ and $m = \frac{c}{G}$



Future Research

- More attacks: adaptation of GRS attack, polynomial annihilator attack
- More family of codes: simple codes, cyclic codes, QC-LRPC, DC-LRPC
- Application of codes to cryptography: GPT and LRPC cryptosystem, authentication