

Course Code: EE 2003	Course Name: Computer Organization and Assembly Language
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Student's Roll No:	Section:

Instructions:

- Except for your Roll No and Section, DO NOT SOLVE anything on this paper.
- Return the question paper.
- Read each question completely before answering it. There are **3 questions on 2 pages (1 Sheet)**.
- In case of any ambiguity, you may make an assumption. But your assumption should not contradict any statement in the question paper.
- All the answers must be solved according to the SEQUENCE given in the question paper, otherwise, points will be deducted.
- This paper is subjective.
- Where asked for values, only provide the **hex-decimal** values.
- Problems needing iterations should be coded using iterative instructions. No points will be awarded otherwise.

Time Allowed: 60 minutes.

Maximum Points: 32 points

Q. No. 1 Briefly answer each of the following:

[8 x 2 = 16 points]

- (i) Explain the difference between the direct-offset and indexed operands with examples.

4.1.8 Direct-Offset Operands

You can add a displacement to the name of a variable, creating a direct-offset operand. This lets you access memory locations that may not have explicit labels. Let's begin with an array of bytes named `arrayB`:

```
arrayB  BYTE  10h,20h,30h,40h,50h
```

If we use `MOV` with `arrayB` as the source operand, we automatically move the first byte in the array:

```
mov al,arrayB           ; AL = 10h
```

We can access the second byte in the array by adding 1 to the offset of arrayB:

```
mov al,[arrayB+1]           ; AL = 20h
```

The third byte is accessed by adding 2:

```
mov al,[arrayB+2] ; AL = 30h
```

4.4.3 Indexed Operands

An *indexed operand* adds a constant to a register to generate an effective address. Any of the 32-bit general-purpose registers may be used as index registers. There are different notational forms permitted by MASM (the brackets are part of the notation):

```
constant[reg]
[constant + reg]
```

The first notational form combines the name of a variable with a register. The variable name is translated by the assembler into a constant that represents the variable's offset. Here are examples that show both notational forms:

arrayB[esi]	[arrayB + esi]
arrayD[ebx]	[arrayD + ebx]

Indexed operands are ideally suited to array processing. The index register should be initialized to zero before accessing the first array element:

```
.data
arrayB BYTE 10h,20h,30h
.code
mov esi,0
mov al,[arrayB + esi]          ; AL = 10h
```

- (ii) In which case, we use MOVSB instruction? Explain one example of MOVSB in which a smaller value is moved into a larger destination.

MOVSB Instruction

The MOVSB instruction (move with sign-extend) copies the contents of a source operand into a destination operand and sign-extends the value to 16 or 32 bits. This instruction is only used with signed integers. There are three variants:

```
MOVSB reg32,reg/mem8
MOVSB reg32,reg/mem16
MOVSB reg16,reg/mem8
```

```
.data
byteVal BYTE 10001111b
.code
movsb ax,byteVal          ; AX = 1111111110001111b
```

- (iii) Discuss the Signed and Unsigned integers with reference to the hardware viewpoint.

Signed/Unsigned Integers: Hardware Viewpoint



- All CPU instructions operate exactly the same on signed and unsigned integers
- The CPU cannot distinguish between signed and unsigned integers
- YOU, the programmer, are solely responsible for using the correct data type with each instruction

(iv) Discuss the term label as an identifier and directive with examples.

Label

A *label* is an identifier that acts as a place marker for instructions and data. A label placed just before an instruction implies the instruction's address. Similarly, a label placed just before a variable implies the variable's address.

Data Labels A data label identifies the location of a variable, providing a convenient way to reference the variable in code. The following, for example, defines a variable named *count*:

```
count  DWORD 100
```

The assembler assigns a numeric address to each label. It is possible to define multiple data items following a label. In the following example, *array* defines the location of the first number (1024). The other numbers following in memory immediately afterward:

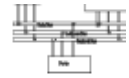
```
array  DWORD 1024, 2048  
        DWORD 4096, 8192
```

Code Labels A label in the code area of a program (where instructions are located) must end with a colon (:) character. Code labels are used as targets of jumping and looping instructions. For example, the following *JMP* (jump) instruction transfers control to the location marked by the label named *target*, creating a loop:

```
target:  
    mov    ax,bx  
    ...  
    jmp    target
```

(v) The *LOOP* instruction creates a counting loop. What happens when a loop instruction is encountered?

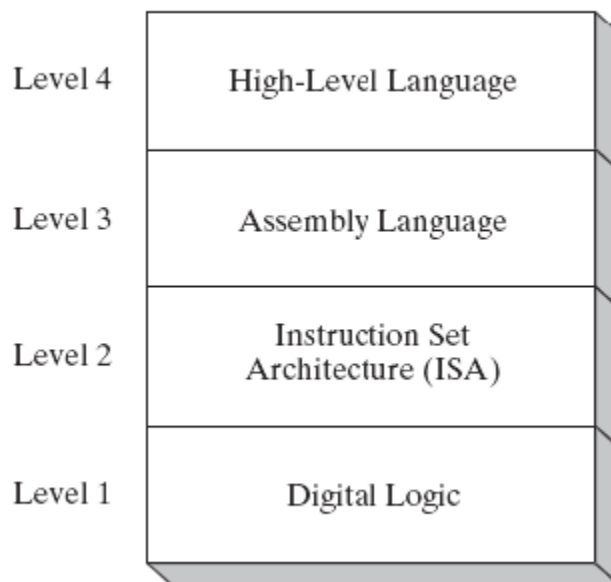
LOOP Instruction



- The `LOOP` instruction creates a counting loop
- Syntax: `LOOP target`
- Logic:
 - $ECX \leftarrow ECX - 1$
 - if `ECX != 0`, jump to *target*

(vi) Why the concept of virtual machines is important in computer organization? At which level does assembly language appear at the virtual machine level? Consider L1 as the lowest level.

- Virtual Machine Concept:
 - An effective way to explain how a computer's hardware and software are related is called the *virtual machine concept*.



(vii) How do we load and store the status flag bits? Give example instructions.

4.1.6 LAHF and SAHF Instructions

The LAHF (load status flags into AH) instruction copies the low byte of the EFLAGS register into AH. The following flags are copied: Sign, Zero, Auxiliary Carry, Parity, and Carry. Using this instruction, you can easily save a copy of the flags in a variable for safekeeping:

```
.data
saveflags BYTE ?
.code
lahf                ; load flags into AH
mov saveflags,ah    ; save them in a variable
```

The SAHF (store AH into status flags) instruction copies AH into the low byte of the EFLAGS register. For example, you can retrieve the values of flags saved earlier in a variable:

```
mov ah,saveflags    ; load saved flags into AH
sahf                ; copy into Flags register
```

- (viii) Compare the following x86 processors' mode of operations:
 - (a) Real address mode
 - (b) Protected address mode

❖ Real-Address mode (original mode provided by 8086)

- ❖ Only 1 MB of memory can be addressed, from 0 to FFFFF (hex)
 - ❖ Programs can access any part of main memory
 - ❖ MS-DOS runs in real-address mode
- ❖ Implements the programming environment of the Intel 8086 processor
- ❖ This mode is available in Windows 98, and can be used to run an MS-DOS program that requires direct access to system memory and hardware devices
- ❖ Programs running in real-address mode can cause the operating system to crash (stop responding to commands)
-
- #### ❖ Protected mode (introduced with the 80386 processor)
- ❖ Each program can address a maximum of 4 GB of memory
 - ❖ The operating system assigns memory to each running program
 - ❖ Programs are prevented from accessing each other's memory (*segments*)
 - ❖ Native mode used by Windows NT, 2000, XP, and Linux

Q. No. 2

Consider the following initialization:

[3 + 3 + 2 = 8 points]

```
X1    WORD    0E342H, 4 DUP(0Eh)
```

- (i) Assign proper physical addresses (using a real address mode) to each byte stored in the data segment, and draw a memory map. (Assume DS= 2FF0h, and the starting offset is 2304h).
- (ii) For the above data definition directives, give the content of the destination register after execution of each of the following instructions:

Address	value
32204	42h
32205	E3h
32206	0Eh
32207	00
32208	0Eh
32209	00
3220A	0E
3220B	00
3220C	0E
3220D	00

```
MOV  EAX, DWORD PTR X1          ; (a) EAX = 000EE342h
```

```
MOV  BL,  SIZEOF X1             ; (b) BL = 0Ah
```

```
MOV  ESI, 4
```

```
MOV  BX,  [X1+ESI]              ; (c) BX = 000Eh
```

- (iii) Where indicated, write down the values of the Carry, Sign, Zero, and Overflow flags after the execution of each instruction below:

```
MOV  AX, 7FF0H
```

```
ADD  AL, 10H          ; (a) CF = 1      SF = 0      ZF = 1  OF = 0
```

ADD AH, 1 ; (b) CF = 0 SF = 1 ZF = 0 OF = 1

Q. No. 3 Write assembly language programs for the following problems: [4 + 4 = 8 points]

- (i) Declare two variables “val1” of type BYTE and “val2” of type WORD, initialized with hexadecimal values 79h and 100h respectively. Find the multiplication of these variables using a loop instruction and store the result into a third variable “val3” of type DWORD.

```
INCLUDE Irvine32.inc
.data
val1 BYTE 79h
val2 WORD 100h
val3 DWORD ?
.code
main PROC
    MOVZX EBX, val2
    MOVZX ECX, val1
    MOV EAX, 0
L1:
    ADD EAX, EBX
    LOOP L1
    MOV val3, EAX
    call dumpregs
exit
main ENDP
END main
```

- (ii) The Lucas sequence has the same recursive relationship as the Fibonacci sequence, where each term is the sum of the previous two terms, but with different starting values. If the starting values are 2 and 1, write an assembly language program that finds and stores the missing elements in the following series into a WORD type variable X1:

2, 1, 3, 4, 7, __, __, __, __, __ .

Hint: X1 word 2, 1, 3, 4, 7, 5 DUP(?)

```
INCLUDE Irvine32.inc
.code
main PROC
    MOV EBX, 2
    MOV EDX, 1
    MOV EAX, 2
    CALL WRITEDEC
    MOV ECX, 10 - 1
L1:
    MOV AL, ','
    CALL WRITECHAR
```

```
        MOV EAX,EDX
        CALL WRITEDEC
        ADD EAX, EBX
        MOV EBX, EDX
        MOV EDX, EAX

    LOOP L1

exit
main ENDP
END main
```

Best of LUCK