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AIM:	Divide and Conquer- Strassen's Matrix Multiplication.
ALGORITHM/ THEORY:	Strassen's Matrix multiplication can be performed only on square matrices where $\bf n$ is a power of $\bf 2$. Order of both of the matrices are $\bf n \times \bf n$. Divide $\bf X$, $\bf Y$ and $\bf Z$ into four $(n/2)\times(n/2)$ matrices as represented below – $ Z = \begin{bmatrix} I & J \\ K & L \end{bmatrix} \qquad X = \begin{bmatrix} A & B \\ C & D \end{bmatrix} \text{and} \qquad Y = \begin{bmatrix} E & F \\ G & H \end{bmatrix} $

Using Strassen's Algorithm compute the following -

$$M_1 := (A+C) \times (E+F)$$

$$M_2 := (B+D) \times (G+H)$$

$$M_3:=(A-D)\times (E+H)$$

$$M_4 := A imes (F - H)$$

$$M_5 := (C+D) \times (E)$$

$$M_6 := (A+B) \times (H)$$

$$M_7 := D \times (G - E)$$

Then,

$$I := M_2 + M_3 - M_6 - M_7$$

$$J := M_4 + M_6$$

$$K := M_5 + M_7$$

$$L := M_1 - M_3 - M_4 - M_5$$

Analysis

$$T(n) = \left\{ egin{array}{ll} c & if \, n = 1 \ 7 \, x \, T(rac{n}{2}) + d \, x \, n^2 & otherwise \end{array}
ight.$$
 where \emph{c} and \emph{d} are constants

Using this recurrence relation, we get $\ T(n) = O(n^{log7})$

Hence, the complexity of Strassen's matrix multiplication algorithm is $\ O(n^{log7})$.

PROGRAM:

```
#include<stdio.h>
int main(){
  int i,j;
  int p,q,r,s,t,u,v;
  int a[2][2],b[2][2],c[2][2];
  printf("\n\t\t******Strassins Matrix Multiplication
Program******\t\t\n");
  printf("\nPlease enter the 4 elements of first matrix: ");
  for(i=0;i<2;i++)</pre>
      for(j=0;j<2;j++)</pre>
           scanf("%d",&a[i][j]);
  printf("Please Enter the 4 elements of second matrix: ");
  for(i=0;i<2;i++)</pre>
      for(j=0;j<2;j++)</pre>
           scanf("%d",&b[i][j]);
  printf("\nThe first matrix is\n");
  for(i=0;i<2;i++){</pre>
      printf("\n");
      for(j=0;j<2;j++)</pre>
           printf("%d\t",a[i][j]);
  printf("\nThe second matrix is\n");
  for(i=0;i<2;i++){</pre>
      printf("\n");
      for(j=0;j<2;j++)</pre>
           printf("%d\t",b[i][j]);
  p=(a[0][0] + a[1][1])*(b[0][0]+b[1][1]);
  q= (a[1][0]+a[1][1])*b[0][0];
  r= a[0][0]*(b[0][1]-b[1][1]);
  s= a[1][1]*(b[1][0]-b[0][0]);
  t= (a[0][0]+a[0][1])*b[1][1];
  u= (a[1][0]-a[0][0])*(b[0][0]+b[0][1]);
  v= (a[0][1]-a[1][1])*(b[1][0]+b[1][1]);
```

```
c[0][0]=p+s-t+v;
c[0][1]=r+t;
c[1][0]=q+s;
c[1][1]=p-q+r+s;
printf("\nMatrix multiplication using Strassins Algorithm: \n");
for(i=0;i<2;i++){</pre>
    printf("\n");
    for(j=0;j<2;j++)</pre>
         printf("%d\t",c[i][j]);
return 0;
```

RESULT:

```
*******Strassins Matrix Multiplication Program******
Please enter the 4 elements of first matrix: 4 5 3 1
Please Enter the 4 elements of second matrix: 7 8 9 1
The first matrix is
The second matrix is
Matrix multiplication using Strassins Algorithm:
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PS C:\Users\maazs\OneDrive\Desktop\Studies\DAA\DAA Coding>
```

CONCLUSION: I understood the Strassen's matrix multiplication algorithm and also understood how its time complexity is less than normal matrix multiplication method.