

CSC 330 / Spring 2023

# OCaml First-class Functions

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*based on Dan Grossman's Lecture materials*

## Example

```
let double x = x * 2 (* old-fashioned function *)
let incr   x = x + 1 (* old-fashioned function *)

let funcs = [ double; incr ] (* list of functions *)

let rec apply_funcs (fs, x) =
  match fs with
  | [] -> x
  | f :: fs' -> apply_funcs (fs', f x)

let foo = apply_funcs (funcs, 100)          (* 201 *)
let bar = apply_funcs (List.rev funcs, 100) (* 202 *)
```

## Code demonstration

# What is *functional programming*?

- ▶ No crisp definition (you get to know it when you see it), but usually includes:
  - Avoiding mutation in most (or all) cases
  - Using functions as values: the topic of today's lecture
- ▶ Definition may also include:
  - Using recursion and recursive variant types
  - Style closer to mathematical definitions
  - Idioms using “laziness” (later topic, briefly)
  - (*bad definition*) Anything not OOP or C?
- ▶ A *functional language* just encourages functional programming
  - Often not a clear yes or no

# First-class functions

- ▶ **First-class** means something is *in the language* of expressions:  
you can compute them, pass them around, put them in data structures, etc.
- ▶ Functions in many languages, including OCaml, are first class
  - I never said they weren't
  - OCaml functions are *almost* values
- ▶ Most common use is as an argument to another function
  - Lets us abstract over *what to compute* in certain situations: caller just passes in code to handle!
  - A function that takes or returns functions is called higher-order

## Function closures

- ▶ Functions can use bindings from the enclosing environment where they were defined
  - Even if function is passed around and called somewhere else
  - Makes first-class functions *much* more powerful and useful
- ▶ Will study this carefully after some simpler examples
  - Will need function *values* to be not just functions, but also the environments where they were defined
  - A *function and its environment* is called a **function closure**
- ▶ In theory, a language could have first-class functions without closures or vice versa, but typically a language with one has the other

## Plan for this lecture

1. How to use first-class functions and function closures?
2. The precise semantics for function closures and function calls
3. Multiple powerful idioms this semantics enables

## Functions as arguments

- ▶ Can pass one function as an argument to another function
  - Again, not really a new feature: we just haven't done this before
- ▶ Elegant way to factor out common code
  - Replace  $n$  similar functions with calls to 1 function with  $n$  different (short) function arguments

```
let rec n_times (f, n, x) =  
  if n = 0 then x else f (n_times (f, n - 1, x))
```

## Types for `n_times`

- ▶ What's the type of `n_times`?

```
let rec n_times (f, n, x) =
  if n = 0
  then x
  else f (n_times (f, n - 1, x))
```

- It's `val n_times : ('a -> 'a) * int * 'a -> 'a`
  - ▶ Types are *inferred* based on how arguments are used!
  - ▶ More useful than `(int -> int) * int * int -> int`

## Relation to types

- ▶ Higher-order functions often “so reusable” that they have polymorphic types
  - i.e., types with **type variables** (e.g., `'a`)
- ▶ There are higher-order functions that are not polymorphic
  - e.g., `val times_until_zero : (int -> int) * int -> int`
- ▶ And there are polymorphic functions that are not higher-order
  - e.g., `val len : 'a list -> int`
- ▶ Always a good idea to understand a function’s type, especially for higher-order functions

## Code demonstration

# Map

```
let rec map (f, xs) =
  match xs with
  | [] -> []
  | x :: xs' -> f x :: map (f, xs')
(* ('a -> 'b) * 'a list -> 'b list *)
```

- ▶ Hall-of-fame higher-order function
- ▶ Name is standard (used for any sort of collection data structure)
  - Generally provided in standard libraries for standard collections
- ▶ Used all the time in functional code, thus *idiomatic*
  - Using it where appropriate isn't just shorter code but *communicates what you are doing*
  - Conversely, not using it when "doing a map" confuses your reader

# Filter

```
let rec filter f xs =
  match xs with
  | [] -> []
  | x :: xs' -> if f x
    then x :: filter f xs'
    else filter f xs'
(* ('a -> bool) * 'a list -> 'a list *)
```

- ▶ Another hall-of-famer
  - Keep only some elements from a collection
  - Again an idiom, so you should use whenever “doing a filter”

# Anonymous functions

# Anonymous functions

- ▶ **Syntax:**

```
fun p -> e
```

where `p` is a pattern and `e` is an expression

- ▶ Function as an expression!

- Type checking, evaluation rules “the same” as function bindings
- Note that use `->` instead of `=` to separate parameters from function body
- No function name!
- *Just an expression:* can appear anywhere expressions allowed!

# Using anonymous functions

- ▶ Usually use anonymous functions as arguments to other functions
  - No need to name a function that you use only once
- ▶ **Limitation:** anonymous functions cannot call themselves
  - No name to make a recursive call with
- ▶ Non-recursive function bindings are really just syntactic sugar
  - These two lines are equivalent!

```
let f p = e  
let f = fun p -> e
```

- Again, doesn't work for recursive functions
- And as usual, better style to use the syntactic sugar available

## A point on style

BAD

```
if e then true else false  
fun x -> f x  
n_times ((fun x -> List.tl x), 3, xs)
```

GOOD

```
e  
f  
n_times (List.tl, 3, xs)
```

# Generalizing

- ▶ Our examples so far have been awfully similar
  - Take one function as argument then process a number or list
- ▶ We can do much more!
  - Pass several functions as arguments
  - Put functions in data structures
  - Return functions as results
  - Write higher-order functions over variant types
- ▶ Useful whenever abstracting over “what to compute with”

# Returning functions

- ▶ Functions are first-class, so they can be returned from (other) functions
- ▶ Silly example:

```
let double_or_triple f =
  if f 7 then fun x -> 2 * x else fun x -> 3 * x
```

- Has type  $(\text{int} \rightarrow \text{bool}) \rightarrow (\text{int} \rightarrow \text{int})$
- But REPL reports  $(\text{int} \rightarrow \text{bool}) \rightarrow \text{int} \rightarrow \text{int} !$
- Well, this is **the same thing**:
  - ▶ REPL never prints unnecessary parentheses
  - ▶  $t_1 \rightarrow t_2 \rightarrow t_3 \rightarrow t_4$  means  $t_1 \rightarrow (t_2 \rightarrow (t_3 \rightarrow t_4))$
  - ▶ We will learn soon why this *precedence* is convenient

## Other data structures

- ▶ Higher-order functions are not just for lists
- ▶ They work great for common traversals over your own data structures
- ▶ See the code demonstration for an example

## Code demonstration

