

# Modern Algebra

## CS 2LC3

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- *Modern algebra* is the study of the structure of certain sets along with operations on them.
- The algebras discussed here are *semigroups*, *monoids*, *groups*, and *boolean algebras*. They are useful throughout computer science and mathematics.
- Semigroups and monoids find application in formal languages, automata theory, and coding theory.
- One boolean algebra is the standard model of the propositional calculus.
- Important in our study is not only the various algebras but their interrelationship.
- Thus, we study topics like *isomorphisms*, *homomorphisms*, and *automorphisms* of algebras.

# The Structure of Algebras

- An algebra consists of two components:
  - 1 A set  $S$  of elements, called the *carrier* of the algebra.
  - 2 Operators defined on the carrier.
- Formally an algebra is a pair  $(S, \Psi)$ , where  $S$  is a carrier and  $\Psi$  is a list of operators.
- Each operator is a total function of type  $S^m \rightarrow S$  for some  $m$ , where  $m$  is called the *arity* of the operator.
- The algebra is *finite* if its carrier  $S$  is finite; otherwise, it is *infinite*.

- Formally an algebra is a pair  $(S, \Psi)$ , where  $S$  is a carrier and  $\Psi$  is a list of operators.
- Each operator is a total function of type  $S^m \rightarrow S$  for some  $m$ , where  $m$  is called the *arity* of the operator.
- Operators of arity 0, called *nullary* operators, are functions of no arguments.
- The nullary operators are interpreted as *constants* in the carrier.
- For example, we consider 1 to be a function that takes no arguments and returns the value one.
- Operators of arity 1 are *unary* operators; of arity 2 , *binary operators*; of arity 3 , *ternary operators*.
- Unary operators are written in prefix form (for example  $-x$ ); binary operators in infix form (for example  $a + b$ ).

# Examples of Algebras

- (a) The set of even integers and the operator  $+$  form an algebra  $(\text{Even}, +)$ .
- (b) The set of even numbers together with the operations multiplication and division is not an algebra, because division is not a total function on the even integers (division by 0 is not defined).
- (c) The set  $\{\text{false}, \text{true}\}$  and operators  $\vee, \wedge$  and  $\neg$ , is an algebra  $(\mathbb{B}, \vee, \wedge, \neg)$ .  
This is a finite algebra, because the set is finite.

- The *signature* of an algebra consists of the name of its carrier and the list of types of its operators.

- For example, the algebra  $(\mathbb{B}, \vee, \wedge, \neg)$  has the signature:

$$(\mathbb{B}, \mathbb{B} \times \mathbb{B} \rightarrow \mathbb{B}, \mathbb{B} \times \mathbb{B} \rightarrow \mathbb{B}, \mathbb{B} \rightarrow \mathbb{B})$$

- Two algebras are said to have the same signature if
  - (i) they have the same number of operators and
  - (ii) corresponding operators have the same types (modulo the name of the carrier).
- For example, algebras  $(\mathbb{B}, \vee, \wedge, \neg)$  and  $(\mathcal{P}(S), \cap, \cup, \sim)$  for some set  $S$  have the same signature.
- Algebra  $(\mathcal{P}(S), \sim, \cup, \cap)$  has a different signature, since  $\sim$  is of arity 1 and  $\vee$  and  $\cap$  are of arity 2 .

## Definition

- An element  $1$  in  $S$  is a *left identity* (or *unit*) of binary operator  $\circ$  over  $S$  if  $1 \circ b = b$  (for  $b \in S$ );
- $1$  is a *right identity* if  $b \circ 1 = b$  (for  $b \in S$ ); and
- $1$  is an *identity* if it is both a left and a right identity.

## Theorem

If  $c$  is a left identity of  $\circ$  and  $d$  is a right identity of  $\circ$ , then  $c = d$ .

## Proof.

$$c = c \circ d = d.$$



## Definition

- An element  $0$  in  $S$  is a *left zero* of binary operator  $\circ$  over  $S$  if  $0 \circ b = 0$  (for  $b \in S$ );
  - $0$  is a *right zero* if  $b \circ 0 = 0$  (for  $b \in S$ ); and
  - $0$  is an *zero* if it is both a left and a right zero.
- 
- An algebra can have more than one left zero. For example, consider algebra  $(\{b, c\}, \circ)$  with operator  $\circ$  defined below.  
$$b \circ b = b \quad c \circ b = c$$
$$b \circ c = b \quad c \circ c = c$$
  - Both  $b$  and  $c$  are left zeros - and both are right identities!

## Theorem

If  $c$  is a left zero of  $\circ$  and  $d$  is a right zero of  $\circ$ , then  $c = d$ .

## Proof.

$$d = c \circ d = c.$$





# One-to-one and Onto Functions

## Definition

- 1 A function  $f : B \rightarrow C$  is *one-to-one* iff  $f(b) = f(c) \implies b = c$  for all  $b, c \in B$ .
- 2 A function  $f : B \rightarrow C$  is *onto* iff  $\text{Ran}(f) = C$ , i.e. for every  $c \in C$  there is  $b \in B$  such that  $f(b) = c$ .

## Definition

Let  $1$  be the identity of binary operator  $\circ$  on  $S$ . Then  $b$  has a **right inverse**  $c$  with respect to  $\circ$  and  $c$  has a **left inverse**  $b$  with respect to  $\circ$  if  $b \circ c = 1$ . Elements  $b$  and  $c$  are called **inverses** of each other if  $b \circ c = c \circ b = 1$ .

Examples of inverses:

- In algebra  $(\mathbb{Z}, +)$ ,  $0$  is an identity. Every element  $b \in \mathbb{Z}$  has an inverse  $-b$ .
- In algebra  $(\mathbb{R}, \cdot)$ ,  $1$  is an identity. Every element  $b \in \mathbb{R}$  except  $0$  has an inverse  $\frac{1}{b}$ .
- Consider the set  $F$  of functions of arity 1 over a set  $S$ , and let
  - be function composition:  $(f \bullet g)(b) = f(g(b))$ . Then the function  $id$  given by  $id(b) = b$  (for all  $b \in S$ ) is an identity.**Every onto function has a right inverse, every one-to-one function has a left inverse. Every one-to-one and onto function has an inverse.**

## Definition

A subset  $T$  of a set  $S$  is *closed* under an operator if applying the operator to elements of  $T$  always produces a result in  $T$ .

### Example of closed operators

- (a) The set of even integers is closed under  $+$  because the sum of two even integers is even.
- (b) Subset  $\{0, 1\}$  of the integers is not closed under  $+$  because  $1 + 1$  is not in this subset.
- (c) Subset  $\{0, 1\}$  of the integers is closed under  $\uparrow$  (maximum) because the maximum of any two of these integers is one of the integers.  $\square$

## Definition

$(T, \Phi)$  is a subalgebra of  $(S, \Phi)$  if

- (a)  $\emptyset \subset T \subseteq S$ , and
- (b)  $T$  is closed under every operator in  $\Phi$ .

## Examples

- (a) Algebra  $\langle \mathbb{N}, + \rangle$  is a subalgebra of  $\langle \mathbb{Z}, + \rangle$  because  $\mathbb{N} \subseteq \mathbb{Z}$  and  $\mathbb{N}$  is closed under  $+$ .
- (b)  $\langle \{0, 1\}, + \rangle$  is not a subalgebra of  $\langle \mathbb{Z}, + \rangle$  because  $\{0, 1\}$  is not closed under  $+$ .
- (c) Algebra  $\langle \{0, 1\}, \cdot \rangle$  is a subalgebra of  $\langle \mathbb{N}, \cdot \rangle$ .
- (d) Any algebra is a subalgebra of itself. □

# Isomorphism

## Definition

Let algebras  $A = (S, \Phi)$  and  $\hat{A} = (\hat{S}, \hat{\Phi})$  have the same signature. A function  $h : S \rightarrow \hat{S}$  is an **isomorphism** from  $A$  to  $\hat{A}$  if:

- (a) Function  $h$  is one-to-one and onto.
- (b) For each pair of corresponding nullary operators (constants)  $c$  in  $\Phi$  and  $\hat{c}$  in  $\hat{\Phi}$ ,  $h(c) = \hat{c}$ .
- (c) For each pair of corresponding unary operators  $\sim$  in  $\Phi$  and  $\hat{\sim}$  in  $\hat{\Phi}$ ,  $h(\sim b) = \hat{\sim}h(b)$  (for  $b$  in  $S$ ).
- (d) For each pair of corresponding binary operators  $\circ$  in  $\Phi$  and  $\hat{\circ}$  in  $\hat{\Phi}$ ,  $h(b \circ c) = h(b)\hat{\circ}h(c)$ .

$A$  and  $\hat{A}$  are **isomorphic**, and  $\hat{A}$  is the **isomorphic image** of  $A$  under  $h$ .

- Property (d) is sometimes depicted as the **commuting diagram**.

$$\begin{array}{ccc} S \times S & \xrightarrow{\circ} & S \\ h \downarrow & & \downarrow h \\ \hat{S} \times \hat{S} & \xrightarrow{\hat{\circ}} & \hat{S} \end{array}$$

# Examples of Isomorphism

- ① Let  $A = (\mathbb{B}, \vee)$  and  $\hat{A} = (\mathbb{B}, \wedge)$ . Clearly  $A$  and  $\hat{A}$  have the same signature. Define  $h : \mathbb{B} \rightarrow \mathbb{B}$  by  $h(b) = \neg b$ . Function  $h$  is one-to-one and onto. Moreover:

$$h(b \vee c) = \neg(b \vee c) = \neg b \wedge \neg c = h(b) \wedge h(c).$$

- ② Let  $A = (\mathbb{N}, +)$  and  $\hat{A} = (\text{even}, +)$ , where *even* is the set of even natural numbers.  $A$  and  $\hat{A}$  have the same signature. Define  $h : \mathbb{N} \rightarrow \text{even}$  by  $h(b) = 2 \cdot b$  (for  $b \in \mathbb{N}$ ). Function  $h$  is one-to-one and onto. Moreover:

$$h(b + c) = 2 \cdot (b + c) = 2 \cdot b + 2 \cdot c = h(b) + h(c).$$

- ③ Let  $A = (\mathbb{R}^+, \cdot)$  and  $\hat{A} = (\mathbb{R}, +)$ , where  $\mathbb{R}^+$  is the set of positive real numbers.  $A$  and  $\hat{A}$  have the same signature. Define  $h : \mathbb{R}^+ \rightarrow \mathbb{R}$  by  $h(r) = \log(r)$  for  $r > 0$ , so that  $h^{-1}(r) = 2^r$ . Function  $h$  is clearly one-to-one and onto. Moreover:

$$h(b \cdot c) = \log(b \cdot c) = \log(b) + \log(c) = h(b) + h(c).$$

## Theorem

- 1 *An isomorphism maps identities to identities, zeros to zeros, and inverses to inverses.*
- 2 *If  $\hat{A}$  is an isomorphic image of  $A$ , then  $A$  is an isomorphic image of  $\hat{A}$ .*
- 3 *Let  $\mathcal{C}$  be a set of algebras. The relation “ $A$  is isomorphic to  $\hat{A}$ ” is an equivalence relation.*

## Definition

An isomorphism from  $A$  to  $A$  is called an **automorphism**.

Examples of automorphism:

- Let  $A = \hat{A} = (S, \Phi)$ . Let  $h$  be the identity function on  $S$ , i.e.  $h(b) = b$  for  $b \in S$ . Here  $h$  is automorphism.
- Let  $A = \hat{A} = (\mathbb{Z}, +)$  and  $h$  be defined as  $h(b) = -b$  for all  $b \in \mathbb{Z}$ . Again  $h$  is automorphism.



## Definition

Let algebras  $A = (S, \Phi)$  and  $\hat{A} = (\hat{S}, \hat{\Phi})$  have the same signature. A function  $h : S \rightarrow \hat{S}$  is a **homomorphism** from  $A$  to  $\hat{A}$  if it satisfies:

- (a) For each pair of corresponding nullary operators  $c$  in  $\Phi$  and  $\hat{c}$  in  $\hat{\Phi}$ , we have  $h(c) = \hat{c}$ .
- (b) For each pair of corresponding unary operators  $\sim$  in  $\Phi$  and  $\hat{\sim}$  in  $\hat{\Phi}$ ,  $h(\sim b) = \hat{\sim}h(b)$  (for  $b \in S$ ).
- (c) For each pair of corresponding binary operators  $\circ$  in  $\Phi$  and  $\hat{\circ}$  in  $\hat{\Phi}$ ,  $h(b \circ c) = h(b)\hat{\circ}h(c)$  (for  $b, c \in S$ ).

- An isomorphism is a homomorphism that is also one-to-one and onto.

## Theorem

Let  $h$  be a homomorphism from  $A = (S, \Phi)$  to  $\hat{A} = (\hat{S}, \hat{\Phi})$ . Then  $(h(S), \hat{\Phi})$  is a subalgebra of  $\hat{A}$ , called the *homomorphic image of  $A$  under  $h$* .

## Proof.

We show that  $(h(S), \hat{\Phi})$  satisfies the definition of a subalgebra.

- (a) Since  $h : S \rightarrow \hat{S}$ ,  $h(S) \subseteq \hat{S}$ .
- (b) We show that  $h(S)$  is closed under each binary operator  $\hat{o}$  in  $\hat{\Phi}$ . Let  $\hat{b}$  and  $\hat{c}$  be in  $h(S)$ . Then there exist values  $b, c$  in  $S$  that satisfy  $h(b) = \hat{b}$  and  $h(c) = \hat{c}$ . Moreover we have:

$$\hat{b}\hat{o}\hat{c} = h(b)\hat{o}h(c) = h(b \circ c).$$

Hence,  $\hat{b}\hat{o}\hat{c}$  is in  $h(S)$  and  $h(S)$  is closed under  $\hat{o}$ . Similarly,  $h(S)$  is closed under all the nullary and unary operators of  $\hat{\Phi}$ .

# Examples of Homomorphism

- (a) Function  $h.b = 5 \cdot b$  is a homomorphism from algebra  $\langle \mathbb{N}, + \rangle$  to itself. There are no unary operators, and  $h(b + c) = 5 \cdot (b + c) = 5 \cdot b + 5 \cdot c = h.b + h.c$  (for  $b$  and  $c$  in  $\mathbb{N}$ ). Actually, for any integer  $k$  (including 0),  $h.b = k \cdot b$  is a homomorphism from  $\langle \mathbb{N}, + \rangle$  to itself.
- (b) Let  $\oplus$  be the function defined by  $b \oplus c = (b + c) \bmod 5$ . Then  $h.b = b \bmod 5$  is a homomorphism from  $\langle \mathbb{N}, \oplus \rangle$  to  $\langle 0..4, \oplus \rangle$ .

# Lattices as Algebras

## Definition

A **lattice** is an algebra  $(S, \sqcup, \sqcap)$ , where  $\sqcup$  and  $\sqcap$ , called *join* and *meet* are two binary, commutative and associative operators that satisfy, for all  $a, b \in S$ :

(a)  $a \sqcup (a \sqcap b) = a$

(b)  $a \sqcap (a \sqcup b) = a$

(c)  $a \sqcup a = a$

(d)  $a \sqcap a = a$

The axioms (a) and (b) are called *absorption laws*, (c) and (d) are called *idempotency laws*.

## Example (Examples of Algebra Lattices)

- $(\mathbb{B}, \vee, \wedge)$  - Boolean Lattice,
- $(\mathcal{P}(S), \cup, \cap)$ , where  $S \neq \emptyset$  is a set - Set Lattice,
- $(\mathcal{R}, \uparrow, \downarrow)$ , where  $\uparrow$  is a maximum, and  $\downarrow$  is a minimum, is a lattice.

## Theorem

Let  $(S, \sqcup, \sqcap)$  be a lattice algebra. Define a relation  $\preceq$  on  $S$  as follows, for all  $a, b \in S$ :  $a \preceq b \iff a \sqcap b = a$ .

The pair  $(S, \preceq)$  is a *partial order lattice*.

## Proof.

First note that  $a = a \sqcap b \implies b = b \sqcup (b \sqcap a) = (a \sqcap b) \sqcup b = a \sqcup b$ , so also  $a \preceq b \iff a \sqcup b = b$ . We will show that  $\preceq$  is a partial order. From  $a \sqcap a = a$  we have  $a \preceq a$ . Consider  $a \preceq b \wedge b \preceq a$ . We have  $a \sqcap b = a \wedge a \sqcap b = b$  so  $a = b$ . Consider  $a \preceq b \wedge b \preceq c$ . Here we have  $a \sqcap c = (a \sqcap b) \sqcap c = a \sqcap (b \sqcap c) = a \sqcap b = a$ , so  $a \preceq c$ . Hence  $\preceq$  is a partial order. We will show  $a \sqcap b = \text{glb}(\{a, b\})$  and  $a \sqcup b = \text{lub}(\{a, b\})$ . We have  $(a \sqcap b) \sqcup a = a \iff a \sqcap b \preceq a$  and  $(a \sqcap b) \sqcup b = a \iff a \sqcap b \preceq b$ , so  $a \sqcap b$  is a lower bound of  $\{a, b\}$ . Consider  $c$  such that  $c \preceq a \wedge c \preceq b$ . Hence  $c \sqcap a = c \wedge c \sqcap b = c$ , i.e.  $c = c \sqcap c = (c \sqcap a) \sqcap (c \sqcap b) = (a \sqcap b) \sqcap c$ , which means  $c \preceq a \sqcap b$ . So  $a \sqcap b = \text{glb}(\{a, b\})$ . Similarly we can show that  $a \sqcup b = \text{lub}(\{a, b\})$ . ■

## Theorem

Let  $(S, \preceq)$  be a partial order lattice. Define an algebra  $(S, \sqcup, \sqcap)$ , where  $\sqcup, \sqcap$  are binary operators defined as follows, for all  $a, b \in S$ :

- $a \sqcap b = \text{lub}(\{a, b\})$ ,
- $a \sqcup b = \text{glb}(\{a, b\})$ .

The algebra  $(S, \sqcup, \sqcap)$  is a *lattice algebra*.

## Proof.

Directly from the definition we have that  $\sqcup, \sqcap$  are binary and symmetric. Moreover  $\text{lub}(\{a, \text{lub}(\{b, c\})\}) = \text{lub}(\{a, b, c\})$  so  $\text{lub}$  is associative. and similarly for  $\text{glb}$ .

Since  $\text{glb}(\{a\}) = \text{lub}(\{a\}) = a$ , then the axioms (c) and (d) are satisfied. Note that if  $a \preceq b$  then  $\text{glb}(\{a, b\}) = a$  and  $\text{lub}(\{a, b\}) = b$ . Since  $c = \text{glb}(\{a, b\}) \preceq a$ , so  $\text{lub}(\{a, c\}) = a$ , i.e.  $\text{lub}(\{a, \text{glb}(\{a, b\})\}) = a$ , or  $a \sqcup (a \sqcap b) = a$ , so the axiom (a) is satisfied. Similarly for (b). ■

- Lattice algebra and partial order lattice are different model of the same concept, they are equivalent, however different applications, one of the models might be more convenient.

**Alphabet:** an *arbitrary* (usually finite) set of elements, often denoted by the symbol  $\Sigma$ .

**Sequence:**

- an element  $x = (a_1, a_2, \dots, a_k) \in \Sigma^k$ , where  $\Sigma^k$  is a Cartesian product of  $\Sigma$ 's.

For convenience we write  $x = a_1 a_2 \dots a_k$ .

- a function  $\phi : \{1, \dots, k\} \rightarrow \Sigma$ , such that  $\phi(1) = a_1, \dots, \phi(k) = a_k$ .

♣ The two above definitions are in a sense identical since:

$$\underbrace{\Sigma \times \dots \times \Sigma}_n \equiv \{f \mid f : \{1, \dots, k\} \rightarrow \Sigma\}.$$

- Frequently a *sequence* is considered as a primitive undefined concept that is understood and does not need any explanation.

- If the elements of  $\Sigma$  are *symbols*, then a *finite* sequence of symbols is often called a *string* or a *word*.
- The *length* of a sequence  $x$ , denoted  $|x|$ , is the number of elements composing the sequence.
- The *empty sequence*,  $\varepsilon$ , is the sequence consisting of zero symbols, i.e.  $|\varepsilon| = 0$ .
- A *prefix* of a sequence is any number of leading symbols of that sequence, and a *suffix* is any number of trailing symbols (any number means 'zero included').



# Concatenation

- *Concatenation* (operation)

Let  $x = a_1 \dots a_k$ ,  $y = b_1 \dots b_l$ . Then

$$x \circ y = a_1 \dots a_k b_1 \dots b_l.$$

We usually write  $xy$  instead of  $x \circ y$ .

- Properties of concatenation:

①  $x(yz) = (xy)z$

②  $\varepsilon x = x\varepsilon = x$

*Fact.* A triple  $(\Sigma, \circ, \varepsilon)$  is a *monoid*, or *semigroup* (a concept discussed later).

- Power operator:  $x^0 = \varepsilon$ ,  $x^1 = x$  and  $x^k = \underbrace{x \dots x}_k$ .

- Recursive definition of power:

$$x^0 = \varepsilon$$

$$x^{k+1} = x^k x.$$

- Function  $h : \Sigma^* \rightarrow \mathbb{N}$  defined by  $h(z) = |z|$  is a *homomorphism* from  $(\Sigma^*, \circ, \varepsilon)$  to  $(\mathbb{N}, +, 0)$ .

- Let  $\Sigma$  be a finite alphabet. Then we define  $\Sigma^*$  as:

$$\Sigma^* = \{a_1 \dots a_k \mid a_i \in \Sigma \wedge k \geq 0\},$$

i.e. the set of all sequences, including  $\varepsilon$ , built from the elements of  $\Sigma$ .

- A (*formal*) language over  $\Sigma$  is any subset of  $\Sigma^*$ , including the empty set  $\emptyset$  and  $\Sigma^*$ .

# Semigroups

## Definition

A **semigroup** is an algebra  $(S, \circ)$  where  $\circ$  is a binary associative operator.

## Example

- $(\Sigma^*, \circ)$ , where  $\circ$  is a string concatenation, is a semigroup (an important one).
- $([0, 1], \cdot)$ , where “ $\cdot$ ” is a multiplication is a semigroup.
- $(S, \uparrow)$ , where  $S$  is any nonempty subset of the real numbers and  $b \uparrow c$  is the maximum of  $b$  and  $c$ , is a semigroup.
- $(\{b, c\}, \circ)$ , where  $\circ$  is defined by  $b \circ b = c \circ b = b$  and  $b \circ c = c \circ c = c$ . This is a finite semigroup (since  $S$  is finite).
- Let  $X$  be a set.  $(\text{Rel}(X), \circ)$ , where  $\text{Rel}(X)$  is the set of all binary relations over  $X$  and  $\circ$  is a composition of relations, is a semigroup.

## Definition

Let  $T$  be a subset of carrier  $S$  of semigroup  $(S, \circ)$ . Suppose  $T$  is closed under  $\circ$ . Then algebra  $(T, \circ)$  is called a **subsemigroup** of  $(S, \circ)$ .

# Monoids

## Definition

- A **monoid**  $(S, \circ, 1)$  is a semigroup  $(S, \circ)$  with an identity 1.
- If  $\circ$  is also symmetric, the monoid is called **Abelian**
- A subalgebra of a monoid that contains the identity of the monoid is called a **submonoid**.

## Example

- $(\Sigma^*, \circ, \epsilon)$ , where  $\circ$  is a string concatenation, is a monoid. This monoid is not Abelian.
- $([0, 1], \cdot, 1)$ , where “ $\cdot$ ” is a multiplication is an Abelian monoid.
- $(S, \uparrow)$ , where  $S$  is any nonempty subset of the real numbers and  $b \uparrow c$  is the maximum of  $b$  and  $c$ , is not a monoid, since  $\uparrow$  has no identity in  $\mathbb{R}$ .
- $(\mathbb{N}, \uparrow, 0)$  is an Abelian monoid. Note that  $0 \uparrow b = b \uparrow 0 = b$  for all  $b \in \mathbb{N}$ .
- Let  $X$  be a set.  $(\text{Rel}(X), \circ, \text{id}_X)$ , where  $\text{Rel}(X)$  is the set of all binary relations over  $X$  and  $\circ$  is a composition of relations, is a monoid. This monoid is not Abelian.

- Any semigroup  $(S, \circ)$  can be made into a monoid  $(S \cup \{c\}, \circ, c)$  for  $c \notin S$  a fresh element that is defined to satisfy  $c \circ b = b \circ c = b$  for all elements of  $S \cup \{c\}$ .
- For example, operator  $\uparrow$  can be extended to  $\mathbb{R} \cup \{\infty\}$  by  $r \uparrow \infty = \infty \uparrow r = r$  for all elements of  $\mathbb{R} \cup \{\infty\}$ , so that  $\uparrow$  has an identity.
- One must be wary of this extension, however, because other properties of the reals  $\mathbb{R}$  may not hold for  $\mathbb{R} \cup \{\infty\}$ . For example,  $1 + b > b$  does not hold for  $b = \infty$ .

## Definition

A **group** is an algebra  $(S, \circ, 1)$  in which

- (a)  $\circ$  is a binary, associative operator,
- (b)  $\circ$  has the identity 1 in  $S$ ,
- (c) Every element  $b \in S$  has an inverse, which we write as  $b^{-1}$ .

A **symmetric**, **commutative**, or **Abelian group** is an Abelian monoid in which every element has an inverse.

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## Examples of groups

- (a) The *additive group of integers*  $\langle \mathbb{Z}, +, 0 \rangle$  is a group. The inverse  $b^{-1}$  of  $b$  is  $-b$ .
- (b) Let  $K$  be the set of multiples of 5. Then  $\langle K, +, 0 \rangle$  is a group. The inverse  $b^{-1}$  of  $b$  is the element  $-b$ .
- (c) Let  $n > 0$  be an integer. Define  $\oplus$  for operands  $b$  and  $c$  in  $0..(n-1)$  by  $b \oplus c = (b + c) \bmod n$ . Then  $M_n = \langle 0..(n-1), \oplus, 0 \rangle$  is a group, called the *additive group of integers modulo  $n$* .
- (d)  $\langle \mathbb{R}, \cdot, 1 \rangle$  has identity 1 but is not a group, because 0 has no inverse.
- (e)  $\langle \mathbb{R}^+, \cdot, 1 \rangle$  is a group. The inverse  $r^{-1}$  of  $r$  in  $\mathbb{R}^+$  is  $1/r$ . □

# Theorems for Groups

$$(18.18) \quad b = (b^{-1})^{-1}$$

$$(18.19) \quad \textbf{Cancellation:} \quad \begin{aligned} b \circ d = c \circ d &\equiv b = c \\ d \circ b = d \circ c &\equiv b = c \end{aligned}$$

$$(18.20) \quad \textbf{Unique solution:} \quad \begin{aligned} b \circ x = c &\equiv x = b^{-1} \circ c \\ x \circ b = c &\equiv x = c \circ b^{-1} \end{aligned}$$

$$(18.21) \quad \textbf{One-to-one:} \quad \begin{aligned} b \neq c &\equiv d \circ b \neq d \circ c \\ b \neq c &\equiv b \circ d \neq c \circ d \end{aligned}$$

$$(18.22) \quad \textbf{Onto:} \quad \begin{aligned} (\exists x | : b \circ x = c) \\ (\exists x | : x \circ b = c) \end{aligned}$$

Proof of 18.18.

$$(b^{-1})^{-1} = 1 \circ (b^{-1})^{-1} = b \circ b^{-1} \circ (b^{-1})^{-1} = b \circ 1 = b \quad \blacksquare$$



## Definition

We define integral powers  $b^n$  of an element  $b$  of a group  $(S, \circ, 1)$  as follows:

$$b^0 = 1$$

$$b^n = b^{n-1} \circ b \quad (\text{for } n > 0)$$

$$b^{-n} = (b^{-1})^n \quad (\text{for } n > 0)$$

## Properties of powers of group elements

$$b^m \circ b^n = b^{m+n} \quad (\text{for } m \text{ and } n \text{ integers})$$

$$(b^m)^n = b^{m \cdot n} \quad (\text{for } m \text{ and } n \text{ integers})$$

$$b^n = b^p \equiv b^{n-p} = 1$$

## Definition

A boolean algebra is an algebra  $(S, \oplus, \otimes, \sim, 0, 1)$  in which

- (a)  $\oplus$  and  $\otimes$  are associative binary operators;
- (b)  $\oplus$  and  $\otimes$  are symmetric;
- (c) 0 and 1 are the identities of  $\oplus$  and  $\otimes$ ;
- (d) unary operator  $\sim$  satisfies  $b \oplus (\sim b) = 1$  and  $b \otimes (\sim b) = 0$  (for all  $b$ );  $\sim b$  is called the **complement** of  $b$ ;
- (e)  $\otimes$  distributes over  $\oplus$ :  $b \otimes (c \oplus d) = (b \otimes c) \oplus (b \otimes d)$ ;
- (f)  $\oplus$  distributes over  $\otimes$ :  $b \oplus (c \otimes d) = (b \oplus c) \otimes (b \oplus d)$ .

- $\oplus$  is often called “sum” or “plus”, while  $\otimes$  is often called “product” or “times”.

## Example

- $(\mathbb{B}, \vee, \wedge, \neg, \text{false}, \text{true})$  is a boolean algebra. It is our model for the propositional calculus, it provides intuition for the general definition.
- $(\mathcal{P}(S), \cup, \cap, \sim, \emptyset, S)$  is a boolean algebra, where  $S$  is any nonempty set. We call this a **power-set algebra**.
- For  $n$  in  $\mathbb{Z}^+$ , let  $F_n$  be the set of functions of type  $\mathbb{B}^n \rightarrow \mathbb{B}$ , i.e. the set of boolean functions of  $n$  boolean arguments. Let  $s$  denote a sequence of  $n$  boolean values. Define  $\oplus, \otimes$  and  $\sim$  by  $(f1 \oplus f2)(s) = f1(s) \vee f2(s)$ ,  $(f1 \otimes f2)(s) = f1(s) \wedge f2(s)$ ,  $(\sim f)(s) = \neg f(s)$ . Then  $(F_n, \oplus, \otimes, \sim, f, t)$  is a boolean algebra. The identity of  $\oplus$  is the function  $f$  that always yields *false*, and the identity  $t$  of  $\otimes$  always yields *true*.

## Definition

- The *greatest common divisor*  $b \mathbf{gcd} c$  of integers  $b$  and  $c$  that are not both zero is the greatest integer that divides both.  
For example  $24 \mathbf{gcd} 30 = 6$ .
- The *least common multiple*  $b \mathbf{lcm} c$  of  $b$  and  $c$  is the smallest positive integer that is a multiple of both  $b$  and  $c$ .  
For example  $12 \mathbf{lcm} 18 = 36$ .

## Example

$(\{1, 2, 3, 6\}, \mathbf{lcm}, \mathbf{gcd}, \sim, 1, 6)$ , where  $\sim x = \frac{6}{x}m$ , is a boolean algebra.

# Theorems for Boolean Algebras

(18.49) **Idempotency:**  $b \oplus b = b, \quad b \otimes b = b$

(18.50) **Zero:**  $b \oplus 1 = 1, \quad b \otimes 0 = 0$

(18.51) **Absorption:**  $b \oplus (b \otimes c) = b, \quad b \otimes (b \oplus c) = b$

(18.52) **Cancellation:**  $(b \oplus c = b \oplus d) \wedge (\sim b \oplus c = \sim b \oplus d) \equiv c = d$   
 $(b \otimes c = b \otimes d) \wedge (\sim b \otimes c = \sim b \otimes d) \equiv c = d$

(18.53) **Unique complement:**  $b \oplus c = 1 \wedge b \otimes c = 0 \equiv c = \sim b$

(18.54) **Double complement:**  $\sim(\sim b) = b$

(18.55) **Constant complement:**  $\sim 0 = 1, \quad \sim 1 = 0$

(18.56) **De Morgan:**  $\sim(b \oplus c) = (\sim b) \otimes (\sim c)$   
 $\sim(b \otimes c) = (\sim b) \oplus (\sim c)$

(18.57)  $b \oplus (\sim c) = 1 \equiv b \oplus c = b, \quad b \otimes (\sim c) = 0 \equiv b \otimes c = b$

(18.58) A homomorphic image of a boolean algebra is a boolean algebra.

# Partial Order Generated by a Boolean Algebra

## Definition

Consider an arbitrary Boolean algebra  $(S, \oplus, \otimes, \sim, 0, 1)$ . Define the relations  $\leq$  and  $<$  on  $S$  as follows:

$$b \leq c \equiv b \otimes c = b$$

$$b < c \equiv b \leq c \wedge b \neq c$$

## Theorem

*Relation  $\leq$  is a partial order.*

## Proof.

Since  $b \otimes b = b$  then  $b \leq b$ , so  $\leq$  is reflexive.

$b \leq c \wedge c \leq b \iff b \otimes c = b \wedge b \otimes c = c \iff b = c$ , so  $\leq$  is antisymmetric.

$b \leq c \wedge c \leq d \iff b \otimes c = b \wedge c \otimes d = c \iff$   
 $b \otimes c \otimes d = b \otimes d \iff b \leq d$ . Hence  $\leq$  is transitive. ■

### Lemma

$$b \otimes c = b \iff b \oplus c = c$$

### Proof.

$$(\Rightarrow) \quad b \oplus c = (b \otimes c) \oplus c = c.$$

$$(\Leftarrow) \quad b \otimes c = b \otimes (b \oplus c) = b. \quad \blacksquare$$

### Theorem

$$b \leq c \iff b \oplus c = c$$

### Proof.

$$b \leq c \iff b \otimes c = b \iff b \oplus c. \quad \blacksquare$$

- **Intuition.** Consider a Boolean algebra  $(\mathcal{P}(S), \cup, \cap, \sim, \emptyset, S)$ . The *singletons*, i.e. sets containing only one element,  $\{b\}$ , for all  $b \in S$ , can be called *atoms*, as they are not divisible and each non empty set can be built from them.
- Note that only  $\emptyset$  is smaller than singletons with respect to the partial order  $\subseteq$ .

## Definition

Consider an arbitrary Boolean algebra  $(S, \oplus, \otimes, \sim, 0, 1)$ .

An element  $a \in S$  is called an **atom** if the following predicate is satisfied;

$$a \neq 0 \wedge (\forall b : S \mid 0 \leq b \leq a : 0 = b \vee b = a).$$

We then write  $atom(a)$ .

## Properties of atoms of a boolean algebra

$$(18.64) \quad atom.a \Rightarrow a \otimes b = 0 \vee a \otimes b = a$$

$$(18.65) \quad atom.a \wedge atom.b \wedge a \neq b \Rightarrow a \otimes b = 0$$

$$(18.66) \quad (\forall a \mid atom.a : a \otimes b = 0) \Rightarrow b = 0$$



- Let  $B$  be a non-empty set. Clearly  $B = \bigcup \{\{b\} \mid b \in B\}$ , or equivalently  $B = \bigcup_{b \in B} \{b\}$ .
- For example,  $\{a, b, c, d\} = \{a\} \cup \{b\} \cup \{c\} \cup \{d\}$ ,  
 $\mathbb{N} = \{1\} \cup \{2\} \cup \{3\} \dots$ , etc.
- Hence each set is a union of all singletons that it contains.
- Singletons are atoms of the Boolean algebra of sets, so each set is a union of all its atoms.
- We can extend this property to all finite Boolean algebras.

## Theorem

Any element  $b$  of a finite Boolean different than 0 is equal to its "sum" of atoms, i.e. for every  $b \in S$

$$b = \bigoplus(a \mid \text{atom}(a) \wedge a \oplus b \neq 0 : a).$$

## Definition

For every  $b \in B$ , we define  $y(b) = \bigoplus(a \mid \text{atom}(a) \wedge a \oplus b \neq 0 : a)$ .

## Lemma

$$b \otimes y(b) = y(b)$$

## Proof.

Clearly  $y(b) \in B$ . We have:

$$\begin{aligned} b \otimes y(b) &= b \otimes \bigoplus(a \mid \text{atom}(a) \wedge a \oplus b \neq 0 : a) = \\ &\bigoplus(a \mid \text{atom}(a) \wedge a \oplus b \neq 0 : b \otimes a) = \\ &\bigoplus(a \mid \text{atom}(a) \wedge a \oplus b \neq 0 : a) = y(b), \text{ so we are done.} \quad \blacksquare \end{aligned}$$

## Lemma

$$b \otimes \sim y(b) = 0$$

## Proof.

It suffices to show that for each atom  $s$ ,  $(b \otimes \sim y(b)) \otimes a = 0$ . Then by property (18.66), page 37 of this Lecture Notes, we have  $b \otimes \sim y(b) = 0$ . We have to consider two cases;  $b \otimes a = 0$  and  $b \otimes a \neq 0$ .

Case  $b \otimes a = 0$ :  $(b \otimes \sim y(b)) \otimes a = 0 \otimes \sim y(b) = 0$ .

Case  $b \otimes a \neq 0$ : We have  $(b \otimes \sim y(b)) \otimes a = b \otimes a \otimes \sim \bigoplus (c \mid \text{atom}(c) \wedge b \oplus b \neq 0 : c) = b \otimes a \otimes (\sim a) \otimes \bigotimes (c \mid c \neq a \wedge \text{atom}(c) \wedge b \oplus b \neq 0 : \sim c) = b \otimes 0 \otimes \bigotimes (c \mid c \neq a \wedge \text{atom}(c) \wedge b \oplus b \neq 0 : \sim c) = 0. \quad \blacksquare$

## Theorem

*Any element  $b$  of a finite Boolean different than 0 is equal to its “sum” of atoms, i.e. for every  $b \in S$*

$$b = y(b) = \bigoplus (a \mid \text{atom}(a) \wedge a \oplus b \neq 0 : a).$$

*Moreover  $y(b) = \bigoplus (a \mid \text{atom}(a) \wedge a \leq b : a)$ , and this representation is unique.*

## Proof.

We have proven:  $b \otimes y(b) = y(b)$  and  $b \otimes \sim y(b) = 0$ .

We will show that  $b = y(b)$ .

$$\begin{aligned} \text{We have } b &= b \otimes 1 = b \otimes (y(b) \oplus (\sim y(b))) = \\ &= (b \otimes y(b)) \oplus (b \otimes (\sim y(b))) = y(b) \oplus 0 = y(b). \end{aligned}$$

Clearly  $\text{atom}(a) \wedge a \oplus b \neq 0 \equiv \text{atom}(a) \wedge a \leq b$ , so  $y(b) = \bigoplus (a \mid \text{atom}(a) \wedge a \leq b : a)$ , which immediately implies uniqueness. ■

## Theorem

*A boolean algebra with  $n$  atoms has  $2^n$  elements.*

## Proof.

Since if  $|X| = n$  then  $|\mathcal{P}(X)| = 2^n$ . ■

## Theorem

A finite boolean algebra  $A = (S, \oplus, \otimes, \sim, 0, 1)$  with  $n$  atoms is isomorphic to algebra  $\hat{A} = (\mathcal{P}(\{1, \dots, n\}), \cup, \cap, \sim, \emptyset, \{1, \dots, n\})$ .

## Proof.

Let label the elements of  $S$  with  $\{1, \dots, n\}$ , i.e. assume  $a_1, \dots, a_n$  are atoms of  $S$ . Clearly  $A$  and  $\hat{A}$  have the same signature. Define the function

$h: S \rightarrow \mathcal{P}(\{1, \dots, n\})$  by, for each  $b \in S$ ,

$$h(b) = \{i \mid \text{atom}(a_i) \wedge a_i \leq b\}.$$

Since the representation of each element of  $A$  as a sum is unique, this mapping is well defined, one-to-one, and onto. Clearly  $h(0) = \emptyset$  and  $h(1) = \{1, \dots, n\}$ , so  $h$  preserves constants.

Since  $b = \bigoplus (i \mid \text{atom}(a_i) \wedge a_i \leq b)$  we actually have

$h(\bigoplus (i \mid \text{atom}(a_i) \wedge a_i \leq b)) = \{i \mid \text{atom}(a_i) \wedge a_i \leq b\}$ ! Hence

$$h(b \oplus c) = h(\bigoplus (i \mid \text{atom}(a_i) \wedge a_i \leq b) \oplus \bigoplus (i \mid \text{atom}(a_i) \wedge a_i \leq c)) =$$

$$h(\bigoplus (i \mid (\text{atom}(a_i) \wedge a_i \leq b) \vee (\text{atom}(a_i) \wedge a_i \leq c))) =$$

$$\{i \mid (\text{atom}(a_i) \wedge a_i \leq b) \vee (\text{atom}(a_i) \wedge a_i \leq c)\} =$$

$$\{i \mid \text{atom}(a_i) \wedge a_i \leq b\} \cup \{i \mid \text{atom}(a_i) \wedge a_i \leq c\} = h(b) \cup h(c).$$

Similarly we can show  $h(b \otimes c) = h(b) \cap h(c)$ .

Define  $c = \sim b$ . Hence  $b \oplus c = 1$  and  $b \otimes c = 0$ , i.e.

$$h(b \oplus c) = h(b) \oplus h(c) = h(b) \cup h(c) = \{1, \dots, n\} \text{ and}$$

$$h(b \otimes c) = h(c) \otimes h(c) = h(b) \cap h(c) = \emptyset. \text{ Hence } h(b) = \{1, \dots, n\} \setminus h(c), \text{ i.e.}$$

$h(b) = \sim h(c)$ . This means  $h$  is an isomorphism. ■