

MATH 465 - INTRODUCTION TO COMBINATORICS

LECTURE 10

1. TRIANGULATIONS OF AN $(n + 2)$ -GON

Let \mathcal{P}_{n+2} be a convex $(n + 2)$ -gon. We consider the edge as a convex 2-gon which has $C_0 = 1$ triangulation (the empty one). For example, the pentagon has $C_3 = 5$ triangulations (Figure [1](#)).

Theorem 1.1. *The number of triangulations of \mathcal{P}_{n+2} is C_n .*

Proof. Let e be an edge of \mathcal{P}_{n+2} and let T be a triangulation of \mathcal{P}_{n+2} . If we remove e , we obtain two triangulated polygons \mathcal{Q}_1 and \mathcal{Q}_2 (in clockwise order starting from e) with a common vertex (see Figure [2](#)). If \mathcal{Q}_1 is a $(k + 2)$ -gon and \mathcal{Q}_2 is an $(l + 2)$ -gon, we must have $k + l = n - 1$. This procedure is clearly invertible, so if h_n denotes the number of triangulations of an $(n + 2)$ -gon, we obtain

$$h_n = \sum_{k+l=n-1} h_k h_l, \quad h_0 = 1,$$

the recurrence for Catalan numbers. □

2. PLANE TREES

A *plane tree* or *ordered rooted tree* P on $n + 1$ vertices is a set with $n + 1$ elements, called *vertices*, defined recursively as follows:

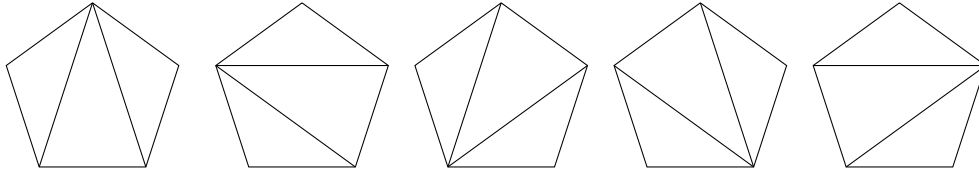


FIGURE 1. Triangulations of a pentagon.

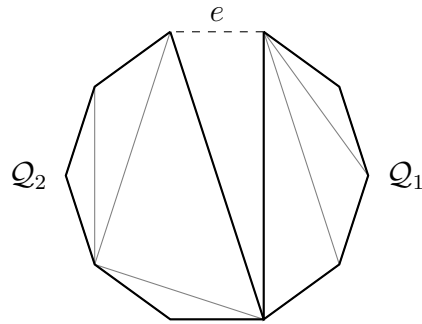


FIGURE 2. Removing an edge from a triangulation of \mathcal{P}_{n+2} .

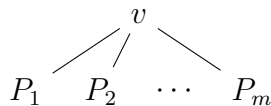


FIGURE 3. Drawing a plane tree.

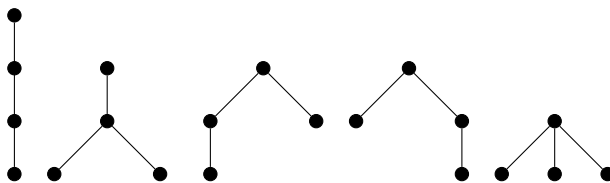


FIGURE 4. Plane trees with 4 vertices.

- (1) P has a distinguished vertex v called the *root*. In particular, $P \neq \emptyset$.
- (2) If $n = 0$, then $P = \{v\}$. If $n > 0$, P has a sequence (P_1, \dots, P_m) of sub-plane trees, where each $P_i, 1 \leq i \leq m$ is a plane tree with $\sum_{i=1}^m \#P_i = n$.

The roots of P_i are called the *children* of v , and P is drawn by placing the root v at the top, writing P_1, \dots, P_m from left to right below v and drawing edges from v to each of its children (Figure 3).

Lemma 2.1. *A plane tree with $n + 1$ vertices has n edges.*

Proof. We prove this by strong induction. If $n = 0$, then $P = \{v\}$ and so P has 0 edges. Suppose the statement is true for all plane trees with fewer than $n + 1$ vertices. Suppose P is a plane tree with $n + 1$ vertices. Let P_1, \dots, P_m denote the set of subtrees of the root of P . Each P_i has fewer than $n + 1$ vertices, so by the induction hypothesis, P_i has $\#P_i - 1$ edges. Therefore, P has $\sum_{i=1}^m (\#P_i - 1) + m = \#P - 1$ edges. \square

Theorem 2.2. *The number of plane trees with $n + 1$ vertices is C_n .*

Proof. We will define a bijection between plane trees and ballot sequences. We first define an order called *preorder* on the vertices of a plane tree P . The definition is recursive. Let v be the root with subtrees P_1, \dots, P_m . Define

$$\text{ord}(P) = v, \text{ord}(P_1), \dots, \text{ord}(P_m).$$

Traverse P in preorder, and each time we take a down step (i.e., away from the root), record a 1 and each time we take an up step (i.e., towards the root), record a -1 . For the five plane trees in Figure 4, the respective sequences are

$$111 ---, \quad 11 - 1 ---, \quad 11 - -1 -, \quad 1 - 11 ---, \quad 1 - 1 - 1 -.$$

By Lemma 2.1, there are n edges. Since each edge is traversed twice, the length of the sequence is $2n$. At any instant, the difference between the number of 1s and -1 s is the depth below the root, so the ballot sequence condition is satisfied. This is a bijection between plane trees with $n + 1$ vertices and ballot sequences of length $2n$. \square

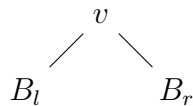


FIGURE 5. Drawing a binary tree.

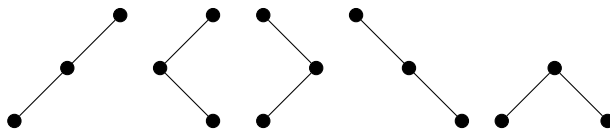


FIGURE 6. Binary trees with 3 vertices.

3. BINARY TREES

A *binary tree* on n vertices is a set B with $\#B = n$ defined recursively by the following rule:

- (1) If $n = 0$, then $B = \emptyset$.
- (2) If $n \neq 0$, then it has a root v , a left subtree B_l and a right subtree B_r with $\#B_l + \#B_r = n - 1$.

The root of B_l (resp., of B_r) is called the left (resp., right) child of v . We draw B as shown in Figure 5.

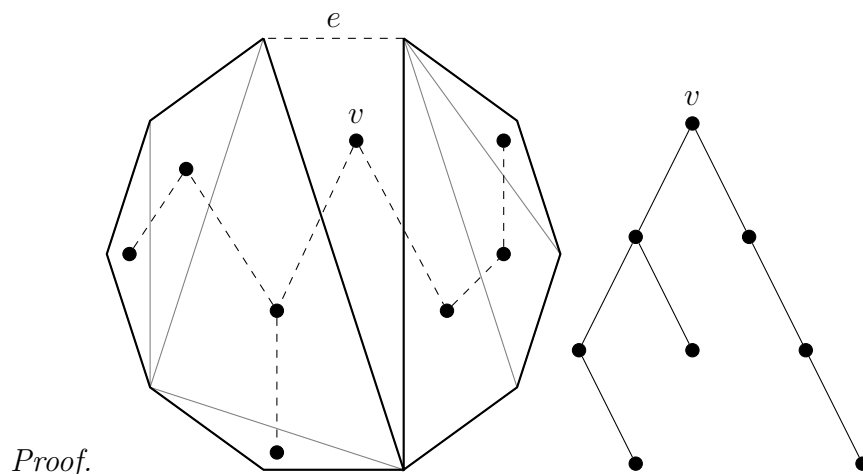
Theorem 3.1. *There are C_n binary trees with n vertices.*

Proof. There is a bijection between binary trees with $n \geq 1$ vertices and pairs (B_l, B_r) of left and right subtrees with k, l vertices respectively such that $k + l = n - 1$. Therefore, if h_n is the number of binary trees with n vertices, we have

$$h_n = \sum_{k+l=n-1} h_k h_l, \quad h_0 = 1,$$

the Catalan recurrence. □

Theorem 3.2. *Binary trees with n vertices are in bijection with triangulations of $(n + 2)$ -gons.*



□

4. BRACKETINGS

A bracketing w of $n + 1$ *s ($n \geq 0$) is defined recursively as follows:

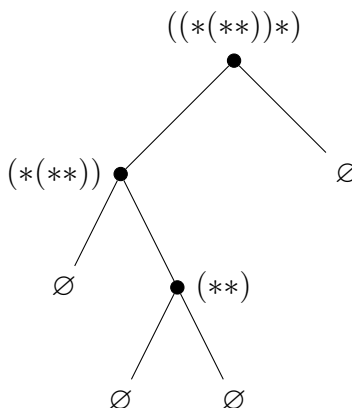
- (1) If $n = 0$, then $w = *$.
- (2) If $n > 0$, then $w = (st)$ where s and t are bracketings of $k + 1$ and $l + 1$ *s respectively for some k, l such that $k + l = n - 1$.

Theorem 4.1. *The number bracketings of a sequence of $n + 1$ *s is C_n .*

$$(*(*(**))), \quad *((**)*), \quad ((**)(**)), \quad ((*(**))*), \quad (((**)*)*).$$

Proof. The parse tree construction provides a bijection between bracketings of $n + 1$ *s and binary trees with n vertices. Given a binary tree B , we define the bracketing w_B recursively as follows:

- (1) If $B = \emptyset$, let $w_B = *$.
- (2) If v is the root with left and right subtrees B_l and B_r respectively, let $w_B = (w_{B_l}w_{B_r})$.

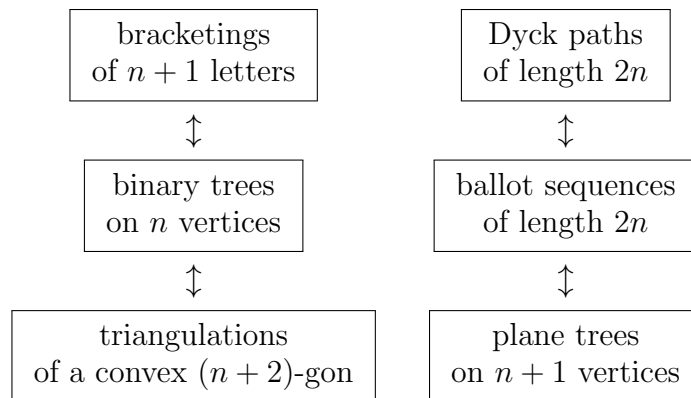


This procedure is invertible. Given a bracketing w , define the binary tree B_w recursively as follows: If $n = 0$, let $B_w = \emptyset$. Otherwise w is a product (st) . Define the right and left subtrees of B_w to be B_s and B_t respectively. For example, if $w = (**)$, we have $s = *$ and $t = *$. Therefore, $B_s = B_t = \emptyset$, so B is the tree with one vertex.

□

5. BIJECTIONS BETWEEN CATALAN OBJECTS

We constructed bijections between the following families of combinatorial objects counted by the Catalan numbers C_n :



Is there a bijection linking the families on the left and on the right?

Theorem 5.1. *There is a bijection between bracketings of $n + 1$ letters and ballot sequences of length $2n$.*

Exercise. Prove that the following construction is a bijection. Remove all $)$'s and replace all $($'s with a 1 and each letter with -1 except for the last letter. For example, for $n = 3$, we have

$$\begin{array}{ccccc}
 (*(*(**))) & (*((**)*)) & ((**)(**)) & ((*(**))*) & (((**)*))* \\
 1 - 1 - 1 - & 1 - 11 - - & 11 - -1 - & 11 - 1 - - & 111 - - - .
 \end{array}$$

□