

# RATIONALMAPS, A PACKAGE FOR MACAULAY2

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ABSTRACT. This paper describes the `RationalMaps` package for Macaulay2. This package provides functionality for computing several aspects of rational maps.

## 1. INTRODUCTION

This package aims to compute a number of things about rational maps between varieties. In particular, this package will compute

- The base locus of a rational map.
- Whether a rational map is birational.
- The inverse of a birational map.
- Whether a map is a closed embedding.
- And more!

Our functions have numerous options which allow them to run much more quickly in certain examples if configured correctly. The `Verbose` option gives hints as to the best way to apply these.

A rational map  $\mathfrak{F} : X \subseteq \mathbb{P}^n \dashrightarrow Y \subseteq \mathbb{P}^m$  between projective varieties is presented by  $m + 1$  forms  $\mathbf{f} = \{f_0, \dots, f_m\}$  of the same degree in the coordinate ring of  $X$ , denoted by  $R$ . The idea of looking at the syzygies of the forms  $\mathbf{f}$  to detect the geometric properties of  $\mathfrak{F}$  goes back at least to [HKS92] in the case where  $X = \mathbb{P}^n$ ,  $Y = \mathbb{P}^m$  and  $m = n$  (see also [ST69]). In [RS01] this method was developed by Russo and Simis to handle the case  $X = \mathbb{P}^n$  and  $m \geq n$ . Simis pushed the method further to the study of general rational maps between two integral projective schemes in arbitrary characteristic by an extended ideal-theoretic method emphasizing the role of the Rees algebra associated to the ideal generated by  $\mathbf{f}$  [Sim04]. Recently, Doria, Hassanzadeh, and Simis applied these Rees algebra techniques to study the birationality of  $\mathfrak{F}$  [DHS12]. Our core functions, in particular the functions related to computing inverse maps, rely heavily on this work.

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## 2. BASE LOCI

We begin with the problem of computing the base locus of a map to projective space. Let  $X$  be a projective variety over any field  $k$  and let  $\mathfrak{F} : X \rightarrow \mathbb{P}_k^m$  be a rational map from  $X$  to projective space. Then we can choose some representative  $(f_0, \dots, f_m)$  of  $\mathfrak{F}$ , where each  $f_i$  is the  $i^{\text{th}}$  coordinate of  $\mathfrak{F}$ . A priori, each  $f_i$  is in  $K = \text{frac } R$ , where  $R$  is the coordinate ring of  $X$ . However, we can get another representative of  $\mathfrak{F}$  by clearing denominators. (Note this does not enlarge the base locus of

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$\mathfrak{F}$  since  $\mathfrak{F}$  is undefined whenever the denominator of any of the  $f_i$  vanishes.) Thus we can assume that  $f_i \in R$  for all  $i$ , and that all the  $f_i$  are homogeneous of the same degree.

In this setting, one might naively think that the map  $\mathfrak{F}$  is undefined exactly when all of the  $f_i$  vanish, and thus the base locus is the vanishing set of the ideal  $(f_0, \dots, f_m)$ . However, this can yield a base locus that's too big. Indeed, to find the base locus of a rational map, we must consider all possible representatives of the map and find where none of them are defined. To do this, we use the following result.

**Proposition 2.1.** [Sim04, Proposition 1.1] *Let  $\mathfrak{F} : X \dashrightarrow \mathbb{P}^m$  be a rational map and let  $\mathbf{f} = \{f_0, \dots, f_m\}$  be a representative of  $\mathfrak{F}$  with  $f_i \in R$  homogeneous of degree  $d$  for all  $i$ . Set  $I = (f_0, \dots, f_m)$ . Then the set of such representatives of  $\mathfrak{F}$  corresponds bijectively to the homogeneous vectors in the rank 1 graded  $R$ -module  $\text{Hom}_R(I, R) \cong (R :_K I)$ .*

The bijection comes from multiplying our fixed representative  $\mathbf{f}$  of  $\mathfrak{F}$  by  $h \in (R :_K I)$ . Now, in the setting of Proposition 2.1, let

$$\bigoplus_s R(-d_s) \xrightarrow{\varphi} R(-d)^{m+1} \xrightarrow{[f_0, \dots, f_m]} I \rightarrow 0$$

be a free resolution of  $I$ . Then we get

$$0 \rightarrow \text{Hom}_R(I, R) \rightarrow (R(-d)^{m+1})^\vee \xrightarrow{\varphi^t} \left( \bigoplus_s R(d_s) \right)^\vee$$

where  $\varphi^t$  is the transpose of  $\varphi$  and  $R^\vee$  is the dual module of  $R$ . Thus, we get that  $\text{Hom}_R(I, R) \cong \ker \varphi^t$ , and so each representative of  $\mathfrak{F}$  corresponds to a vector in  $\ker \varphi^t$ . The correspondence takes a representative  $(hf_0, \dots, hf_m)$  to the map that multiplies vectors in  $R^{m+1}$  by  $[hf_0, \dots, hf_m]$  on the left.

The base locus of  $\mathfrak{F}$  is the intersection of the sets  $V(f_0^i, \dots, f_m^i)$  as  $\mathbf{f}^i = (f_0^i, \dots, f_m^i)$  ranges over all the representatives of  $\mathfrak{F}$ . The above implies that this is the same as the intersection of the sets  $V(w_0^i, \dots, w_m^i)$  as  $\mathbf{w}^i = (w_0^i, \dots, w_m^i)$  ranges over the vectors in  $\ker \varphi^t$ . Now, given any  $a, f, g \in R$ , we have  $V(af) \supseteq V(f)$  and  $V(f+g) \supseteq V(f) \cap V(g)$ . Thus, it's enough to take a generating set  $\mathbf{w}^1, \dots, \mathbf{w}^n$  of  $\ker \varphi^t$  and take the intersection over this generating set.

The base locus of  $\mathfrak{F}$  is then the variety cut out by the ideal generated by all the entries of all of the  $\mathbf{w}^i$ . Our function `baseLocusOfMap` returns this ideal.

```
i1 : loadPackage "RationalMaps";
i2 : R = QQ[x,y,z];
i3 : f = {x^2*y, x^2*z, x*y*z};
i4 : baseLocusOfMap(f);
o4 = ideal (y*z, x*z, x*y)
o4 : Ideal of R
```

If the `SaturateOutput` option is set `true`, our function will return the saturation of this ideal.

### 3. BIRATIONALITY AND INVERSE MAPS

Again, a rational map  $\mathfrak{F} : X \subseteq \mathbb{P}^n \dashrightarrow Y \subseteq \mathbb{P}^m$  between projective spaces is defined by  $m+1$  forms  $\mathbf{f} = \{f_0, \dots, f_m\}$  of the same degree in the coordinate ring of  $X$ , denoted by  $R$ .  $R$  is a standard graded ring in  $n+1$  variables. Here we assumed the varieties over a field  $k$  and  $\dim R \geq 1$ . The idea to find a ring theoretic criterion for birationality and on top of that to find the inverse of a rational map, is to study the Rees algebra of the ideal  $I = (\mathbf{f})$  in  $R$ . So that let  $R \simeq k[x_0, \dots, x_n] = k[\mathbf{X}]/\mathfrak{a}$  with  $k[\mathbf{X}] = k[X_0, \dots, X_n]$  and  $\mathfrak{a}$  a homogeneous ideal. The Rees algebra is defined by the polynomial relations among  $\{f_0, \dots, f_m\}$  in  $R$ . To this end, we consider the polynomial extension  $R[\mathbf{Y}] = R[Y_0, \dots, Y_m]$ . To keep track of the variables by degrees, we set the standard bigrading  $\deg(X_i) = (1, 0)$  and  $\deg(Y_j) = (0, 1)$ . Mapping  $Y_j \mapsto f_j t$  yields a presentation  $R[\mathbf{Y}]/\mathcal{J} \simeq \mathcal{R}_R((\mathbf{f}))$ ,

with  $\mathcal{J}$  a bihomogeneous *presentation ideal*.  $\mathcal{J}$  is a bigraded ideal depends only on the rational map defined by  $\mathbf{f}$  and not on this particular representative.

$$\mathcal{J} = \bigoplus_{(p,q) \in \mathbb{N}^2} \mathcal{J}_{(p,q)},$$

where  $\mathcal{J}_{(p,q)}$  denotes the  $k$ -vector space of forms of bidegree  $(p, q)$ . Every pieces of this ideal contains information about the rational map. For example  $\mathcal{J}_{0,*}$  determines the dimension of the image of the map. For birationality, the following bihomogeneous piece is important:

$$\mathcal{J}_{1,*} := \bigoplus_{r \in \mathbb{N}} \mathcal{J}_{1,q}$$

with  $\mathcal{J}_{1,q}$  denoting the bigraded piece of  $\mathcal{J}$  spanned by the forms of bidegree  $(1, q)$  for all  $q \geq 0$ . Now, a form of bidegree  $(1, *)$  can be written as  $\sum_{i=0}^n Q_i(\mathbf{Y}) x_i$ , for suitable homogeneous  $Q_i(\mathbf{Y}) \in R[\mathbf{Y}]$  of the same degree.

One then goes to construct a matrix which can measure the birationality of the map. The first step is to lift the polynomials  $Q_i(\mathbf{Y}) \in R[\mathbf{Y}]$  into  $k[\mathbf{X}, \mathbf{Y}]$ . Since the  $\{y_0, \dots, y_m\}$  are indeterminates over  $R$ , each pair of such representations of the same form gives a syzygy of  $\{x_0, \dots, x_n\}$  with coefficients in  $k$ . This is where one must take into attention whether  $X \subseteq \mathbb{P}^n$  is minimally embedded or not. To measure this one can easily check the vector space dimension of  $\mathfrak{a}_1$ , the degree-1 part of  $\mathfrak{a}$ ; if it is zero then  $X \subseteq \mathbb{P}^n$  is non-degenerated.

Next, one can pick a minimal set of generators of the ideal  $(\mathcal{J}_{1,*})$  consisting of a finite number of forms of bidegree  $(1, q)$ , for various  $q$ 's. Let's assume  $X \subseteq \mathbb{P}^n$  is non-degenerated. Let  $\{P_1, \dots, P_s\} \subset k[\mathbf{X}, \mathbf{Y}]$  denote liftings of these bifurms, consider the Jacobian matrix of the polynomials  $\{P_1, \dots, P_s\}$  with respect to  $\{x_0, \dots, x_n\}$ . This is a matrix with entries in  $k[\mathbf{Y}]$ . Write  $\psi$  for the corresponding matrix over  $S = k[\mathbf{Y}]/\mathfrak{b}$ , the coordinate ring of  $Y$ . This matrix is called the *weak Jacobian dual matrix* associated to the given set of generators of  $(\mathcal{J}_{1,*})$ . Note that a weak Jacobian matrix  $\psi$  is not uniquely defined due to the lack of uniqueness in the expression of an individual form and to the choice of bihomogeneous generators. However, it is shown in [DHS12, Lemma 2.13] that if the weak Jacobian matrix associated to one set of bihomogeneous minimal generators of  $(\mathcal{J}_{1,*})$  has rank over  $S$  then the weak Jacobian matrix associated to any other set of bihomogeneous minimal generators of  $(\mathcal{J}_{1,*})$  has rank over  $S$  and the two ranks coincide.

The following criterion is [DHS12, Theorem 2.18 ]. In the package we consider only the cases where  $\mathbf{X}$  is irreducible i.e.  $R$  is a domain.

**Theorem 3.1.** *Let  $X \subseteq \mathbb{P}^n$  be non-degenerate. Then  $\mathfrak{F}$  is birational onto  $\mathbf{Y}$  if and only if  $\text{rank}(\psi) = \text{edim}(R) - 1 (= n)$ . Moreover*

- (i) *We get a representative for the inverse of  $\mathfrak{F}$  by taking the coordinates of any homogeneous vector of positive degree in the (rank one) null space of  $\psi$  over  $S$  for which these coordinates generate an ideal containing a regular element.*
- (ii) *If, further,  $R$  is a domain, the representative of  $\mathfrak{F}$  in (i) can be taken to be the set of the (ordered, signed)  $(\text{edim}(R) - 1)$ -minors of an arbitrary  $(\text{edim}(R) - 1) \times \text{edim}(R)$  submatrix of  $\psi$  of rank  $\text{edim}(R) - 1$ .*

As expected, the most expensive part of applying this theorem is computing the Rees ideal  $\mathcal{J}$ . In the package `RationalMaps` we use `ReesStrategy` to compute the Rees equations. The algorithm is the standard elimination technique. However we do not use the `ReesAlgebra` package, since verifying birationality according to Theorem 3.1 only requires computing a small part of the Rees ideal, namely elements of first-degree 1. This idea is applied in the `SimisStrategy`. More precisely, if the given map  $\mathfrak{F}$  is birational, then the Jacobian dual rank will attain its maximum

value of  $\text{edim}(R) - 1$  after computing the Rees equations up to degree  $(1, N)$  for  $N$  sufficiently large. This allows us to compute the inverse map. The downside of `SimisStrategy` is that if  $\mathfrak{F}$  is not birational, the desired number  $N$  cannot be found and the process never terminates. To provide a definitive answer for birationality, we use `HybridStrategy`, which is a hybrid of `ReesStrategy` and `SimisStrategy`. The default strategy is `HybridStrategy`.

`HybridLimit` is an option to switch `SimisStrategy` to `ReesStrategy`, if the computations up to degree  $(1, \text{HybridLimit})$  do not lead to  $\text{rank}(\psi) = \text{edim}(R) - 1$ . The default value for `HybridLimit` is 15. The change from `SimisStrategy` to `ReesStrategy` is done in such a way that the generators of the Rees ideal computed in the `SimisStrategy` phase are not lost; the program computes other generators of the Rees ideal while keeping the generators it found before attaining `HybridLimit`.

There is yet another method for computing the Rees ideal called `SaturationStrategy`. In this option the whole Rees ideal is computed by saturating the defining ideal of the symmetric algebra with respect to a non-zero element in  $R$  (we assume  $R$  to be a domain). This strategy appears to be slower in some examples, though one might be able improve this option in the future by stopping the computation of the saturation at a certain step.

Computing inverse maps is the most important functionality of this package, and is done by the function `inverseOfMap`. According to Theorem 3.1, there are two ways to compute the inverse of a map: (1) by finding any syzygy of the Jacobian dual matrix, and (2) by finding a sub-matrix of  $\psi$  of rank  $\text{edim}(R) - 1$ . Each way has its own benefits. Method (1) is quite fast in many cases, however method (2) is very useful if the rank of the Jacobian dual matrix  $\psi$  is relatively small compared to the degrees of the entries of  $\psi$ . Our function `inverseOfMap` starts by using the second method and later switches to the first method if the second method didn't work. The timing of this transition from the first method to the second method is controlled by the option `MinorsCount`. Setting `MinorsCount` to zero will mean that no minors are checked and the inverse map is computed just by looking at the syzygies of  $\psi$ . If `MinorsCount` is left as null (the default value), the program will try to make an educated guess as to how big to set this option, depending on varieties the user is working with.

In addition, to improve the speed of the function `inverseOfMap`, we have two other options, `AssumeDominant` and `CheckBirational`. If `AssumeDominant` is set to be `true`, then `inverseOfMap` assumes that the map from  $X$  to  $Y$  is dominant and does not compute the image of the map; this is time consuming in certain cases as it computes the kernel of a ring map. However, this function goes through a call to `idealOfImageOfMap` which first checks whether the ring map is injective (at least if the target is a polynomial ring) using the method described in [Sim03, Proposition 1.1]. Similarly, if `CheckBirational` set `false`, `inverseOfMap` will not check birationality although it still computes the Jacobian dual matrix. The option `QuickRank` is available to many functions. At various points, the rank of a matrix is computed, and sometimes it is faster to compute the rank of an interesting looking submatrix (using the tools of the package `FastLinAlg`). Turning `QuickRank` off will make showing that certain maps are birational slower, but will make showing that certain maps are *not* birational faster.

In general, as long as `Verbose` is `true`, the function will make suggestions as to how to run it more quickly. For example.

```
i1 : loadPackage "RationalMaps";
i2 : Q=QQ[x,y,z,t,u];
i3 : phi=map(Q,Q,matrix{{x^5,y*x^4,z*x^4+y^5,t*x^4+z^5,u*x^4+t^5}});
o3 : RingMap Q <--- Q
i4 : time inverseOfMap(phi, AssumeDominant=>true,CheckBirational=>false);
Starting inverseOfMapSimis(SimisStrategy or HybridStrategy)
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree {1, 1}.
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree {1, 2}.
getSubmatrixOfRank: Trying to find a submatrix of rank at least: 4 with attempts = 4
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree {1, 4}.
getSubmatrixOfRank: Trying to find a submatrix of rank at least: 4 with attempts = 4
```

```

inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree {1, 7}.
getSubmatrixOfRank: Trying to find a submatrix of rank at least: 4 with attempts = 4
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree {1, 11}.
getSubmatrixOfRank: Trying to find a submatrix of rank at least: 4 with attempts = 4
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree {1, 16}.
inverseOfMapSimis: We give up. Using the previous computations, we compute the whole
Groebner basis of the rees ideal. Increase HybridLimit and rerun to avoid this.
inverseOfMapSimis: Found Jacobian dual matrix (or a weak form of it), it has 5 columns and about 20 rows.
inverseOfMapSimis: Looking for a nonzero minor.
If this fails, you may increase the attempts with MinorsCount => #
getSubmatrixOfRank: Trying to find a submatrix of rank at least: 4 with attempts = 20
getSubmatrixOfRank: found one, in 3 attempts
inverseOfMapSimis: We found a nonzero minor.
-- used 0.703125 seconds

o4 : RingMap Q <--- Q

```

#### 4. EMBEDDINGS

Our package also checks whether a rational map  $\mathfrak{F} : X \rightarrow Y$  is a closed embedding. The strategy is quite simple.

- (a) We first check whether  $\mathfrak{F}$  is regular (by checking if its base locus is empty).
- (b) We next invert the map (if possible).
- (c) Finally, we check if the inverse map is also regular.

If all three conditions are met, then the map is a closed embedding and the function returns `true`. Otherwise `isEmbedding` returns false. The following exam illustrates this. We begin with an example where we take a plane quartic, choose a point  $Q$  on it and take map associated to the divisor  $12Q$ . This map must be an embedding, which we now verify.

```

i1 : needsPackage "Divisor"; --used to quickly define a map
i2 : C = ZZ/101[x,y,z]/(x^4+x^2*y*z+y^4+z^3*x);
i3 : Q = ideal(y,x+z);
o3 : Ideal of C
i4 : f2 = mapToProjectiveSpace(12*divisor(Q));
      ZZ
o4 : RingMap C <--- ---[YY , YY , YY , YY , YY , YY , YY , YY , YY , YY ]
      101 1 2 3 4 5 6 7 8 9 10
i5 : loadPackage "RationalMaps";
i6 : time isEmbedding(f2)
isEmbedding: About to find the image of the map.
If you know the image, you may want to specify that and set AssumeDominant=>true option if this is slow.
isEmbedding: Checking to see if the map is a regular map
isEmbedding: computing the inverse map
Starting inverseOfMapSimis(SimisStrategy or HybridStrategy)
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree {1, 1}.
getSubmatrixOfRank: Trying to find a submatrix of rank at least: 2 with attempts = 4
getSubmatrixOfRank: found one, in 1 attempts
inverseOfMapSimis: We computed enough of the Groebner basis.
inverseOfMapSimis: Found Jacobian dual matrix (or a weak form of it), it has 3 columns and about 17 rows.
inverseOfMapSimis: MinorsCount => 0, so we now compute syzygies instead.
If this doesn't terminate quickly, you may want to try increasing the option MinorsCount.
isEmbedding: checking if the inverse map is a regular map
-- used 0.53125 seconds
o6 = true

```

Notice that `MinorsCount => 0` by default for `isEmbedding`. This is because the expressions defining the inverse map obtained from an appropriate minor frequently are more complicated than the expressions for the inverse map obtained via the syzygies. Complicated expressions can sometimes slow down the checking of whether the inverse map is regular.

## 5. FUNCTIONALITY OVERLAP WITH OTHER PACKAGES

We note that our package has some overlaps in functionality with other packages.

While the `Parametrization` package [Boe10] focuses mostly on curves, it also includes a function called `invertBirationalMap` which has the same functionality as `inverseOfMap`. On the other hand, these two functions were implemented somewhat differently and so sometimes one function can be substantially faster than the other.

The package `Cremona` [Sta16] focuses on very fast probabilistic computation in general cases and very fast deterministic computation for special kinds of maps from projective space. In particular, in `Cremona`,

- `isBirational` gives a probabilistic answer to the question of whether a map between varieties is birational. Furthermore, if the source is projective space, then `degreeOfRationalMap` with `MathMode=>true` can give a deterministic answer that is frequently faster than what our package can provide with `isBirationalMap`.
- `inverseMap` gives a very fast computation of the inverse of a birational map if the source is projective space and the map has maximal linear rank. If you pass this function a map not from projective space, then it calls a modified, faster version of `invertBirationalMap` originally from `Parametrization`. Even in some cases with maximal linear rank, our `inverseOfMap` function appears to be quite competitive however.

## 6. COMMENTS AND COMPARISONS ON FUNCTION SPEEDS

We begin with a comparison using examples with maximal linear rank where `Cremona` excels. These examples were run using version 4.3 of `Cremona` and version 0.3 of `RationalMaps`.

Indeed, in this example (taken from `Cremona`'s documentation), `Cremona` is faster.

```
i1 : loadPackage "Cremona"; loadPackage "RationalMaps";
i3 : ringP20=QQ[t_0..t_20];
i4 : phi=map(ringP20,ringP20,{t_10*t_15-t_9*t_16+t_6*t_20,t_10*t_14-t_8*t_16+t_5*t_20,t_9*t_14-t_8*t_15+t_4*t_20,
t_6*t_14-t_5*t_15+t_4*t_16,t_11*t_13-t_16*t_17+t_15*t_18-t_14*t_19+t_12*t_20,t_3*t_13-t_10*t_17+t_9*t_18-t_8*t_19
+t_7*t_20,t_10*t_12-t_2*t_13-t_7*t_16-t_6*t_18+t_5*t_19,t_9*t_12-t_1*t_13-t_7*t_15-t_6*t_17+t_4*t_19,t_8*t_12
-t_0*t_13-t_7*t_14-t_5*t_17+t_4*t_18,t_10*t_11-t_3*t_16+t_2*t_20,t_9*t_11-t_3*t_15+t_1*t_20,t_8*t_11-t_3*t_14
+t_0*t_20,t_7*t_11-t_3*t_12+t_2*t_17-t_1*t_18+t_0*t_19,t_6*t_11-t_2*t_15+t_1*t_16,t_5*t_11-t_2*t_14+t_0*t_16,
t_4*t_11-t_1*t_14+t_0*t_15,t_6*t_8-t_5*t_9+t_4*t_10,t_3*t_6-t_2*t_9+t_1*t_10,t_3*t_5-t_2*t_8+t_0*t_10,t_3*t_4
-t_1*t_8+t_0*t_9,t_2*t_4-t_1*t_5+t_0*t_6});
i5 : time inverseOfMap(phi, AssumeDominant=>true, Verbose=>false)-- Function from "RationalMaps"
-- used 1.04688 seconds
i6 : time inverseMap phi -- Function from "Cremona"
-- used 0.046875 seconds
i7 : isSameMap(o5, o6) -- Function from "RationalMaps"
o7 = true
```

Increasing the option `MinorSize` makes `inverseOfMap` much slower.

However, sometimes the `RationalMaps` function is faster, even in examples with maximal linear rank (see the first example from `Cremona`'s documentation). We now include an example where the map does not have maximal linear rank.

```
i1 : loadPackage "Cremona"; loadPackage "RationalMaps";
i3 : Q=QQ[x,y,z,t,u];
i4 : phi=map(Q,Q,matrix{{x^5,y*x^4,z*x^4+y^5,t*x^4+z^5,u*x^4+t^5}});
o4 : RingMap Q <--- Q
i5 : (time g = inverseOfMap(phi, AssumeDominant=>true, CheckBirational=>false, Verbose=>false));
-- Function from "RationalMaps"
-- used 0.265625 seconds
i6 : (time f = inverseOfMap(phi, AssumeDominant=>true, CheckBirational=>false, Verbose=>false, MinorsCount=>0));
-- used 125.141 seconds
i7 : (time h = inverseMap(phi)); -- Function from "Cremona"
-- used 141.25 seconds
i8 : isSameMap(f, h)
```

```
o8 = true
i9 : isSameMap(g, h)
o9 = true
```

In this final example, for instance, setting `MinorsCount=>0` makes `inverseOfMap` much slower – approximately the same speed as corresponding command from `Cremona`. The takeaway for the user should be that changing the options `Strategy`, `HybridLimit`, `MinorSize` and `QuickRank`, can make a large difference in performance.

We conclude with discussions of the limits of this package.

Work of O. Gabber shows that if  $f : \mathbb{P}^n \rightarrow \mathbb{P}^n$  is defined by forms of degree  $d$ , then its inverse can be defined by forms of degree  $d^{n-1}$ , [1]. This bound is sharp, as the map

$$(x_0^d : x_1 x_0^{d-1} : x_2 x_0^{d-1} - x_1^d : \dots : x_n x_0^{d-1} - x_{n-1}^d)$$

has inverse given by forms of degree  $d^{n-1}$ , see [1]. Thus we might expect that this family of maps would be good to explore to see the limits of `RationalMaps`. We ran these examples with the following code.

```
R = ZZ/101[x_0..x_n];
L = {x_0^d, x_1*x_0^(d-1)} | toList(apply(2..n, i -> (x_i*x_0^(d-1) + x_(i-1)^d)));
psi = map(R, R, L);
time inv = inverseOfMap(psi, AssumeDominant=>true, CheckBirational=>false, Verbose=>false);
```

When  $n = 3$  (we are working on  $\mathbb{P}^3$ ) here is a table showing the computation time for the inverse for various  $d$  at least one one system. The degrees are those we would expect in this example (when  $d = 100$ , the degree of the forms in the inverse is 10000). Note `Cremona` has nearly identical performance for these examples in  $\mathbb{P}^3$  ( $n = 3$ ), but has slower as we increase the dimension.

$d$	5	10	20	40	60	80	100
seconds	0.0625	0.09375	0.203125	2.10937	14.1406	43.1875	138.641

However, as the size of projective space increases, this becomes much slower. Here is a table when  $n = 4$ .

$d$	5	8	10	11	12	13	14	15
seconds	0.125	1.92187	11.8281	26.9375	53.5156	106.313	211.484	396.703

We conclude with a table when  $n = 5$ .

$d$	3	4	5	6
seconds	0.3125	9.45312	335.172	15257.9

Note the  $d = 6$  case took more than 4 hours.

Finally, we mention that Ana-Maria Castravet and Zhuang He communicated to us that they used `RationalMaps` to compute the inverse of a rational map from  $\mathbb{P}^3$  to  $\mathbb{P}^3$ . This map was given by 4 degree 13 forms, with 485, 467, 467, and 467 terms respectively. Computing the inverse of this map took several hours to compute, but it was successful. [1]

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