Homework Assignment 01

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1 Problem 1

Let e = (10001111) be the exponent. Illustrate the addition chains produced by each one of the following algorithms. Compute the length of each addition chain.

1.1 Binary method

i	e_i	Step 2a	Step 2b
6	0	$(M)^2 = M^2$	M^2
5	0	$(M^2)^2 = M^4$	M^4
4	0	$(M^4)^2 = M^8$	M^8
3	1	$(M^8)^2 = M^{16}$	$M^{16}.M = M^{17}$
2	1	$(M^{17})^2 = M^{34}$	$M^{34}.M = M^{35}$
1	1	$(M^{35})^2 = M^{70}$	$M^{70}.M = M^{71}$
0	1	$(M^{71})^2 = M^{142}$	$M^{142}.M = M^{143}$

Addition chain = 1 2 4 8 16 17 34 35 70 71 142 143, Length = 12.

1.2 m-ary method

1.2.1 For d=2

bits	w	M^w
00	0	1
01	1	M
10	2	$M.M = M^2$
11	3	$M^2.M = M^3$

i	F_i	Step 4a	Step 4b
2	00	$(M^2)^4 = M^8$	M^8
1	11	$(M^8)^4 = M^{32}$	$M^{32}.M^3 = M^{35}$
0	11	$(M^{35})^4 = M^{140}$	$M^{140}.M^3 = M^{143}$

Addition chain = 1 2 3 4 8 16 32 35 70 140 143, Length = 11.

1.2.2 For d=4

bits	w	M^w
0000	0	1
0001	1	M
0010	2	M.M = M^2
0011	3	$M^2.M = M^3$
0100	4	$M^3.M = M^4$
0101	5	$M^4.M = M^5$
0110	6	$M^5.M = M^6$
0111	7	$M^6.M = M^7$
1000	8	$M^7.M = M^8$
1001	9	$M^8.M = M^9$
1010	10	$M^9.M = M^{10}$
1011	11	$M^{10}.M = M^{11}$
1100	12	$M^{11}.M = M^{12}$
1101	13	$M^{12}.M = M^{13}$
1110	14	$M^{13}.M = M^{14}$
1111	15	$M^{14}.M = M^{15}$

i	F_i	Step 4a	Step 4b
0	1111	$(M^8)^{16} = M^{128}$	$M^{128}.M^{15} = M^{143}$

Addition chain = 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 32 64 128 143, Length = 20.

1.3 Factor Method

Compute: $M \rightarrow M^2$
$Assign : a = M^2$
$Compute : a \rightarrow a^2 \rightarrow a^4 \rightarrow a^5$
$Assign : b = a^5 = M^{10}$
$Compute : b.M \rightarrow M^{11}$
$Assign : c = M^{11}$
$Compute : c \rightarrow c^2 \rightarrow c^3$
$Assign : d = c^3 = (M^{11})^3$
$Compute : d \rightarrow d^2 \rightarrow d^4$
$Assign : e = d^4 = (M^{11})^{12}$
$Compute : e.c \rightarrow (M^{11})^{13} = M^{143}$
Addition chain = 1 2 4 8 10 11 22 33 66 132 143, Length = 11.

1.4 Power Tree Method

A tree of height 11 leads to 143 as its leaf node. Refer code `mk_tree.py` for the details. The path from the root is: 1 2 3 5 7 14 21 35 70 140 143. Addition chain is same as this path.
Addition chain = 1 2 3 5 7 14 21 35 70 140 143, Length = 11.

1.5 Canonical Recording for d=1

Truth table for canonical recording is as shown below:

c_i	e_{i+1}	e_i	c_{i+1}	f_i
0	0	0	0	0
0	0	1	0	1
0	1	0	0	0
0	1	1	1	$\overline{1}$
1	0	0	0	1
1	0	1	1	0
1	1	0	1	$\overline{1}$
1	1	1	1	0

Using truth table, recording for given number is:

c_i	e_{i+1}	e_i	c_{i+1}	f_i
0	1	1	1	$\overline{1}$
1	1	1	1	0
1	1	1	1	0
1	0	1	1	0
1	0	0	0	1
0	0	0	0	0
0	1	0	0	0
0	0	1	0	1

$1001000\overline{1} = 2^7 + 2^4 - 2^0 = 143$

Computing the exponent:

i	f_i	Step 2a	Step 2b
7	1	M	M
6	0	$(\mathrm{M})^2 = M^2$	M^2
5	0	$(\mathrm{M}^2)^2 = M^4$	M^4
4	1	$(\mathrm{M}^4)^2 = M^8$	$\mathrm{M}^8.M = M^9$
3	0	$(\mathrm{M}^9)^2 = M^{18}$	M^{18}
2	0	$(\mathrm{M}^{18})^2 = M^{36}$	M^{36}
1	0	$(\mathrm{M}^{36})^2 = M^{72}$	M^{72}
0	$\overline{1}$	$(\mathrm{M}^{72})^2 = M^{144}$	$\mathrm{M}^{144}.\mathrm{M}^{-1} = \mathrm{M}^{143}$

Addition chain: 1 2 4 8 9 18 36 72 144 143, Length = 10.

2 Illustrate the steps of the standard multiplication algorithm for computing c =a * b = 456 * 555

i	j	Step	(C,S)	Partial t
0	0	$t_0 + a_0 b_0 + C$	(0,*)	000000
		0 + 6*5 + 0	(3,0)	00000**0**
	1	$t_1 + a_1 b_0 + C$		
		0 + 5*5 + 3	(2,8)	0000**8**0
	2	$t_2 + a_2 b_0 + C$		
		0 + 4*5 + 2	(2,2)	000**2**80
				00**2**280
1	0	$t_1 + a_0 b_1 + C$	(0,*)	
		8 + 6*5 + 0	(3,8)	0022**8**0
	1	$t_2 + a_1 b_1 + C$		
		2 + 5*5 + 3	(3,0)	002**0**80
	2	$t_3 + a_2 b_1 + C$		
		2 + 4*5 + 3	(2,5)	00**5**080
				0**2**5080
2	0	$t_2 + a_0 b_2 + C$	(0,*)	
		0 + 6*5 + 0	(3,0)	025**0**80
	1	$t_3 + a_1 b_2 + C$		
		5 + 5*5 + 3	(3,3)	02**3**080
	2	$t_4 + a_2 b_2 + C$		
		2 + 4*5 + 3	(2,5)	0**5**3080
				253080

3 Illustrate the steps of the standard squaring algorithm for computing c =a * a = 456 * 456

i	j	Step	(C,S)	Partial t
0	1	$t_0 + a_0 a_0$		000000
		0 + 6*6	(3,6)	00000**6**
		$t_1 + 2a_1 a_0 + C$	(3,*)	000006
		0 + 2*5*6 + 3	(6,3)	0000**3**6
0	2	$t_2 + 2a_2 a_0 + C$	(6,*)	000036
		0 + 2*4*6 + 6	(5,4)	000**4**36
				00**5**436
1	2	$t_2 + a_1 a_1$		005436
		4 + 5*5	(2,9)	005**9**36
		$t_3 + 2a_2 a_1 + C$	(2,*)	005936
		5 + 2*4*5 + 2	(4,7)	00**7**936
				0**4**7936
2	2	$t_4 + a_2 a_2$		047936
		4 + 4*4	(2,0)	0**0**7936
				207936

4 Let r = 32, n = 21, a = 13, and b = 15. Compute $c = a * b * r^{-1}$ mod n using the standard Montgomery multiplication algorithm. Illustrate the steps and give all temporary results

iteration	q	g_0	g_1	u_0	u_1	v_0	v_1
0	-	32	21	1	0	0	1
1	1	21	11	0	1	1	-1
2	1	11	10	1	-1	-1	2
3	1	10	1	-1	2	2	-3
4	10	1	0	2	-21	-3	32
		•		•		•	

From Table **??**, GCD = 1, $r^{-1} = 2$, $n' = 3$.
Consider $\overline{x}$ = a = 13, such that x = $\overline{x} * r^{-1}$ mod n = 5. $\overline{y}$ = b = 15, such that y = $\overline{y} * r^{-1}$ mod n = 9.
Now, $a * b * r^{-1}$ mod n is same as $\overline{x} * \overline{y} * r^{-1}$ mod n.
So, $a * b * r^{-1}$ mod n = $\overline{x} * \overline{y} * r^{-1}$ mod n = MonPro($\overline{x}$ = 13,$\overline{y}$ = 15).

4.1 function MonPro($\overline{a}$ = 13 , $\overline{b}$ = 15)

1. t = 13*15 = 195
2. m = (195*3) mod 32 = 9
3. u = (195 + 9 * 21) / 32 = 384/32 = 12
4. 12 < 21 **return** 12

Final result is : 12.

5 Let p = 29, a = 23, and g = 10. Compute g^a (mod p) using the binary method of exponentiation and the Montgomery multiplication where r = 32. Show the steps and temporary values.

Consider n=p=29, M=g=10, e=a=23.

iteration	q	g_0	g_1	u_0	u_1	v_0	v_1
0	-	32	29	1	0	0	1
1	1	29	3	0	1	1	-1
2	9	3	2	1	-9	-1	10
3	1	2	1	-9	10	10	-11
4	2	1	0	10	-29	-11	32
		•		•		•	

From Table **??**, GCD = 1, $r^{-1} = 10$, $n' = 11$.

- Step 2

  $\overline{M}$ = M * r mod n = 10 * 32 mod 29 = 1

- Step 3

  $\overline{C}$ = 1 * r mod n = 1 * 32 mod 29 = 3

- Step 4

e_i	Step 5	Step 6
1	MonPro(3,3) = 3	MonPro(1,3) = 1
0	MonPro(1,1) = 10	
1	MonPro(10,10) = 14	MonPro(1,14) = 24
1	MonPro(24,24) = 18	MonPro(1,18) = 6
1	MonPro(6,6) = 12	MonPro(1,12) = 4

$MonPro(3,3)$
t = 3 * 3 = 9
m = 9 * 11 mod 32 = 3
u = (9 + 3 * 29) / 32 = 3

$MonPro(1,3)$
t = 1 * 3 = 3
m = 3 * 11 mod 32 = 1
u = (3 + 1 * 29) / 32 = 1

$MonPro(1,1)$
t = 1 * 1 = 1
m = 1 * 11 mod 32 = 11
u = (1 + 11 * 29) / 32 = 10

$MonPro(10,10)$
t = 10 * 10 = 100
m = 100 * 11 mod 32 = 12
u = (100 + 12 * 29) / 32 = 14

$MonPro(1,14)$
t = 1 * 14 = 14
m = 14 * 11 mod 32 = 26
u = (14 + 26 * 29) / 32 = 24

$MonPro(24,24)$
t = 24 * 24 = 576
m = 576 * 11 mod 32 = 0
u = (576 + 0 * 29) / 32 = 18

$MonPro(1,18)$
t = 1 * 18 = 18
m = 18 * 11 mod 32 = 6
u = (18 + 6 * 29) / 32 = 6

$MonPro(6, 6)$
t = 6 * 6 = 36
m = 36 * 11 mod 32 = 12
u = (36 + 12 * 29) / 32 = 12

$MonPro(1, 12)$
t = 1 * 12 = 12
m = 12 * 11 mod 32 = 4
u = (12 + 4 * 29) / 32 = 4

- Step 7
  C = MonPro(4,1) = 11

Result of 10^{23} (mod 29) = **11**

6 Let an RSA key be determined by the parameters {p,q,n,e,d} = {17,23,391,29,85}. Compute S = M^d (mod n) for M = 175 with and without the Chinese remainder theorem and the binary exponentiation.

6.1 Chinese remainder theorem

d_1 = 85 mod (16) = 5
d_2 = 85 mod (22) = 19

iteration	quotient	g_0	g_1	u_0	u_1	v_0	v_1
0	-	23	17	1	0	0	1
1	1	17	6	0	1	1	-1
2	2	6	5	1	-2	-1	3
3	1	5	1	-2	3	3	-4
4	5	1	0	3	-17	-4	23
		•		•		•	

In Table **??**, Initial values of g_0 = q = 23 and g_1 = p = 17. This p^{-1} = -4 and q^{-1} = 3.
$M_1 = M^{d_1}$ mod p = 175^5 mod 17 = 14.
$M_2 = M^{d_2}$ mod q = 175^{19} mod 23 = 10.
S = M_1 + p *((M_2 - M_1) * p^{-1} mod q) = 14 + 17*((10-14)*-4 mod 23) = 14 + 272 = **286**.

6.2 Binary Exponentiation.

Representing 85 in binary results in: 1 0 1 0 1 0 1
Following table shows computation of RSA decryption using binary exponentiation.

i	e_i	Step 2a	Step 2b
6	0	$((175)^2)$ mod 391 = 127	127
5	1	$((127)^2)$ mod 391 = 98	98*175 mod 391 = 337
4	0	$((337)^2)$ mod 391 = 179	179
3	1	$((179)^2)$ mod 391 = 370	370*175 mod 391 = 235
1	0	$((235)^2)$ mod 391 = 94	94
0	1	$((94)^2)$ mod 391 = 234	234*175 mod 391 = **286**